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1 Tips

1. pow 返回浮点数可能会导致误差，慎用!!! 自己写。
2. 样例数用 long long 可能会 TLE
3. 多组数据记得初始化
4. 小心过程量 overflow
5. Type conversions
6. 在取余的情况下，要避免减法运算的结果出现负数 $(... + MOD) \% MOD$;
7. fill 注意初始范围为从 1 到 n 的情况
8. 提交换行
9. c++11: long long int abs (long long int n);
10. 输入读取判断其中一个为零
11. 搜索 flag 标记可能要加上方案数
12. 数据范围特别大时考虑不可能情况
13. SegmentTree 开大点以防 RE ($1e5 \rightarrow 1 \ll 17$)
14. 滚动数组未全部清零

2 Basic Algorithm

2.1 Sort

```
1 // Bubble Sort
2 for (int i = 0; i < N; i++)
3     for (int j = 0; j < N - i - 1; j++)
4         if (A[j] > A[j + 1]) swap(A[j], A[j + 1]);
```

```
1 // Insertion Sort
2 for (int i = 1; i < N; i++) {
3     int tmp = A[i], j;
4     for (j = i - 1; j >= 0 && A[j] > tmp; j--)
5         A[j + 1] = A[j];
6     A[++j] = tmp;
7 }
```

```
1 // Selection Sort
2 for (int i = 0; i < N; i++)
3     for (int j = i + 1; j < N; j++)
4         if (A[i] > A[j]) swap(A[i], A[j])
```

```
1 // Merge Sort
2 void merge_sort(int A[], int l, int r) {
3     if (l >= r) return ;
4     int mid = (l + r) / 2;
5     merge_sort(A, l, mid);
6     merge_sort(A, mid + 1, r);
7     int i, j, k;
8     i = l; j = mid + 1; k = l;
9     while (i <= mid && j <= r) {
10         if (A[i] <= A[j]) B[k] = A[i++];
11         else B[k] = A[j++];
12         k++;
13     }
14     while (i <= mid) {
15         B[k] = A[i++];
16         k++;
17     }
18     while (j <= r) {
19         B[k] = A[j++];
20         k++;
21     }
```

```

21     }
22     memcpy(A, B, r - l + 1);
23     return ;
24 }

1 // Quick Sort
2 void quicksort(int A[], int l, int r) {
3     int i = l, j = r, mid = A[(r - l) / 2 + l];
4     while (i <= j) {
5         while (A[i] < mid) i++;
6         while (A[j] > mid) j--;
7         if (i <= j) {
8             swap(A[i], A[j]);
9             ++i; --j;
10        }
11    }
12    if (i < r) quicksort(A, i, r);
13    if (l < j) quicksort(A, l, j);
14    return ;
15 }

```

- Heap Sort (见堆的内容)

2.2 DP

2.2.1 LIS

```

1 int A[MAX_N];
2 long lis(int n) {
3     int dp[MAX_N];
4     fill(dp, dp + n, INF);
5     for (int i = 0; i < n; ++i)
6         *lower_bound(dp, dp + n, A[i]) = A[i]; // lds: -A[i]; ln:
           upper_bound
7     return lower_bound(dp, dp + n, INF) - dp;
8 }

```

2.2.2 Knapsack Problem

- 0/1 背包

$$f[i, j] = \max(f[i-1, j], f[i-1, j-w[i]] + v[i])$$

```

1 for (int i = 0; i < N; ++i)
2     for (int j = W; j >= w[i]; --j)
3         f[j] = max(f[j - w[i]] + c[i], f[j]);

```

- 完全背包

$$f[i, j] = \max(f[i-1, j], f[i-1, j-w[i]] + j[i])$$

```

1 for (int i = 0; i < N; ++i)
2     for (int j = w[i]; j <= W; ++j)
3         f[j] = max(f[j - w[i]] + c[i], f[j]);

```

注意循环顺序的不同背后思路。

- 一个简单的优化：若两件物品 i 、 j 满足 $w[i] \leq w[j]$ 且 $c[i] \geq c[j]$ ，则讲物品 j 去掉，不用考虑。
- 转化为 01 背包问题求解：
 - 第 i 种物品转化为 $\frac{V}{w[i]}$ 件费用于价值均不变的物品。
 - 第 i 种物品拆成费用为 $w[i] * 2^k$ ，价值为 $c[i] * 2^k$ 的若干件物品其中 k 满足 $w[i] * 2^k < V$
- 多重背包

$$f[i, j] = \max(f[i-1, j-w[i]*k] + v[i]*k | 0 \leq k \leq m[i])$$

- 优化：转化为 01 背包问题
 - 将第 i 件物品分成若干件物品，每件物品的系数分别为：1, 2, 4, ..., $2^{(k-1)}$, $n[i] - 2^k$
 - ** 根据 w , v 范围改变 DP 对象，可以考虑针对不同价值计算最小的重量。（ $f[i][j]$, 其中 j 代表价值总和）**

```

1 for (int i = 0; i < N; ++i) {
2     int num = m[i];
3     for (int k = 1; num > 0; k <= 1) {
4         int mul = min(k, num);
5         for (int j = W; j >= w[i] * mul; --j) {
6             f[j] = max(f[j - w[i] * mul] + v[i] * mul, f[j]);
7         }
8         num -= mul;
9     }
}

```

10 }

- 混合三种背包

弄清楚上面三种背包后分情况就好

- 超大背包

$$- 1 \leq n \leq 40, 1 \leq w_i, v_i \leq 10^{15}, 1 \leq W \leq 10^{15}$$

```

1  int n;
2  ll w[MAX_N], v[MAX_N], W;
3  Pll ps[1 << (MAX_N / 2)]; // (w, v);
4
5  void solve() {
6      int n2 = n / 2;
7      for (int i = 0; i < 1 << n2; ++i) {
8          ll sw = 0, sv = 0;
9          for (int j = 0; j < n2; ++j)
10             if (i >> j & 1) {
11                 sw += w[j];
12                 sv += v[j];
13             }
14             ps[i] = Pll(sw, sv);
15     }
16     sort(ps, ps + (1 << n2));
17     int m = 1;
18     for (int i = 1; i < 1 << n2; ++i)
19         if (ps[m - 1].second < ps[i].second)
20             ps[m++] = ps[i];
21
22     ll res = 0;
23     for (int i = 0; i < 1 << (n - n2); ++i) {
24         ll sw = 0, sv = 0;
25         for (int j = 0; j < n - n2; ++j)
26             if (i >> j & 1) {
27                 sw += w[n2 + j];
28                 sv += v[n2 + j];
29             }
30         if (sw <= W) {
31             ll tv = (lower_bound(ps, ps + m, make_pair(W - sw, INF)) -
32                 1) -> second;
33             res = max(res, sv + tv);
34     }
35 }
```

```

35     printf("%lld\n", res);
36 }

```

- 二维费用背包

$$f[i, j, k] = \max(f[i - 1, j, k], f[i - 1, j - a[i], k - b[i]] + c[i])$$

二维费用可由最多取 m 件等方式隐蔽给出。

- 分组背包

$$f[k, j] = \max(f[k - 1, j], f[k - 1, j - w[i]] + v[i] | i \in K)$$

```

1  for (int k = 0; k < K; ++k)
2      for (j = W; j >= 0; --j)
3          for (int i = 0; i <= m[k]; ++i)
4              f[j] = max(f[j - w[i]] + v[i], f[j]);

```

显然可以对每组中物品应用完全背包中“一个简单有效的优化”

