# **Assignment 3**

## COS30023 - Languages in Software Development

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August 31, 2014

- 1. Problem 1
- 1.1. String specification

$$S ::= \begin{array}{cc} \epsilon \\ & aS \end{array}$$

1.2. String Length Specification

$$length(\epsilon) = 0$$
  
$$length(aS) = 1 + length(S)$$

1.3. String Concatenation Specification

$$s_1 = \epsilon$$
:  $\epsilon \oplus s_2 = s_2$   
 $s_1 = as'_1$ :  $as'_1 \oplus s_2 = a(s'_1 \oplus s_2)$ 

## 1.4. a)

Show that if  $s \in S$ , then  $s \oplus \epsilon = s$ 

Base Case 1

$$s \equiv \epsilon$$

$$\epsilon \oplus \epsilon = \epsilon$$

The first base case is when s is structurally equivalent to the empty string  $\epsilon$  and it's a fact that concatenating two empty strings will yield another empty string.

#### Base Case 2

if length(s) = 1 then

$$s \oplus \epsilon = s$$

In this second base case we take s to be any string of length 1 and again it is a fact that any string concatenated with the empty string  $\epsilon$  will be the original first string. In this case the string will still be of length 1.

#### **Inductive Step**

We assume that the base case holds for all lengths of s greater than 1, therefore:

if length(s) > 1

$$s \oplus \epsilon = s$$

Q.E.D.

## 1.5. b)

Show that if  $s_1, s_2 \in S$  then  $length(s_1 \oplus s_2) = length(s_1) + length(s_2)$ 

Base Case 1

$$s_1, s_2 \equiv \epsilon$$

Given that  $length(\epsilon) = 0$  then it goes to say that  $length(\epsilon \oplus \epsilon) = 0$ .

#### Base Case 2

 $s_1 \neq \epsilon$ , where  $length(s_1) = 1$  and  $s_2 \equiv \epsilon$ .

$$length(s_1 \oplus \epsilon) = length(s_1)$$

For the case where  $s_1$  is a string of length 1 and  $s_2$  is an empty string, the concatenated length will be the length of  $s_1$ .

#### Base Case 3

```
s_1, s_2 \neq \epsilon, where length(s_1) = 1 and length(s_2) = 1.
```

$$length(s_1 \oplus s_2) = length(s_1) + length(s_2) = 2$$

For the case where both  $s_1$  and  $s_2$  are strings of length 1, the concatenated length will be the sum of the individual lengths, which is 2.

### **Inductive Step**

 $length(s_1) > 1$  and  $length(s_2) > 1$ 

$$length(s_1 \oplus s_2) = length(s_1) + length(s_2)$$

We assume the base case 3 holds for strings  $s_1, s_2$  of any length greater than 1.

Q.E.D.

## 2. Problem 2

```
% List length predicate
list_length([], 0).
list_length([_|T], N) :- list_length(T,M), N is M+1.
% List concatenation predicate
list_concatenation([],Xs,Xs).
list\_concatenation([H|Xs],Ys,[H|Zs]) :- list\_concatenation(Xs,Ys,Zs).
% List equality predicate
equal([],[]) :- true.
equal([Xhead|Xtail],[Yhead|Ytail]) :- Xhead =:= Yhead, equal(Xtail, Ytail).
% Check that concatenating the empty set to Xs equals Xs
checkA(Xs) :- list_concatenation([],Xs,R), equal(R,Xs).
% Check that concatenating two strings equals the length of the sum
% of individual lengths of the substrings
checkB(Xs, Ys) :-
                        list_concatenation(Xs,Ys,Rs),
                                        list_length(Rs, R1),
                                        list_length(Xs, X1),
                                        list_length(Ys, Yl),
                                        R1 = := X1 + Y1.
```