

⇒ Graph Travels :-

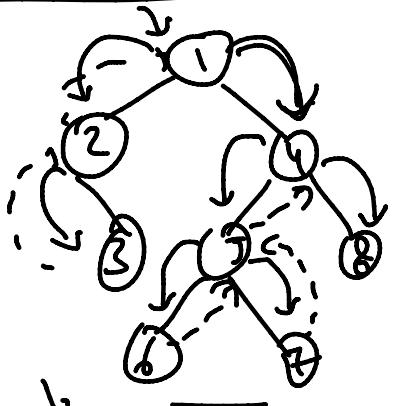
→ Visiting all the nodes in some order once.

DFS

BFS

① DFS - Depth First Search :-

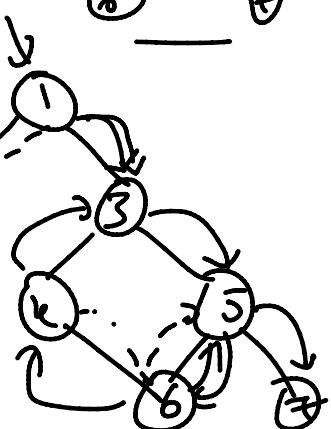
$$\begin{aligned} n &= 8 \\ l &= n-1 \\ &= 7 \end{aligned}$$



$\begin{array}{l} 1, 2, 3, 4, 5, 6, 7, 8 \\ 1, 4, 8, 5, 6, 7, 2, 3 \\ 1, 4, 5, 6, 7, 8, 2, 3 \end{array} \quad \left\{ \right.$

→ Tree.

$\begin{array}{l} 1, 4 \\ 1, 2 \\ 1 \rightarrow 4, 2 \\ 1, 2 \\ 1, 4 \\ 1 \rightarrow 2, 4 \end{array}$



Multiple DFS are possible.

$\begin{array}{l} 1, 2, 3, 4, 6, 5, 7 \\ 1, 2, 3, 5, 6, 4, 7 \end{array}$

`bool vis[8];
vis[i] = true;`

⇒ Recursive :- `dfs(1);`

`void dfs(int i) {
 vis[i] = true;
 print(i);
 for x in adj[i] {
 if (vis[x] == false) {`

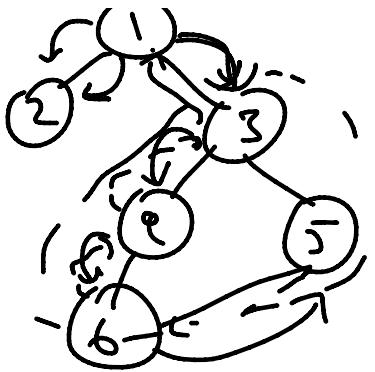
$1 \rightarrow 3, 2$

`dfs(x);
 }
}`



for
if (vis(x) == F)
 dfs(x);

<u>3</u>	<u>3</u>
<u>1, 3, 4, 6, 5, 2</u>	<u>1, 2, 3, 4, 5, 6</u>
<u>F</u>	<u>F F F F F X</u>
<u>T</u>	<u>T T T T T T</u>



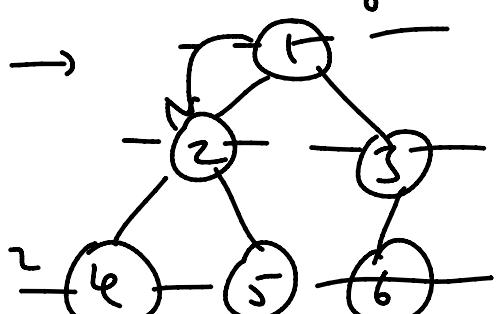
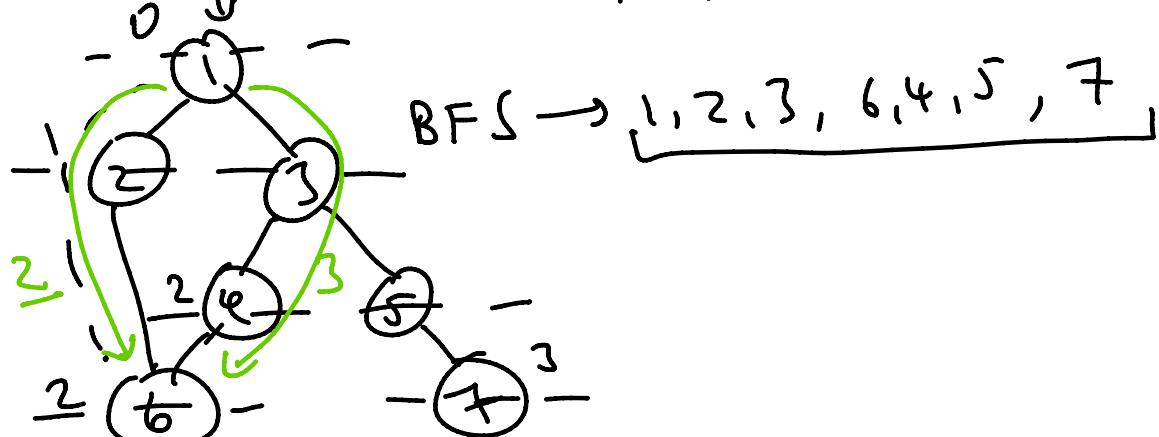
$$TC = O(V+E), \quad AS = O(V)$$