- 有依赖背包

由 NOIP2006 金明的预算方案引申，对每个附件先做一个 01 背包，再与组件得到一个 $V - w[i] + 1$ 个物品组。更一般问题，依赖关系由「森林」形式给出，涉及到树形 DP 以及泛化物品，这里不表。

- 背包问题方案总数

$$f[i, j] = \text{sum}(f[i - 1, j], f[i - 1, j - w[i]] + v[i]), f[0, 0] = 0$$

更多内容详见「背包九讲」

2.2.3 Maximum Subarray Sum

```

1  int max_subarray_sum(int A[], int n) {
2      int res, cur;
3      if (!A || n <= 0) return 0;
4      res = cur = a[0];
5      for (int i = 0; i < n; ++i) {
6          if (cur < 0) cur = a[i];
7          else cur += a[i];
8          res = max(cur, res);

```

```
9     }
10    return res;
11 }
```

2.3 Set

```
1 // 子集枚举
2 int sub = sup;
3 do {
4     sub = (sub - 1) & sup;
5 } while (sub != sup); // -1 & sup = sup;
6
7 // 势为 k 的集合枚举
8 int comb = (1 << k) - 1;
9 while (comb < 1 << n) {
10     int x = comb & -comb, y = comb + x;
11     comb = ((comb & ~y) / x >> 1) | y;
12 }
13
14 // 排列组合
15 do {
16
17 } while (next_permutation(A, A + N)); // prev_permutation
```

2.4 Bit operation

```
1 int __builtin_ffs (unsigned int x)
2 //返回x的最后一位1的是从后向前第几位，比如7368 (1110011001000) 返回4。
3 int __builtin_clz (unsigned int x)
4 // 返回前导的0的个数。
5 int __builtin_ctz (unsigned int x)
6 // 返回后面的0的个数，和__builtin_clz相对。
7 int __builtin_popcount (unsigned int x)
8 // 返回二进制表示中1的个数。
9 int __builtin_parity (unsigned int x)
10 // 返回x的奇偶校验位，也就是x的1的个数模2的结果。
```

3 Data Structure

```
1 // Heap
2 int heap[MAX_N], sz = 0;
3 void push(int x) {
4     int i = sz++;
5
6     while (i > 0) {
7         int p = (i - 1) / 2;
8         if (heap[p] <= x) break;
9         heap[p] = heap[i];
10        i = p;
11    }
12    heap[i] = x;
13 }
14 int pop() {
15     int ret = heap[0];
16     int x = heap[--sz];
17     int i = 0;
18     while (i * 2 + 1 < sz) {
19         int a = i * 2 + 1, b = i * 2 + 2;
20         if (b < sz && heap[b] < heap[a]) a = b;
21         if (heap[a] >= x) break;
22         heap[i] = heap[a];
23         i = a;
24     }
25     heap[i] = x;
26     return ret;
27 }
```

```
1 // Binary Search Tree
2 struct node {
3     int val;
4     node *lch, *rch;
5 };
6
7 node *insert(node *p, int x) {
8     if (p == NULL) {
9         node *q = new node;
10        q->val = x;
11        q->lch = q->rch = NULL;
12        return q;
13    }
```



```
13     } else {
14         if (x < p->val) p->lch = insert(p->lch, x);
15         else p->rch = insert(p->rch, x);
16         return p;
17     }
18 }
19 bool find(node *p, int x) {
20     if (p == NULL) return false;
21     else if (x == p->val) return true;
22     else if (x < p->val) return find(p->lch, x);
23     else return find(p->rch, x);
24 }
25 node *remove(node *p, int x) {
26     if (p == NULL) return NULL;
27     else if (x < p->val) p->lch = remove(p->lch, x);
28     else if (x > p->val) p->rch = remove(p->rch, x);
29     else if (p->lch == NULL) {
30         node *q = p->rch;
31         delete p;
32         return q;
33     } else if (p->lch->rch == NULL) {
34         node *q = p->lch;
35         q->rch = p->rch;
36         delete p;
37         return q;
38     } else {
39         // 把左儿子子孙中最大的节点提到需要删除的节点上
40         node *q;
41         for (q = p->lch; q->rch->rch != NULL; q = q->rch);
42         node *r = q->rch;
43         q->rch = r->lch;
44         r->lch = p->lch;
45         r->rch = p->rch;
46         delete p;
47         return r;
48     }
49     return p;
50 }
```

```
1 // Union-find Set
2 int par[MAX_N];
3 int rnk[MAX_N];
4 void init(int n) {
```

```
5     for (int i = 0; i < n; ++i) {
6         par[i] = i;
7         rnk[i] = 0;
8     }
9 }
10 int find(int x) {
11     return par[x] == x? x : par[x] = find(par[x]);
12 }
13 bool same(int x, int y) {
14     return find(x) == find(y);
15 }
16 void unite(int x, int y) {
17     x = find(x);
18     y = find(y);
19     if (x == y) return;
20     if (rnk[x] < rnk[y]) {
21         par[x] = y;
22     } else {
23         par[y] = x;
24         if (rnk[x] == rnk[y]) rnk[x]++;
25     }
26 }
```

当然，更快捷简单的做法，是使用 C++ 的 container。

```
1 // Segment Tree
2 const int MAX_N = 1 << 17;
3 int n, dat[2 * MAX_N - 1];
4 void init(int _n) {
5     n = 1;
6     while (n < _n) n <= 1;
7     for (int i = 0; i < 2 * n - 1; ++i)
8         dat[i] = INF;
9 }
10 void update(int k, int a) {
11     k += n - 1;
12     dat[k] = a;
13     while (k > 0) {
14         k = (k - 1) / 2;
15         dat[k] = min(dat[2 * k + 1], dat[2 * k + 2]);
16     }
17 }
18 // query [a, b), index k in [l, r)
```

```

19 // query(a, b, 0, 0, n)
20 int query(int a, int b, int k, int l, int r) {
21     if (r <= a || b <= l) return INF;
22     if (a <= l && r <= b) return dat[k];
23     else {
24         int v1 = query(a, b, k * 2 + 1, l, (l + r) / 2);
25         int v2 = query(a, b, k * 2 + 2, (l + r) / 2, r);
26         return min(v1, v2);
27     }
28 }

```

```

1 // RMQ
2 int n, dat[2 * MAX_N - 1];
3 void init(int _n) {
4     n = 1;
5     while (n < _n) n <= 1;
6     for (int i = 0; i < 2 * n - 1; ++i)
7         dat[i] = INF;
8 }
9 void update(int k, int a) {
10    k += n - 1;
11    dat[k] = a;
12    while (k > 0) {
13        k = (k - 1) / 2;
14        dat[k] = min(dat[2 * k + 1], dat[2 * k + 2]);
15    }
16 }
17 // query [a, b), index k in [l, r)
18 // query(a, b, 0, 0, n)
19 int query(int a, int b, int k, int l, int r) {
20     if (r <= a || b <= l) return INF;
21     if (a <= l && r <= b) return dat[k];
22     else {
23         int v1 = query(a, b, k * 2 + 1, l, (l + r) / 2);
24         int v2 = query(a, b, k * 2 + 2, (l + r) / 2, r);
25         return min(v1, v2);
26     }
27 }

```

```

1 //Sparse Table
2 const int MAX_N = 1e5 + 10;
3 const int MAX_K = 31 - __builtin_clz(MAX_N);
4

