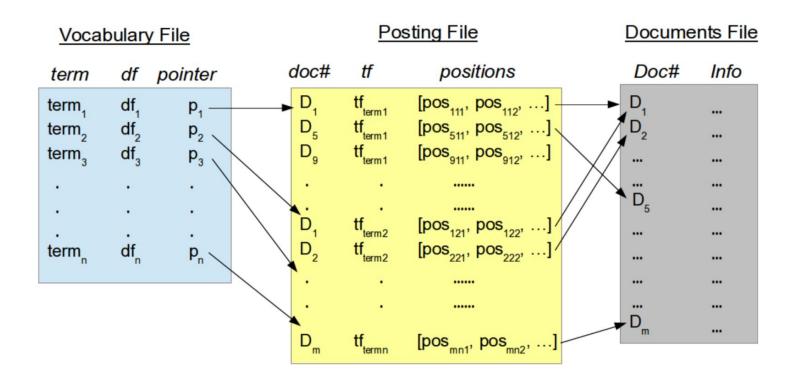
Partial Indexing and Merging

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The Inverted File



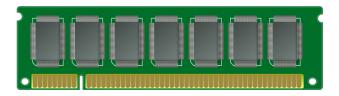


Why Partial Indexing?



- In memory index construction does not scale
- We should be able to construct indexes for very large collections
- Posting file occupies the largest amount of memory
- Think a collection of 5 gb size
 - For an inverted file 3 times the collection size, we would need 15 gb!
 - Solution: Write chunks of postings to disk and merge them afterwards to a single posting file





The Two Extremes



Write partially the inverted file and merge it with other partial inverted files to construct the final one

→ Fast and can adapt on the available resources

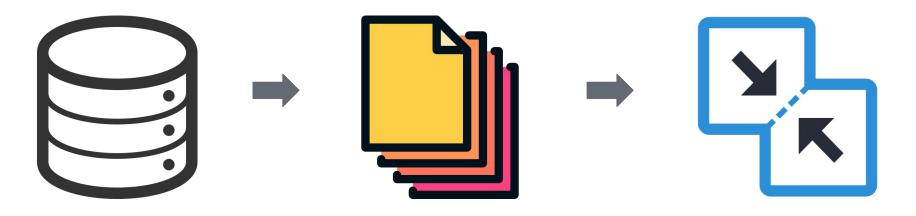
Write each record immediately when is indexed Too many I/Os

→ Low need in resources but requires lot of time to finish

Keep all inverted file to main memory and write it directly after reading the collection Requires too much resources → Extremely fast but requires huge amount of main memory

An Overview





Read the text collection once

Write the partial indices sorted when RAM is going to exceed or the inverted file reaches a certain number of records

Merge the partial indices by keeping the sorted order to construct the final inverted file

Increasing Heap Size



 In order to utilize better the main memory during indexing, you should increase the heap size of JVM

- Increase heap size from IDE
 - Intellij: https://www.jetbrains.com/help/idea/increasing-memory-heap.html
 - Netbeans:
 https://stackoverflow.com/questions/15460779/how-to-increase-the-java-h
 eap-size-in-netbeans
 - Eclipse:http://www.planetofbits.com/eclipse/increase-jvm-heap-size-in-eclipse
- Increase heap size through terminal
 - https://stackoverflow.com/questions/1565388/increase-heap-size-in-java

Partial Indexing Process - Introduction



- Choosing when the partial indices should be written
 - Option 1: Write the inverted file each time heap size is going to exceed
 - Periodically check if heap size is going to exceed. E.g. every 5 seconds check if heap size reaches 80% usage
 - https://stackoverflow.com/questions/12807797/java-get-available-memory
 - This approach could be slow so choose a period of time to ask the remaining available memory

Partial Indexing Process - Introduction



- Choosing when the partial indices should be written
 - Option 2: Choose how many records in the inverted file will be kept in memory
 - Define a good threshold Assume that:
 - Posting file contains: Tf (double = 8 bytes) and a pointer to document (long = 8 bytes)
 - Vocabulary file contains: word (string = 2 bytes/char), df (long = 8 bytes), a pointer to posting file (long = 8 bytes) and an offset (int = 4 bytes) indicating how many records to read in posting file

