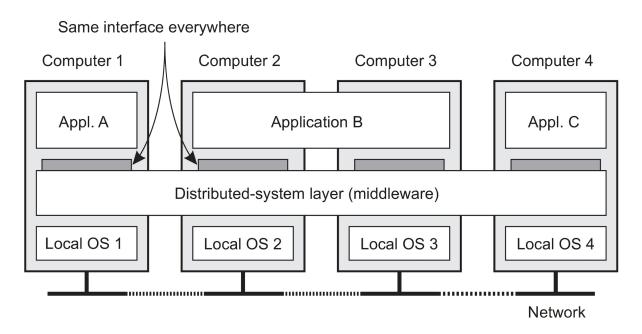
Communication

Agenda

- Remote Procedure Call
- Message-oriented Middleware

References: Chapter 4 of "Distributed Systems by M. van Steen and A. S. Tanenbaum"

Middleware



A software layer that abstracts from heterogeneity of networks, hardware, operating systems and programming languages, and provides a uniform computational and communication model

Services

A middleware offers services to applications, e.g.,

- Communication
- Security mechanism
- Transactions
- Error-recovery
- Management of shared resources

Services are independent of the specific application

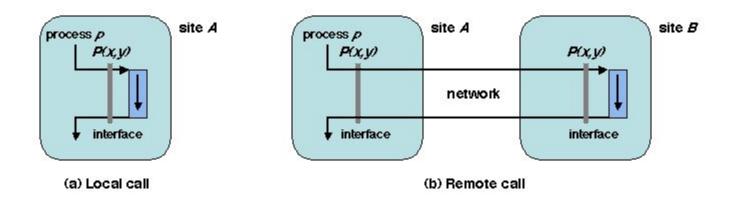
Direct Coordination

A kind of communication where components are

- Referentially coupled: during the communication components use an explicit reference to the communication partner (an address, a name)
- Temporally coupled: both components must be up and running to communicate

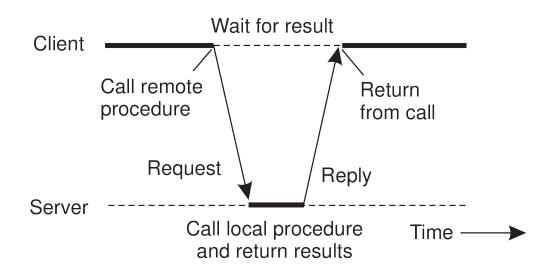
Remote Procedure Call

The idea: Programs can call procedures on other machines



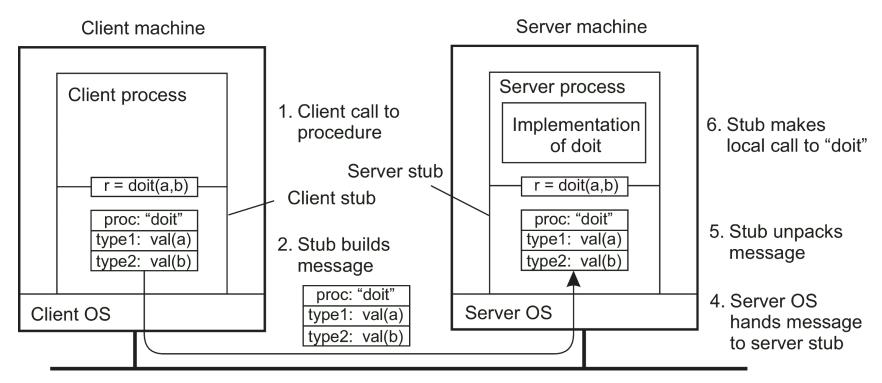
A Glimpse

- When a client calls a remote procedure, its execution is suspended
- A request is send to the server
- The server calls the corresponding local procedure
- The results of the call is returned to the client
- The client resumes the execution



Challenges: two different address spaces, serialization of parameters & return value, management of possible failures

How RPC works



3. Message is sent across the network

The Steps For Calling a Remote Procedure

- 1. The client calls the procedure stub
- 2. The stub builds a message pass it to OS
- 3. The client OS sends the message to the server
- 4. The server receives the message and calls the server stub
- 5. The stub extracts the parameter from the message
- 6. The stub calls the procedure and gets the result
- 7. The stub builds a message and pass it to OS
- 8. The server sends the message to the client
- 9. The client stub gets the message and unpacks the result
- 10. The client resumes the execution

