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CS590 Homework Assignment  
11: Application Exercises

Due Date: April 10, 2022

### **Problem 13.7.39:**

A customized DFS search can be performed in this case for each vertex or station, to find the stations reachable by at 4 or less links from the vertex. The DFS algorithm will be:

**Algorithm** CustDFS( $G, v, \text{depth}, S$ )

Inputs: graph  $G$ , begin vertex  $v$ , depth to traverse to,  $S$  sequence to store the result in

Output: a sequence  $S$  of vertices reachable from  $v$  in at most depth hops

count  $\leftarrow 0$

mark  $v$  as visited

if (depth = 0) then

    return  $S$

for each edge in  $G.\text{incidentEdges}()$  do

    tov  $\leftarrow G.\text{opposite}(\text{edge}, v)$

    if (! isVisited(tov)) then

$S.\text{insertLast}(\text{tov})$

$S \leftarrow \text{CustDFS}(G, \text{tov}, \text{depth} - 1, S)$

return  $S$

## **Algorithm** ComputeFourSets(G)

Input: a graph G

Output: a dictionary keyed on a vertex with values being sequences of nodes 4-reachable from that vertex

D is a blank dictionary

for each v in G.vertices() do

    D.insert(v, CustDFS(G,v,4, a new sequence))

For a graph where every vertex is at most 4 edges away from others, n complete DFS traversals will be performed practically. So the worst-case run time is  $O(n(n + m))$ . But the running time is  $O(nm)$  if the network is connected (when  $m \geq n - 1$ ).