Devin Gendron CS225-402 Midterm 11/12/17 – 6:50PM EST Exam Time

Logic Symbols:  $\geq \leq \neq \neg \land \lor \oplus \equiv \rightarrow \longleftrightarrow \exists \forall$ Set Symbols:  $\in \notin \subseteq \subset \supseteq \supset \varnothing \cup \cap \times$ 

1.

a)

P=to log on to the server

Q=to have a valid password

"If we want to log on to the server, then a valid password is necessary"

The contrapositive =  $\sim Q \rightarrow \sim P$ , so:

~P=to not log on to the server

~Q=to not have a valid password

Thus, the contrapositive is: "If there is not a valid password, then you cannot log on to the server"

b)

P=x is nonnegative

Q=x is positive OR x is zero

The contrapositive =  $\sim Q \rightarrow \sim P$ , so:

~P=x is not nonnegative

 $\sim$ Q=x is not positive and x is not zero

Thus, the contrapositive is: "If x is not positive and x is not zero, then x is not nonnegative."

1.2.

a)

P=you pay the bill

Q=you lose the company's services

$$\sim (P \lor Q) = (\sim P \land \sim Q)$$

Thus the negation is: "You do not pay the bill and you do not lose the company's services."

b)

P=A positive integer is prime

~Q=it has no divisors other than 1 and itself

 $\sim (P \rightarrow \sim Q) = P \cap Q$  Thus the negation is:

"A positive integer is prime and it does have divisors other than 1 or itself."

2.

P(x): x is a mathematics problem

Q(x): x is time consuming

R(x): x is easy

S(x): x is solvable

- a) For all problems in science, such that if it is a mathematics problem, then it is solvable or it is not easy.
- b) Not all problems are solvable if they are easy. -Or- There exists a problem that is solvable and is not easy.
- c) There exists a problem in science such that it is a mathematics problem and it is time consuming.
- d) There does not exist a problem in science such that it is time consuming and it is solvable.

3.

a) "For all people, if they are young, then they are not good athletes."

$$\forall x \mid S(x) \rightarrow \neg W(x)$$

b) "There exists a person, such that, they are female and not good athletes, but they are young"

$$\exists x \mid B(x) \rightarrow \neg W(x) \land S(x)$$

c) Not all young people are good athletes.

$$\neg \forall x | S(x) \rightarrow W(x)$$

d) There exists a person who is not a good athlete and is not a female.

$$\exists x | \neg W(x) \land \neg B(x)$$

Р	Q	R	~R	(Q v R)	(P ∧ ¬R)	$(P \rightarrow (Q \ V \ R))$	$((P \land \neg R) \rightarrow Q)$
Т	Т	Т	F	Т	F	Т	Т
Т	Т	F	Т	Т	Т	Т	Т
Т	F	Т	F	T	F	Т	Т
Т	F	F	Т	F	T	F	F
F	Т	Т	F	T	F	Т	Т
F	Т	F	Т	T	F	T	Т
F	F	Т	F	T	F	T	T
F	F	F	Т	F	F	Т	Т

Since  $(P \rightarrow (Q \lor R))$  and  $((P \land \neg R) \rightarrow Q)$  have the same truth table values, they are both logically equivalent.

5. Given two rational numbers r and s with r< s, we must find if there is another rational number between r and s. By definition of a rational number, r=a/b and s=c/d where a, b, c, and d are integers and b and d l=0. To find if there is a number between r and s that is rational, we add r and s together (r+s) and then divide it by two ((r+s)/2). We then plug in our integers a, b, c, d:

$$(((a/b) + (c/d))/2)$$
 -we then solve to get:  
= $(ad + cb)/2bd$ 

Since a, b, c, and d are integers, by the definition of summation, and by the definition of a product, (ad + cb)/2bd is rational and the original statement is proven. An easy way to view this is: r < ((r+s)/2) < s

6. Show that if k is an integer and 5k + 4 is odd then k is odd. Suppose not. Suppose that 5k+4 is odd and k is even. By contradiction, P=k is an integer and 5k+4 is odd, Q=k is odd, then  $P^{\sim}Q=k$  is an integer and 5k+4 is odd and k is even.

So, if k is even, then by definition of an even number, k=2m where m is an arbitrary integer that is even. If we plug in k, we get 5(2m)+4 which is even and not odd. Thus by contradiction we prove that the supposition is false and the original statement is true.

7.

By the definition of a set,  $B \subseteq A$ . For some arbitrary integer x, x is an element in B and x is an element in A by the previously defined set  $B \subseteq A$ . To determine if  $B \subseteq A$ , then 18b-2 and 6(random integer)+4 will produce a random integer a that is in the set B.

3b-1=a is an integer by definition of a difference. So if we plug it back in to 6a+4 we get:

Thus  $B \subseteq A$  is proven.

8. 
$$(B^c \cup (B^c - A))^c = B$$

$$(B^c \cup (B^c - A))^c = (B^c \cup (B^c \cap A^c))^c$$
 by the set difference law  $= ((B^c)^c \cap ((B^c)^c \cup (A^c)^c)$  by De Morgan's Laws for all sets A and B by the double complement law by the absorption laws

9.

9.1:

$$\sum i=1 \rightarrow 7$$
:  $4i - \sum i=1 \rightarrow 2$ :  $4i + \sum i=0 \rightarrow 7$ :  $(-1)^i - \sum i=0 \rightarrow 2$ :  $(-1)^i =$   
= $(7(7+1)/2) - (2(2+1)/2) + ((-1^i(7+1)-1)/2) - ((-1^i(2+1)-1)/2)$   
= $28 - 3 + 0 + (-1) = 24(ans)$   
9.2:

$$\sum i=0 \rightarrow 4$$
:  $(4^{j+3}+6) = 64*\sum i=0 \rightarrow 4$ :  $(4^{j}) + \sum i=0 \rightarrow 4$ :  $6 = 64*((4^{4+1}-1)/4-1) + 6*4 = (ans)$ 

$$64*((4^{5}-1)/3) + 24 = (ans)$$

10.

$$P(n) = 3^0 + 3^1 + \dots + 3^n = \frac{1}{2}(3^{n+1} - 1)$$

Base Case P(0):

LHS:  $3^0 = 1$ 

RHS:  $\frac{1}{2}(3^{0+1}-1) = 1$ 

P(0) is true.

Inductive Case: If P(k) is true, then show that P(k+1) is also true.

$$P(k) = 3^0 + 3^1 + \dots + 3^k = \frac{1}{2}(3^{(k+1)} - 1)$$

Then:

$$P(k + 1) = 3^0 + 3^1 + \dots + 3^(k+1) = \frac{1}{2}(3^(k+2) - 1)$$

Thus to prove that P(k+1) is true solve:

$$\frac{1}{2}(3^{k+1}-1) + 3^{k+1}$$

$$= \frac{1}{2}(3^{k+1}-1) + (2/2)*3^{k+1}$$

$$=\frac{1}{2}(3^{k+1}-1+2*3^{k+1})$$

$$=\frac{1}{2}(1^{1*} 3^{k+1}-1+2*3^{k+1})$$

$$=\frac{1}{2}(3^{k+2}-1)$$

Thus by mathematical induction, P(k+1) is proven to be true.