\Rightarrow BFS - Breadth First Search :-

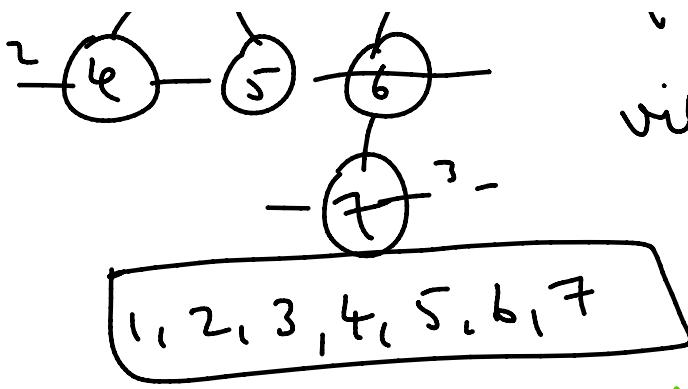
↑ level by level traversal.

DFS \rightarrow 1, 2, 4, 5, 3, 6

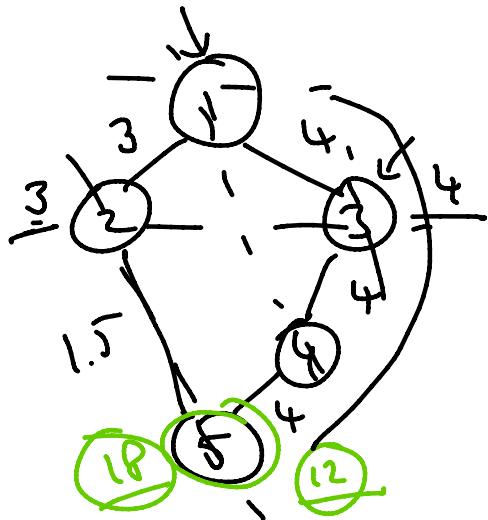
BFS \rightarrow 1, 2, 3, 4, 5, 6
1, 3, 2, 6, 5, 4



\Downarrow [2, 3], [4, 5, 6], 7
0 1 2 3
9 \rightarrow X, 2, 3, X, 4, 5, 6, 7
 \therefore 1, 2, 3, 4, 5, 6, 7



$\text{vis} \rightarrow 1, 2, 3, 4, 5, 6, 7$
 F F F F F F F
 $T T T T T T T$



When we push an element in queue, we mark it as true, assuming uniform edge weight.

```

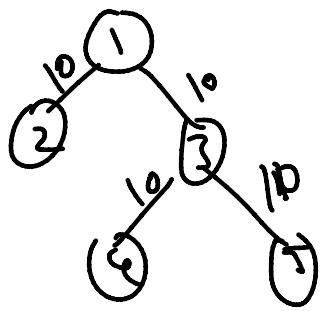
→ queue<int> q;
q.push(1);
vis[1] = T;
while (q.empty() == F)
{
    i = q.front();
    q.pop();
    print(i);
    for xl in adj[i];
        if (vis[xl] == F)
            q.push(xl);
            vis[xl] = T;
}