If a word has approximately 10 chars (=20 bytes) then we need 56 bytes for each record in vocabulary and posting file. If the heap size is 4 GB a good threshold could be 50 million records utilizing the 80% of available heap size



Partial Indexing Process - Introduction



- A partial index may contain
 - A partial vocabulary file that is sorted lexicographically based on each word and stores
 - Word
 - DF of the word
 - Pointer to the partial posting file
 - A partial posting file
 - o TF
 - A document id
 - Pointer to the document file

Partial Indexing Process - Example



Partial Vocabulary 1

| Word | DF | Pointer |
|--------------------------------------|-------------|-------------|
| Brutus Caesar Julius Killed | 1 1 1 | P P P |

Partial Posting File 1

| | TF | #Doc | Pointer |
|---|------|------|---------|
| | 1.5 | D2 | Р |
| | 2.33 | D1 | Р |
| 4 | 5.3 | D1 | P |
| | 4 | D2 | Р |
| | | | |

Partial Vocabulary 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 1 | Р |
| Caesar | 1 | Р |
| Noble | 1 | Р |
| With | 1 | Р |

Partial Posting File 2

| | TF | #Doc | Pointer |
|---|------|------|---------|
| | 2 | D3 | Р |
| | 3.33 | D4 | Р |
| | 6.3 | D3 | Р |
| • | 4.5 | D4 | Р |

Inverted Index 2

Inverted Index

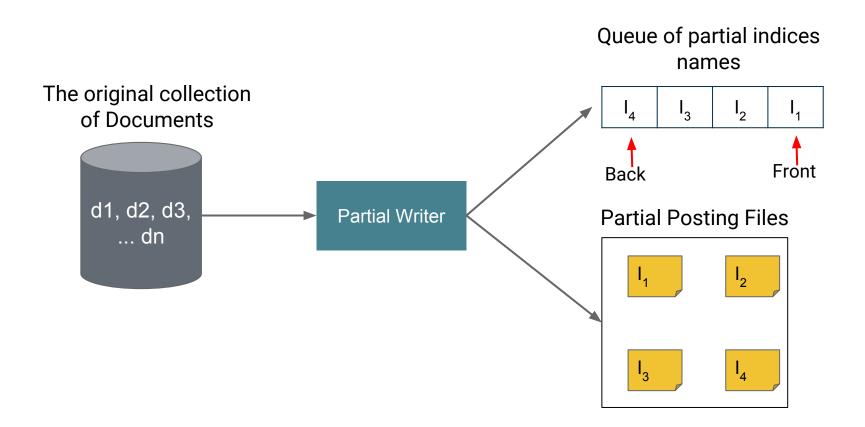
Partial Indexing Process - Pseudocode



- 1: for each document **d** in collection
- 2: for each word **w** of d
- 3: if inverted file size reaches the threshold **t** and the <u>current document ends</u>
- 4: Sort the records in current vocabulary file **lexicographically**
- 5: Start writing the current posting file and for each word keep the bytes to its first record in posting file
- 6: Write the partial vocabulary file and use the stored bytes to point in the first record of each word in its first record in posting file
- 7: Keep a data structure (e.g. a queue) in main memory to store the file names of the partial indexes
- 8: Clear the inverted file from memory
- 9: endif
- 10: Insert the word **w** to the inverted file
- 11; 12: endfor; endfor

