



ECOO 2012

Programming Contest

Solutions and Notes

Final Competition (Round 3)

May 12, 2012

Problem 1: The Word Garden

Recommended Approach

Trying to generate and draw the trees on the screen all at once will be quite difficult. A better idea is to create a two-dimensional array of characters (24 rows, 40 columns), “draw” the trees into this array, then draw the array to the screen. Once the array is created, you can fill it with spaces and draw the ground in the final row immediately. Then you need an algorithm (preferably in a method, function, or procedure) for drawing a tree with its trunk at a given location. Test that algorithm to make sure you can draw a single word tree anywhere you want. It’s a good idea to put it in a method, procedure, or function.

The next part is figuring out where to draw each tree. The first one is not hard. If the word is 5 characters long, its trunk should be in the 5th column of the array. But figuring out where to draw each of the next trees is harder. You could try to do some mathematical or logical reasoning to figure out where each one should go, but once you have the drawing algorithm done, you can copy it and modify it so that instead of actually drawing the tree, it checks all the array locations the tree would need for a “collision” with other trees. This second algorithm can be used to test possible locations for the next tree. If you know where the trunk of the last tree was, you can start at that location plus 1 and then keep moving to the right until your tree checking algorithm says it’s ok. Then draw the tree there.

The Test Cases

The text for both data sets is from Martin Luther King Jr.’s famous “I have a dream” speech.

Solution to DATA11.txt

```
      d
    drd
  h   drerd   t
hah dreaerd tht
havahdreamaerdthaht o   d   t
havevah d   thatahtono dad tht
  h   r   t onenodayadotheht
  a   e   h   o   d ono t
I va   a   a   n   a o h
I ea   m   t   e   y n e
=====
```

```

      G
    GeG
  GeoeG      f
    GeoroeG   fof
      h   GeorgroeG   forof
    hih   GeorgigroeG   formrof
  hilihGeorgiaigroeG s   formemrof
    hilllih   G   sosformeremrof
  rhillsllih   e t   sonos   f
rer   h   o t t s o n s n o s   o
reder i o   r t h e h t   s o   r
r   lofo   g t   oofo   m
e   l o   i h   n o   e
d   s f   a e   s f   r
=====

```

```

      s      f
    sls      fof
  slals      forof      s
slavals      formrof   sls
slavevals      s   formemrof   slals
slavevals      sosformeremrofslavals
  s a   t   sonos   f   slavevals
l ana t h t s o n s n o s   o   s
a a n d n a t h e h t   s o   r   l
v a   t   oofo   m   a
e n   h   n o   e   v
s d   e   s f   r   e
=====

```

```

      o
    owo
  ownwo
ownenwo
ownerenwo w   a      d
ownersrenwowi w   aba   dod
  o   wiliw   ablba s   dowod
w   willliwablelbasisdowod
n   w b   a t s i t i s   d
e   ibeb   b t o t s   o
r   l b   l t i   w
s   l e   e o t   n
=====

```

```

b
brb
brorb
brotorb
brothtorb
brothehtorb
t    brotherehtorb
tot  brotherhrehtorb
togat brotherhohrehtorb
togeget brotherhooohrehtorb
togetegotbrotherhoodoohrehtorb
togethtegot    b
togethehtegot  t  r
togetherhtegot tat o
t    tabat t
o    tablbat h
g    ttablelbate
e  tht  t  r
t atheht  a o h
hata t    bofo o
e a h    l o o
r t e    e f d
=====

```

Solution to DATA12.txt

```

l
lil
d    litil
drd  litttil
h    drerd  t    f  littltil
hah  dreaerd  tht  foflitttleltil
havahdreamaerdthaht  fouof  l
havevah  d  thatahtfouruof  i
h    r    t m  f  t
a    e    hmym  o  t
I va    a    a m  u  l
I ea    m    t y  r  e
=====

```

```

c
chc
chihc
chilihc
childlihc
childrdlihc
childrerddlihc
childrenerddlihc
n
nan
natan
natitan
c  w    l  natioitan
h  wiw    lilnationoitan
i  wiliw  o  d  livil  n
lwillliwono  dadlivevil  a
d  w  onenodayad  l i  t
r  i  o  d  iini  i
e  l  n  a  v  ia  o
n  l  e  y  e  na  n
=====

```

```

j
juj
juduj
judgduj
judgegduj
judgedgduj
wherehw tht wiw judgedegduj
wherehwtheht wiliw n j t
w theyehtwillliwnon u tht
h t w notonb d btheht
e h i n bebgbyb t
r e l o b e b h
e y l t e d y e
=====

```

```

c
coc
coloc t
cololoc tht
colouoloc theht s
colouruoloctheieht sks
c theiriehtskiks b t
o t skiniksbub tht
l o h s butubbtheht
ofo e k b byb t
u o i i u b h
r f r n t y e
=====

```

```

c
chc
chahc
charahc
chararahc
characarahc
charactcarahc
charactetcarahc
tcharacteretcarahc
contentnoc tht c
contentnetnoctheht h
c theieht a
o theirieht r
n t a
t o h c
eifo e t
n o i e
t f r r
=====

```

Problem 2: Jewelry Tips

Recommended Approach

The best way to solve this is just a top to bottom, left to right sweep of the entire board. For each jewel, simulate swaps in all 4 directions (in order of priority – LT, UP, RT, DN) and check for lines, then undo the swap and continue. Keep track of the first “good” and “excellent” tips found, but quit immediately if you find a “normal” tip and return that. If you get to the end of the board without any normal tips, return the good tip you found, or the excellent tip if there was no good tip, or “game over” if no tips of any kind were found.