```

```

5  int n, ST[MAX_N][MAX_K + 1], A[MAX_N];
6  void build(int N) {
7      for (int i = 0; i < N; ++i)
8          ST[i][0] = A[i];
9      int k = 31 - __builtin_clz(N);
10     for (int j = 1; j <= k; ++j)
11         for (int i = 0; i <= N - (1 << j); ++i)
12             ST[i][j] = min(ST[i][j - 1], ST[i + (1 << (j - 1))][j - 1]);
13 }
14 int query(int l, int r) {
15     if (l >= r) return 0;
16     int ans = INF, k = 31 - __builtin_clz(r - l);
17     for (int j = k; j >= 0; --j)
18         if (l + (1 << j) - 1 <= r) {
19             ans = min(ans, ST[l][j]);
20             l += 1 << j;
21         }
22     return ans;
23 }
24 int RMQ(int l, int r) {
25     if (l >= r) return 0;
26     int k = 31 - __builtin_clz(r - l);
27     return min(ST[l][k], ST[r - (1 << k)][k]);
28 }

```

4 Graph

```

1  struct edge {
2      int from;
3      int to, dis;
4  };
5  vector<edge> G[MAX_V];
6  vector<edge> es;
7  bool vis[MAX_V];
8  int V, E, pre[MAX_V], dist[MAX_V];
9  // int cost[MAX_V][MAX_V];

```

```

1  // Shortest Way

```

```
2 void floyed() {
3     for (int k = 1; k <= n; ++k)
4         for (int i = 1; i <= n; ++i)
5             for (int j = 1; i <= n; ++j)
6                 dis[i][j] = min(dis[i][k] + dis[k][j], dis[i][j]);
7 }
8 void dijkstra(int s) {
9     priority_queue<Pii, vector<Pii>, greater<Pii> > que; // first 是最短
10    // 距离, second 是顶点编号
11    fill(dist, dist + V, INF);
12    dist[s] = 0; que.push(Pii(0, s));
13    while (!que.empty()) {
14        Pii p = que.top(); que.pop();
15        int v = p.second;
16        if (dist[v] < p.first) continue;
17        for (int i = 0; i < G[v].size(); i++) {
18            edge e = G[v][i];
19            if (dist[e.to] > dist[v] + e.dis) {
20                dist[e.to] = dist[v] + e.dis;
21                que.push(Pii(dist[e.to], e.to));
22            }
23        }
24    }
25    void bellman_ford(int s) {
26        fill(dist, dist + V, INF);
27        dist[s] = 0;
28        while (true) {
29            bool update = false;
30            for (int i = 0; i < E; ++i) {
31                edge e = es[i];
32                if (dist[e.from] != INF && dist[e.from] + e.dis < dist[e.to]) {
33                    update = true;
34                    dist[e.to] = dist[e.from] + e.dis;
35                }
36            }
37            if (!update) break;
38        }
39    }
40    bool find_negative_loop() {
41        memset(dist, 0, sizeof dist);
42        for (int i = 0; i < V; ++i)
```

```

43     for (int j = 0; j < E; ++j) {
44         edge e = es[j];
45         if (d[e.to] > d[e.from] + e.dis) {
46             d[e.to] = d[e.from] + e.dis;
47             if (i == V - 1) return true;
48         }
49     }
50     return false;
51 }
52 void spfa(int s) {
53     queue<int> que;
54     fill(dist, dist + V, INF);
55     fill(vis, vis + V, false);
56     dist[s] = 0; que.push(s); vis[s] = true;
57     while (!que.empty()) {
58         int v = que.front(); que.pop();
59         vis[v] = false;
60         for (int i = 0; i < G[v].size(); ++i) {
61             int u = G[v][i].to;
62             if (dist[u] > dist[v] + G[v][i].dis) {
63                 dist[u] = dist[v] + G[v][i].dis;
64                 if (!vis[u]) {
65                     que.push(u);
66                     vis[u] = true;
67                 }
68             }
69         }
70     }
71 }

```

```

1 // Spanning Tree
2 int prime() {
3     /*
4     fill(dist, dist + V, INF);
5     fill(vis, vis + V, false);
6     dist[0] = 0;
7     int res = 0;
8     while (true) {
9         int v = -1;
10        for (int u = 0; u < V; ++u) {
11            if(!vis[u] && (v == -1 || dist[u] < dist[v])) v = u;
12        }
13        if (v == -1) break;

```

```
14     vis[v] = true;
15     res += dist[v];
16     for (int u = 0; u < V; u++)
17         dist[u] = min(dist[u], cost[v][u]);
18 }
19 /**/
20 priority_queue<Pii, vector<Pii>, greater<Pii> > que;
21 int res = 0;
22 fill(dist, dist + V, INF);
23 fill(vis, vis + V, false);
24 dist[0] = 0;
25 que.push(Pii(0, 0));
26 while (!que.empty()) {
27     Pii p = que.top(); que.pop();
28     int v = p.second;
29     if (vis[v] || dist[v] < p.first) continue;
30     res += dist[v]; vis[v] = true;
31     for (int i = 0; i < G[v].size(); ++i) {
32         edge e = G[v][i];
33         if (dist[e.to] > e.dis) {
34             dist[e.to] = e.dis;
35             que.push(Pii(dist[e.to], e.to));
36         }
37     }
38 }
39 return res;
40 }
41
42 bool cmp(const edge e1, const edge e2) {
43     return e1.dis < e2.dis;
44 }
45 int kruskal() {
46     sort(es.begin(), es.end(), cmp);
47     init(V);
48     int res = 0;
49     for (int i = 0; i < E; ++i) {
50         edge e = es[i];
51         if (!same(e.from, e.to)) {
52             unite(e.from, e.to);
53             res += e.dis;
54         }
55     }
56     return res;
```

57 }

```

1 // SCC
2 int V, cmp[MAX_V];
3 vector<int> G[MAX_V], rG[MAX_V], vs;
4 bool used[MAX_V];
5
6 void add_edge(int from, int to) {
7     G[from].push_back(to); rG[to].push_back(from);
8 }
9 void dfs(int v) {
10     used[v] = true;
11     for (int i = 0; i < G[v].size(); ++i)
12         if (!used[G[v][i]]) dfs(G[v][i]);
13     vs.push_back(v);
14 }
15 void rdfs(int v, int k) {
16     used[v] = true;
17     cmp[v] = k;
18     for (int i = 0; i < rG[v].size(); ++i)
19         if (!used[rG[v][i]]) rdfs(rG[v][i], k);
20 }
21 int scc() {
22     memset(used, 0, sizeof used);
23     vs.clear();
24     for (int v = 0; v < V; ++v)
25         if (!used[v]) dfs(v);
26     memset(used, 0, sizeof used);
27     int k = 0;
28     for (int i = vs.size() - 1; i >= 0; --i)
29         if (!used[vs[i]]) rdfs(vs[i], k++);
30     return k;
31 }

```

```

1 // Bipartite Matching
2 void add_edge(int u, int v) {
3     G[u].push_back(v); G[v].push_back(u);
4 }
5 bool dfs(int v) {
6     used[v] = true;
7     rep(i, 0, G[v].size()) {
8         int u = G[v][i], w = match[u];
9         if (w < 0 || (!used[w] && dfs(w))) {

```



```

10         match[v] = u; match[u] = v;
11         return true;
12     }
13 }
14 return false;
15 }
16 int bipartite_matching() {
17     int res = 0;
18     memset(match, -1, sizeof match);
19     rep (v, 0, V)
20         if (match[v] < 0) {
21             memset(used, false, sizeof used);
22             if (dfs(v)) ++res;
23         }
24     return res;
25 }