A Conceptual Example (Client-side)

```
# Stub on the client
class Client:
    def append(self, data, dbList):
        msglst = (APPEND, data, dbList)
        self.chan.sendTo(self.server, msglst)
        msgrcv =
self.chan.recvFrom(self.server)
    return msgrcv[1]
```

```
# Remote invocation on the client
dbHandle = client.append(newTable,
dbHandle)
```

A Conceptual Example (Server-side)

```
# Actual data structure on the server
                                                # Main loop of the server
                                                while True:
class DBList:
 def append(self, data):
                                                  msgreq = self.chan.recvFromAny()
   self.value = self.value + [data]
                                                  client = msgreq[0]
   return self
                                                  msgrpc = msgreq[1]
                                                  if APPEND == msgrpc[0]:
# Stub on the server
                                                    result = self.append(msgrpc[1],
class Server:
                                                msgrpc[2])
 def append(self, data, dbList):
                                                    self.chan.sendTo(client, result)
   return dbList.append(data)
```

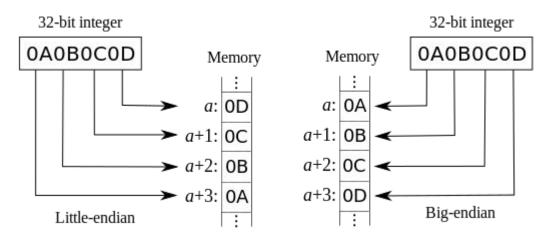
Parameter Passing

- The operation of packing parameters into a message is called parameter marshaling
- In the end, a message is a sequence of bytes containing parameters + other information useful for the server
- There are two challenges to solve in parameter passing:
 - 1. Client and server may have **different data representation**
 - 2. They have to use the **same encoding**

Data representation

The order of byte in memory may differ between machines

Little-endian vs Big-endian



Data representation

Solution: transform data into a machine-independent format

Network Byte Order

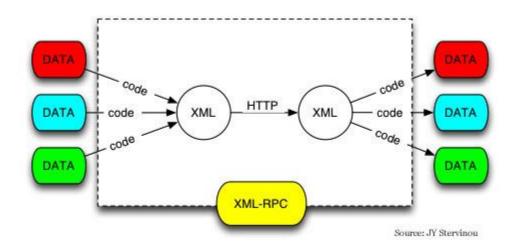
```
#include <netinet/in.h>
uint32_t htonl(uint32_t hostlong);
uint16_t htons(uint16_t hostshort);
uint32_t ntohl(uint32_t netlong);
uint16_t ntohs(uint16_t netshort);
```

Data Encoding

Client and server have to agree on how basic/complex data values are represented, e.g.,

- Layout of multi-dimensional arrays: row-major vs column-major
- Layout of records/objects

Example: XML-RPC



XML-RPC uses HTTP as the transport protocol and XML as the encoding format

XML-RPC: Data Encoding

```
<br/>
<br/>
doolean>1</boolean>
<array>
<data>
 <value><i4>1404</i4></value>
                                            <dateTime.iso8601>
 <value><string>
                                              19980717T14:08:55
   Something here
                                            </dateTime.iso8601>
  </string></value>
 <value><i4>1</i4></value>
                                            <double>-12.53</double>
</data>
                                            <string>Hello world!</string>
</array>
```

XML-RPC: Message Encoding

```
<?xml version="1.0"?>
<methodCall>
<methodName>
 examples.getStateName
</methodName>
<params>
 <param>
   <value><i4>40</i4></value>
 </param>
</params>
</methodCall>
```

```
<?xml version="1.0"?>
<methodResponse>
<params>
 <param>
   <value><string>
    South Dakota
   </string></value>
 </param>
</params>
</methodResponse>
```

Passing References/Pointers

```
# Stub on the client
class Client:
    def append(self, data, dbList):
        msglst = (APPEND, data, dbList)
        self.chan.sendTo(self.server, msglst)
        msgrcv =
self.chan.recvFrom(self.server)
    return msgrcv[1]
```

