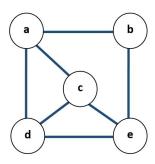
What is Graph?

- A graph consists of a number of data items, each of which is called a vertex. Any vertex may be connected to any other, these connections are called edges.
- A graph G = (V,E) is composed of:

V: set of vertices

E: set of edges connecting the vertices in V

- An edge e = (u,v) is a pair of vertices
- Example:



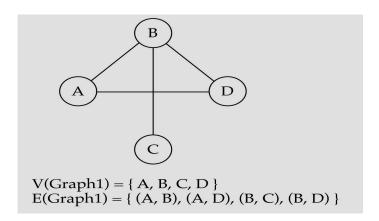
 $V= \{a,b,c,d,e\}$ $E= \{(a,b),(a,c),(a,d),$ $(b,e),(c,d),(c,e),(d,e)\}$

Main Application of Graphs

- Electronic circuits
- · Printed circuit board
- Integrated circuit
- Transportation networks
- Highway network
- Flight network
- Computer networks
- Local area network
- Wide area network
- Internet
- Databases
- Entity-relationship

Directed And Undirected Graph

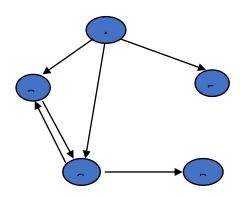
Undirected graph When the edges in a graph have no direction, the graph is called *undirected*.



• When the edges in a graph have a direction, the graph is called directed graph. Also known as digraph. This kind of graph contains ordered pair of vertices i.e. if the graph is directed, the order of the vertices in each edge is important.

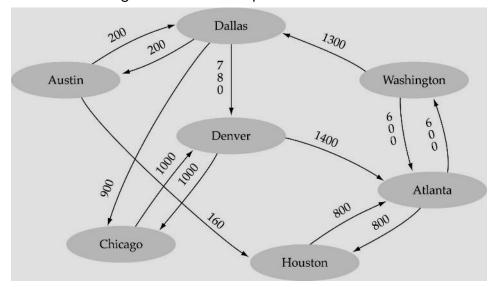
$$V = {A, B, C, D, E}$$

 $E = {(A,B), (B,C), (C,B), (A,C), (A,E), (C,D)}$

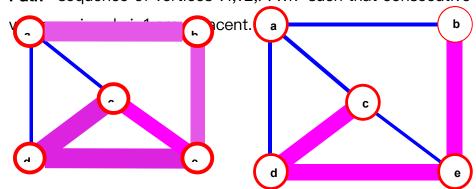


Basic Terminologies

• **Weighted Graph**— A graph in which every edge is assigned some value which is greater than or equal to zero.



- Adjacent Nodes
 — When there is an edge from one node to another
 then these nodes are called adjacent nodes.
- Path- sequence of vertices v1,v2,...vk such that consecutive

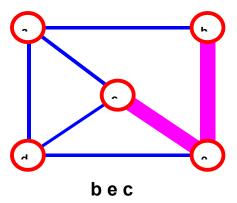


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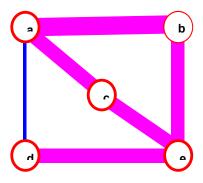
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Length of a Path
 — Length of a path is nothing but the total number
 of edges included in the path from source to destination node.

• Simple Path- path such that all its vertices and edges are distinct.

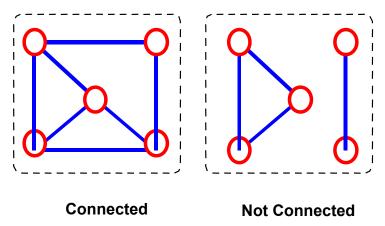


- Cycle- simple path, except that the last vertex is the same as the first vertex.
- Acyclic Graph- A graph without cycle is called acyclic Graph.



Connected Graph

— An undirected graph is connected if, for any
pair of vertices, there is a path between them.