```

```

1 // Network Flow
2 struct edge{
3     int to, cap, rev;
4 };
5 vector<edge> G[MAX_V];
6 int level[MAX_V], iter[MAX_V];
7 void add_edge(int from, int to, int cap) {
8     G[from].push_back((edge){to, cap, static_cast<int>(G[to].size())});
9     G[to].push_back((edge){from, 0, static_cast<int>(G[from].size() - 1)});
10 }
11 // Ford-Fulkerson
12 int dfs(int v, int t, int f) {
13     if (v == t) return f;
14     flag[v] = true;
15     rep(i, 0, G[v].size()) {
16         edge &e = G[v][i];
17         if (!flag[e.to] && e.cap > 0) {
18             int d = dfs(e.to, t, min(f, e.cap));
19             if (d > 0) {
20                 e.cap -= d;
21                 G[e.to][e.rev].cap += d;
22                 return d;
23             }
24         }
25     }
26 }

```

```
26     return 0;
27 }
28 int max_flow(int s, int t) {
29     int flow = 0;
30     for(;;) {
31         memset(flag, false, sizeof flag);
32         int f = dfs(s, t, INF);
33         if (!f) return flow;
34         flow += f;
35     }
36 }
37 // Dinic
38 void bfs(int s) {
39     memset(level, -1, sizeof(level));
40     queue<int> que;
41     level[s] = 0; que.push(s);
42     while (!que.empty()) {
43         int v = que.front(); que.pop();
44         for (int i = 0; i < G[v].size(); ++i) {
45             edge &e = G[v][i];
46             if (e.cap > 0 && level[e.to] < 0) {
47                 level[e.to] = level[v] + 1;
48                 que.push(e.to);
49             }
50         }
51     }
52 }
53 int dfs(int v, int t, int f) {
54     if (v == t) return f;
55     for (int &i = iter[v]; i < G[v].size(); ++i) {
56         edge &e = G[v][i];
57         if (e.cap > 0 && level[v] < level[e.to]) {
58             int d = dfs(e.to, t, min(f, e.cap));
59             if (d > 0) {
60                 e.cap -= d;
61                 G[e.to][e.rev].cap += d;
62                 return d;
63             }
64         }
65     }
66     return 0;
67 }
68 int max_flow(int s, int t) {
```

```

69     int flow = 0;
70     for (;;) {
71         bfs(s);
72         if (level[t] < 0) return flow;
73         memset(iter, 0, sizeof iter);
74         int f;
75         while ((f = dfs(s, t, INF)) > 0) {
76             flow += f;
77         }
78     }
79 }

```

```

1 // min_cost_flow
2 void add_edge(int from, int to, int cap, int cost) {
3     G[from].push_back((edge){to, cap, cost, (int)G[to].size()});
4     G[to].push_back((edge){from, 0, -cost, (int)G[from].size() - 1});
5 }
6 int min_cost_flow(int s, int t, int f) {
7     int res = 0;
8     fill(h, h + V, 0);
9     while (f > 0) {
10         priority_queue<Pii, vector<Pii>, greater<Pii> > que;
11         fill(dist, dist + V, INF);
12         dist[s] = 0; que.push(Pii(0, s));
13         while (!que.empty()) {
14             Pii p = que.top(); que.pop();
15             int v = p.second;
16             if (dist[v] < p.first) continue;
17             rep(i, 0, G[v].size()) {
18                 edge &e = G[v][i];
19                 if (e.cap > 0 && dist[e.to] > dist[v] + e.cost + h[v] -
20                     h[e.to]) {
21                     dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
22                     prevv[e.to] = v;
23                     preve[e.to] = i;
24                     que.push(Pii(dist[e.to], e.to));
25                 }
26             }
27             if (dist[t] == INF) return -1;
28             rep(v, 0, V) h[v] += dist[v];
29             int d = f;
30             for (int v = t; v != s; v = prevv[v])

```

```
31         d = min(d, G[prevv[v]][preve[v]].cap);
32         f -= d;
33         res += d * h[t];
34         for (int v = t; v != s; v = prevv[v]) {
35             edge &e = G[prevv[v]][preve[v]];
36             e.cap -= d;
37             G[v][e.rev].cap += d;
38         }
39     }
40     return res;
41 }
```

```
1 // stoer_wagner 全局最小割
2 void search() {
3     memset(vis, false, sizeof vis);
4     memset(wet, 0, sizeof wet);
5     S = T = -1;
6     int imax, tmp;
7     rep(i, 0, V) {
8         imax = -INF;
9         rep(j, 0, V)
10             if (!cmb[j] && !vis[j] && wet[j] > imax) {
11                 imax = wet[j];
12                 tmp = j;
13             }
14         if (T == tmp) return;
15         S = T; T = tmp;
16         mc = imax;
17         vis[tmp] = true;
18         rep(j, 0, V)
19             if (!cmb[j] && !vis[j])
20                 wet[j] += G[tmp][j];
21     }
22 }
23 int stoer_wagner() {
24     memset(cmb, false, sizeof cmb);
25     int ans = INF;
26     rep(i, 0, V - 1) {
27         search();
28         ans = min(ans, mc);
29         if (ans == 0) return 0;
30         cmb[T] = true;
31         rep(j, 0, V)
```

```

32         if (!cmb[j]) {
33             G[S][j] += G[T][j];
34             G[j][S] += G[j][T];
35         }
36     }
37     return ans;
38 }

```

```

1 // LCA--Doubling
2 const int MAX_LOG_V = 32 - __builtin_clz(MAX_V);
3
4 vector<int> G[MAX_V];
5 int root, parent[MAX_LOG_V][MAX_V], depth[MAX_V];
6
7 void dfs(int v, int p, int d) {
8     parent[0][v] = p;
9     depth[v] = d;
10    for (int i = 0; i < G[v].size(); i++)
11        if (G[v][i] != p) dfs(G[v][i], v, d + 1);
12 }
13 void init(int V) {
14     dfs(root, -1, 0);
15     for (int k = 0; k + 1 < MAX_LOG_V; k++)
16         for (int v = 0; v < V; v++)
17             if (parent[k][v] < 0) parent[k + 1][v] = -1;
18             else parent[k + 1][v] = parent[k][parent[k][v]];
19 }
20 int lca(int u, int v) {
21     if (depth[u] > depth[v]) swap(u, v);
22     for (int k = 0; k < MAX_LOG_V; k++)
23         if ((depth[v] - depth[u]) >> k & 1)
24             v = parent[k][v];
25     if (u == v) return u;
26     for (int k = MAX_LOG_V - 1; k >= 0; k--)
27         if (parent[k][u] != parent[k][v])
28             u = parent[k][u], v = parent[k][v];
29     return parent[0][u];
30 }
31 // LCA--RMQ
32 vector<int> G[MAX_V];
33 int root, vs[MAX_V * 2 - 1], depth[MAX_V * 2 - 1], id[MAX_V];
34
35 int ST[2 * MAX_V][MAX_K];

