CS61A NOTE12 Scheme

Primitives and Defining Variables 初始和定义变量

Scheme has a set of **atomic** primitive expressions. Atomic means that these expressions cannot be divided up.Scheme 有原子性的原始表达式。原子:这些表达式不能被分割。

Unlike in Python, the only primitive in Scheme that is a false value is #f and its equivalents, false and False. This means that 0 is not false.

与之前不同, Scheme 中唯一假值是 #f /false /False

In Scheme, we can use the **define** special form to bind values to symbols, which we can then use as variables. When a symbol is defined this way, the **define** special form returns the symbol.在 Scheme 中可以用 define 特殊形式将值 绑定到 symbols 上,然后我们可以将其作为变量使用。当 symbols 以这种方式被定义时,define 返回该 symbol。

(define <variable name> <value>)

Evaluates <value> and binds the value to <variable name> in the current environment.算值并绑定到变量名(当前环境都用这个)直接输出是输出变量名

WWSD

```
纯文本
scm> (define a 1)
a
scm> a
1
```

```
scm> (define b a)
b
scm> b
1 #不是a,a已经是1
scm> (define c 'a)
c
scm> c
a #'a 不是1,是字符
```

Call Expressions

Call expressions apply a procedure to some arguments.

```
纯文本
(<operator> <operand1> <operand2> ...)
```

Call expressions in Scheme work exactly like they do in Python. To evaluate them:

- 1. Evaluate the operator to get a procedure.用 operator 操作符
- 2. Evaluate each of the operands from left to right.从左到右用 operands
- 3. Apply the value of the operator to the evaluated operands.应用 operator 在 operands 上

For example, consider the call expression (+ 1 2). First, we evaluate the symbol 符号 + to get the built-in addition procedure 内置加法. Then we evaluate the two operands 1 and 2 to get their corresponding atomic values 原子值. Finally, we apply the addition procedure to the values 1 and 2 to get the return value 3.应用操作符在 12 上

Operators may be symbols, such as + and *, or more complex expressions, as long as they evaluate to procedure values.操作可以是符号,如 + 和*, 也可以是更复杂的表达式,只要它们评估为过程值。

WWSD

What would Scheme display? As a reminder, the built-in quotient function performs floor division.内置的商函数可以执行地板除法 | | //。

```
xcm> (define a (+ 1 2)) #define直接输出为本身
a
scm> a
3
scm> (define b (- (+ (* 3 3) 2) 1))
b
scm> (+ a b)
13
scm> (= (modulo b a) (quotient 5 3)) #前者mod 1,后者//,1
#t
```

Special Forms

Special form expressions contain a **special form** as the operator. Special form expressions *do not* follow the same rules of evaluation as call expressions. Each special form has its own rules of evaluation — that's what makes them special! Here's the <u>Scheme Specification</u> to reference the special forms we will cover in this class.特殊形式表达式包含一个特殊形式作为运算符。特殊形式表达式不遵循与调用表达式相同的评估规则。每个特殊形式都有自己的求值规则

It is important to note that everything in Scheme is either an **atomic** or an **expression**, so although these special forms look and operate similarly to Python, they are evaluated differently.需要注意的是,Scheme 中的所有东西要么是**原子**,要么是**表达式**,所以尽管这些特殊形式的外观和操作与 Python 相似,但它们的求值方式却不同。

Special forms like if, cond, and, or in Python direct the control flow of a program and allow you to evaluate specific expressions under some condition. In Scheme, however, these special forms are expressions that take in a set amount of parameters and return some value based on the condition passed in.像 Python 中的 if, cond, and, or 这样的特殊形式可以引导程序的控制流,并允许你在某些条件下评估特定的表达式。但在 Scheme 中,这些特殊形式是表达式,它接收了一定数量的参数,并根据传入的条件返回一些值。

If Expression(选前面/后面)

An if expression looks like this:

```
(if cate> <if-true> [if-false])
```

and <if-true> are required expressions and [if-false] is
optional.前两者必须,后一个可选

The rules for evaluation are as follows:

