

# Build X : Algorithms

## Week 3

- Week 2 - Winner && Challenges
- Lee's Algorithm
- Mouse Labyrinth
- Minesweeper Game
- Subarray of Maximum Sum



## Week 2 - Challenges

- The Dragon (Not the Android One) – 100<sub>pts</sub>
  - Solved with a simple formula :  $3 * D + H$ 
    - D = the number of days our hero fights the dragon
    - H = the number of heads the dragon initial had
  - Another solution would be to use a loop to pass through each day and calculate the number of heads the dragon will have after each day.



## Week 2 - Challenges

- Let's First Sort – 100<sub>pts</sub>
  - Input : 5 6 4 7 2 1
    - 5 elements : 6 4 7 2 1
  - Sort the elements using any sorting algorithm learned.
    - New list : 1 2 4 6 7
  - Select and print the elements that are on even positions.
    - Elements to print : 2 6

# And the winner is ...Amandine Jala

- Alex Clapa
- Jerome De Chillaz
- Tania Copocean
- Amandine Jala
- Razvan-Gabriel Ceangu
- Ruxandra Anghel

# Lee's Algorithm - Short Description

- Identifies the minimum number of steps we need to make in a matrix.
  - Example of usage:
- Labyrinth Problem – Determine the number of steps we need to make to find an exit

# Lee's Algorithm - Characteristics

- Identical with BF but used on a 2D data structure such as a matrix.
- Considered to be efficient:
  - Time Complexity :  $O(m*n)$
- Used quite often

# Lee's Algorithm - Representation

Row/ Column	0	1	2	3	4
0	X	X	X	X	X
1	X	0	0	0	-1
2	X	0	0	X	X
3	X	0	0	X	X
4	X	0	X	X	X

# Labyrinth Problem - Description

- Given a labyrinth find a path to the solution.
- Steps :
  - Represent your labyrinth;
  - Specify the possible moves;
  - Apply the algorithm;



# Labyrinth Problem - Representation

[illegible]

**Time to code**

**-Matrix Representation-**

# Labyrinth Problem - Possible Moves

- a. N, S, E, V (4 moves)
  - Use 2 lists for xOy representations
    - $dx = \{-1, 0, 1, 0\};$
    - $dy = \{0, 1, 0, -1\};$
- b. N, S, E, V, NE, NV, etc. (8 moves)
  - Use the same 2 lists but with more coordinates:
    - $dx = \{-1, -1, 0, 1, 1, 1, 0, -1\};$
    - $dy = \{0, -1, -1, -1, 0, 1, 1, 1\};$

# Time to code

-Lee Algorithm-

# Different variations

- Add obstacles
- Add more solutions
- Use a Queue structure to expand only certain nodes

**Break**  
**-10 min-**

# Queue

- It is an abstract data structure which works after the FIFO technique
- Elements are kept are kept in order
- Has only 2 methods
  - Enqueue :
    - addition of elements to the rear of the queue;
  - Dequeue :
    - Removal of elements from the front of the queue;