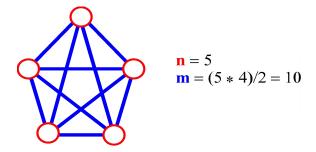
- A directed graph
 is connected if, for any pair of vertices, there is
 a path between them.
- Complete Graph
 — A graph in which all pairs of vertices are adjacent
 OR

A graph in which every vertex is directly connected to every other vertex

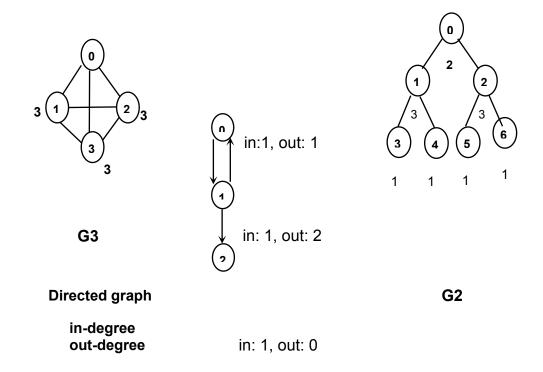
Let n = Number of vertices, and

m = Number of edges

For a complete graph with n vertices, the number of edges is n(n-1)/2. A graph with 6 vertices needs 15 edges to be complete.



- Degree of a node— the total number of edges linked to a node in an Undirected graph, For directed graph there are two degree for every node:
- Indegree
 The indegree of a node is the total number of edges coming to that node.
- Outdegree
 The outdegree of a node is the total number of edges going out from that node.



Graph Representation

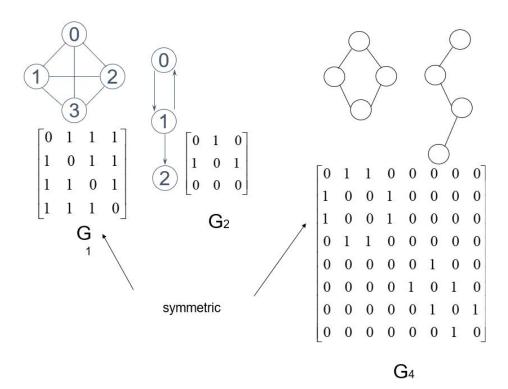
There are two methods to represent a Graph

- Adjacency Metrix (Array Implementation)
- Adjacency List (Linked List Implementation)

Adjacency Matrix

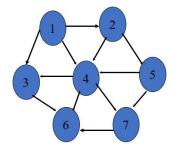
- Let G=(V,E) be a graph with n vertices.
- The adjacency matrix of G is a two-dimensional n by n array, say adj_mat
- If the edge (vi, vj) is in E(G), adj_mat[i][j]=1
- If there is no such edge in E(G), adj_mat[i][j]=0
- The adjacency matrix for an undirected graph is symmetric; since adj_mat[i][j]=adj_mat[j][i]
- The adjacency matrix for a digraph may not be symmetric

Examples for Adjacency Matrix



Adjacency Matrix

- From the adjacency matrix, to determine the connection of vertices is easy
- The degree of a vertex is
- We can also represent weighted graph with Adjacency Matrix. Now the contents of the matrix will not be 0 and 1 but the value is substituted with the corresponding weight.



	1	2	3	4	5	6	7
1	0	1	1	1	0	0	0
2	0	0	0	1	1	0	0
3	0	0	0	0	0	1	0
4	0	0	1	0	0	1	1
5	0	0	0	1	0	0	1
6	0	0	0	0	0	0	0
7	0	0	0	0	0	1	0