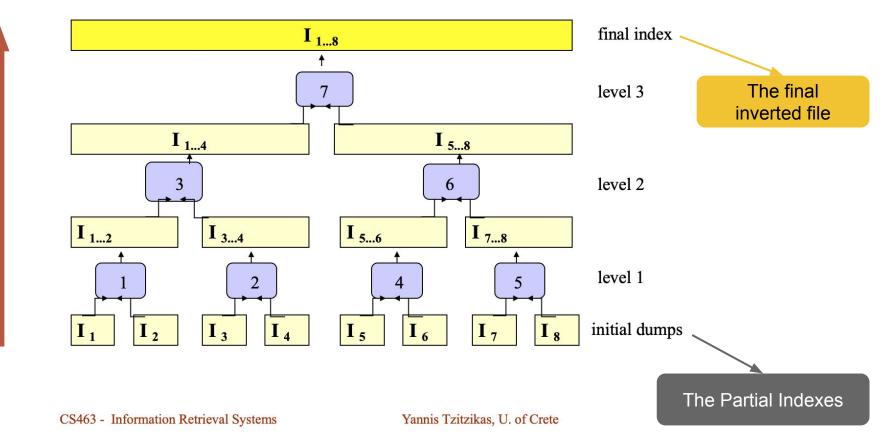
Partial Indexing Process - Output





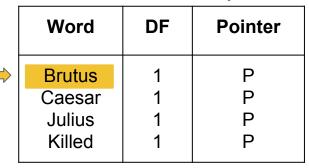
Merging Process - Overview







Partial Vocabulary 1



Partial Vocabulary 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 1 | Р |
| Caesar | 1 | Р |
| Noble | 1 | Р |
| With | 1 | Р |

Merged Vocabulary from 1 and 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 2 | P' |

| TF | #Doc | Pointer |
|--------------|----------|---------|
| 2.33 3.33 | D1 D4 | P P |



Partial Vocabulary 1

| | Word | DF | Pointer |
|---|--------|----|---------|
| | Brutus | 1 | Р |
| > | Caesar | 1 | Р |
| | Julius | 1 | Р |
| | Killed | 1 | Р |

Partial Vocabulary 2

| | Word | DF | Pointer |
|---|--------|----|---------|
| | Brutus | 1 | Р |
| • | Caesar | 1 | Р |
| | Noble | 1 | Р |
| | With | 1 | Р |

Merged Vocabulary from 1 and 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 2 | P' |
| Caesar | 2 | P' |

| TF | #Doc | Pointer |
|------|------|---------|
| 2.33 | D1 | Р |
| 3.33 | D4 | Р |
| 1.5 | D2 | Р |
| 2 | D3 | Р |
| | | |



Partial Vocabulary 1

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 1 | Р |
| Caesar | 1 | Р |
| Julius | 1 | Р |
| Killed | 1 | Р |

Partial Vocabulary 2

| | Word | DF | Pointer |
|---|--------|----|---------|
| | Brutus | 1 | Р |
| | Caesar | 1 | Р |
| • | Noble | 1 | Р |
| | With | 1 | Р |

Merged Vocabulary from 1 and 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 2 | P' |
| Caesar | 2 | P' |
| Julius | 1 | P' |

| TF | #Doc | Pointer |
|------|------|---------|
| 2.33 | D1 | Р |
| 3.33 | D4 | Р |
| 1.5 | D2 | Р |
| 2 | D3 | Р |
| 5.3 | D1 | Р |



Partial Vocabulary 1

| DF | Pointer |
|----|------------------|
| 1 | Р |
| 1 | Р |
| 1 | Р |
| 1 | Р |
| | 1 1 1 1 |