- 2. If continue = continue =

if is a special form as not all of its operands will be evaluated. The value of the first operand determines whether the second or the third operator is evaluated.if 是一个特殊的形式,因为不是所有的操作数都会被评估。第一个操作数的值决定了第二个或第三个操作数是否被评估。Only #f is a false-y value in Scheme; everything else is truth-y, including 0.在 Scheme 中,只有#f 是假 Y 值,其他都是真 Y 值,包括 0

Cond Expression 条件表达式(直到找到真值)

A cond expression looks like this:

```
(cond (<pred1> <if-pred1>) (<pred2> <if-pred2>) ... (<predn> <if-
predn>) [(else <else-expression>)])
```

The rules for evaluation are as follows:

1. Evaluate the predicates <pred1> , <pred2> , ..., <predn> in order until you reach one that evaluates to a truth-y value.直到找到真值

- 2. If you reach a c that evaluates to a truth-y value, evaluate and return the corresponding expression in the clause.
- 3. If none of the predicates are truth—y and there is an else clause, evaluate and return <else-expression> .如果没真,返回<else-expression>

cond is a special form because it does not evaluate its operands in their entirety; the predicates are evaluated separately from their corresponding return expression. In addition, the expression short circuits 短路 upon reaching the first predicate that evaluates to a truth—y value, leaving the remaining predicates unevaluated.第一个真值短路,剩下不 value

Let Expressions(返回最后值)

```
A let expression looks like this: (let ([binding_1] ... [binding_n]) <body> ...)

Each binding corresponds to expressions of the form (<name> <expression>).
```

Scheme evaluates a let expression using the following steps:

1. Create a new local frame that extends the current environment (in other words, it creates a new child frame whose parent is the current frame).创建一个新的本地框架/它创建了创建一个新的本地框架一个新的子框架,其父级是当前框架)

- 2. For each binding provided, bind each name to its corresponding evaluated expression .绑定
- 3. Finally, the body expressions are evaulated in order in this new frame, returning the result of evaluating the last expression.按顺序算, 返回最后值

Note that bindings are optional within a let statement, but we typically include them

Note that (-x y) in the body of this let expression *does* get evaluated, but the result doesn't get returned by the let expression because only the value of the *last* expression in the body, (+x y), gets returned. Thus, the interpreter does *not* display -5 (the result of (-x y)).(-x y)% with (-x y)% determined by the let (-x y)0.(-x y)8 with (-x y)9 with (

However, we see that 5 and 10 are displayed out by the interpreter. This is because printing 5 and printing 10 were side effects of evluating the expressions (print x) and (print y), respectively. 5 and 10 are not the return values of (print x) and (print y).打印 5 和打印 10 是打印 x 和打印 y 的副作用,5 和 10 不是 print x 和 print y 的返回值

Begin Expressions(返回最后值)

A begin expression looks like this: (begin <body 1> ... <body n>)

Scheme evaluates a begin expression by evaluating each body in order in the current environment, returning the result of evaluating the last body 方案通过在 当前环境中依次评估每个体来 value,返回最后一个 body 的结果

(+1 2)确实被计算了,但是结果 3 并没有被 begin 表达式返回(因此也没有被解释器显示),因为它不是最后一个主体表达式

Boolean operators

Like Python, Scheme has the boolean operators and, or, and not. and and or are special forms because they are short-circuiting operators, while not is a builtin procedure.短路运算符: and or; 内置过程: not

- and takes in any amount of operands and evaluates these operands from left to right until one evaluates to a false-y value. It returns that first false-y value or the value of the last expression if there are no false-y values.返回第一个假 Y 值,如果没有假 Y 值返回最后一个表达式的值。
- or also evaluates any number of operands from left to right until one evaluates to a truth-y value. It returns that first truth-y value or the value of the last expression if there are no truth-y values.返回第一个真-Y 值,如果没有真-Y 值返回最后一个表达式的值
- not takes in a single operand, evaluates it, and returns its opposite truthiness value.接收一个操作数,评估,并返回相反值