I found it helpful to create a special counting method (a.k.a. procedure or function) that takes the board and a location as parameters, then counts the jewels of the same colour vertically and horizontally in both directions. If there is only one line ≥ 3 , it returns the length of it. If there are two, it adds their lengths and returns that. If there are none, it returns 0. So 0 means no line, 3 or 4 means a single line, and 5 or more means either more than one line or a line of five.

Some Further Notes

When you swap two jewels, you have to check that the other jewel in the swap doesn’t end up creating lines as well. If it does, you’ve got an excellent tip, so you shouldn’t return it as a normal tip.

There are not many situations in which you would return a good tip, since when you move one jewel to create a line of 4, there is always another jewel you could move to make a line of 3 instead. It is only when this move creates multiple lines (making it excellent rather than normal) that you can return a good tip.

The Test Cases

The test cases were generated automatically by randomly generating thousands of boards. The boards were thrown away if they already contained a line of 3 or more. If not, they were checked using the jewelry tip routine written above to see what kind of tips they contained and kept if they met the criteria I was looking for.

I went with 10 cases because there are so many different kinds of errors you can make in programming this. I thought that the more test cases there were, the more errors I would shake out of the code being tested. But because rule 5 was added at the last minute to cover an ambiguity in the rules, none of the test cases required the use of that rule.

Solution to DATA21.txt

Norm: Y.UP@7,3
Good: W.UP@6,2
Excl: O.RT@4,3
Game Over
Norm: W.RT@3,4
Norm: Y.UP@7,0
Good: W.RT@6,2
Excl: B.DN@6,2
Norm: O.LT@6,3
Excl: O.RT@6,2

Solution to DATA22.txt

Excl: P.DN@4,5
Excl: G.DN@2,1
Excl: P.RT@7,3
Excl: P.RT@4,0
Good: W.DN@3,6
Good: G.RT@6,4
Norm: R.LT@0,6
Norm: Y.DN@2,2
Norm: G.DN@2,4
Norm: G.RT@3,0

Problem 3: Steam Arithmetic

The name “Steam” is a spoof of “Scheme”, an actual dialect of Lisp that is currently in use and is taught as a first programming language at Waterloo and some other universities. You might think that programming this in a Lisp dialect would give an advantage, but because of the requirement that the program read the data from a file and interpret it, it is not necessarily so easy.

Recommended approach

Recursion is the best approach here. Write a method (or procedure or function) that takes a 3-element list, separates the operator and the operands and applies the operator to the operands. If an operand is a list, it should be recursively evaluated before being applied.

The restrictions on spacing and the restriction of operators and operands to a single character make it easy to process the list by pulling apart the string one character at a time. There is no need for any further tokenizing, except in the case of lists as operands.

To isolate list operands, you could use a bracket counting method. That is, start to the right of the first open bracket with the counter set to 1. Step through the string one character at a time. Increment the counter at an open bracket, and decrement it at a closed bracket. When the counter hits 0 you’ve found the end of the list.

The Test Cases

In the test cases, I tried to hit a couple of simple cases, then progressively harder ones. I tried for a few cases with very deep nesting, some with tail recursion, and some with head recursion

(http://en.wikipedia.org/wiki/Tail_call)

Solution to DATA31.txt

```
1
0
-35
2
75
77
130
-19683
0
61
```

Solution to DATA32.txt

```
0
10
157
8
-20
160
6
-3456
1152
-9
```


Problem 4: Splitsville

This problem is inspired by the field of Machine Learning, specifically Concept Learning Systems. In a concept learning system, you have a list of training objects (e.g. satellite pictures of possible oil spills, data from medical scanners, etc.) with category labels (e.g. is/is not an oil spill, is/is not a tumor, etc.). Each object is described using a list of properties (usually numbers) and the task is to try to form a theory that separates the objects into their different categories. This theory can be viewed as a partitioning of a multi-dimensional space.

In this problem the training objects are houses, the labels are A and B, and the list of properties for each object are just the x and y location of the house. The “theory” is the set of fences that separate the A’s from the B’s. In machine learning, the theory you form should eventually do well at categorizing previously unseen objects, but that would only work in this case if A and B families tend to live close to other families of the same type.

There are a number of ways to partition spaces like this, including “nearest neighbour” classification and neural network classifiers. The algorithm used in this question is a simplification of machine learning algorithms that produce decision trees or decision rules as their output.

Recommended approach

This problem lends itself well to a recursive solution in which you partition a region of space into two parts, then recursively subdivide the two new regions created. The base case for the recursion is a “pure” region (i.e. a region with only A or B houses in it). Within each region the simplest strategy is to try all possible partitions and count the number of houses out of place in each one, where the number of houses out of place is just the sum of the minima of the counts of A and B houses on each side. So if the left side has 1 A and 3 B’s while the right side has 10 A’s and 2 B’s, the total number of out of place houses is 3.

There are at least two ways to represent a region in this space. The first is to think of it as a 2-dimensional space bounded by integer coordinates above and below and on the left and right. To process it, you try all partitions half way between each pair of integer coordinates. (To keep life simpler and avoid floating point numbers, you could also double the original coordinates and then only consider splits on odd numbered coordinates.) This approach is not very efficient but it’s fast enough for problems of this size. One potential pitfall is that you have to rule out splits that would not separate at least one house from the others.

The second, more efficient, method is to think of a region as just a list of houses. If there are 12 houses in the region, you only need to try at most 11 different cuts on each dimension, even if the houses are separated by thousands of units. To do this, you would have to sort the houses first, then find a split between each pair. Because of the sorting step, and the problems of creating and managing lists, this solution is more difficult to implement.

The Test Cases

The last test case in each set requires a few seconds of processing in Java when using the less efficient representation described above.

Solution to DATA41.txt

3
1
60
0
132

Solution to DATA42.txt

25
0
1
65
156