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CS 590 Homework Assignment
9: Divide-and-Conquer
Reinforcement Exercises

Due Date: April 3, 2022

Problem 11.6.1:

To characterize the recurrence equations using the master theorem, we must apply check which case apply as follows. The master theorem is used for recurrence equations of the form:

$$T(n) = \begin{cases} c & \text{if } n < d \\ aT\left(\frac{n}{b}\right) + f(n) & \text{if } n \geq d \end{cases}$$

Where $d \geq 1$ is an integer constant, $a \geq 1$, $c > 0$, and $b > 1$ are real constants, and $f(n)$ is a function that is positive for $n \geq d$.

The Master Theorem is defined as below:

Let $f(n)$ and $T(n)$ be defined as above.

1. If there is a small constant $\epsilon > 0$, such that $f(n)$ is $O(n^{\log_b(a) - \epsilon})$, then $T(n)$ is $\Theta(n^{\log_b(a)})$.
2. If there is a constant $k \geq 0$, such that $f(n)$ is $\Theta(n^{\log_b(a)} \log^k n)$, then $T(n)$ is $\Theta(n^{\log_b(a)} \log^{(k+1)} n)$.
3. If there are small constants $\epsilon > 0$ and $\delta < 1$, such that $f(n)$ is $\Omega(n^{\log_b(a) + \epsilon})$ and $af\left(\frac{n}{b}\right) \leq \delta f(n)$, for $n \geq d$, then $T(n)$ is $\Theta(f(n))$.

Case 1 characterizes the situation where $f(n)$ is polynomially smaller than the special function, $n^{\log_b(a)}$. Case 2 characterizes the situation when $f(n)$ is asymptotically close to the special function, and Case 3 characterizes the situation when $f(n)$ is polynomially larger than the special function.

Part A:

Let's characterize the recurrence equation below using the master theorem:

$$T(n) = \left(2 * T \left(\frac{n}{2} \right) \right) + \log (n)$$

From the given above,

$$T(n) = \left(a * T \left(\frac{n}{b} \right) \right) + f(n) \text{ for } n \geq d$$

From the above, we can deduce the following:

$$f(n) = \log(n), a = 2, \text{ and } b = 2$$

Let's find what the special function gets us:

$$n^{\log_b(a)} = n^{\log_2(2)} = n^1 = n$$

From the analysis, we can see that:

$$n > \log(n)$$

Which means that:

$$n^{\log_b(a)} > f(n)$$

This characterizes case 1 above. This means that $T(n)$ is $\Theta(n^{\log_b(a)})$ based on case 1 above.

So,

$$**$T(n)$ is $\Theta(n^{\log_2(2)})$**$$

Which means that

$$***T(n) is \Theta(n)***$$

Part B:

Let's characterize the recurrence equation below using the master theorem:

$$T(n) = 8T\left(\frac{n}{2}\right) + n^2$$

From the given above,

$$T(n) = \left(a * T\left(\frac{n}{b}\right)\right) + f(n) \text{ for } n \geq d$$

From the above, we can deduce the following:

$$f(n) = n^2, a = 8, \text{ and } b = 2$$

Let's find what the special function gets us:

$$n^{\log_b(a)} = n^{\log_2(8)} = n^3$$

From the analysis, we can see that:

$$n^3 > n^2$$

Which means that:

$$n^{\log_b(a)} > f(n)$$

This characterizes case 1 above. This means that $T(n)$ is $\Theta(n^{\log_b(a)})$ based on case 1 above.

So,

$$T(n) \text{ is } \Theta(n^{\log_2(8)})$$

Which means that

$$T(n) \text{ is } \Theta(n^3)$$

Part C:

Let's characterize the recurrence equation below using the master theorem:

$$T(n) = 16T\left(\frac{n}{2}\right) + (n \log(n))^4$$

From the given above,

$$T(n) = \left(a * T\left(\frac{n}{b}\right)\right) + f(n) \quad \text{for } n \geq d$$

From the above, we can deduce the following:

$$f(n) = (n \log(n))^4, a = 16, \text{ and } b = 2$$

Let's find what the special function gets us:

$$n^{\log_b(a)} = n^{\log_2(16)} = n^4$$

From the analysis, we can see that the special function is asymptotically close to $f(n)$, which means that we are in case 2.

Since we are in case 2 above, this means that $T(n)$ is $\Theta(n^{\log_b(a)} \log^{(k+1)} n)$ based on case 2 above.

We need to determine k, we can determine it as follows:

From case 2 definition:

$$f(n) \text{ is } \Theta(n^{\log_b(a)} \log^k n)$$

We know from the equation given in the Exercise that f(n) is:

$$f(n) = n^4 \log^4(n)$$

And since we have a and b we know that:

$$\log_b(a) = 4$$

This means that

$$f(n) \text{ is } \Theta(n^4 \log^k n)$$

This means that k = 4

Now, back to where we were at:

Since we are in case 2 above, this means that $T(n)$ is $\Theta(n^{\log_b(a)} \log^{(k+1)} n)$ based on case 2 above.