```

```

36 void rmq_init(int* A, int N) {
37     for (int i = 0; i < N; i++)
38         ST[i][0] = i;
39     int k = 31 - __builtin_clz(N);
40     for (int j = 1; j <= k; j++)
41         for (int i = 0; i <= N - (1 << j); ++i)
42             if (A[ST[i][j - 1]] <= A[ST[i + (1 << (j - 1))][j - 1]])
43                 ST[i][j] = ST[i][j - 1];
44             else ST[i][j] = ST[i + (1 << (j - 1))][j - 1];
45 }
46 int query(int l, int r) {
47     if (l >= r) return -1;
48     int k = 31 - __builtin_clz(r - l);
49     return (depth[ST[l][k]] <= depth[ST[r - (1 << k)][k]]) ? ST[l][k] :
50         ST[r - (1 << k)][k];
51 }
52 void dfs(int v, int p, int d, int &k) {
53     id[v] = k;
54     vs[k] = v;
55     depth[k++] = d;
56     for (int i = 0; i < G[v].size(); i++) {
57         if (G[v][i] != p) {
58             dfs(G[v][i], v, d + 1, k);
59             vs[k] = v;
60             depth[k++] = d;
61         }
62     }
63 }
64 void init(int V) {
65     int k = 0;
66     dfs(root, -1, 0, k);
67     rmq_init(depth, 2 * V - 1);
68 }
69 int lca(int u, int v) {
70     return vs[query(min(id[u], id[v]), max(id[u], id[v]) + 1)];
71 }

```

5 Computational Geometry

```

1  const double EPS = 1e-8;
2  int sgn(double x) { return x < -EPS ? -1 : x > EPS ? 1 : 0; }
3  struct Point {
4      double x, y;
5      Point(double x = 0, double y = 0) : x(x), y(y) {}
6      Point operator + (Point p) { return Point(x + p.x, y + p.y); }
7      Point operator - (Point p) { return Point(x - p.x, y - p.y); }
8      Point operator * (double d) { return Point(x * d, y * d); }
9      bool operator < (Point p) {
10         return x != p.x ? x < p.x : y < p.y;
11     }
12     double dot(Point p) { return add(x * p.x, y * p.y); } // 内积
13     double det(Point p) { return add(x * p.y, -y * p.x); } // 外积
14 };
15 bool on_seg(Point p1, Point p2, Point q) {
16     return (p1 - q).det(p2 - q) == 0 && (p1 - q).dot(p2 - q) <= 0;
17 }
18 P intersection(Point p1, Point p2, Point q1, Point q2) {
19     return p1 + (p2 - p1) * ((q2 - q1).det(q1 - p1) / (q2 - q1).det(p2 - p1));
20 }
21 // 凸包
22 int convex_hull(Point *ps, int n, Point *ch) {
23     sort(ps, ps + n);
24     int k = 0;
25     for (int i = 0; i < n; ++i) {
26         while (k > 1 && (ch[k - 1] - ch[k - 2]).det(ps[i] - ch[k - 1]))
27             <= 0) k--;
28         ch[k++] = ps[i];
29     }
30     for (int i = n - 2, t = k; i >= 0; --i) {
31         while (k > t && (ch[k - 1] - ch[k - 2]).det(ps[i] - ch[k - 1]))
32             <= 0) k--;
33         ch[k++] = ps[i];
34     }
35     return k - 1;
36 }

```

Simpson 公式——二次函数近似原函数积分:

$$\int_a^b f(x)dx \approx \frac{b-a}{6} * \left(f(a) + 4f\left(\frac{a+b}{2}\right) + f(b) \right)$$

6 Math Problem

```

1 // returning count of nk in range [l, r], from Infinity
2 template<typename T> T mps(T l, T r, T k) {
3     return ((r - (r % k + k) % k) - (l + (k - l % k) % k)) / k + 1;
4 }
5 template<typename T> T gcd(T a, T b) {
6     //return (b)? gcd(b, a % b) : a;
7     while (b) { T t = a % b; a = b; b = t; } return a;
8 }
9 template<typename T> T lcm(T a, T b) {
10    return a / gcd(a, b) * b;
11 }
12 // find (x, y) s.t. a x + b y = gcd(a, b) = d
13 template<typename T> T exgcd(T a, T b, T &x, T &y) {
14     T d = a;
15     if (b) {
16         d = exgcd(b, a % b, y, x);
17         y -= a / b * x;
18     } else {
19         x = 1; y = 0;
20     }
21     return d;
22 }
23 template<typename T> T modular_linear(T a, T b, T n) {
24     T d, e, x, y;
25     d = exgcd(a, n, x, y);
26     if (b % d)
27         return -1;
28     e = x * (b / d) % n + n;
29     return e % (n / d);
30 }
31 template<typename T> T mod_mult(T a, T b, T mod) {
32     T res = 0;
33     while (b) {
34         if (b & 1) {
35             res = (res + a) % mod;
36             // res += a;
37             // if (res >= mod) res -= mod;

```



```

38     }
39     a = (a + a) % mod;
40     // a <= 1;
41     // if (a >= mod) a -= mod;
42     b >>= 1;
43 }
44 return res;
45 }
46 template<typename T> T mod_pow(T x, T n, T mod) {
47     T res = 1;
48     while (n) {
49         if (n & 1) res = mod_mult(res, x, mod);
50         x = mod_mult(x, x, mod);
51         n >>= 1;
52     }
53     return res;
54     // return b ? mod_pow(a * a % mod, b >> 1, mod) * (b & 1 ? a : 1) %
55     mod : 1;
56 }
57 template<typename T> T mod_inverse(T a, T m) {
58     T x, y;
59     exgcd(a, m, x, y);
60     return (m + x % m) % m;
61 }
62 // 线性求逆元 小心乘法溢出?
63 void init_inverse() {
64     inv[1] = 1;
65     for (int i = 2; i < MAX_N; i++)
66         inv[i] = (MOD - (MOD / i) * inv[MOD % i] % MOD) % MOD;
67 }
68 // A[i] * x % M[i] = B[i];
69 std::pair<int, int> linear_congruence(const std::vector<int> &A, const
70 std::vector<int> &B, const std::vector<int> &M) {
71     // wa 了把中间量开大? * 溢出
72     int x = 0, m = 1;
73     for(int i = 0; i < A.size(); i++) {
74         int a = A[i] * m, b = B[i] - A[i] * x, d = gcd(M[i], a);
75         if(b % d != 0) return std::make_pair(0, -1); // no
76         solution
77         int t = b / d * mod_inverse(a / d, M[i] / d) % (M[i] / d
78 );
79         x = x + m * t;
80         m *= M[i] / d;

```

```

77     }
78     while (x < m) x += m;
79     return std::make_pair(x % m, m);
80 }
81 ll CRT(vector<ll> &a, vector<ll> &m) {
82     ll M = 1LL, res = 0;
83     for (int i = 0; i < m.size(); ++i)
84         M *= m[i];
85     for (int i = 0; i < m.size(); ++i) {
86         ll Mi, Ti;
87         Mi = M / m[i]; Ti = mod_inverse(Mi, m[i]);
88         res = (res + a[i] * (Mi * Ti % M) % M) % M;
89     }
90     return res;
91 }
92 // only for MOD < 1e6;
93 ll fact[MOD + 10];
94 void init() {
95     fact[0] = 1;
96     for (int i = 1; i <= MOD; ++i)
97         fact[i] = fact[i - 1] * i % MOD;
98 }
99 int mod_fact(int n, int p, int &e) {
100     e = 0;
101     if (n == 0) return 1;
102     int res = mod_fact(n / p, p, e);
103     e += n / p;
104     if (n / p % 2 != 0) return res * (p - fact[n % p]) % p;
105     return res * fact[n % p] % p;
106 }
107 int mod_comb(int n, int k, int p) {
108     if (n < 0 || k < 0 || n < k) return 0;
109     if (n == 0) return 1;
110     int e1, e2, e3;
111     int a1 = mod_fact(n, p, e1), a2 = mod_fact(k, p, e2), a3 = mod_fact(
        n - k, p, e3);
112     if (e1 > e2 + e3) return 0;
113     return a1 * mod_inverse(a2 * a3 % p, p) % p;
114 }
115 ll lucas(ll n, ll k, const ll &p) {
116     if (n < 0 || k < 0 || n < k) return 0;
117     if (n == 0) return 1;
118     return lucas(n / p, k / p, p) * mod_comb(n % p, k % p, p) % p;