# Queue - Advantages

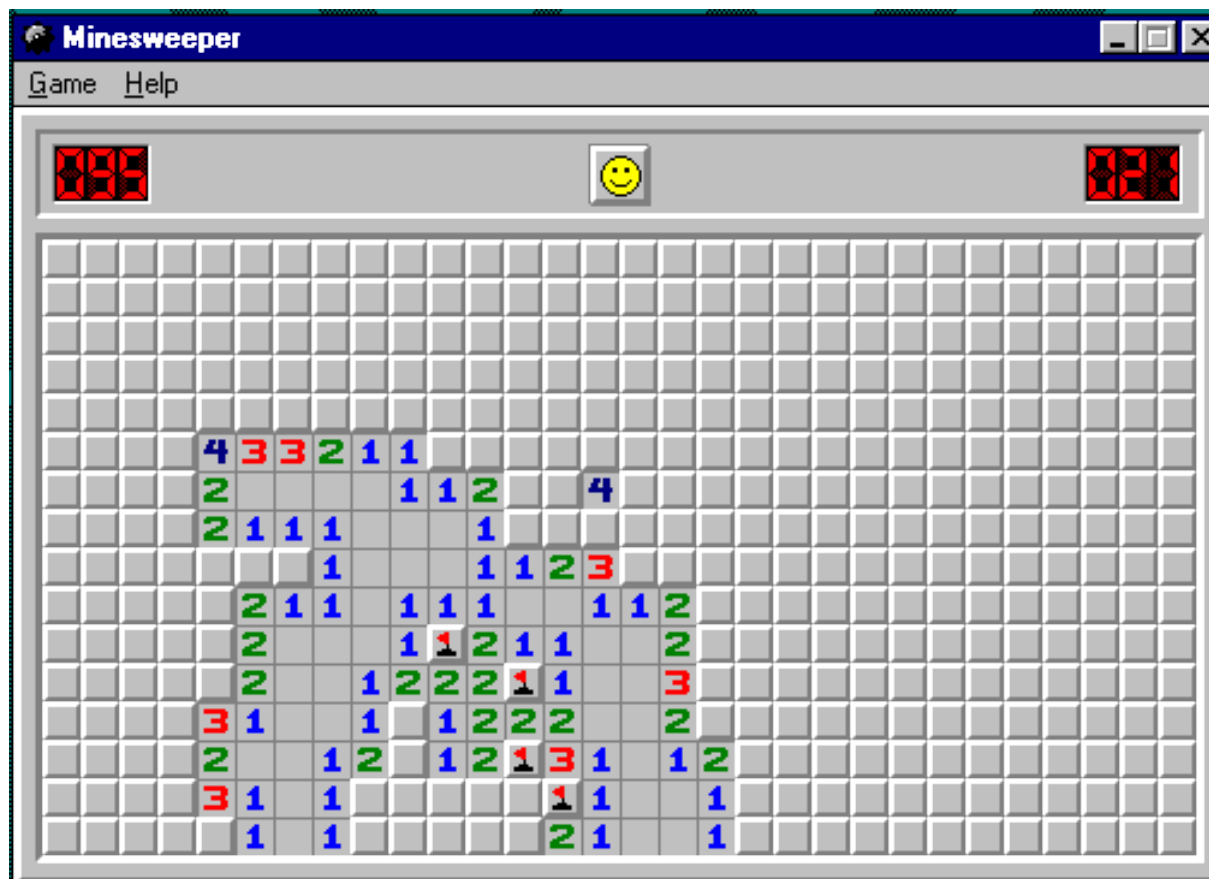
- No real limit on the memory
- Used when we need to use the elements in the same order we got them
  - E.g.
    - Manage print commands from multiple users connected to a printer
    - Hotel Booking Systems



# Queue in Lee's Algorithm

- We can use Queues in order to keep track of the nodes we need to expand next or nodes that have already been expanded

# Minesweeper



# Minesweeper - Steps

- Create the field
- Border it
- Place the bombs
- Initialize the rest of the places with the number of neighbours that contains a bomb
- Use Lee's Expansion Algorithm to interact with the user's click

# Maximum Subarray - Examples

- E.g.
  - 5 4 1 -10 2
    - 5 4 1
  - 1 2 3 -4 10
    - 1 2 3 -4 10
  - -2 -1 -3 4 -1 2 1 -5 4
    - 4 -1 2 1

# Maximum Subarray - Description

- Given a list of elements find the sub-list that have the maximum sum
- Let's take:
  - $S[] = (s_1, s_2, s_3, \dots, s_n)$ ,  $S$  contains  $n$  elements
  - A sub-list of  $S$  is a list of the form :
    - $(s(i), s(i+1), s(i+2), \dots, s(j)); 0 \leq i \leq j < n$
  - Sum of the sub-list is
    - $\text{Sum} = s(i) + s(i+1) + s(i+2) + \dots + s(j);$

# Maximum Subarray - Solutions

- 1.  $O(n^3)$ 
  - Set 2 indexes  $i$  and  $j$  and iterate between them
- 2.  $O(n^2)$ 
  - Use the same 2 indexes but calculate the sum while you iterate
- 3.  $O(n * \log(n))$ 
  - Use the Divide and Conquer technique.

# Maximum Subarray - Solutions

- 4.  $O(n)$ 
  - Using a new structure to calculate all the sums up to the position  $i$
- 5.  $O(n)$ 
  - Dynamic Programming

# Next week

Special Guest : **Luca Vigano**  
-Algorithms in Security-

