

For faster operation, the subkey should be pre-computed and stored in cache for faster encryption.

International Data encryption Algorithms (IDEA)

(James Neasey, Lai - 1990)

- Block cipher (64 bit)
- Symmetric (Sender & Receiver same key)
- Reversible like DES (Decry is just Reverse of Encryption)

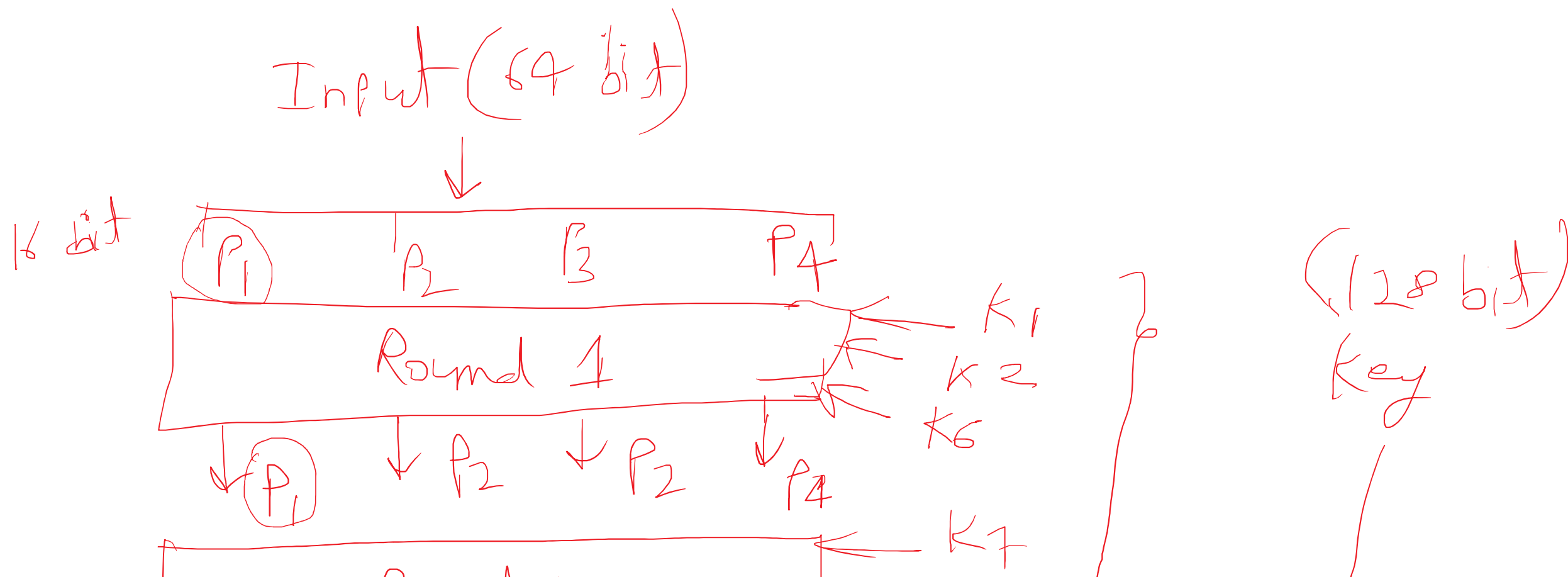
Design principle behind IDEA is the mixing of Arithmetic operation from different algebraic

