Projekt 2: "fast" Matrix-Matrix multiplication

The standard algorithm to multiply two matrices $A, B \in \mathbb{R}^{n \times n}$ has complexity $\mathcal{O}(n^3)$. Volker Strassen published in 1969 the *Strassen-Algorithm* that reduced this complexity (i.e., the number of arithmetic operations) to $\mathcal{O}(n^{\log_2(7)})$.

The Strassen algorithm is as follows. For simplicity, let $n=2^m$. The goal is to compute

$$C = AB. (1)$$

The matrices are decomposed into 4 blocks of size 2^{m-1}

$$A = \begin{pmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{pmatrix}, B = \begin{pmatrix} B_{11} & B_{12} \\ B_{21} & B_{22} \end{pmatrix} \text{ und } C = \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix}$$
 (2)

The standard multiplication can be written as

$$C = \begin{pmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{pmatrix} = \begin{pmatrix} A_{11}B_{11} + A_{12}B_{21} & A_{11}B_{12} + A_{12}B_{22} \\ A_{21}B_{11} + A_{22}B_{21} & A_{21}B_{12} + A_{22}B_{22} \end{pmatrix}.$$
(3)

Another, equivalent, representation is obtained in terms of the matrices

$$M_1 := (A_{11} + A_{22})(B_{11} + B_{22}) \tag{4a}$$

$$M_2 := (A_{21} + A_{22})B_{11} \tag{4b}$$

$$M_3 := A_{11}(B_{12} - B_{22}) \tag{4c}$$

$$M_4 := A_{22}(B_{21} - B_{11}) \tag{4d}$$

$$M_5 := (A_{11} + A_{12})B_{22} \tag{4e}$$

$$M_6 := (A_{21} - A_{11})(B_{11} + B_{12}) \tag{4f}$$

$$M_7 := (A_{12} - A_{22})(B_{21} + B_{22}) \tag{4g}$$

and the observation that

$$C_{11} = M_1 + M_4 - M_5 + M_7,$$

$$C_{12} = M_3 + M_5,$$

$$C_{21} = M_2 + M_4, \text{ und}$$

$$C_{22} = M_1 - M_2 + M_3 + M_6.$$

The multiplication in (4) is done recursively until only 1×1 matrices have to be multiplied.

- 1. Check the correctness of the above algorithm and prove its complexity.
- 2. Program the following functions for square matrices of size 2^m of type float in C or C++:
 - a) the standard algorithm for the matrix-matrix multiplication
 - b) the Strassen algorithm for the matrix-matrix multiplication

The function for the Strassen algorithm is recursive and it shouldn't allocate memory. Therefore, it should have a fourth argument with *work space*, e.g.,

```
float * A = new float [n*n]; FillWithSomething (A);
float * B = new float [n*n]; FillWithSomething (B);
float * C = new float [n*n];
float * Work = new float [n*n]; // don't use 'new' again
ComputeStrassen(n,n,A,B,C,Work); // A*B=C using 'Work'
```

Document your source code sufficiently. Test your code with matrices for small m until you are convinced that your code works properly. Test first the function for the standard multiplication using, if necessary, matlab. Test your implementation of the Strassen algorithm by comparing with the standard multiplication. Document your test examples and results.

3. The complexity and the run-time behavior of the two algorithms are to be compared. Write a program that measure the CPU-time needed for the algorithms of Problem 2. Use the function clock(),

```
#include <ctime>
...
  clock_t start = clock();
  ComputeStrassen(...);
  clock_t end = clock();

took = difftime(end, start)*1000.0/CLOCKS_PER_SEC;
  std::cout << "Took" << took << " [millisec]" << std::endl;</pre>
```

Write a program that inputs m from the command line and generates two matrices A and B of size $n = 2^m$ with random entries in [0, 1]. Measure the times needed:

Call your program with m = 6, ..., 11 (or larger if your memory permits it). Collect the timings in a table. Plot the timings using matlab. What do you observe?

hint: Do not use the optimization flags of the compiler as this may improve the timings significantly for certain values.