WWSD

```
纯文本
scm> (if (or #t (/ 1 0)) 1 (/ 1 0)) #看清括号,最后两个是前/后,or后是两个只执行第
1
scm > ((if (< 4 3) + -) 4 100)
-96
                                                                纯文本
scm> (cond #到真
       ((and (- 4 4) (not #t)) 1)
       ((and (or (< 9 (/ 100 10)) (/ 1 0)) #t) -1) #or执行前者, and全对, 输出
       (else (/ 1 0))
   )
-1
                                                                纯文本
scm> (let (
           (a (- 3 2))
           (b (+ 5 7))
         )
       (* a b) 不输出
       (if (< (+ a b) b)
          (/ a b)
           (/ b a) 只输出最后
       )
   )
>>>12
                                                                纯文本
scm> (begin
       (if (even? (+ 2 4)) #对
           (print (and 2 0 3)) #执行此行, 3
           (/10)
       (+ 2 2) #begin不return
       (print 'lisp) #lisp, side effect
       (or 2 0 3) #begin只return此行2
```

```
3
lisp
2
```

Defining Functions

All Scheme procedures are constructed as lambda procedures.

One way to create a procedure is to use the lambda special form.

```
(lambda (<param1> <param2> ...) <body>)
```

This expression creates a lambda function with the given parameters and body, but does not evaluate the body. As in Python, the body is not evaluated until the function is called and applied to some argument values. The fact that neither the parameters nor the body is evaluated is what makes lambda a special form.

We can also assign the value of an expression to a name with a define special form:

```
    (define (<name> <param> ...) <body> ...)
    (define <name> (lambda (<param> ...) <body> ...))
```

These two expressions are equivalent; the first is a concise version of the second.

```
#文本

scm> ; Bind lambda function to square

scm> (define square (lambda (x) (* x x)))

square

scm> (define (square x) (* x x)) ; 直接lambda写成square

square

scm> square
(lambda(x) (*x x)) #注意最外括号

scm> (square 4)

16
```

lab11 WWSD

disc Q1: Virahanka-Fibonacci

Write a function that returns the n -th Virahanka-Fibonacci number.

```
(define (vir-fib n)
    'YOUR-CODE-HERE
    (if (<= n 1)
        n #不用加括号时别加,且不是1, n=1时1, n=0时0, 不用定义n=2
        (+ (vir-fib (- n 1)) (vir-fib (- n 2)))
    )
)
```

```
(expect (vir-fib 10) 55)
(expect (vir-fib 1) 1)
```

Q2: fib2 (Su21 Final #5a) Solution

Each element of the **fibonacci2** sequence is defined as twice the absolute value of the difference between the previous two elements. Assume that the 0th element of the **fibonacci2** sequence is 0, and the 1st element is 1. Implement the function **fib2**, which takes in one parameter **n**, a non-negative integer, and returns the nth element of the **fibonacci2** sequence. Reminder: Scheme has a built in procedure **abs** which returns the absolute value of the argument that is passed in.

lab11 Q2 over or under

Define a procedure over-or-under which takes in a number num1 and a number num2 and returns the following:

- -1 if num1 is less than num2
- 0 if num1 is equal to num2
- 1 if num1 is greater than num2

lab11 Q3: Make Adder

Write the procedure make-adder which takes in an initial number, num, and then returns a procedure. This returned procedure takes in a number inc and returns the result of num + inc .提示用 lambda 函数 inc 是输入的

Hint: To return a procedure, you can either return a lambda expression or define another nested procedure. Remember that Scheme will automatically return the last clause in your procedure.

You can find documentation on the syntax of <u>lambda</u> expressions in <u>the 61A</u> scheme specification!

```
纯文本
(define (make-adder num)
    (lambda(x)(+ x num))
)
```

lab11 Q4: Compose

Write the procedure composed, which takes in procedures f and g and outputs a new procedure. This new procedure takes in a number x and outputs the result of calling f on g of x.返回 f(g(x))没给 x 提示用 lambda

```
纯文本
(define (composed f g)
    (lambda(x) (f(g x))
    )
)
```

lab11 Q5: Repeat

Write the procedure $\ repeat$, which takes in a procedure $\ f$ and a number $\ n$, and outputs a new procedure. This new procedure takes in a number $\ x$ and outputs the result of applying $\ f$ to $\ x$ a total of $\ n$ times. For example:

```
scm> (define (square x) (* x x))
square
```

```
scm> ((repeat square 2) 5); (square (square 5))
625
scm> ((repeat square 3) 3); (square (square (square 3)))
6561
scm> ((repeat square 1) 7); (square 7)
49
```

Hint: The **composed** function you wrote in the previous problem might be useful.

```
(define (repeat f n)
  (if (< n 1)
        (lambda(x) x)
        (composed f (repeat f (- n 1) ))
)</pre>
```

lab11 Q6: Greatest Common Divisor

The GCD is the the greatest common divisor of two numbers.