```

$$TC = O(V+E)$$

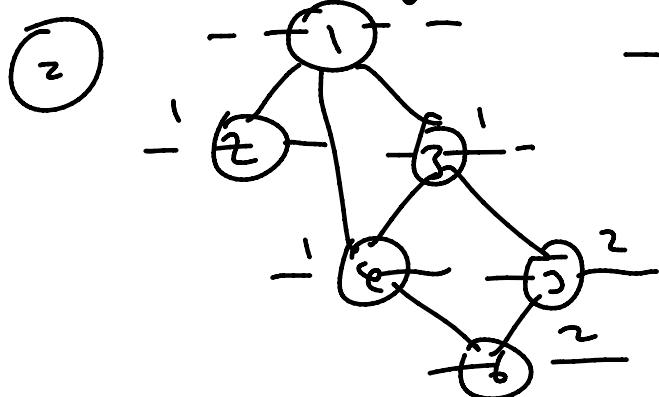
$$AS = O(V)$$

3 3 3



→ Application :-

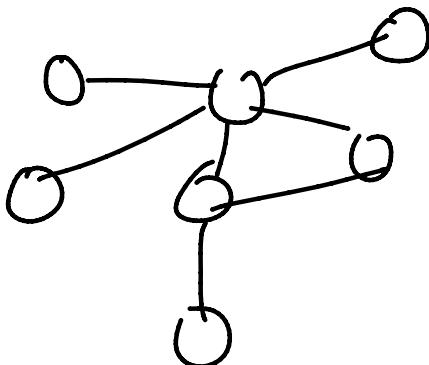
① Find min distance b/w 2 nodes
when edge weights are uniform.



→ Find shortest distance
from source node to
all other nodes.

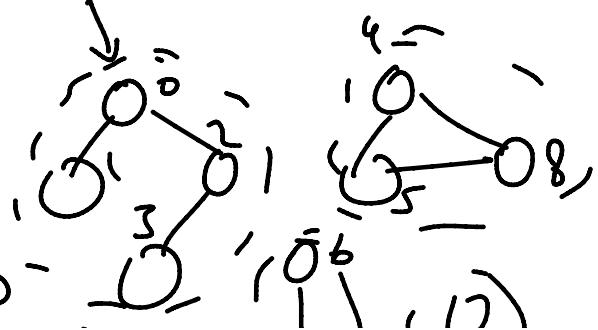
⇒ Connected graph :-

There is a path b/w
each and every pair
of vertices.

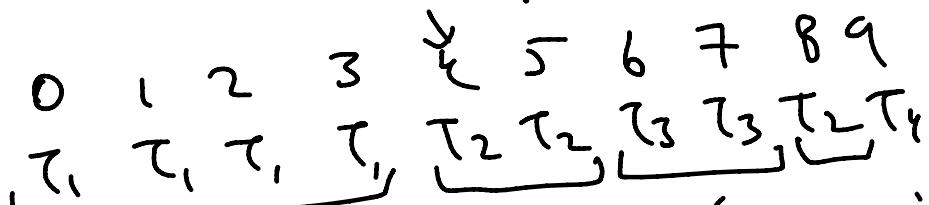


⇒ Unconnected graph

(4) connected entl.



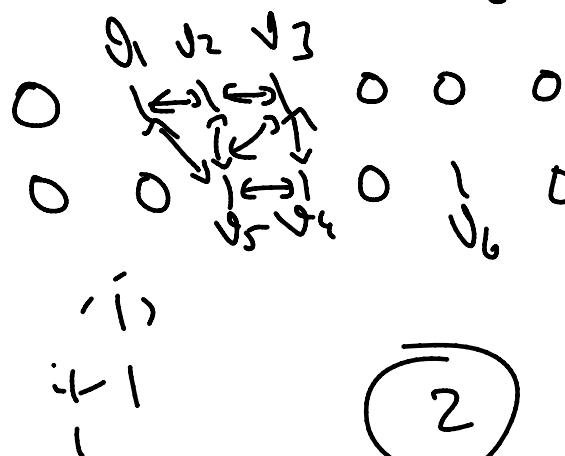
(+) connected components \rightarrow $\{O_1, O_2, O_3\}$