Partial Vocabulary 2

| DF | Pointer |
|----|------------------|
| 1 | P P |
| 1 | P D |
| | 1 1 1 1 |

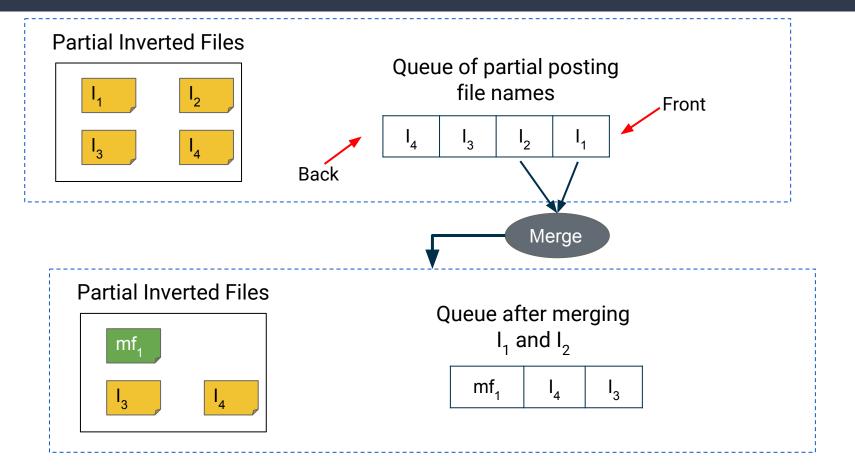
Merged Vocabulary from 1 and 2

| Word | DF | Pointer |
|--------|----|---------|
| Brutus | 2 | P' |
| Caesar | 2 | P' |
| Julius | 1 | P' |
| | | |

| TF | #Doc | Pointer |
|------|------|---------|
| 2.33 | D1 | Р |
| 3.33 | D4 | Р |
| 1.5 | D2 | Р |
| 2 | D3 | Р |
| 5.3 | D1 | Р |
| | | |

Merging Process - Schema

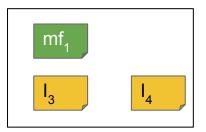




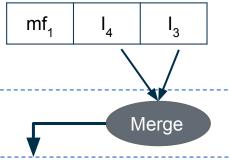
Merging Process - Schema



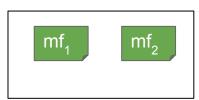
Partial Inverted Files



Queue of partial posting file names



Partial Inverted Files



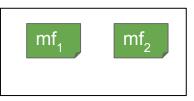
Queue after merging I_3 and I_4

| mf ₂ | mf ₁ |
|-----------------|-----------------|
|-----------------|-----------------|

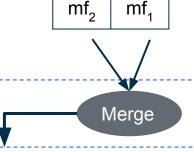
Merging Process - Schema







Queue of partial posting file names



Partial Inverted Files



Queue after merging mf₁ and mf₂

Final Inverted File

Updating Vocabulary File



- Case 1: If two words are not equal
 - Find the smaller word lexicographically
 - Write its data from partial posting in the common posting file
 - Write the data from partial vocabulary file in the common vocabulary file and move the file pointer from file with the smaller word
- Case 2: If two words are equal
 - Merge two lines and update word's DF. Add DF of the first line with DF of the second
 - Keep the word's data in main memory until a different word is discovered
 - Write the word's data in the common vocabulary file and move the file pointers of both files

Merging Process - Pseudocode



- 1: Let **Q** be a queue that keeps every file name from partial inverted indices
- 2: for each pair of partial indices **I1**, **I2** in Q
- 3: Pop I1 and I2 from Q and get the partial vocabulary and posting file names
 - for each word **w**_i of file vocab1 and **w**_i of file vocab2
- 5: if $w_i < w_i$

6:

10:

- Write posting data from the pointer to posting file of w_i vocab1 record, write w_i to the merged vocab file and move file pointer of file vocab1
- elseif w_i > w_i
 - Write posting data from the pointer to posting file of w_j vocab2 record, write w_j to the merged vocab file and move file pointer of file vocab2
- else ▷ The two words are equal
 - Keep the record in memory, update DF and write the word's data when a new word is discovered. Move the file pointers from both files
- 11: endif
- 12: Close and **delete** the files used for merging
- 13: Add the merged file name to Q
- 14; 15: endfor; endfor

Useful Tips



- Build indexes in a small collection for testing (check the partial indexing using a small threshold e.g. 10 records). Then, scale it to the full collection.
- Give the JVM at least 4 gb of heap space (if available)
- Perform plenty of tests
 - Each word in vocabulary file must have its posting data
 - Each record in posting file should correspond to a document
- Use random access file API from java to obtain the bytes and seek into posting and document file



Conclusion



- ✓ Partial indexing is a good technique to index large collections
- ✓ It is preferable to define a threshold and tune it according to the available resources.
- ✓ Keep a data structure (e.g. a queue) in main memory that stores the
 file names of partial indices and use it in the merging phase
- ✓ Be careful while merging. The lexicographic order should be preserved!