

5. Single source shortest path - Dijkstra's Algorithm(a) Pseudocodeprocedure dijkstra (G, l, s)

Input: graph $G = (V, E)$, directed or undirected;
 positive edge lengths $\{l_e; e \in E\}$; vertex $s \in V$.

Output: For all vertices u reachable from s , $\text{dist}(u)$ is set to the distance from s to u .

for all $u \in V$: $\text{dist}(u) = \infty$ $\text{prev}(u) = \text{nil}$ $\text{dist}(s) = 0$ $H = \text{makequeue}(V)$ (using dist-values as keys)while H is not empty: $u = \text{deletemin}(H)$ for all edges $(u, v) \in E$:if $\text{dist}(v) > \text{dist}(u) + l(u, v)$: $\text{dist}(v) = \text{dist}(u) + l(u, v)$

→ Relaxation

 $\text{prev}(v) = u$ $\text{decreasekey}(H, v)$

(b) Compute the shortest dist from vertex 1 to all other vertices using Dijkstra's algorithm.

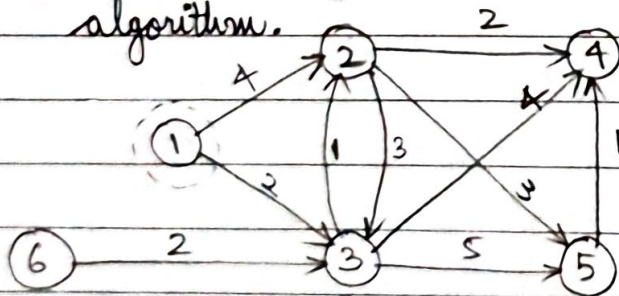
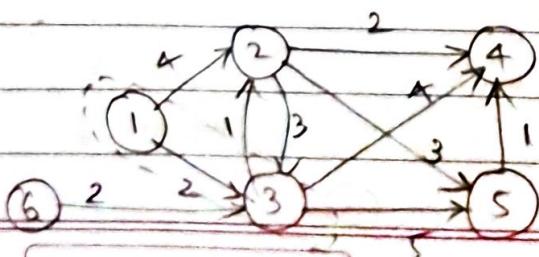
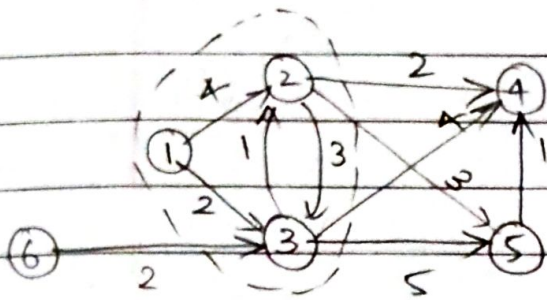


table with node:dist values

1: 0	4: ∞
2: 4	5: ∞
3: 2	6: ∞



1: 0	4: $\infty > 2 + 4$ 6
2: 3 ($4 > 2 + 1$)	5: $\infty > 2 + 5$ 7
3: 2	6: ∞



1: 0

2: 3

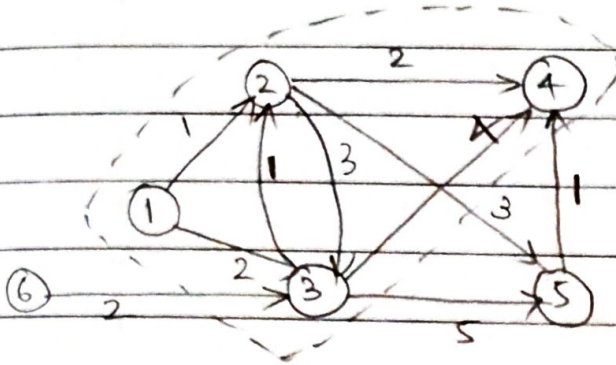
3: 2

$6 > 3+2$

4: 5

5: $7 > 3+3$
6

6: ∞



1: 0

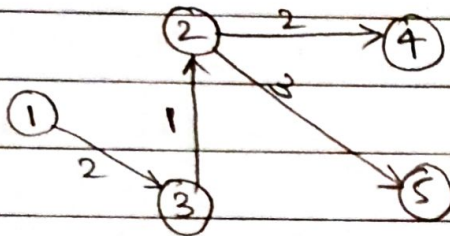
2: 3

3: 2

4: 5

5: 6

6: ∞



shortest distances from 1 to the other vertices:

1: 0

2: 3

3: 2

4: 5

5: 6

6: ∞ (no path)