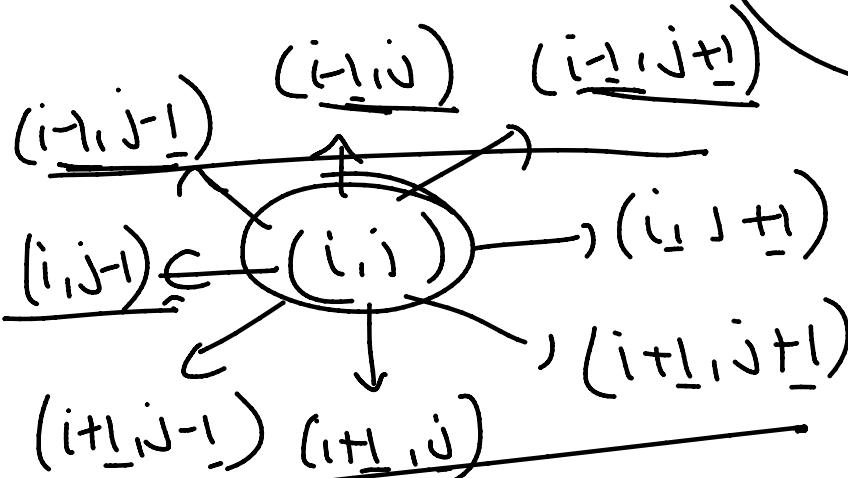
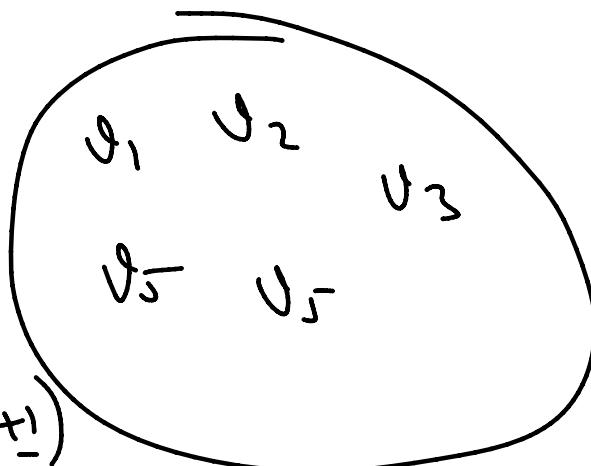
```

arb = 0;
for (i=0; i < v; i++) {
    if (vis[i] == -F)
        dfs(i); arb++;
}
    
```

$arb = 1$
 $arb = 2$
 $arb = 3$
 $arb = 4$



all \rightarrow vertical.

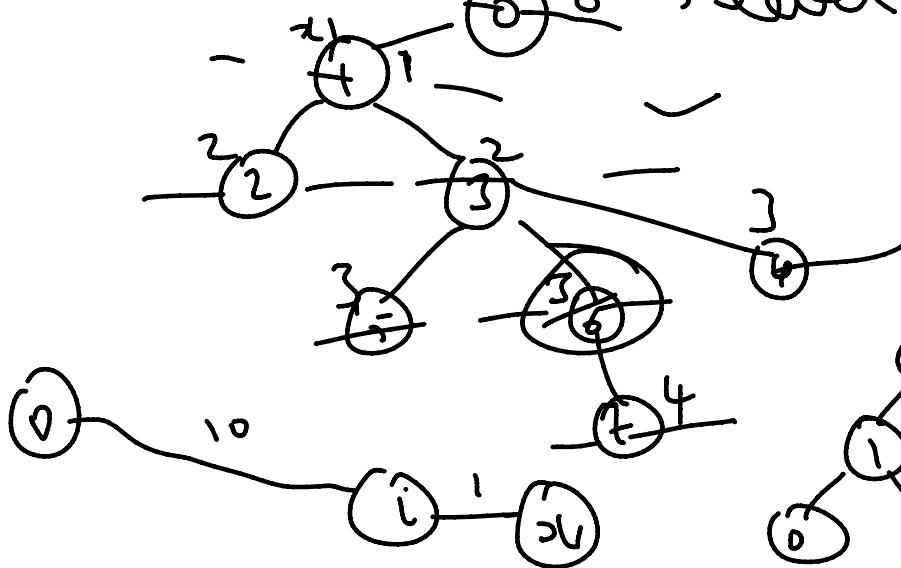


```

for (ii=-1; ii<=1; ii++)
    for (jj=-1; jj<=1; jj++)
        (i+ii, j+jj)
    
```

Q1 graph. $x \rightarrow$ node.