Write the procedure gcd, which computes the GCD of numbers a and b. Recall that Euclid's Algorithm tells us that the GCD of two values is either of the following:

- the smaller value if it evenly divides the larger value, or
- the greatest common divisor of the smaller value and the remainder of the larger value divided by the smaller value

```
In other words, if a is greater than b and a is not divisible by b, then gcd(a, b) = gcd(b, a % b)
```

```
纯文本
(define (max a b)
  (if (> a b)
```

```
a
b))

(define (min a b)
    (if (> a b)
        b
        a))

(define (gcd a b)
    (cond
        ((zero? a) b)
        ((zero? b) a)
        ((= (modulo (max a b) (min a b)) 0) (min a b))
        (else (gcd (min a b) (modulo (max a b) (min a b) )))
)

#別多写括号
```

lab12 Q6: Count

Implement count, which takes in an element x and a list s and returns the number of times that x is contained in s.

Scheme Lists - Useful Functions

```
(length <list>)
```

- Returns the number of elements in **list**.
- (length (list 1 2 3 4 5)) -> 5
- (length (list 1 2 (list 3 4) 5)) -> 4

```
(append <list1> <list2>)
```

- Returns a new list containing the elements of **list1**, then the elements of **list2** in that order.
- (append (list 1 2 3 4) (list 5 6 7 8)) -> (1 2 3 4 5 6 7 8)

Scheme Lists - Useful Functions

```
(map <func> <list>)
```

- Returns a new list containing the elements of list in the same order but with func applied to them
- (map even? (list 1 2 3 4)) -> (#f #t #f #t)

```
(filter <func> <list>)
```

- Returns a new list that only contains the elements of **list** that return true when **func** is called on them
- (filter odd? (list 1 2 3 4)) -> (1 3)

lambda 提醒用 map+append

```
Q4: join (Fa20 Final #5b)
```

The **join** procedure takes two lists of lists s and t. It returns a list of lists that has one element for each possible pairing of an element of s with an element of t. Each element of the result is a list that has all the elements of a list from s followed by all the elements of a list from t.

Q4: join (Fa20 Final #5b) Solution

The **join** procedure takes two lists of lists \mathbf{s} and \mathbf{t} . It returns a list of lists that has one element for each possible pairing of an element of \mathbf{s} with an element of \mathbf{t} . Each element of the result is a list that has all the elements of a list from \mathbf{s} followed by all the elements of a list from \mathbf{t} .

```
(define (join s t)
    (if (null? s) nil
        (append (map (lambda (v) (append (car s) v)) t)
        ; (a) (b) (c) (d) (e)
        (join (cdr s) t))))
        ; (f)
```

cons 递归

Q5: countdown (Su21 #5b) Solution

lab12 Q7: Unique

Implement unique, which takes in a list s and returns a new list containing the same elements as s with duplicates removed.

```
scm> (unique '(1 2 1 3 2 3 1))
(1 2 3)
scm> (unique '(a b c a a b b c))
(a b c)

(define (unique s)
   (if (null? s) nil
        (cons (car s))
```

```
(unique (filter (lambda (x) (not (eqv? (car s) x))) (cdr s)))))
```

lab12 Q8: Tally

Implement tally, which takes a list of names and returns a list of 2-element lists, one for each unique name in names. Each 2-element list should contain a name and the number of times that the name appeared in names. Each name should appear only once in the output, and the names should be ordered by when they first appear in names.

Hint: You might want to use count and unique procedures in your solution. You may also want to use map.

```
(tally '())
()
scm> (tally '(kevin irene ashley irene irene ashley brandon irene irene bra
((kevin 1) (irene 5) (ashley 2) (brandon 2))

(define (tally names)
   (map (lambda (name) (list name (count name names))) (unique names)))
)
```