We cannot simply pass a reference to the server

The Problem

- Client and server have two different address spaces
- References and pointers refer to memory locations that have meaning locally only to the calling process
- A reference may refer to something different in the client and in the server

A First Solution

We copy the entire data structure the pointer is referring to

- Quite easy for flat data structure, arrays
- Hard for nested data types, e.g., trees, graphs, etc.
 - Copying large and nested data structure may be inconvenient because of the overhead due to the (un)marshalling and transmission processes
 - Many OO languages provide an automatic support for (un)marshalling objects

Our Conceptual Example Revisited

```
# Stub on the client
                                               # Main loop of the server
class Client:
                                               while True:
 def append(self, data, dbList):
                                                 msgreq = self.chan.recvFromAny()
   msglst = (APPEND, data, dbList)
                                                 client = msgreq[0]
   msgsnd = pickle.dumps(msglst)
                                                 msgrpc = pickle.loads(msgreq[1])
   self.chan.sendTo(self.server, msglst)
                                                 if APPEND == msgrpc[0]:
                                                    result = self.append(msgrpc[1],
   msgrcv =
self.chan.recvFrom(self.server)
                                               msgrpc[2])
                                                    msgres = pickle.dumps(result)
   return msgrcv[1]
                                                    self.chan.sendTo(client, result)
```

Example: XML-RPC Request

```
<?xml version="1.0"?>
<methodCall>
<methodName> list.multiply </methodName>
<params> <param>
   <array> <data>
    <value><i4>1404</i4></value>
    <value><i4>10</i4></value>
    <value><i4>1</i4></value>
   </data></array>
 </param></params>
</methodCall>
```

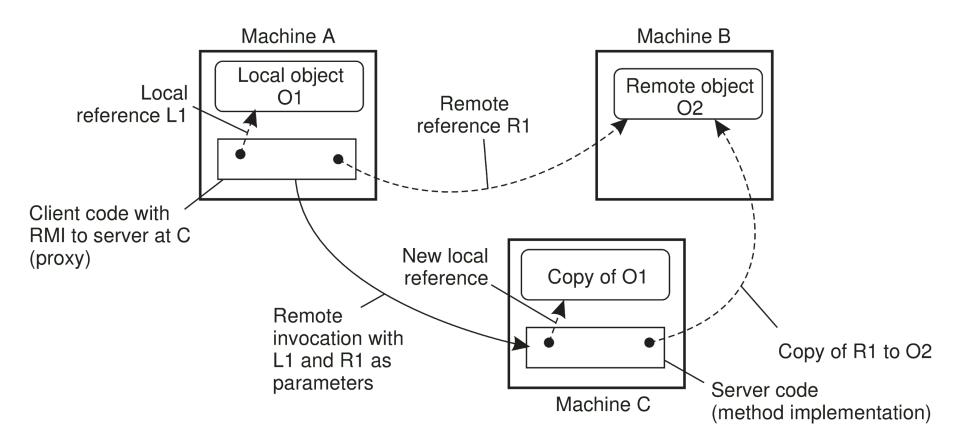
Global References

- References are meaningful to the client and the server
- When they are passed as parameters to a remote procedure, they can be simply copied by one machine to the other
- They are typically used in object-based systems

Parameter Passing in Object-based Systems

- There are two kinds of objects, local and remote, that are treated differently:
 - Local objects are copied and transmitted entirely
 - o For remote objects, only the stub is copied and transmitted
- In Java local and remote objects have different types, i.e., remote objects implements Remote interface

Example



RPC Support

There are two ways in which a RPC mechanism can be provided to developers:

- 1. As a framework/library together with suitable tools, e.g., a compiler
- 2. As a programming language construct

RPC Framework

- Programmers need to specify what is remotely exported by providing an interface of the service
- Typically, the interface is written in ad hoc language, called
 Interface Definition Language (IDL), that needs to be compiled
 to generate the stubs for the client and the server

