| 3.04.2010 | |
|---|-------|
| 15. Silyle Source Shortest Poth & Bellman-Ford Algorithm. | |
| - Algorithm Bellman-Ford (G, w, 8) | |
| Input: Directed graph G= (V, E); | |
| Edge weights è we : e E E 3 | |
| vertex s E V. | |
| Output: For all vertices v reachable from 8, vod is set to the | |
| distance from sto v. | |
| The state of the state of the property of the second of the second of | |
| # Initialize-Stugle-Source: | |
| for each vertex v E & · V: | |
| v.d = 00 | |
| V.TZ NIL | |
| 8.d=0 | |
| | |
| for i=1 to 1G.VI-1: (repeat WI-1 time) | |
| for each edge (u,v) € G.E: | |
| RELAX (u,v,w) | |
| for each edge (u,v) & G.E: | |
| if v.d > u.d + w(u,v) | |
| return FALSE | |
| return TRUE | |
| | 1 |
| Procedure RELAX (u, v, w) | 1 |
| if v.d > u.d + w(u,v): | |
| $v \cdot d = u \cdot d + w(u_1 v)$ | |
| v. A = u | |
| | |
| . The algorithm solves the single source shortest path problem of a | |
| directed graph G= (V, E) in which the edge weights may be negation | ve. |
| · Moreover, this algorithm can be applied to find the shortest if there does not exist any negative weighted eycle. | sath, |
| If there does not exist any negative weighted eyels | |