

sob sob

① Să dem că $2^{n+1} \in O(2^n)$

Def. $f \in O(g)$ dacă $\exists c, n_0 > 0$ aî. $\forall n \geq n_0$
avem $f(n) \leq c \cdot g(n)$

$$f = 2^{n+1} \quad g = 2^n$$

Algem $c=2, n_0=1 \Rightarrow 2^{n+1} \leq 2 \cdot 2^n, \forall n \geq 1$

$$2^{2n} \in O(2^n)$$

$$\exists c \text{ aî } 2^{2n} \leq c \cdot 2^n \quad \forall n > n_0 \Leftrightarrow 2^n \leq c \quad \forall n \geq n_0$$

Fals deoarece $2n \rightarrow \infty$ și c e o constantă

Teoremă $f \in \Theta(g) \Leftrightarrow f \in O(g)$ și $f \in \Omega(g)$

Dem: " \Leftarrow "

$$f \in O(g) \Rightarrow \exists c_1 \text{ aî } f(n) \leq c_1 \cdot g(n) \quad \forall n \geq n_{01}$$

$$f \in \Omega(g) \Rightarrow \exists c_2 \text{ aî } f(n) \geq c_2 \cdot g(n) \quad \forall n \geq n_{02}$$

$$\text{Algem } n_0 = \max(n_{01}, n_{02})$$

$$\Rightarrow c_1 \cdot g(n) \geq f(n) \geq c_2 \cdot g(n), \quad \forall n \geq n_0$$

" \Rightarrow "

$$f \in \Theta(g) \Rightarrow \exists c_1, n_0 > 0 \text{ aî. } c_1 \cdot g(n) \geq f(n) \geq c_2 \cdot g(n), \quad \forall n \geq n_0$$

$$\Rightarrow 1) c_1 \cdot g(n) \geq f(n) \quad \forall n \geq n_0 \Rightarrow f \in O(g)$$

$$2) f(n) \geq c_2 \cdot g(n) \quad \forall n \geq n_0 \Rightarrow f \in \Omega(g)$$

ex 2 $f(n) + g(n) \in \Theta(\min \{f(n), g(n)\})$

nu e corect. contra ex:

$$\left. \begin{array}{l} f(n) = 1 \\ g(n) = n^2 \end{array} \right\} \Rightarrow n^2 + 1 \in \Theta(1)$$

$$c_1 \cdot 1 \geq n^2 \geq c_2 \cdot 1 \rightarrow \text{fals} \quad 1 + n^2 \notin \Theta(1)$$

$$f(n) + g(n) \in \Theta(\max \{f(n), g(n)\})$$

$$f(n) + g(n) \in \Theta(\max \{f(n), g(n)\}) \Leftrightarrow$$

$$\Leftrightarrow \exists c_1, c_2, n_0 \text{ at } \forall n \geq n_0 \Rightarrow$$

$$\Rightarrow c_1 \cdot \max(\dots) \geq f(n) + g(n) \geq c_2 \cdot \max(\dots)$$

Alegem $c_1 = 2, c_2 = 1$ și inegalitatea este adevărată
(pt $f(n) = 1, g(n) = n^2$)

ex $n! \quad o(n^n)$

$$\Theta(n^n)$$

$$n \in o(n^2)$$

$$\Leftrightarrow \exists c > 0 \quad c \cdot n^2 > n \quad \forall n \geq n_0 \Rightarrow \frac{1}{n} < c \Rightarrow n > \frac{1}{c}$$

$\exists n_0 > 0$ aș.
 $\forall n \geq n_0$

$$n_0 = \frac{1}{c} + 1$$

$$n! \in o(n^n)$$

$$\Leftrightarrow \exists c > 0 \text{ aș. } \exists n_0 > 0 \text{ aș. } \forall n \geq n_0 \Rightarrow c \cdot n^n > n!$$

$$(\nexists) \forall c > 1 \text{ dar cum dem?}$$

$$n! < c \cdot n^n \quad | : n^{n-1}$$

$$\frac{n!}{n^{n-1}} < c \cdot n \quad | : c$$

$$\frac{n!}{n^{n-1}} < n \Rightarrow n > \frac{n!}{n^{n-1}}$$

$$n_0 = \frac{1}{c} \text{ satisfacă}$$

n_0 trebuie să fie în funcție de c

ex $\log_2 n! \in \Theta(n \log_2 n)$
0 sau Θ ?

$$\Leftrightarrow \exists c_1, c_2, n_0 > 0 \text{ aș. } \forall n \geq n_0 \text{ și}$$

$$c_1 \cdot g(n) \leq f(n) \leq c_2 \cdot g(n)$$

$$\Rightarrow c_1 \cdot \log_2 n^n \leq \log_2 n! \leq c_2 \cdot \log_2 n^n \quad \forall n \geq n_0$$

$$\log_2 n! = \log_2 (1 \cdot 2 \cdot 3 \cdot \dots \cdot n) = \log_2 1 + \log_2 2 + \dots + \log_2 n =$$

$$= \log_2 1 + \log_2 2 + \dots + \log_2 \left(\frac{n}{2} - 1\right) + \log_2 \frac{n}{2} + \dots + \log_2 n \geq$$

$$\geq \log_2 \frac{n}{2} + \dots + \log_2 n \geq \frac{n}{2} \cdot \log_2 \frac{n}{2} \rightarrow$$

$$\rightarrow \frac{n}{2} \log_2 n - \frac{n}{2} \log_2 2 \geq \frac{n \log_2 n}{4}$$

ex $f(n) + o(f(n)) \in \Theta(f(n))$

$$c_1 \cdot f(n) \geq f(n) + o(f(n)) \geq c_2 \cdot f(n)$$

$$g \in o(f(n)) \text{ d.h. } \forall c > 0 \exists n_0 > 0 \text{ a. } \forall n \geq n_0,$$

$$g(n) < c \cdot f(n)$$

Alagem $c=1$ - d.h. $f(n) + g < 2f(n) \quad \forall n \geq n_0$
 $\left[\exists n_0 \text{ a. } g(n) < f(n) \quad \forall n \geq n_0 \right]$

$$\forall n \geq n_0 \quad f(n) + o(f(n)) \geq f(n)$$

$$f(n) + o(f(n)) \leq 2f(n)$$