Pros: programming language independent

Cons: not fully transparent to programmers

Some Frameworks

Apache Thrift ™







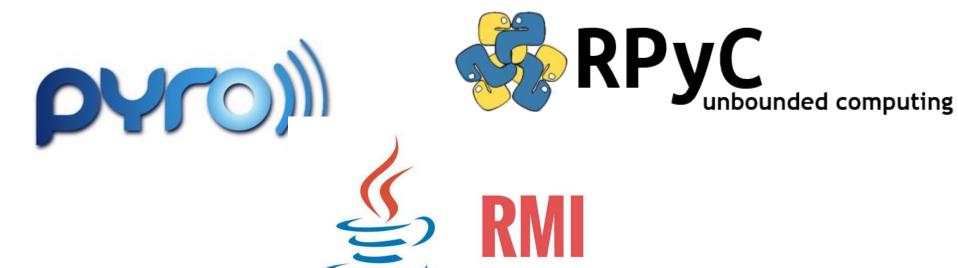
Programming Language-based Support

- The programming language provides constructs or mechanisms to denote a procedure/object as remote
- The compiler of the language takes care of generating the stubs during the compilation

Pros: quite transparent for the programmer

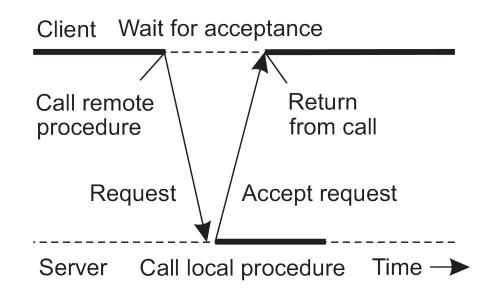
Cons: client & server must be written in the same language

Some Examples



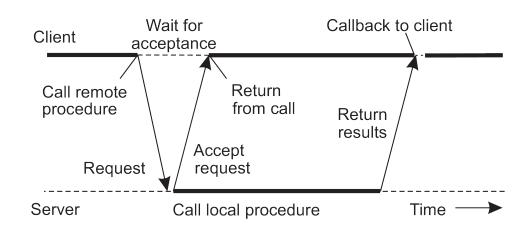
Asynchronous RPC

- After the request the client does not wait for the procedure results but continues its execution and it will wait the response of the server later
- The server immediately sends an acknowledgement as soon as it receives the request from the client



Callbacks

- The client implement a remote procedure to be invoked by the server to return the response
- The client sends the request that include a reference to the procedure, receives the acknowledgement and continues its computation
- The server sends the ack, processes the request and finally invokes the the callback to return the result to the client



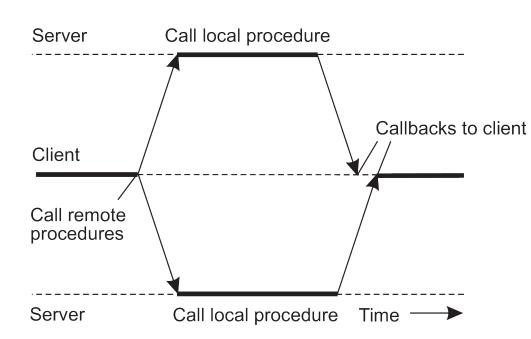
Multicast RPC

One-way RPC: the client does not wait for the acknowledgment of the server

The basic idea: the client sends a request to several servers that process it independently and in parallel

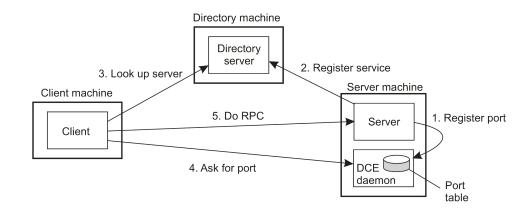
What about the responses?

- Only the first response that arrives is accepted
- 2. The responses from all servers are merged together



Binding a Client to a Server

- In real applications there is a preliminary step called "binding" that allows a client to take a "reference" to the server
- The registry or directory server stores information about the servers, usually mapping a symbolic name to the network address of the server
- The server registers to the registry
- The client looks the server up in the registry by using its symbolic name



