





# (Compiler Design)

**Derivation Tree or Parse Tree**







## Derivation process

There are two types of derivation process,

1. Left most derivation (LMD)
2. Right most derivation (RMD)

### **Left most derivation (LMD):**

A derivation  $S \Rightarrow w$  is called a leftmost derivation if we apply a production only to the leftmost variable at every step.

### **Right most derivation (RMD):**

A derivation  $S \Rightarrow w$  is called a rightmost derivation if we apply a production to the right most variable at every step.



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**Example:** Consider a grammar  $G$  for the language,  $L = \{a^{2n}b^m, m, n \geq 0\}$

$G: S \rightarrow AB, A \rightarrow aaA \mid \epsilon, B \rightarrow bB \mid \epsilon$ , find the LMD and RMD for string  $w = aab$



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**Example:**  $G: S \rightarrow S+S \mid S*S \mid a \mid b \mid c$ , find the LMD and RMD for string  $w = a + b * c$









## Ambiguous Grammar

A grammar  $G$ , is ambiguous grammar if  $\exists w \in L(G)$ , such that  $w$  has  $>$  one Parse tree either using 2 LMD or 2 RMD i.e., a grammar  $G$  is ambiguous if there is more than one Parse Tree or LMD / RMD for a string  $w \in L(G)$ .

**Example:**  $G: S \rightarrow S+S \mid S*S \mid a \mid b \mid c$ , Grammar  $G$  is ambiguous because there exist two different LMD for a string  $w = a + b * c$



**NOTE:**

- 1- A grammar  $G$  is unambiguous if there exist exactly one parse tree or LMD / RMD for all the string,  $w \in L(G)$
- 2- If a grammar  $G$  is ambiguous, it doesn't mean language  $(L)$  is ambiguous.

**Example:** State whether given grammar is ambiguous or not.  $G: E \rightarrow E+E \mid E-E \mid id$









## NOTE:

- Regular Grammar:
  - Ambiguous (NFA)
  - Unambiguous (DFA)
- Regular language is always unambiguous.
- DCFL is always unambiguous language
- Ambiguity starts from CFL



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**Inherent ambiguous language:**





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**Example:** Check whether the given language  $L = \{a^n \mid n \geq 0\}$  is an inherently ambiguous or not



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**Example:** Given language  $L = \{a^m b^n c^k \mid m, n, k \geq 1, \text{ either } m = n \text{ or } n = k\}$  is Inherent Ambiguous?



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**Note:** If grammar is ambiguous then it is not suitable for any kind of parsing technique except backtracking and operator precedence parsing



## Remove Ambiguity:

It is not always possible to convert an ambiguous grammar into an unambiguous grammar because ambiguity finding & removal both are undecidable





## Remove Ambiguity by precedence & associativity rules:

An ambiguous grammar may be converted into an unambiguous grammar by implementing:

- Precedence Constraints
- Associativity Constraints

These constraints are implemented using the following rules:

**Rule 1:** The precedence constraint is implemented using the following rules:

- The level at which the production is present defines the priority of the operator contained in it.
- The top level of the production, the lower the priority of the operator.
- The bottom level of production, the higher the priority of the operator.

**Rule 2:** The associativity constraint is implemented using the following rule:

- If the operator is left associative, induce left recursion in its production.

If the operator is right associative, induce right recursion in its production

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**Example:**  $G: E \rightarrow E+E \mid E * E \mid id$ . Ambiguous grammar converts into unambiguous grammar



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**Example:** Convert the following ambiguous grammar into unambiguous grammar.

$G: E \rightarrow E+E \mid E * E \mid E^E \mid id$

$E \rightarrow E+E$

$E \rightarrow E * E$

$E \rightarrow E^E$

$E \rightarrow id$

**Unambiguous Grammar**

**Apply Rule 1 & 2**

$E \rightarrow E + T \mid T$

(+) left associative

$T \rightarrow T * F \mid F$

(\*) left associative

$F \rightarrow G ^ F \mid G$

(^) Right associative

$G \rightarrow id$

The priority order:  $id > ^ > * > +$

Associativity: + and \* are left associative and ^ operator is right associative



**Example:** Find the precedence and associativity

i)  $G_1: A \rightarrow A \$ B \mid B$   
 $B \rightarrow B \# C \mid C$   
 $C \rightarrow D @ C \mid D$   
 $D \rightarrow d$

ii)  $G_2: E \rightarrow E * F \mid F + E \mid F$   
 $F \rightarrow F - F \mid id$





Consider the grammar defined by the following production rules, with two operators \* and +

```
S --> T * P
T --> U | T * U
P --> Q + P | Q
Q --> Id
U --> Id
```

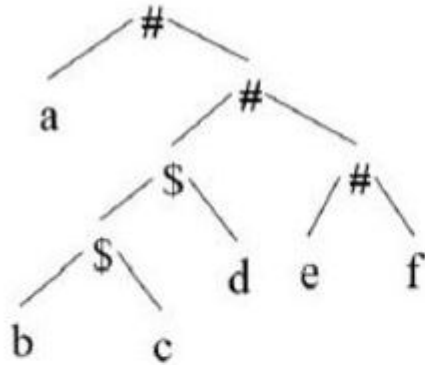
Which one of the following is TRUE?

- (A) + is left associative, while \* is right associative
- (B) + is right associative, while \* is left associative
- (C) Both + and \* are right associative
- (D) Both + and \* are left associative

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Consider the following parse tree for the expression  $a\#b\$c\$d\#e\#f$ , involving two binary operators  $\$$



and  $\#$ .  
parse tree?

Which one of the following is correct for the given

- (A)  $\$$  has higher precedence and is left associative;  $\#$  is right associative
- (B)  $\#$  has higher precedence and is left associative;  $\$$  is right associative
- (C)  $\$$  has higher precedence and is left associative;  $\#$  is left associative
- (D)  $\#$  has higher precedence and is right associative;  $\$$  is left associative

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Recursive CFG:



















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**Left Factoring:**















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**Parsing Technique:**









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**Brute Force Technique:**





Thank You !

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**First and Follow Function:**

**First Function:**











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**Follow Function:**









Thank You !