

25.03.2020

### Kruskal's Algorithm

- The algorithm's state is modelled as a collection of disjoint sets, each of which contains the nodes of a particular component. ~~Initially~~
- Initially, each node is <sup>an</sup> a component by itself.

$\text{makeset}(x)$ : create a singleton set containing just  $x$ .

- We repeatedly test pairs of nodes to see if they belong to the same set.

$\text{find}(x)$ : to which set does  $x$  belong?

- Whenever we add <sup>an</sup> merging edge, we are merging 2 components.

$\text{union}(x, y)$ : merge the sets containing  $x$  &  $y$ .

- The algo uses  $|V|$  makeset,  $2|E|$  find, and  $|V|-1$  union operations.

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## Algorithm

procedure kruskal ( $G, w$ )

Input: A connected undirected graph  $G = (V, E)$  with edge weights  $w_e$

Output: A min. spanning tree defined by the edges  $x$ .

for all  $u \in V$ :

$\text{makeset}(u)$

$x = \emptyset$

Sort the edges  $E$  in non-decreasing order of weight

for all edges  $\{u, v\} \in E$ , in increasing order of weight:

    if  $\text{find}(u) \neq \text{find}(v)$ :

        add edge  $\{u, v\}$  to  $x$

$\text{union}(u, v)$

procedure makeset( $x$ )

$\pi(x) = x$

$\text{rank}(x) = 0$

function find( $x$ )

while  $x \neq \pi(x)$ :  $x = \pi(x)$

return  $x$

procedure union( $x, y$ )

$r_x = \text{find}(x)$

$r_y = \text{find}(y)$

if  $r_x = r_y$ : return

if  $\text{rank}(r_x) > \text{rank}(r_y)$

$\pi(r_y) = r_x$

else:

$\pi(r_x) = r_y$

    if  $\text{rank}(x) = \text{rank}(y)$ :  $\text{rank}(y) = \text{rank}(y) + 1$

One way to store a set is as a directed tree.  
Nodes of the tree are elems of the set, arranged in no particular order, and each has parent pointers that eventually lead up to the root of the tree.

Root elem: convenient representative / name of the set.  
Its parent ptr: self loop

Parent ptr:  $\pi$

rank  $r$ : height of the subtree hanging from that of a node node.

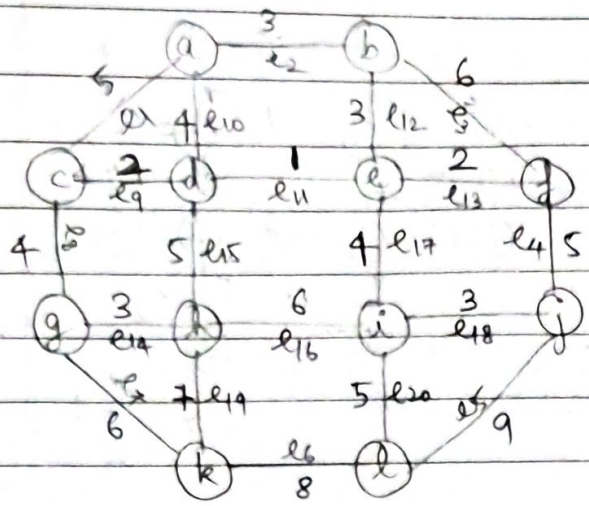
Merging: make the root of the shorter tree point to the root of the taller tree.

Instead of explicitly computing hts of trees, we will use the rank nos. of their root nodes  $\rightarrow$  union by rank

Use the algorithm to find min. cost spanning tree for the graphs given below:



(b)



a b c d e f g h i j k l

$\pi = [a|b|c|d|e|f|g|h|i|j|k|l]$

rank:  $[0|0|0|0|0|0|0|0|0|0|0|0]$

Ein non-decreasing order of wts

$e_{11}, e_9, e_{13}, e_2, e_{12}, e_{14}, e_{18}, e_8, e_{10}, e_{17},$   
 1 2 2 3 3 3 3 4 4 4

$e_1, e_4, e_{15}, e_{20}, e_3, e_7, e_{16}, e_{19}, e_6, e_5$   
 5 5 5 5 6 6 6 7 8 9