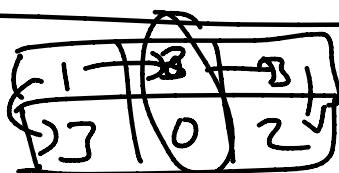
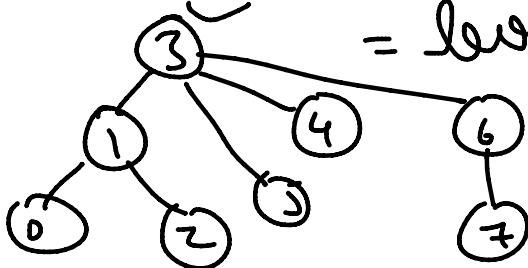
Q1. graph : $x \rightarrow$ node
- $\circ \rightarrow$ label.



$$x = 6$$

level (x_1)

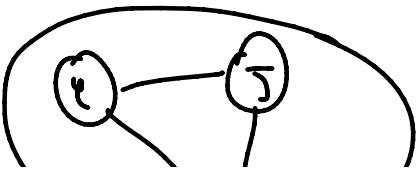
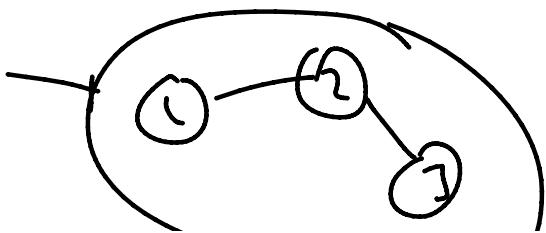
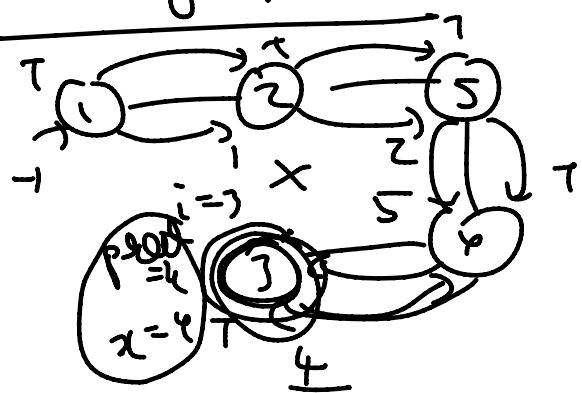
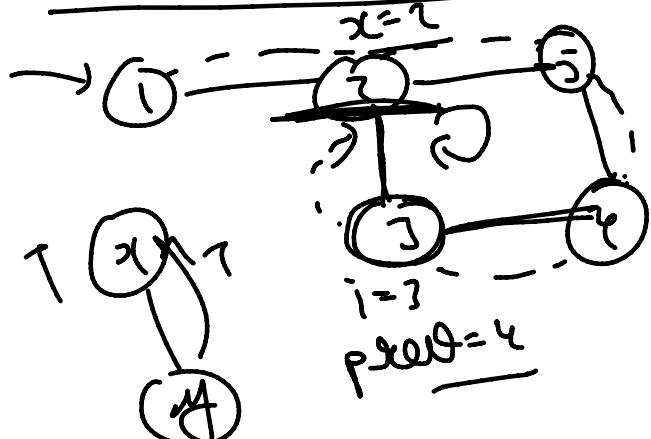
= level[i] + 1;