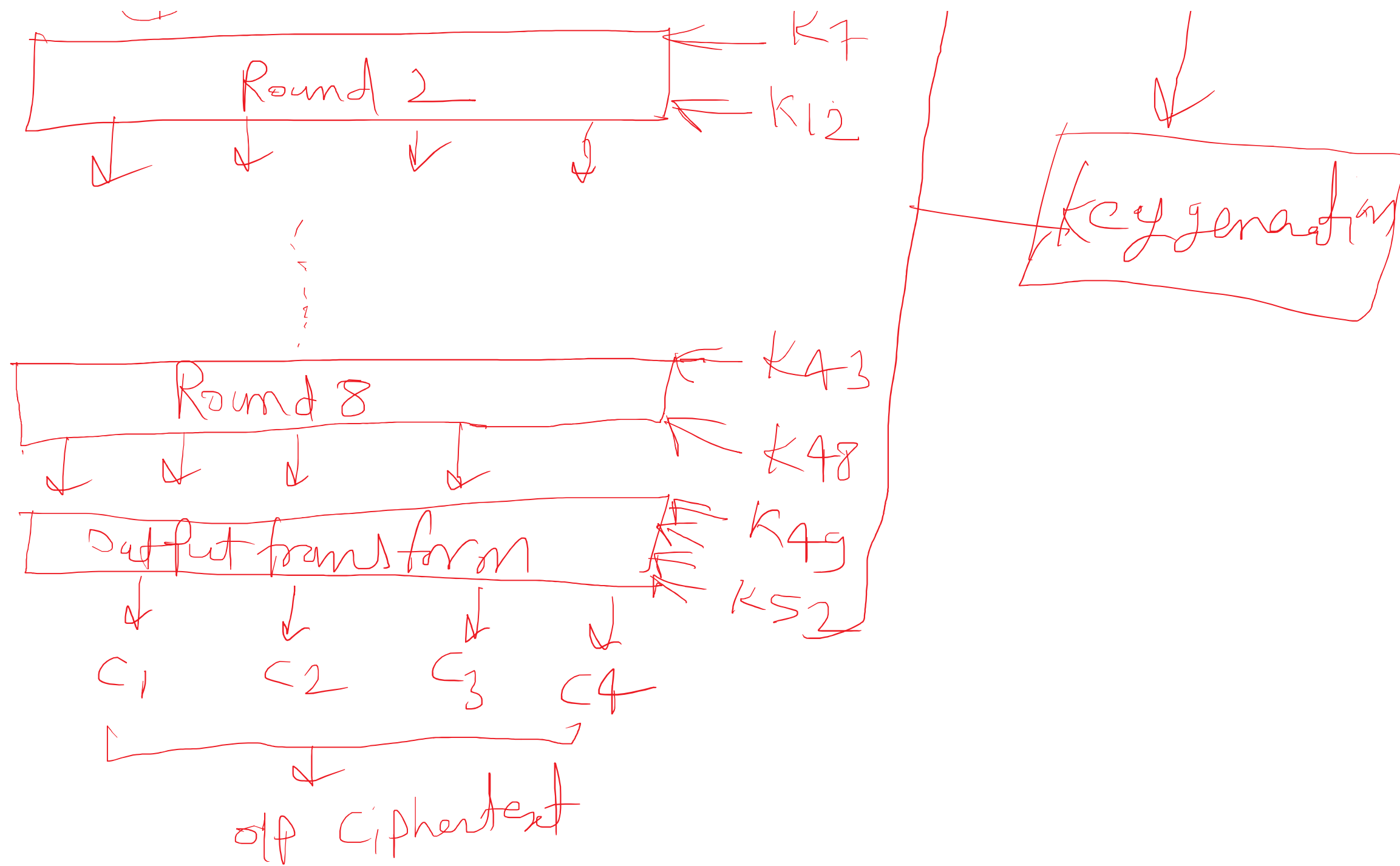
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groups, which are easily implementable
on HW and SW.

- Underline operation are XOR, add, mult.
- It does not use S-Box (explicitly)
- Key size 128 \rightarrow sub key (16 bit)

Structure IDEA : 8 Round





Round functionality:

each Round (P_1, P_2, P_3, P_4)

$$\left[\begin{array}{ll} 1. & \text{multiply} \quad P_1 * K_1 \rightarrow S_1 \\ 2. & \text{add} \quad P_2 + K_2 \rightarrow S_2 \\ 3. & \text{add} \quad P_3 + K_3 \rightarrow S_3 \\ 4. & \text{multiply} \quad P_4 * K_4 \rightarrow S_4 \end{array} \right] \text{ 4key}$$

$$\left[\begin{array}{ll} 5. & \text{XOR} \quad S_1 \oplus S_3 \rightarrow S_5 \\ 6. & \text{XOR} \quad S_2 \oplus S_4 \rightarrow S_6 \end{array} \right] \text{ No Key}$$

$$\left[\begin{array}{ll} 7. & \text{mult} \rightarrow S_5 * K_5 \rightarrow S_7 \\ 8. & \text{add} \quad S_6 + S_7 \rightarrow S_8 \\ 9. & \text{mult} \rightarrow S_8 * K_6 \rightarrow S_9 \\ 10. & \text{add} \quad S_7 + S_9 \rightarrow S_{10} \end{array} \right] \text{ 2 Key}$$

$$\left. \begin{array}{ll}
 i1 & \text{xOR} \\
 i2 & \text{xOR} \\
 i3 & \text{xOR} \\
 i4 & \text{xOR}
 \end{array} \right\} \begin{array}{l}
 S_1 \oplus S_9 \rightarrow S_{11} \rightarrow \text{New } P_1 \\
 S_2 \oplus S_9 \rightarrow S_{12} \rightarrow \text{New } P_2 \\
 S_3 \oplus S_{10} \rightarrow S_{13} \rightarrow \text{New } P_3 \\
 S_4 \oplus S_{10} \rightarrow S_{14} \rightarrow \text{New } P_4
 \end{array} \left. \begin{array}{l}
 \text{Q/P} \\
 \\
 \text{Round 1}
 \end{array} \right\}$$

