1. B05202068

2.8 f: x5, g: x6, h: x7, i: x28, j: x29

&A[0]: x10, &B[0]: x11

```
addi  x30, x10, 8   // x30 = &A[0]+8 (i.e.,&A[1])
addi  x31, x10, 0   // x31 = &A[0]+0 (i.e., &A[0])
sd    x31, 0(x30)   // x30 = x31 (i.e., A[1] = &A[0])
ld    x30, 0(x30)   // x30 = A[1] = &A[0]
add   x5 , x30, x31 // f = &A + &A = 2*(&A).
```

2.9

	opcode, funct3, 7	rs1	rs2	rd	Imm
addi x30, x10, 8 : I-type	0x13, 0x0, -	10	-	30	8
addi x31, x10, 0 : I-type	0x13, 0x0, -	10	-	31	0
sd x31, 0(x30) : S-type,	0x3, 0x3, -	31	30	-	0
ld x30, 0(x30) : I-type	0x3, 0x3, -	30	-	30	0
add x5, x30, x31 : R-type	0x33, 0x0, 0x0	30	31	5	-

2.16

RISC-V : 32 regs. → 128 registers

Instructions increase 4 times

5 → 7 (2 regs.)

① The rs1, rs2, and rd fields : 5 bits → 7 bits. (R-type)

the opcode field: 7 bits → 9 bits.

② the rs1, rd : 5 bits → 7 bits.

imm won't change by itself, but may be shorten to suit the need of overall instruction size. (I-type.

the opcode : 7 bits → 12 bits.

③ Increasing the size of each bit field → instruction longer → may increase code size. ※

But, the increase in # of registers → less use among lots of registers
→ reduce the total # of instructions → may reduce code size.

Report on matrix multiplication:

53195612

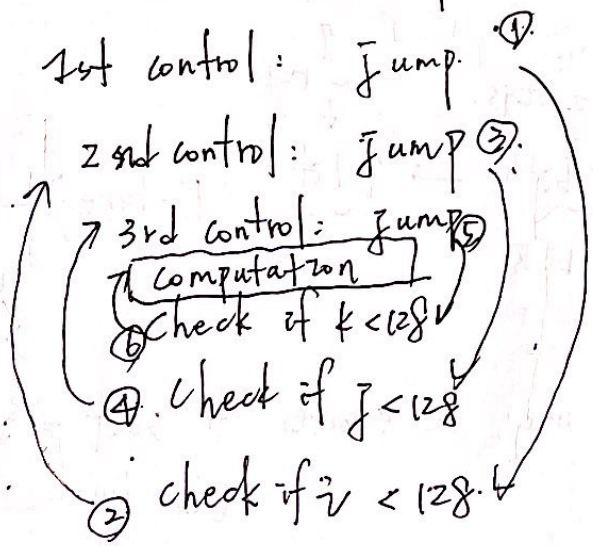
-① naive multiplication : 115558044 cycles / block multiplication:

-② # of load & store : $\underbrace{4\times128}_{\text{1st loop}} + \underbrace{4\times128^2}_{\text{2nd loop}} + \underbrace{(4+17)\times128^3}_{\text{3rd loop}}$
(naive compute)

$= 41011Z$

-④ There are three loops to do the multiplication

the # of loop control for each loop is 2 ⇒ 6 in total.



- ③ the use : C[i][j] += A[i][k] * B[k][j] → result matrix at one entry in the a time

⇒ Can compute a small block of entries simultaneously

eg.
```
for (i=0; i<128; i+=2){
    for (j=0; j<128; j+=2){
        a_00 = a_01 = a_10 = a_11 = 0;
        for (k=0; k<128; k++){
            a_00 += B[k][j+0] * A[i+0][k];
            a_01 += B[k][j+1] * A[i+0][k];
            a_10 += B[k][j+0] * A[i+1][k];
            a_11 += B[k][j+1] * A[i+1][k];
        }
        C[i+0][j+0] = a_00;
        C[i+0][j+1] = a_01;
        C[i+1][j+0] = a_10;
        C[i+1][j+1] = a_11;
    }
}
```

each load values twice ⇓ 4 loads 4 multiply-adds.