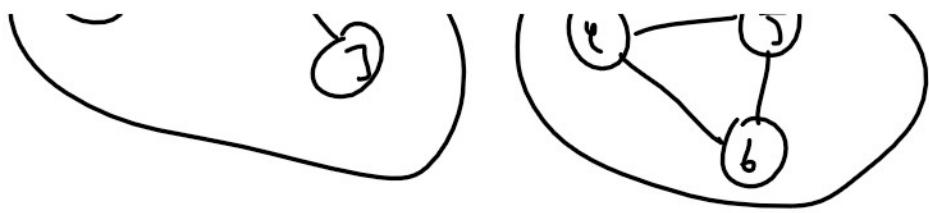


False

Turk

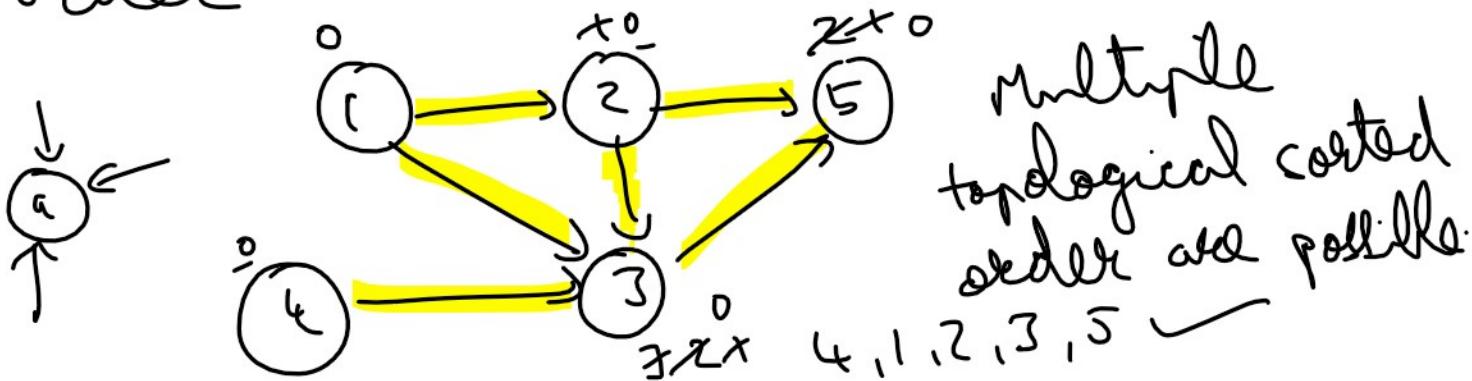
Q1. Cycle in undirected graph :-





\Rightarrow Topological sort :-

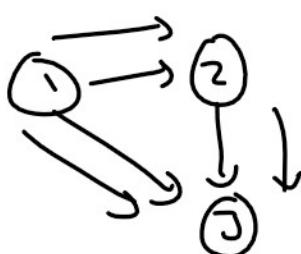
It is an ordering of graph s.t if $a \rightarrow b$, then in topological sorted order a must come before b.



1, 2, 4, 3, 5 ✓

1, 4, 2, 3, 5

only defined for DAG.



1, 2, 3

queue<int> q;

```

for(i=0; i< v; i++)
    if(indegree[i] == 0)
        queue.push(i);
    }
}

```

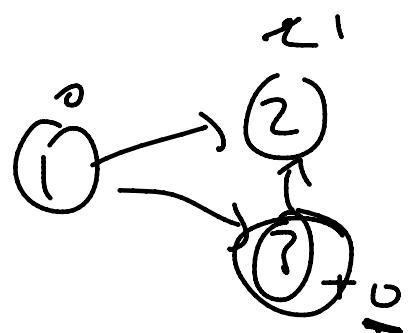
```

for (i=0; i< v; i++)
    if (in-degree(i) == 0)
        q.push(i);
    }
}

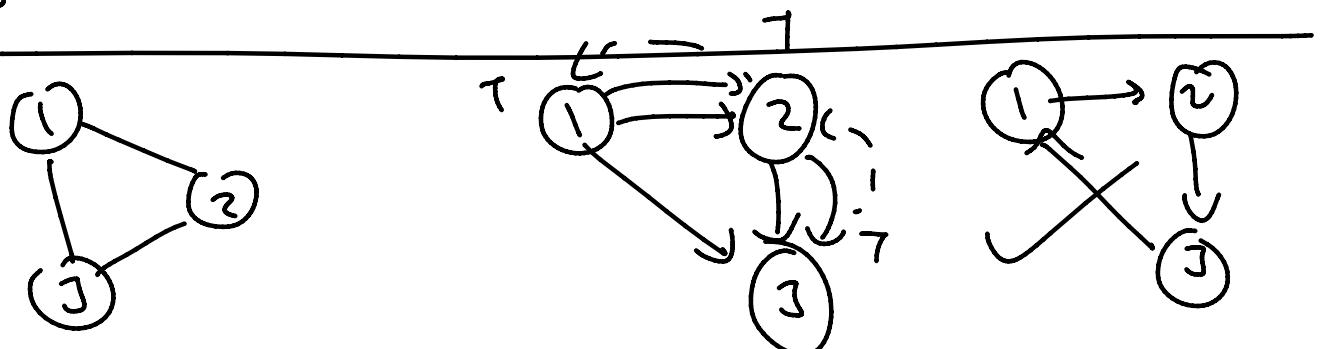
while (!q.empty())
{
    i = q.front();
    q.pop();
    print(i);

    for (x in adj[i])
        if (in-degree(x) -- > 0)
            q.push(x);
}

```



}



=> Doubt :-

1, 2, 3, 1, 2, 1, 2 1, 1, 1, 1, 2, 3, 2

$n=2$

