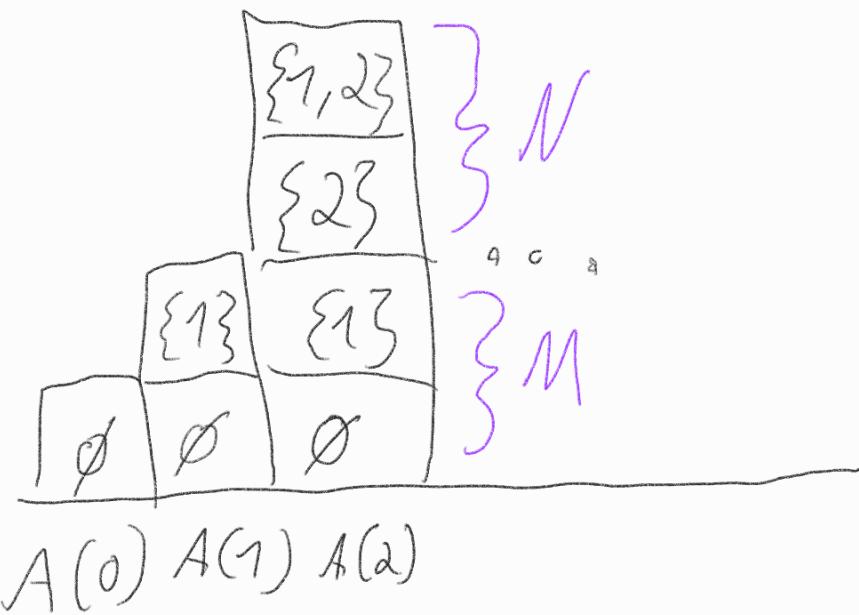


Def.

1) $A(0) = \{\emptyset\}$

$A(n+1) = \underset{\text{def}}{\cup} A(n) \cup \{X \mid X = X' \cup \{n+1\}$
 $, X' \in A(n)\}$



Satz: $\#(A(n)) = 2^n$ (IV)

(IA) $\#(A(0)) = 1 = 2^0$

(IS) $n \rightarrow n+1$

$$\begin{aligned}
 A(n+1) &= P(\{1, 2, \dots, n+1\})^N \\
 &= P(\{1, 2, \dots, n\}) \cup \{X \mid X = X' \cup \{n+1\}\} \\
 &\quad \text{with } X' \in P(\{1, \dots, n\}) \\
 M \cap N &= \emptyset
 \end{aligned}$$

$$\begin{aligned}
 &= 2 \cdot \# P(\{1, \dots, n\}) = \\
 &= 2 \cdot A(n)
 \end{aligned}$$

$$A(n+1) = 2 \cdot A(n)$$

$$\stackrel{(IV)}{=} 2 \cdot 2^n = 2^{n+1}$$

2)