```

119 }

```

1 // 矩阵快速幂
2 typedef vector<int> vec;
3 typedef vector<vec> mat;
4 mat G(MAX_N);
5
6 mat mat_mul(mat &A, mat &B) {
7     mat C(A.size(), vec(B[0].size()));
8     for (int i = 0; i < A.size(); ++i)
9         for (int k = 0; k < B.size(); ++k)
10            for (int j = 0; j < B[0].size(); ++j)
11                C[i][j] = (C[i][j] + A[i][k] % MOD * B[k][j] % MOD + MOD
12                    ) % MOD;
13     return C;
14 }
15
16 mat mat_pow(mat A, ll n) {
17     mat B(A.size(), vec(A.size()));
18     for (int i = 0; i < A.size(); ++i)
19         B[i][i] = 1;
20     while (n > 0) {
21         if (n & 1) B = mat_mul(B, A);
22         A = mat_mul(A, A);
23         n >>= 1;
24     }
25     return B;
26 }

```

```

1 // prime number
2 bool is_prime(int n) {
3     for (int i = 2; i * i <= n; ++i)
4         if (n % i == 0) return false;
5     return n != 1;
6 }
7 vector<int> divisor(int n) {
8     vector<int> res;
9     for (int i = 1; i * i <= n; ++i) {
10         if (n % i == 0) {
11             res.push_back(i);
12             if (i != n / i) res.push_back(n / i);
13         }
14     }
15     return res;
16 }

```

```
16 }
17 map<int, int> prime_factor(int n) {
18     map<int, int> res;
19     for (int i = 2; i * i <= n; ++i) {
20         while (n % i == 0) {
21             ++res[i];
22             n /= i;
23         }
24     }
25     if (n != 1) res[n] = 1;
26     return res;
27 }
28 int prime[MAX_N];
29 bool isPrime[MAX_N + 1];
30 int seive(int n) {
31     int p = 0;
32     fill(isPrime, isPrime + n + 1, true);
33     isPrime[0] = isPrime[1] = false;
34     for (int i = 2; i <= n; ++i)
35         if (isPrime[i]) {
36             prime[p++] = i;
37             for (int j = 2 * i; j <= n; j += i) isPrime[j] = false;
38         }
39     return p;
40 }
41 // the number of prime in [L, r)
42 // 对区间 [l, r) 内的整数执行筛法, prime[i - l] = true <=> i 是素数
43 bool segPrimeSmall[MAX_L];
44 bool segPrime[MAX_SQRT_R];
45 void segment_sieve(ll l, ll r) {
46     for (int i = 0; (ll)i * i < r; ++i) segPrimeSmall[i] = true;
47     for (int i = 0; i < r - l; ++i) segPrime[i] = true;
48     for (int i = 2; (ll)i * i < r; ++i) {
49         if (segPrimeSmall[i]) {
50             for (int j = 2 * i; (ll)j * j <= r; j += i) segPrimeSmall[j]
                    = false;
51             for (ll j = max(2ll, (l + i - 1) / i) * i; j < r; j += i)
                    segPrime[j - l] = false;
52         }
53     }
54 }
55 // Miller_Rabin
56 bool check(ll a, ll n, ll x, ll t) {
```

```
57     ll res = mod_pow(a, x, n);
58     ll last = res;
59     for (int i = 1; i <= t; ++i) {
60         res = mod_mult(res, res, n);
61         if (res == 1 && last != 1 && last != n - 1) return true;
62         last = res;
63     }
64     if (res != 1) return true;
65     return false;
66 }
67 bool Miller_Rabin(ll n) {
68     if (n < MAX_N) return isPrime[n]; // small number may get wrong
69     // answer?!
70     if (n < 2) return false;
71     if (n == 2) return true;
72     if ((n & 1) == 0) return false;
73     ll x = n - 1, t = 0;
74     while ((x & 1) == 0) {
75         x >>= 1;
76         ++t;
77     }
78     for (int i = 0; i < S; ++i) {
79         ll a = rand() % (n - 1) + 1;
80         if (check(a, n, x, t))
81             return false;
82     }
83     return true;
84 }
85 // find factors
86 vector<ll> factor;
87 ll Pollard_rho(ll x, ll c) {
88     ll i = 1, k = 2;
89     ll x0 = rand() % x;
90     ll y = x0;
91     while (true) {
92         ++i;
93         x0 = (mod_mult(x0, x0, x) + c) % x;
94         ll d;
95         if (y == x0) d = 1;
96         else
97             if (y > x0)
98                 d = gcd(y - x0, x);
99             else
100                 d = gcd(x0 - y, x);
```

```

99         if (d != 1 && d != x) return d;
100        if (y == x0) return x;
101        if (i == k) {
102            y = x0;
103            k += k;
104        }
105    }
106 }
107 void find_factor(ll n) {
108     if (n == 1) return ;
109     if (Miller_Rabin(n)) {
110         factor.push_back(n);
111         return ;
112     }
113     ll p = n;
114     while (p >= n) p = Pollard_rho(p, rand() % (n - 1) + 1);
115     find_factor(p);
116     find_factor(n / p);
117 }

```

```

1  #include<bits/stdc++.h>
2  //Meisell-Lehmer
3  const int MAX_N = 5e6 + 2;
4  bool np[MAX_N];
5  int prime[MAX_N], pi[MAX_N];
6  int getprime()
7  {
8      int cnt = 0;
9      np[0] = np[1] = true;
10     pi[0] = pi[1] = 0;
11     for(int i = 2; i < MAX_N; ++i)
12     {
13         if(!np[i]) prime[++cnt] = i;
14         pi[i] = cnt;
15         for(int j = 1; j <= cnt && i * prime[j] < MAX_N; ++j)
16         {
17             np[i * prime[j]] = true;
18             if(i % prime[j] == 0) break;
19         }
20     }
21     return cnt;
22 }
23 const int M = 7;

```

```

24 const int PM = 2 * 3 * 5 * 7 * 11 * 13 * 17;
25 int phi[PM + 1][M + 1], sz[M + 1];
26 void init() {
27     getprime();
28     sz[0] = 1;
29     for(int i = 0; i <= PM; ++i) phi[i][0] = i;
30     for(int i = 1; i <= M; ++i) {
31         sz[i] = prime[i] * sz[i - 1];
32         for(int j = 1; j <= PM; ++j) phi[j][i] = phi[j][i - 1] - phi[j] /
            prime[i][i - 1];
33     }
34 }
35 int sqrt2(ll x) {
36     ll r = (ll)sqrt(x - 0.1);
37     while(r * r <= x) ++r;
38     return int(r - 1);
39 }
40 int sqrt3(ll x) {
41     ll r = (ll)cbrt(x - 0.1);
42     while(r * r * r <= x) ++r;
43     return int(r - 1);
44 }
45 ll getphi(ll x, int s)
46 {
47     if(s == 0) return x;
48     if(s <= M) return phi[x % sz[s]][s] + (x / sz[s]) * phi[sz[s]][s];
49     if(x <= prime[s]*prime[s]) return pi[x] - s + 1;
50     if(x <= prime[s]*prime[s]*prime[s] && x < MAX_N) {
51         int s2x = pi[sqrt2(x)];
52         ll ans = pi[x] - (s2x + s - 2) * (s2x - s + 1) / 2;
53         for(int i = s + 1; i <= s2x; ++i) ans += pi[x / prime[i]];
54         return ans;
55     }
56     return getphi(x, s - 1) - getphi(x / prime[s], s - 1);
57 }
58 ll getpi(ll x) {
59     if(x < MAX_N) return pi[x];
60     ll ans = getphi(x, pi[sqrt3(x)]) + pi[sqrt3(x)] - 1;
61     for(int i = pi[sqrt3(x)] + 1, ed = pi[sqrt2(x)]; i <= ed; ++i) ans
        -= getpi(x / prime[i]) - i + 1;
62     return ans;
63 }
64 ll lehmer_pi(ll x) {

