

Devin Gendron
CS225 Final Exam

Logic Symbols: $\geq \leq \neq \neg \wedge \vee \oplus \equiv \rightarrow \leftrightarrow \exists \forall$ Set Symbols: $\in \notin \subseteq \subset \supseteq \supset \emptyset \cup \cap \times$

Q1:

$P = a$ is rational and ab is irrational

$Q = b$ is irrational

By contradiction:

$P = a$ is rational and ab is irrational

$\neg Q = b$ is rational

Thus: "If a is rational and ab is irrational, then b is rational".

So by definition of a rational integer, $a = m/n$ and $b = p/q$. We can then test this:

$$ab = (m/n)(p/q)$$

$$= mp/nq$$

As we can see, mp/nq is rational and proves by contradiction that the original statement, "If a is rational and ab is irrational, then b is irrational", is true.

Q2:

$P(n)$: For every integer n greater than 1, $n! < n^n$

Base Case: Test $P(n)$ with an n greater than 1. So,

$$P(2): 2! = 2 * 1 = 2$$

$$2^2 = 4$$

This proves the base case $P(2)$.

Inductive Case: If our base case is true, then $P(k)$ is also true. If $P(k)$ is true, then we can check to see if $P(k+1)$ is true.

$(k+1) < (k+1)^{(k+1)}$. As can be seen, any integer $(k+1)$ will be less than $(k+1)^{(k+1)}$ for any integer $n > 1$. For example, $P(2)$:

$$(2+1) < (2+1)^{(2+1)} = 3 < 3^3 = 3 < 27$$

Thus, proven by mathematical induction.

Q3:

Let A and B be sets with $A \subseteq B$. Let x be an element in $A \cap B$. By the definition of an intersection, x is an element in A and x is an element in B. Since A intersects with B in $A \cap B$, then x is also an element in A on the right-hand side. Thus, proving that $A \cap B = A$.

Q4:

$d_k = 3d_{k-1} + 2$, for all integers $k \geq 2$

$$d_1 = 3$$

$$d_2 = 3(d_1) + 2 = 3(3) + 2 = 3^2 + 2 = 11$$

$$d_3 = 3(d_2) + 2 = 3(3^2 + 2) + 2 = 3^3 + 3*2 + 2 = 35$$

$$d_4 = 3(d_3) + 2 = 3(3^3 + 3*2 + 2) + 2 = 3^4 + 3^2*2 + 3*2 + 2 = 107$$

Guess:

$$3^4 + 3^2*2 + 3*2 + 2$$

$$= 3^n + 3^{n-2}*2 + 3^{n-3}*2 + 2$$

$$= 3^n + (\sum_{k=0, n-2} 3^i)$$

$$= 3^n + ((3^{n-2+1} - 1)/(3-1))$$

$$= 3^n + ((3^{n-1} - 1)/2)$$

$$= ((2*3^n)/2) + ((3^{n-1} - 1)/2)$$

Q5:

Base step: According to I, $a \in S$. By II, $s \in S$ so sa and sb will be elements of S .

Inductive step: According to (I), $a \in S$, so by recursion, (II.a) sa will begin the string, followed by (II.b) sb . Since both sa and sb are elements of S , and a is an element of s , every string will begin with a and then followed by b . This is within the means of the restrictions, thus proving that every string S begins with an a .

Q6:

- a There are 900 integers between 100 and 999. To determine how many 3-digit integers between 100 and 999 are divisible by 5, we check:

[] [] [] = box 1 = 1-9 since it cannot be a zero = 9!

Box 2 = 0-10 = 10!

Box 3 = 0-2 since only two integers are divisible by 5 = 2!

Thus: $9! * 10! * 2! =$ the number of 3-digit integers divisible by 5.

- b $\{3, 12\}, \{4, 11\}, \{5, 10\}, \{6, 9\}, \{7, 8\} = 5$ sets of integers that could equal 15.

By the pigeon theorem, we can check if there must be two integers whose sum is 15.

So, there are 6 objects (pigeons) and 5 sets equaling 15 (holes). Thus for at least one hole, there will be 2 pigeons, so there must be two integers whose sum is 15.

Q7: "MANUSCRIPT" = 10 letters. If the letters US must be next to each other, then their new ordering will be: "MAN[US]CRIPT" = 9 letters.

$9! = \text{ans.}$

Q8:

7 types

40 batteries

including 4 batteries of a7b

$40 - 4 = 36$

$C(n + r - 1, n)$

$C(36 + 7 - 1, 36) = C(42, 36) = 42! / 36!(42 - 36)! = \text{ans}$

Q9:

There cannot be a simple graph that have differing degrees. For example, if $n=4$, and the degree of each vertex is: 1, 2, 3, and 4, then the additive of them is 10e. When divided by 2, this equals

5 which is odd thus proving that a simple graph that has n vertices all of different degrees cannot be.

Q10:

This graph is proven to have a Euler Circuit. No edges are repeated, but vertices are repeated. This is a closed circuit. One method of the circuit can be completed as:

$F \rightarrow G \rightarrow A \rightarrow B \rightarrow C \rightarrow E \rightarrow F \rightarrow A \rightarrow D \rightarrow C \rightarrow F$

This shows that we begin and end at the same vertex without repeating any edges, only vertices, and is thus a Euler Circuit.