Example of a List Server: The DB Stub

• • • •

Example of a List Server: The DB Stub

```
class DBClient:
  def sendrecv(self, message):
  def create(self):
    self.listID = self.sendrecv([CREATE])
    return self.listID
  def getValue(self):
    return self.sendrecv([GETVALUE, self.listID])
  def appendData(self, data):
    return self.sendrecv([APPEND, data, self.listID])
```

```
class Server:
  def __init__(self, port=PORT):
    # set up the connection with socket, bind, listen
    self.setOfLists = {} # init: no lists to manage
 def run(self):
   while True:
     (conn, addr) = self.sock.accept() # accept incoming call
     data = conn.recv(1024) # fetch data from client
     request = pickle.loads(data) ....
```

```
class Server:
  def __init__(self, port=PORT):
    . . .
 def run(self):
     if request[0] == CREATE: # create a list
        listID = len(self.setOfLists) + 1 # allocate listID
       self.setOfLists[listID] = [] # initialize to empty
        conn.send(pickle.dumps(listID)) # return ID
```

```
class Server:
  def __init__(self, port=PORT):
    . . .
 def run(self):
     elif request[0] == APPEND: # append request
       listID = request[2] # fetch listID
       data = request[1] # fetch data to append
       self.setOfLists[listID].append(data) # append it to the list
       conn.send(pickle.dumps(OK)) # return an OK
```

```
class Server:
  def __init__(self, port=PORT):
    . . .
 def run(self):
     elif request[0] == GETVALUE: # read request
       listID = request[1] # fetch listID
       result = self.setOfLists[listID] # get the elements
       conn.send(pickle.dumps(result)) # return the list
       conn.close() # close the connection
```

Example of a List Server: 2 Clients

```
client1 = Client(CLIENT1)  # create client
dbClient1 = DBClient(HOST,PORT)  # create reference
dbClient1.create()  # create new list
dbClient1.appendData('Client 1')  # append some data
client1.sendTo(HOSTCL2,CLIENT2,dbClient1) # send to other client
```

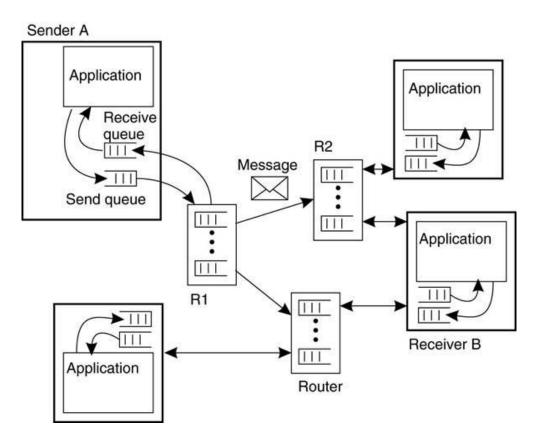
```
client2 = Client(CLIENT2) # create a new client
data = client2.recvAny() # block until data is sent
dbClient2 = pickle.loads(data) # receive reference
dbClient2.appendData('Client 2') # append data to same list
```

Message-oriented Middleware

The idea: components communicates by exchanging messages with each other

Persistent asynchronous communication: sender & receiver do not need to be active during the transmission of messages, the middleware provide storage services

Message-queueing Model



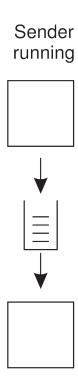
Message-queueing Model

- Each application is equipped with a local queue that contains the sent and received messages
- A queue may be shared by multiple applications
- Messages from the sender queue are forwarded over a series of communication servers until the destination queue is reached

Communication Properties

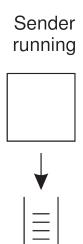
- Communication happens only by inserting/removing messages from queues
- Communication is loosely coupled in time
 - There is no need for the receiver to be running when a message is sent
 - A message remains in the queue until it is explicitly removed

- The sender & the receiver are running at the same time
- The receiver can start receiving the message while the sender is sending it



Receiver running

- The receiver is not ready for receiving messages
- The sender can send messages that will be stored in the queue

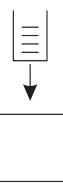




Receiver passive

- The sender is not sending, but there are messages in the queue
- The receiver get the messages from the queue





Receiver running

- The sender & receiver are not involved in communication
- The queue is "running" and stores messages sent until now