①  $e_1 (d,e)$

$\text{find}(d) = d$

$\text{find}(e) = e$

add(d,e) to x

$\text{union}(d,e) \Rightarrow$

$\hookrightarrow r_x = d$

$r_y = e$

$\pi(r_x) = r_y$

$x = \{d, e\}$

a b c d e f g h i j k l

$[a|b|c|d|e|f|g|h|i|j|k|l]$

$[0|0|0|0|1|0|0|0|0|0|0|0]$



$\text{rank}(d) = \text{rank}(e) \Rightarrow \text{rank}(e)++$

②  $e_9 (c,d)$

$x = \{c, d, e\}$

$\text{union}(c,d) : r_x = c$

$r_y = e$



$\text{rank}(r_x) = 0$   
 $\text{rank}(r_y) = 1$   
 $\Rightarrow \text{rank}(c) = e$   
 $\text{rank}(e)++$

a b c d e f g h i j k l

$[0|0|0|0|2|0|0|0|0|0|0|0]$

③  $e_{13} (e,f)$

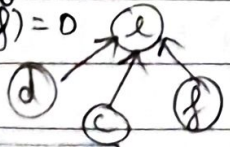
$x = \{c, d, e, f\}$

$\text{union}(e,f) : r_x = e$

$r_y = f$

$\text{rank}(e) = 2, \text{rank}(f) = 0$

$\therefore \pi(f) = e$



a b c d e f g h i j k l

$[0|0|0|0|2|0|0|0|0|0|0|0]$

④  $e_2 (a,b)$

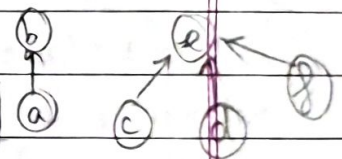
$x = \{a, b, c, d, e, f\}$

$\text{union}(a,b) : r_x = a, r_y = b. \text{rank}(a) = 0, \text{rank}(b) = 0$

$\therefore \pi(b) = a; \text{rank}(b)++$

b c d e f g h i j k l

$[0|1|0|0|2|0|0|0|0|0|0|0]$



⑤  $e_{12} (b,e)$

$\text{find}(b) = b$

$x = \{a, b, c, d, e, f\}$

$\text{find}(e) = e$

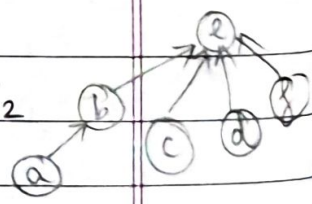
$\text{union}(b,e)$

$r_x = b, r_y = e. \text{rank}(b) = 1, \text{rank}(e) = 2$

$\therefore \pi(b) = e; \text{rank}(e)++$

b c d e f g h i j k l

$[0|1|0|0|3|0|0|0|0|0|0|0]$



⑥  $e_{14} (g,h)$

$\text{find}(g) = g$

$x = \{a, b, c, d, e, f, g, h\}$

$\text{find}(h) = h$

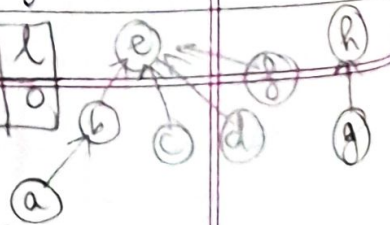
$\text{union}(g,h)$

$r_x = g, r_y = h$

$\pi(g) = h. \text{rank}(h)++$

b c d e f g h i j k l

$[0|1|0|0|3|0|0|1|0|0|0|0]$



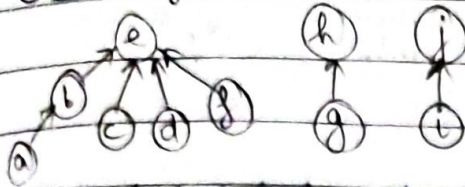
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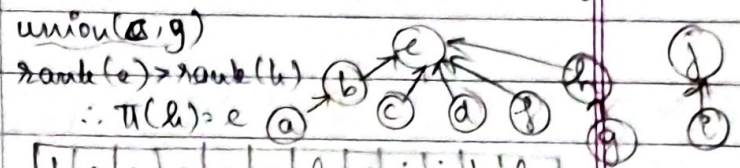


⑦  $e_8(i, j)$



b	e	e	e	e	e	h	h	j	j	k	l
0	1	0	0	3	0	0	1	0	1	0	0

⑧  $e_3(c, g)$   $\text{find}(c) = e \leftarrow rx$   
 $\text{find}(g) = h \leftarrow ry$   
 $x = \{a, b, c, d, e, f, g, h, i, j\}$



b	e	e	e	e	e	h	e	j	j	k	l
0	1	0	0	3	0	0	1	0	1	0	0

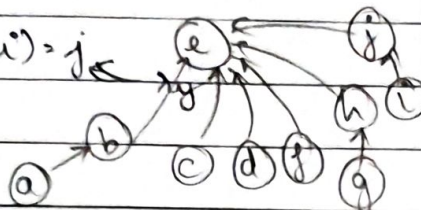
⑨  $e_{10}(a, d)$

$\text{find}(a) = e$ ,  $\text{find}(d) = e$ .