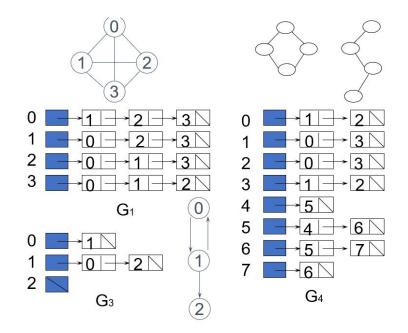
Another Example Adjacency

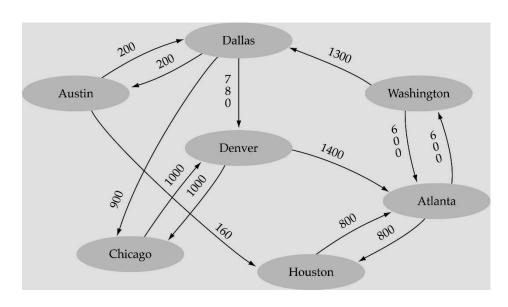
Adjacency List

 A Single Dimension array of Structure is used to

 represent the vertices. A Linked list is used for each vertex V which

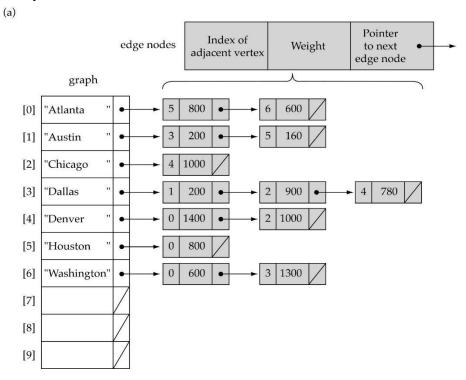
 contains the vertices which are adjacent from V (adjacency list)





Another Example of Adjacency List

Linked-list Implementation



Graph Traversal

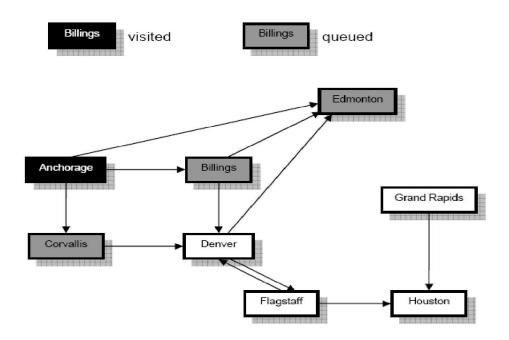
- Traversal is the facility to move through a structure visiting each of the vertices once.
- We looked previously at the ways in which a binary tree can be traversed. Two of the traversal methods for a graph are breadth first and depth—first.

Breadth-First Graph Traversal

- This method visits all the vertices, beginning with a specified start
 vertex. It can be described roughly as "neighbours-first".
- No vertex is visited more than once, and vertices are visited only if they can be reached — that is, if there is a path from the start vertex.

- Breadth-first traversal makes use of a queue data structure. The
 queue holds a list of vertices which have not been visited yet but
 which should be visited soon.
- Since a queue is a first-in first-out structure, vertices are visited in the order in which they are added to the queue.
- Visiting a vertex involves, for example, outputting the data stored in that vertex, and also adding its neighbours to the queue.
- Neighbours are not added to the queue if they are already in the queue, or have already been visited.

Breadth-First Traversal Example

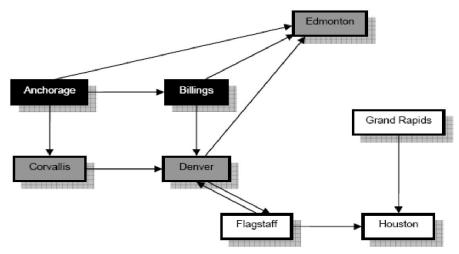


Visited: Anchorage

Queue: Billings, Corvallis, Edmonton

visit Billings next

 Neighbors are added to the queue in alphabetical order. Visited and queued vertices are shown as follows:

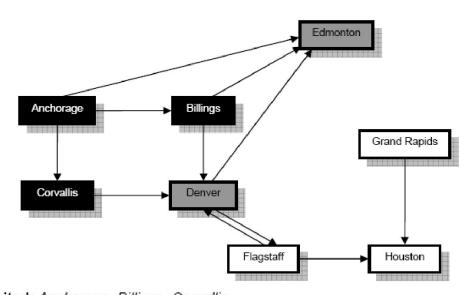


Visited: Anchorage, Billings

Queue: Corvallis, Edmonton, Denver

visit Corvallis next

 Note that we only add Denver to the queue as the other neighbours of Billings are already in the queue.