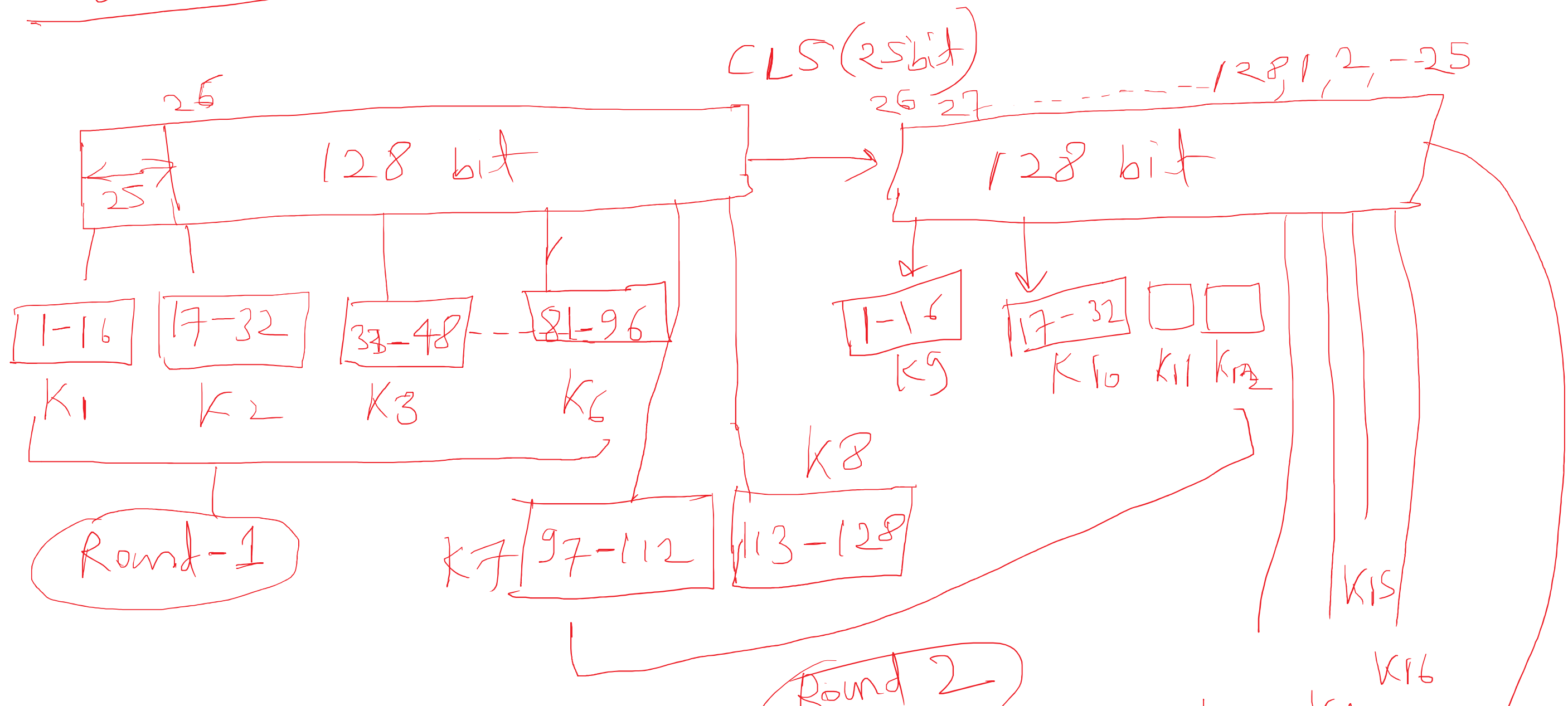
Repeat same up to Round 8, so total 48
 sub key is used till now.
 after 8th Round 4 more operation is need
 to apply in output transform step.