i) $\Sigma = \{a, \dots, z, A, \dots, Z, 0, \dots, 9\}$

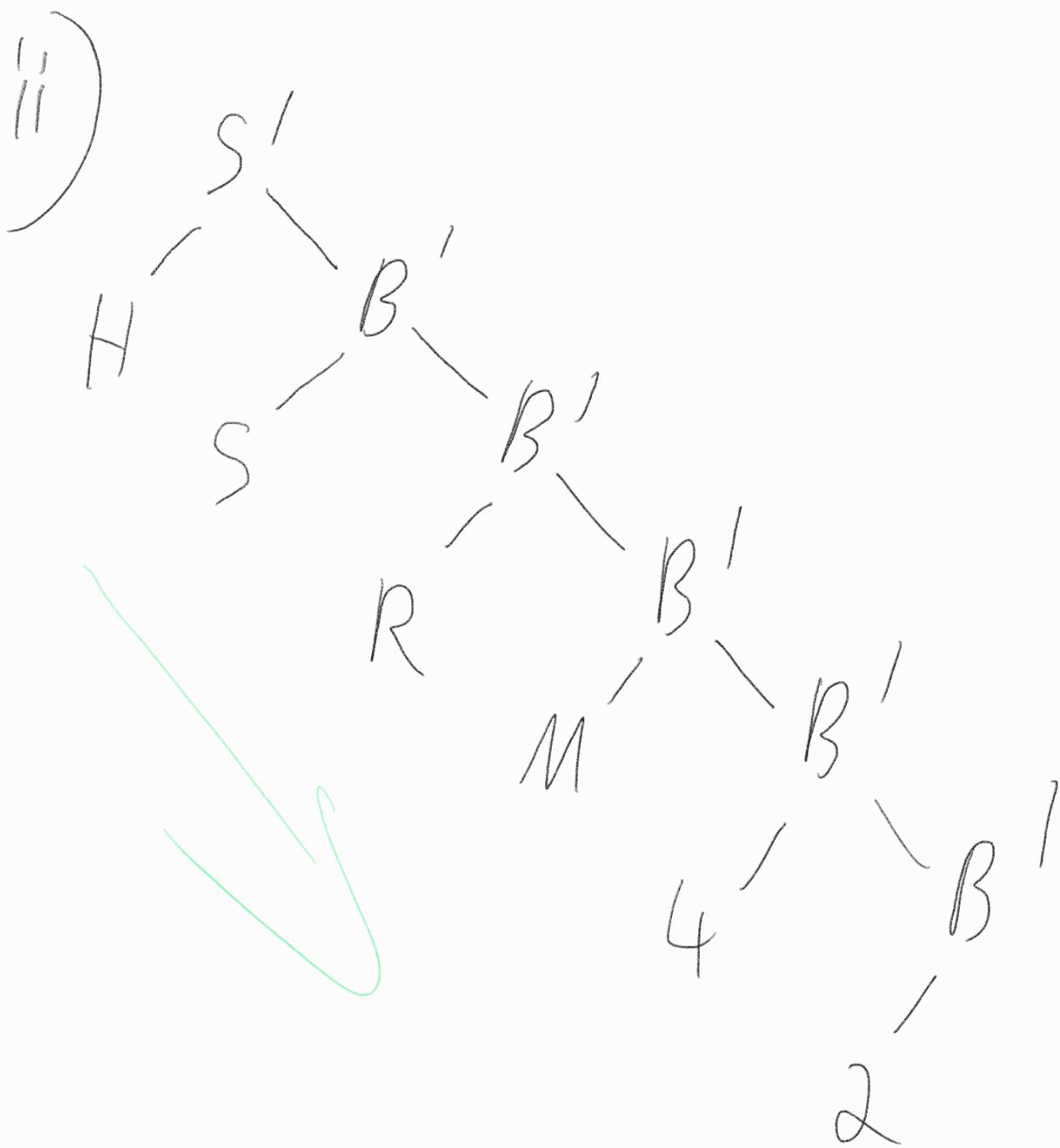
$$N = \{B'\} \quad S = \{S'\}$$

$$P' = \bigcup \{S' \rightarrow x B',$$

$$S' \rightarrow x \}$$

$$x \in \{a, \dots, z, A, \dots, Z\}$$

$$\cup \{B' \rightarrow x B', B' \rightarrow x\}$$
$$x \in \Sigma$$



3) ; - neues Nichtterminalsymbol
 R_A

- Neue Regel $A \rightarrow w$,
wobei w ein Wort mit
einem Terminal Symbol $a \in \Sigma$
- In w das a wird durch
 R_A
- Am Ende brauchen wir noch
eine Regel $R_A \rightarrow a$

$$G = (\Sigma, N, P, S)$$

ii) $L(G)$:

$P = \{A \rightarrow B_1, B_2, B_3, \dots, B_k\}, k \geq 3$

$N = \{A, B_1, B_2, \dots, B_k\}$

$L(G'')$:

$P'' = \{C \rightarrow DE\}$

$N'' = \{C, D, E\}$

dadurch wird λ^0

$A \rightarrow B_1 C_1$

$C_1 \rightarrow B_2 C_2$

$C_2 \rightarrow B_3 C_3$

\vdots

\vdots

\vdots

$C_{k-1} \rightarrow B_C$ C_k, C_1, \dots, C_{k-1}

neue Nicht-terminal-Symbole

$C_1, \dots, C_{k-1} \in N$

aber $C_1, \dots, C_k \in N'$

(4)

$L = \det \{a^n b^m c^l \mid n, m, l \geq 0\}$

$S \rightarrow \epsilon$

$S \rightarrow a L' D'$

$L' \rightarrow a L' D'$

$L' \rightarrow F' D'$, $S \rightarrow b F' D'$

$F' \rightarrow b F' D'$, $S \rightarrow F'$

$$D' \rightarrow c \quad F' \rightarrow \epsilon$$
$$F' \rightarrow b c$$

2) $S \rightarrow \#, S \rightarrow a Sa$
 $S \rightarrow b S b$