```

```

65     if(x < MAX_N)    return pi[x];
66     int a = (int)lehmer_pi(sqrt2(sqrt2(x)));
67     int b = (int)lehmer_pi(sqrt2(x));
68     int c = (int)lehmer_pi(sqrt3(x));
69     ll sum = getphi(x, a) +(ll)(b + a - 2) * (b - a + 1) / 2;
70     for (int i = a + 1; i <= b; i++) {
71         ll w = x / prime[i];
72         sum -= lehmer_pi(w);
73         if (i > c) continue;
74         ll lim = lehmer_pi(sqrt2(w));
75         for (int j = i; j <= lim; j++) sum -= lehmer_pi(w / prime[j]) -
              (j - 1);
76     }
77     return sum;
78 }
79 int main() {
80     init();
81     ll n;
82     while(~scanf("%lld",&n))
83     {
84         printf("%lld\n",lehmer_pi(n));
85     }
86     return 0;
87 }

```

```

1 // 欧拉函数
2 int euler_phi(int n) {
3     int res = n;
4     for (int i = 2; i * i <= n; ++i) {
5         if (n % i == 0) {
6             res = res / i * (i - 1);
7             for (; n % i == 0; n /= i);
8         }
9     }
10    if (n != 1) res = res / n * (n - 1);
11    return res;
12 }
13 int euler[MAX_N];
14 void euler_phi_sieve() {
15     for (int i = 0; i < MAX_N; ++i) euler[i] = i;
16     for (int i = 2; i < MAX_N; ++i)
17         if (euler[i] == i)
18             for (int j = i; j < MAX_N; j += i) euler[j] = euler[j] / i *

```



```

19     (i - 1);
    }

```

- Moebius 如果

$$F(n) = \sum_{d|n} f(d)$$

, 则

$$f(n) = \sum_{d|n} \mu(d) F\left(\frac{n}{d}\right)$$

对于 $\mu(d)$ 函数, 有如下性质:

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & n=1 \\ 0 & n>1 \end{cases}$$

$$\sum_{d|n} \frac{\mu(d)}{d} = \frac{\phi(n)}{n}$$

```

1  int mu[MAX_N];
2  void moebius() {
3      int cnt = 0; mu[1] = 1;
4      memset(vis, 0, sizeof vis);
5      for (int i = 2; i < MAX_N; ++i) {
6          if (!vis[i]) {
7              prime[cnt++] = i;
8              mu[i] = -1;
9          }
10         for (int j = 0; j < cnt && i * prime[j] < MAX_N; ++j) {
11             vis[i * prime[j]] = true;
12             if (i % prime[j])
13                 mu[i * prime[j]] = -mu[i];
14             else
15                 mu[i * prime[j]] = 0, break;
16         }
17     }
18 }

19
20 map<int, int> moebius(int n) {
21     map<int, int> res;
22     vector<int> primes;
23     for (int i = 2; i * i <= n; ++i) {
24         if (n % i == 0) {

```

```

25         primes.push_back(i);
26         while (n % i == 0) n /= i;
27     }
28 }
29 if (n != 1) primes.push_back(n);
30
31 int m = primes.size();
32 for (int i = 0; i < (1 << m); ++i) {
33     int mu = 1, d = 1;
34     for (int j = 0; j < m; ++j) {
35         if (i >> j & 1) {
36             mu *= -1;
37             d *= primes[j];
38         }
39     }
40     res[d] = mu;
41 }
42 return res;
43 }

```

```

1 // Gauss_jordan
2 const double eps = 1e-8;
3 typedef vector<double> vec;
4 typedef vector<vec> mat;
5
6 vec gauss_joedan(const mat &A, const vec& b) {
7     int n = A.size();
8     mat B(n, vec(n + 1));
9     for (int i = 0; i < n; ++i)
10         for (int j = 0; j < n; ++j) B[i][j] = A[i][j];
11     for (int i = 0; i < n; ++i) B[i][n] = b[i];
12
13     for (int i = 0; i < n; ++i) {
14         int pivot = i;
15         for (int j = i; j < n; ++j)
16             if (abs(B[j][i]) > abs(B[pivot][i])) pivot = j;
17         if (i != pivot) swap(B[i], B[pivot]);
18
19         if (abs(B[i][i]) < eps) return vec();
20
21         for (int j = i + 1; j <= n; ++j) B[i][j] /= B[i][i];
22         for (int j = 0; j < n; ++j) if (i != j)
23             for (int k = i + 1; k <= n; ++k) B[j][k] -= B[j][i] * B[i][k];

```

```

24         ];
25     }
26     vec x(n);
27     for (int i = 0; i < n; ++i) x[i] = B[i][n];
28     return x;
29 }
30
31 vec gauss_joedan_xor(const mat& A, const vec& b) {
32     int n = A.size();
33     mat B(n, vec(n + 1));
34     for (int i = 0; i < n; ++i)
35         for (int j = 0; j < n; ++j) B[i][j] = A[i][j];
36     for (int i = 0; i < n; ++i) B[i][n] = b[i];
37
38     for (int i = 0; i < n; ++i) {
39         int pivot = i;
40         for (int j = i; j < n; ++j)
41             if (B[j][i]) {
42                 pivot = j;
43                 break;
44             }
45         if (pivot != i) swap(B[i], B[pivot]);
46
47         for (int j = 0; j < n; ++j) if (i != j && B[j][i])
48             for (int k = i + 1; k <= n; ++k) B[j][k] ^= B[i][k];
49     }
50 }
51
52 vec x(n);
53 for (int i = 0; i < n; ++i) x[i] = B[i][n];
54 return x;
55 }

```

7 String

1. Hash
2. KMP
3. Extend KMP

4. trie 树 poj2001 2503 3630 1056 hdu 1075 1251 1247 1298 1671
5. Manacher 算法
6. AC 自动机
7. 后缀数组
8. 后缀树
9. 后缀自动机
10. 回文自动机

```
1 // 最小最大表示法:
2 int getMinString(const string &s) {
3     int len = (int)s.length();
4     int i = 0, j = 1, k = 0;
5     while(i < len && j < len && k < len) {
6         int t = s[(i + k) % len] - s[(j + k) % len];
7         if(t == 0) k++;
8         else {
9             if(t > 0) i += k + 1; // getMaxString: t < 0
10            else j += k + 1;
11            if(i == j) j++;
12            k = 0;
13        }
14    }
15    return min(i, j);
16 }
```

```
1 // KMP
2 int nxt[MAX_N];
3 void getNext(const string &str) {
4     int len = str.length();
5     int j = 0, k;
6     k = nxt[0] = -1;
7     while (j < len) {
8         if (k == -1 || str[j] == str[k])
9             nxt[++j] = ++k;
10        else k = nxt[k];
11    }
12 }
13 int kmp(const string &tar, const string &pat) {
14     getNext(pat);
15     int num, j, k;
16     int lenT = tar.length(), lenP = pat.length();
17     num = j = k = 0;
```

```

18     while (j < lenT) {
19         if(k == -1 || tar[j] == pat[k])
20             j++, k++;
21         else k = nxt[k];
22         if(k == lenP) {
23             // res = max(res, j - lenP);
24             k = nxt[k];
25             ++num;
26         }
27     }
28     return num;//lenP - res - 1;
29 }