$$\left. \begin{array}{ll}
 P_1 * K_{49} \rightarrow C_1 \\
 P_2 + K_{50} \rightarrow C_2 \\
 P_3 + K_{51} \rightarrow C_3 \\
 P_4 * K_{52} \rightarrow C_4
 \end{array} \right\} \begin{array}{l}
 \text{Final} \\
 \text{Ciphertext} \\
 \hline
 (4 \text{ bit})
 \end{array}$$

$$P4 * K52 \rightarrow 504 \quad (64 \text{ bit})$$

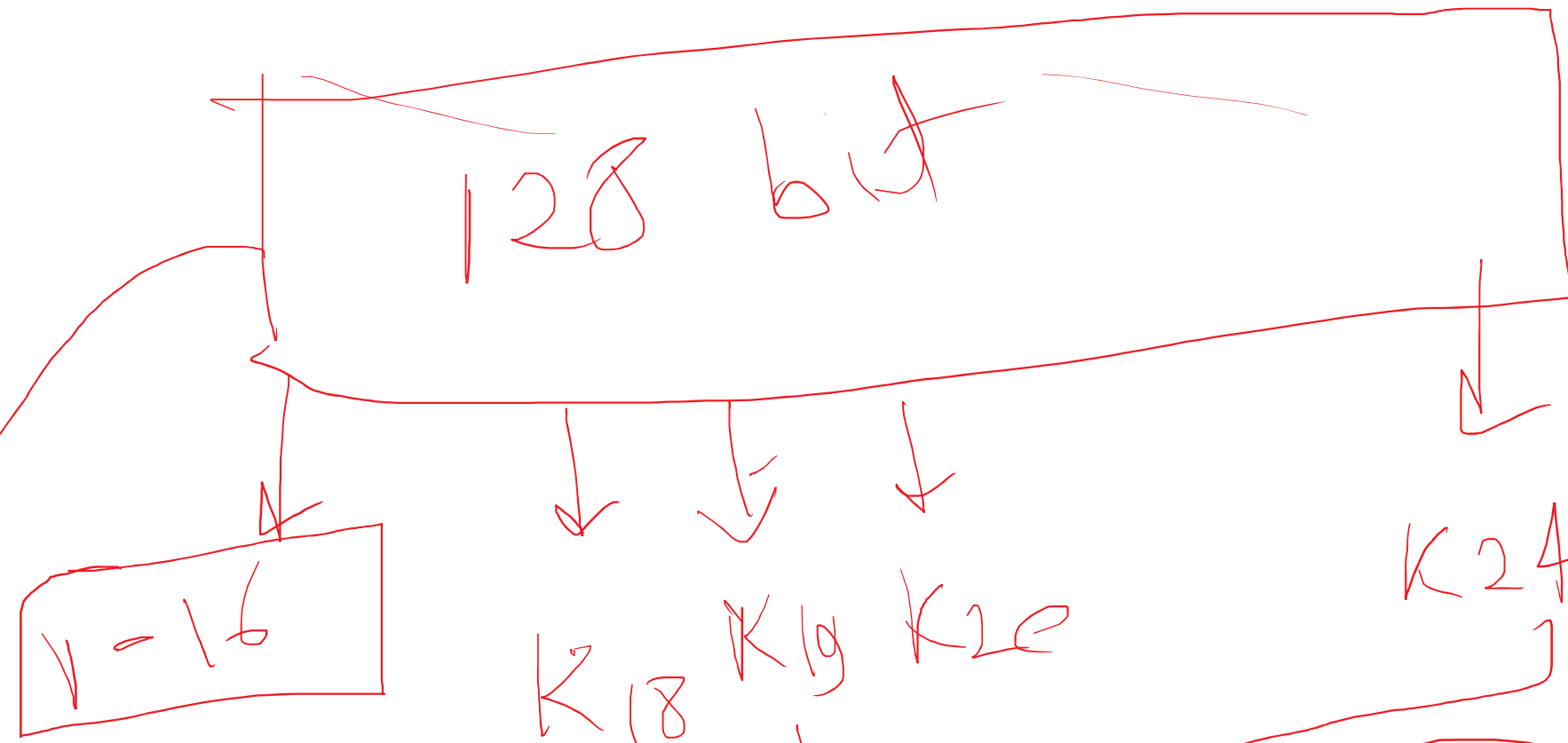
$$\text{Total sub subkey} = 48 + 4 = 52 \text{ key.}$$

key generation: given key k : 128 bit



Round 2

K_{13} K_{14} K_{16}
Round 3



CLS
by
25 bit

K_{17}
Round(3)

Round 4

Round(s)

CLS by 25 bit and So on to generate all 52 sub key.

CLS (3 bit) (10110101) = 11010101

011110101

CLS 11010101010101010101010101010101

CLS
by 3 bit

		↖	1 1 1 0 1 0 1 0
			1 1 0 1 0 1 0 1