So,

$$**$T(n) \text{ is } \Theta(n^4 \log^5(n))$**$$

Part D:

Let's characterize the recurrence equation below using the master theorem:

$$T(n) = \left(7 * T\left(\frac{n}{3}\right) \right) + n$$

From the given above,

$$T(n) = \left(a * T\left(\frac{n}{b}\right) \right) + f(n) \text{ for } n \geq d$$

From the above, we can deduce the following:

$$f(n) = n, a = 7, \text{ and } b = 3$$

Let's find what the special function gets us:

$$n^{\log_b(a)} = n^{\log_3(7)} \cong n^{1.771}$$

From the analysis, we can see that:

$$n^{1.771} > n$$

Which means that:

$$n^{\log_b(a)} > f(n)$$

This characterizes case 1 above. This means that $T(n)$ is $\Theta(n^{\log_b(a)})$ based on case 1 above.

So,

$$T(n) \text{ is } \Theta(n^{\log_3(7)})$$

Part E:

Let's characterize the recurrence equation below using the master theorem:

$$T(n) = \left(9 * T\left(\frac{n}{3}\right) \right) + (n^3 \log(n))$$

From the given above,

$$T(n) = \left(a * T\left(\frac{n}{b}\right) \right) + f(n) \text{ for } n \geq d$$

From the above, we can deduce the following:

$$f(n) = n^3 \log(n), a = 9, \text{ and } b = 3$$

Let's find what the special function gets us:

$$n^{\log_b(a)} = n^{\log_3(9)} = n^2$$

From the analysis, we can see that:

$$n^2 < n^3 \log(n)$$

Which means that:

$$n^{\log_b(a)} < f(n)$$

This characterizes case 3 above. This means that $T(n)$ is $\Theta(f(n))$ based on case 3 above.

So,

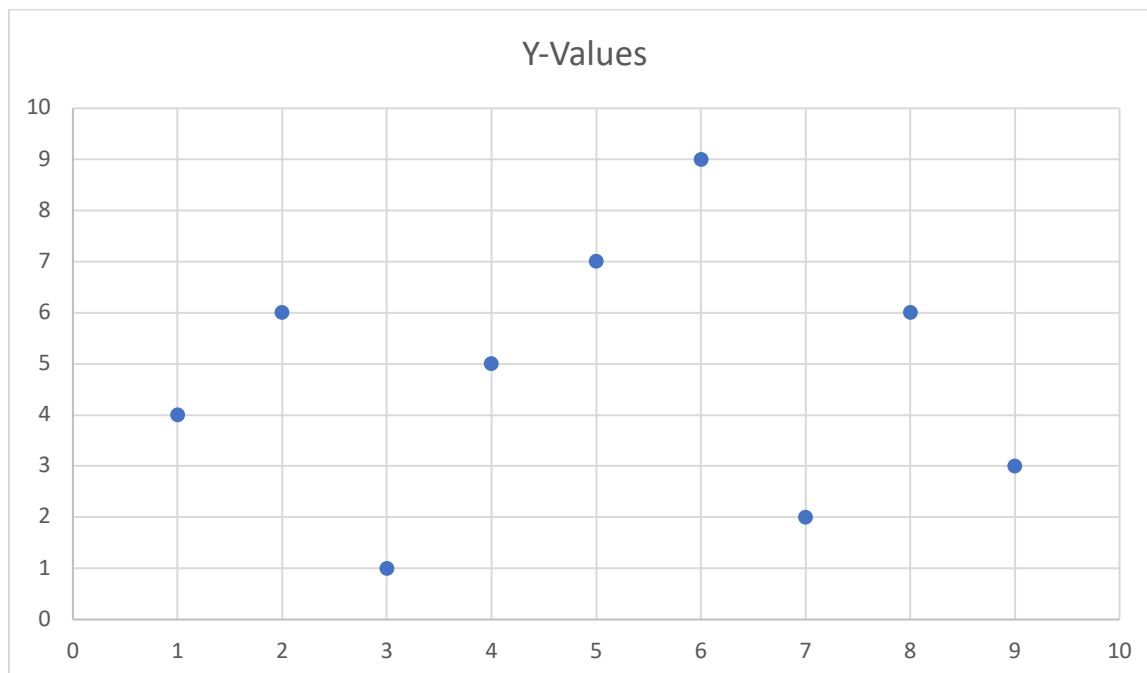
$$**$T(n)$ is $\Theta(n^3 \log(n))$**$$

Problem 11.6.5:

Let's get the maxima set from the following set of points:

$\{(7,2),(3,1),(9,3),(4,5),(1,4),(6,9),(2,6),(5,7),(8,6)\}$

We will use the knowledge learned from Section 11.1 to determine the maxima. Let's first graph the points in the set above:



From the above graph, we can see the following:

Both x and y values of Point (1,4) are less than the x and y values of Point (2,6). This means that Point (1,4) cannot be in the maxima set and will be eliminated.

Both x and y values of Point (6,9) are higher than both x and y values of Points (2,6), (3,1), (4,5), and (5,7). This means that Points (2,6), (3,1), (4,5), and (5,7) cannot be in the maxima set and will be eliminated.

Both x and y values of Points (8,6) and (9,3) are higher than both x and y values of Point (7,2). This means that Point (7,2) cannot be in the maxima set and will be eliminated.

So, the maxima set is as follows:

$\{(6,9), (8,6), (9,3)\}$