```

```

1 // Suffix Array & LCP Array
2 int n, k;
3 int lcp[MAX_N], sa[MAX_N];
4 int rnk[MAX_N], tmp[MAX_N];
5
6 bool compare_sa(int i, int j) {
7     if (rnk[i] != rnk[j]) return rnk[i] < rnk[j];
8     else {
9         int ri = i + k <= n? rnk[i + k] : -1;
10        int rj = j + k <= n? rnk[j + k] : -1;
11        return ri < rj;
12    }
13 }
14 void construct_sa(string S, int *sa) {
15     n = S.length();
16     for (int i = 0; i <= n; i++) {
17         sa[i] = i;
18         rnk[i] = i < n? S[i] : -1;
19     }
20     for (k = 1; k <= n; k *= 2) {
21         sort(sa, sa + n + 1, compare_sa);
22         tmp[sa[0]] = 0;
23         for (int i = 1; i <= n; i++)
24             tmp[sa[i]] = tmp[sa[i - 1]] + (compare_sa(sa[i - 1], sa[i])
25                 ? 1 : 0);
26         memcpy(rnk, tmp, sizeof(int) * (n + 1));
27     }
28 }
29 void construct_lcp(string S, int *sa, int *lcp) {
30     int n = S.length();

```

```
30     for (int i = 0; i <= n; i++) rnk[sa[i]] = i;
31     int h = 0;
32     lcp[0] = 0;
33     for (int i = 0; i < n; i++) {
34         int j = sa[rnk[i] - 1];
35         if (h > 0) h--;
36         for (; j + h < n && i + h < n; h++)
37             if (S[j + h] != S[i + h]) break;
38         lcp[rnk[i] - 1] = h;
39     }
40 }
```

```
1 // AC 自动机
2 int ans[MAX_N], d[MAX_N];
3
4 struct Trie {
5     int nxt[MAX_N][26], fail[MAX_N], end[MAX_N];
6     int root, L;
7     int newnode() {
8         for(int i = 0; i < 26; i++)
9             nxt[L][i] = -1;
10        end[L++] = 0;
11        return L-1;
12    }
13    void init() {
14        L = 0;
15        root = newnode();
16    }
17    void insert(char buf[]) {
18        int len = strlen(buf);
19        int now = root;
20        for(int i = 0; i < len; i++) {
21            if(nxt[now][buf[i]-'a'] == -1)
22                nxt[now][buf[i]-'a'] = newnode();
23            now = nxt[now][buf[i]-'a'];
24        }
25        end[now] = 1;
26        d[now] = len;
27    }
28    void build() {
29        queue<int> Q;
30        fail[root] = root;
31        for(int i = 0; i < 26; i++)
```

```
32         if(nxt[root][i] == -1)
33             nxt[root][i] = root;
34         else {
35             fail[nxt[root][i]] = root;
36             Q.push(nxt[root][i]);
37         }
38     while( !Q.empty() ) {
39         int now = Q.front(); Q.pop();
40         for(int i = 0; i < 26; i++)
41             if(nxt[now][i] == -1)
42                 nxt[now][i] = nxt[fail[now]][i];
43             else {
44                 fail[nxt[now][i]] = nxt[fail[now]][i];
45                 Q.push(nxt[now][i]);
46             }
47     }
48 }
49 void solve(char buf[]) {
50     int cur = root;
51     int len = strlen(buf);
52     int index;
53     for(int i = 0; i < len; ++i) {
54         if(buf[i] >= 'A' && buf[i] <= 'Z')
55             index = buf[i] - 'A';
56         else if(buf[i] >= 'a' && buf[i] <= 'z')
57             index = buf[i] - 'a';
58         else continue;
59         cur = nxt[cur][index];
60         int x = cur;
61         while(x != root) {
62             if(end[x]) {
63                 ans[i + 1] -= 1;
64                 ans[i - d[x] + 1] += 1;
65                 break;
66             }
67             x = fail[x];
68         }
69     }
70 }
71 };
72
73 Trie ac;
```

8 Others

8.1 Divide-and-Conquer Tree

```
1 //uva 12161
2 struct edge {
3     int to, damage, length, next;
4 };
5 int G[MAX_N], En, N, M, T;
6 edge E[MAX_N * 2];
7
8 void add_edge(int from, int to, int damage, int length) {
9     edge e = {to, damage, length, G[from]};
10    E[En] = e;
11    G[from] = En++;
12 }
13
14 int ans, subtree_size[MAX_N];
15 bool flag[MAX_N];
16
17 int s, t;
18 Pii ds[MAX_N];
19
20 int compute_subtree_size(int v, int p) {
21     int c = 1;
22     for (int j = G[v]; ~j; j = E[j].next) {
23         int w = E[j].to;
24         if (w == p || flag[w]) continue;
25         c += compute_subtree_size(w, v);
26     }
27     return subtree_size[v] = c;
28 }
29
30 Pii search_centroid(int v, int p, int t) {
31     Pii res = Pii(INT_MAX, -1);
32     int s = 1, m = 0;
33     for (int j = G[v]; ~j; j = E[j].next) {
34         int w = E[j].to;
35         if (w == p || flag[w]) continue;
```



```

36     res = min(res, search_centroid(w, v, t));
37     m = max(subtree_size[w], m);
38     s += subtree_size[w];
39 }
40 m = max(m, t - s);
41 res = min(res, Pii(m, v));
42 return res;
43 }
44
45 void enumrate_path(int v, int p, int damage, int length) {
46     ds[t++] = Pii(damage, length);
47     for (int j = G[v]; ~j; j = E[j].next) {
48         int w = E[j].to;
49         if (w == p || flag[w]) continue;
50         if (damage + E[j].damage <= M) {
51             enumrate_path(w, v, damage + E[j].damage, length + E[j].
52                 length);
53         }
54     }
55 }
56
57 void remove_useless(int s, int &t) {
58     if (s == t) return;
59     int tt;
60     for (int i = tt = s + 1; i < t; i++) {
61         if (ds[i].first == ds[tt - 1].first) continue;
62         if (ds[i].second <= ds[tt - 1].second) continue;
63         ds[tt++] = ds[i];
64     }
65     t = tt;
66 }
67
68 void solve_sub_problem(int v) {
69     compute_subtree_size(v, -1);
70     int c = search_centroid(v, -1, subtree_size[v]).second;
71     flag[c] = true;
72     for (int j = G[c]; ~j; j = E[j].next) {
73         if (flag[E[j].to]) continue;
74         solve_sub_problem(E[j].to);
75     }
76
77     s = t = 0;
78     for (int j = G[c]; ~j; j = E[j].next) {

```

```
78     int w = E[j].to;
79     if (flag[w]) continue;
80     if (E[j].damage <= M)
81         enumerate_path(w, v, E[j].damage, E[j].length);
82     if (s > 0) {
83         sort(ds + s, ds + t);
84         remove_useless(s, t);
85         for (int l = 0, r = t - 1; l < s && r >= s; l++) {
86             while (r >= s && ds[l].first + ds[r].first > M) r--;
87             if (r >= s)
88                 ans = max(ans, ds[l].second + ds[r].second);
89         }
90     }
91     sort(ds, ds + t);
92     remove_useless(0, t);
93     s = t;
94 }
95
96 flag[c] = false;
97 }
```