Receiver passive

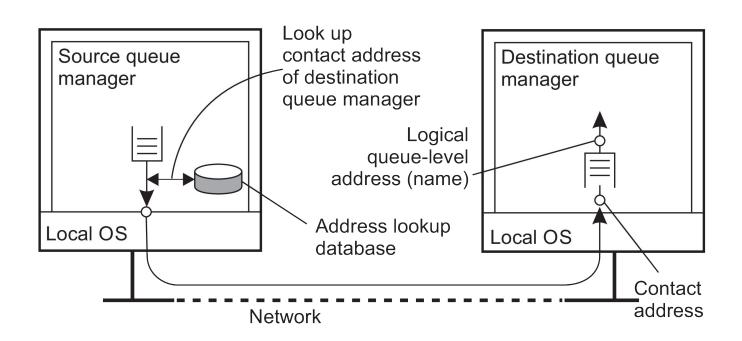
Messages

- Messages can contain any kinds of data and may be of any size within a given limit
- The middleware can support automatically fragmentation of messages
- A destination is identified together with its queue by a systemwide name
- Each message contains the address/name of the destination

Conceptual MoM API

put	Append a msg to a queue
get	Remove the first msg from the queue (blocking)
poll	Remove the first msg from the queue (non-blocking)
notify	Install an handler to be called when a msg arrives

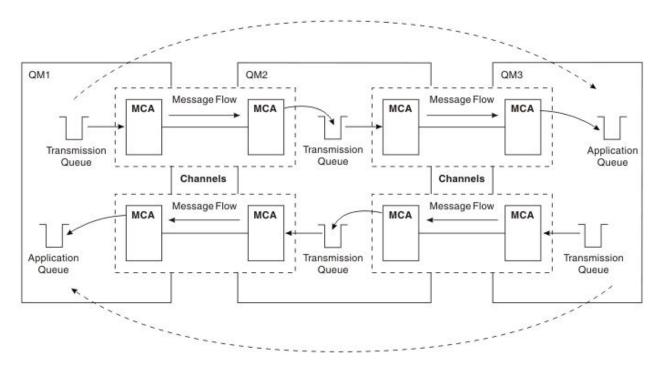
Queue Managers



Queue Managers

- It can be implemented either as a library linked with the application or as a separate process
- Each application is associated with a local queue, hence, it is paired with a queue manager
- An application can send/receive messages to/from its local queue
- Note: to support persistent asynchronous communication queue managers should be implemented as separate processes

Overlay Network



Queue managers form an overlay network used to forward messages in the system

Addresses

- Recall: each queue has a logical name/address that identifies the queue inside the system
- Logical names must be associated to a network address (host, port)
- Each queue manager maintains a translation table mapping names to network addresses
- The table may contain further information

Managing a Messaging Queue System

- An important task in setting up a MQ system is connecting together the queue managers into a overlay network
 - Setting up communication channels
 - Filling and handling the routing and translation tables
- A further effort consists in maintaining the overlay network when the system has to change, e.g., new queue managers

Heterogeneity

- Usually, MoM does not specify the format of messages
- Two different applications can communicate if they use the same protocol (as in network)
- When an application A joins the system, we need to ensure that any other application B interacting with A uses the same protocol
- Idea: provide a protocol for each B (no, it does not scale)

Message Brokers

A component in the MoM system that works as a gateway

- 1. It converts messages so they can be understood by different applications
- 2. It matches receivers in a publish-subscribe kind of interaction