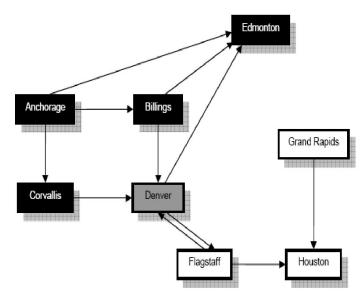


Visited: Anchorage, Billings, Corvallis

Queue: Edmonton, Denver

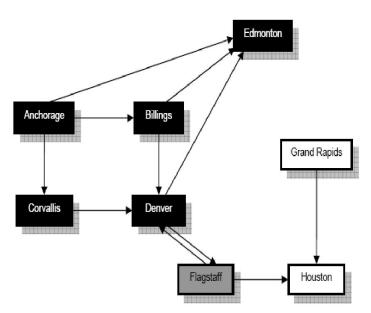
visit Edmonton next

 Note that nothing is added to the queue as Denver, the only neighbour of Corvallis, is already in the queue.



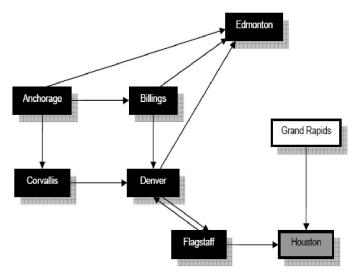
Visited: Anchorage, Billings, Corvallis, Edmonton

Queue: Denver next

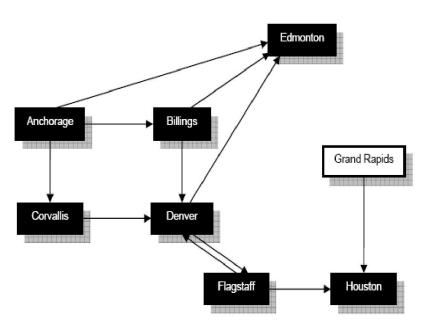


Visited: Anchorage, Billings, Corvallis, Edmonton, Denver

Queue: Flagstaff visit Flagstaff next



Visited: Anchorage, Billings, Corvallis, Edmonton, Denver, Flagstaff **Queue:** Houston visit Houston next



Visited: Anchorage, Billings, Corvallis, Edmonton, Denver, Flagstaff, Houston

Queue: empty

- Note that Grand Rapids was not added to the queue as there is no path from Houston because of the edge direction.
- Since the queue is empty, we must stop, so the traversal is complete.
- The order of traversal was: Anchorage, Billings, Corvallis, Edmonton, Denver, Flagstaff, Houston

Depth-First Search

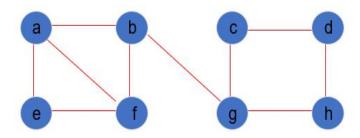
- From the given vertex, visit one of its adjacent vertices and leave others;
- Then visit one of the adjacent vertices of the previous vertex;
- Continue the process, visit the graph as deep as possible until:
 - A visited vertex is reached;
 - An end vertex is reached.

Depth-First Traversal

- Depth-first traversal of a graph:
- Start the traversal from an arbitrary vertex;
- Apply depth-first search;
- When the search terminates, backtrack to the previous vertex of the finishing point,
- Repeat depth-first search on other adjacent vertices, then backtrack to one level up.
- Continue the process until all the vertices that are reachable from the starting vertex are visited.
- Repeat above processes until all vertices are visited.
- This method visits all the vertices, beginning with a specified start vertex.

- This strategy proceeds along a path from vertex V as deeply into the graph as possible.
- This means that after visiting V, the algorithm tries to visit any unvisited vertex adjacent to V.
- When the traversal reaches a vertex which has no adjacent vertex,
 it back tracks and visits an unvisited adjacent vertex.
- Depth-first traversal makes use of a **Stack data structure**.

DFS Traversal of Undirected Graph



DFS traversal stack:

DFS tree:

