

Name: Georges Hatem

CS 590 Homework 1: Creativity
Exercises

Due Date: January 30, 2022

Problem 1.6.39:

We have the following recurrence equation for a function $T(n)$ as follows:

$$T(n) = \begin{cases} 1 & \text{when } n = 0 \\ 2T(n-1) & \text{otherwise} \end{cases}$$

We need to prove by induction that:

$$T(n) = 2^n$$

Base Case:

From the definition of the function, for $n = 0$ we have:

$$T(0) = 1$$

Let's see if we get 1 when we plug $n = 0$ in $T(n) = 2^n$:

$$T(0) = 2^0 = 1$$

This means that $T(n) = 2^n$ works for $n = 0$

Inductive Hypothesis:

For the hypothesis, I assume that $T(n) = 2^n$ is true

Inductive Step:

So, now let's work with the other value of $T(n)$ when $T(n)$ is not equal to zero:

$$T(n) = 2T(n - 1) \text{ when } n \neq 0$$

If $T(n) = 2^n$, then $T(n + 1) = 2^{n+1}$

So, let's plug compute for $T(n+1)$ from the recursion equation when $n \neq 0$:

$$T(n + 1) = 2T(n + 1 - 1) = 2T(n)$$

To check if $T(n) = 2^n$ is true, we will substitute it in $T(n+1)$ above to check what we get as a result:

$$T(n + 1) = 2 * T(n) = 2 * (2^n) = 2^{n+1}$$

Therefore, $T(n) = 2^n$ is true

Conclusion:

$T(n) = 2^n$ satisfies the recursion equation

$$T(n) = \begin{cases} 1 & \text{when } n = 0 \\ 2T(n-1) & \text{otherwise} \end{cases}$$

as you can see from above. **Therefore, by induction**
 $T(n) = 2^n$

Problem 1.6.42:

To show that $\sum_{i=1}^n (i)^2$ is $O(n^3)$ the following below will be done:

If $n = 1$, then $i = n$. Otherwise $i < n$. Therefore,

$$i \leq n$$

So, I can write the following:

$$\sum_{i=1}^n (i)^2 \leq \sum_{i=1}^n (n^2)$$

Let's solve for $\sum_{i=1}^n (n^2)$:

First thing we take the (n^2) out of the sum:

$$\sum_{i=1}^n (n^2) = (n^2) \sum_{i=1}^n (1)$$

Then, $\sum_{i=1}^n (1)$ is as follows:

$$\sum_{i=1}^n (1) = (1 + 1 + 1 + \dots + 1)$$

This means that $\sum_{i=1}^n (1)$ is basically summing 1 n times. So, we can write the following:

$$\sum_{i=1}^n (1) = (1) * (n) = n$$

So, now getting back to solving $\sum_{i=1}^n (n^2)$:

$$\sum_{i=1}^n (n^2) = (n^2) * \sum_{i=1}^n (1) = (n^2) * (n) = n^3$$

Therefore, we can deduce the following:

$$\sum_{i=1}^n (i)^2 \leq \sum_{i=1}^n (n)^2 = (n)^2 \sum_{i=1}^n (1) = n^3$$

And this means that:

$$\sum_{i=1}^n (i^2) \text{ is } O(n^3)$$