Some Brokers

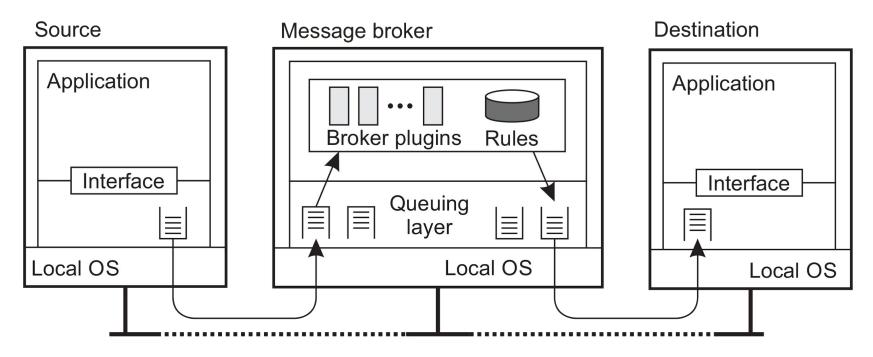






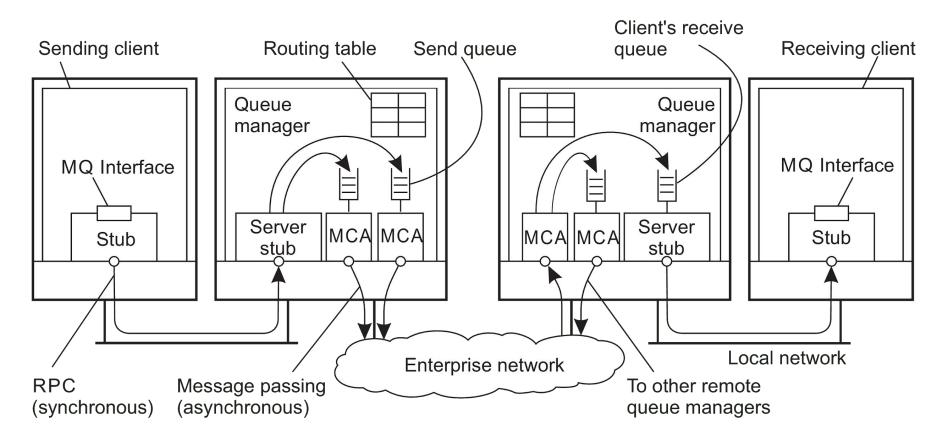


Message Brokers



There are rules that drive the broker to handle messages

Example: IBM WebSphere



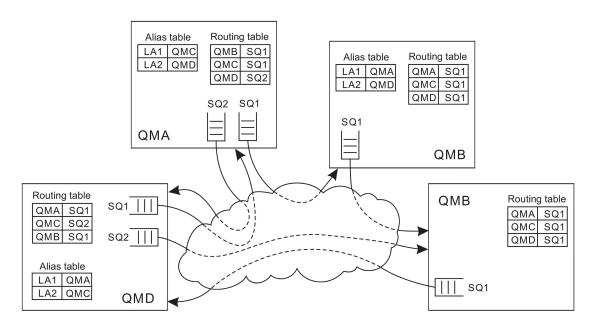
WebSphere Features

- Queue Managers are pairwise connected
- Message channels are unidirectional
- Queue Managers can be:
 - Linked to the application
 - External processes
- The application interacts with its manager through an API that is the same independently of the kind of its Queue Managers

Addressing in WebSphere

- Each message has a header including the name of the destination and the name of the local queue (sender)
- An address is a pair (queue manager, destination)
- Each Queue Manager has a unique name as well as each queue does
- There can be local aliases for denoting Queue Managers and queues

Routing in WebSphere

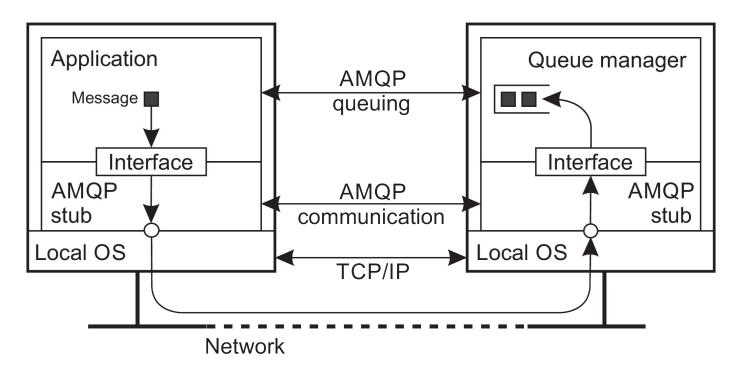


Each Queue Manager has a routing table

Application Interface

MQOPEN	Open a (possible remote) queue
MQCLOSE	Close a queue
MQPUT	Put a message into a opened queue
MQGET	Get a message from a queue

Example: AMQP



An OASIS Standard that allows MoM of different vendors to cooperate

Conclusion

- Remote Procedure Call
- Message-oriented Middleware

References: Chapter 4 of "Distributed Systems by M. van Steen and A. S. Tanenbaum"