⑩  $e_7(e, i)$

$\text{find}(e) = e$ ,  $\text{find}(i) = j$

$\text{rank}(e) > \text{rank}(j)$



⑪  $e_1(a, c)$

$\text{find}(a) = e$ ,  $\text{find}(c) = e$ .

⑫  $e_4(f, j)$

$\text{find}(f) = e$ ,  $\text{find}(j) = e$ .

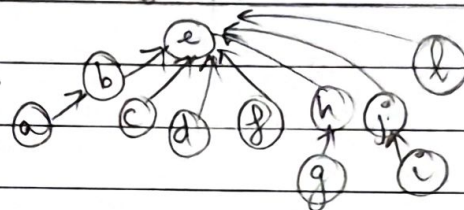
⑬  $e_{15}(d, h)$

$\text{find}(d) = e$ ,  $\text{find}(h) = e$ .

⑭  $e_{20}(i, l)$

$\text{find}(i) = e$ ,  $\text{find}(l) = l$

$x = \{a, b, c, d, e, f, g, h, i, j, l\}$   
 $\text{rank}(e) > \text{rank}(l)$

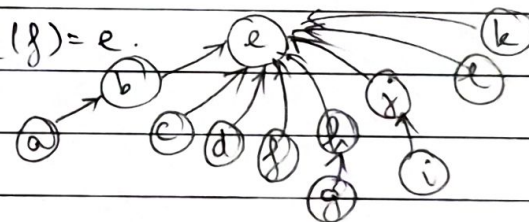


a	b	c	d	e	f	g	h	i	j	k	l
b	e	e	e	e	e	e	h	e	j	e	e
0	1	0	0	3	0	0	1	0	1	0	1

⑮  $e_3(b, f) \rightarrow \text{find}(b) = \text{find}(f) = e$ .

⑯  $e_7(g, k)$

$\text{find}(g) = e$ ,  $\text{find}(k) = k$ .



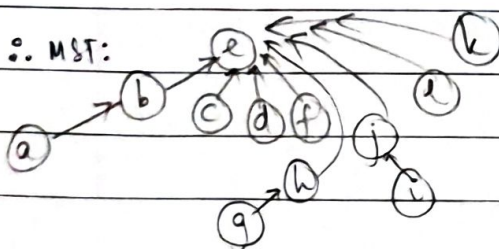
b	e	e	e	e	e	e	h	e	j	e	e
0	1	0	0	3	0	0	1	0	1	0	1

⑰  $e_{16}(h, i)$

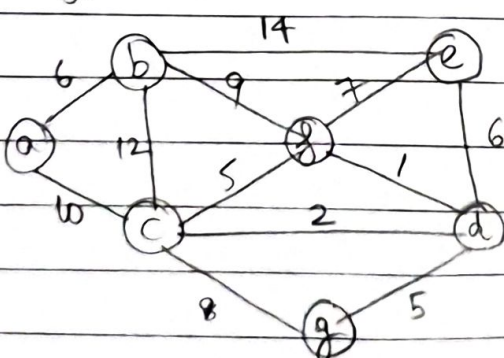
⑱  $e_9(h, k)$

⑲  $e_6(k, l)$

⑳  $(f, l) e_9$



⑵



$\pi$ :	a	b	c	d	e	f	g
rank:	0	0	0	0	0	0	0

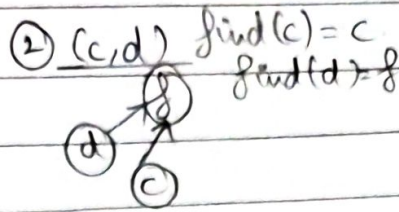
E in non-decreasing order of wts:

(f, d)	(c, d)	(c, f)	(d, g)	(a, b)	(d, e)
1	2	5	5	6	6
(e, f)	(c, g)	(b, f)	(a, c)	(b, d)	(b, e)
7	8	9	10	12	14

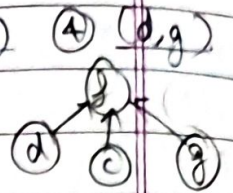
① ~~(f,d)~~ (d,f)



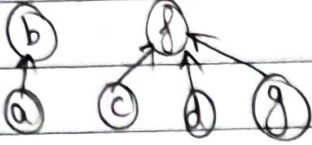
a	b	c	f	e	g
0	0	0	0	0	1



③ (c,f) creates cycle



⑤ (a,b)



⑥ (d,e)



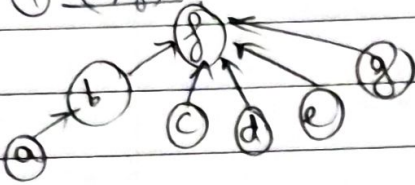
⑦ (e,f)

already present

⑧ (c,g)

will add cycle

⑨ (b,f)



∴ MST:

