

Chapter 4: Intermediate SQL

Database System Concepts, 6th Ed.

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Chapter 4: Intermediate SQL

- Join Expressions
- Views
- Transactions
- Integrity Constraints
- SQL Data Types and Schemas
- Authorization



Joined Relations

- **Join operations** take two relations and return as a result another relation.
- A join operation is a Cartesian product which requires that tuples in the two relations match (under some condition). It also specifies the attributes that are present in the result of the join.
- The join operations are typically used as subquery expressions in the from clause.



Join operations – Example

Relation course

course_id	title	dept_name	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design		4
CS-315	Robotics	Comp. Sci.	3

Relation prereq

course_id	prereq_id
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

Note: prereq information missing for CS-315 and course information missing for CS-437.



Outer Join

- An extension of the join operation that avoids loss of information.
- Computes the join and then adds tuples form one relation that does not match tuples in the other relation to the result of the join.
- Uses null values.



Left Outer Join

course natural left outer join prereq

course_id	title	dept_name	credits	prere_id
CS-190	Game Design		4 4 2	BIO-101 CS-101
	Robotics	Comp. Sci.	3	null

Note: read prere_id as prereq_id



Right Outer Join

course natural right outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101



Full Outer Join

course natural full outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101



Joined Relations

- Join operations take two relations and return as a result another relation.
- These additional operations are typically used as subquery expressions in the **from** clause
- Join condition defines which tuples in the two relations match, and what attributes are present in the result of the join.
- **Join type** defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

inner join left outer join right outer join full outer join

```
Join Conditionsnaturalon < predicate>using (A_1, A_1, ..., A_n)
```



Joined Relations – Examples

course inner join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prere_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

course left outer join prereq on course.course_id = prereq.course_id

course_id	title	dept_name	credits	prere_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null



Joined Relations – Examples

course natural right outer join prereq

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

course right outer join prereq using (course_id)

course_id	title	dept_name	credits	prere_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101



Views

- In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

select *ID*, *name*, *dept_name* **from** *instructor*

- A view provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a "virtual relation" is called a view.



View Definition

A view is defined using the create view statement which has the form

create view v as < query expression >

- where <query expression> is any legal SQL expression. The view name is represented by *v*.
- Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- View definition is not the same as creating a new relation by evaluating the query expression
 - Rather, a view definition causes the saving of an expression;
 the expression is substituted into queries using the view.



Example Views

- A view of instructors without their salary create view faculty as select ID, name, dept_name from instructor
- Find all instructors in the Biology department select name from faculty where dept_name = 'Biology'
- Create a view of department salary totals create view departments_total_salary(dept_name, total_salary) as select dept_name, sum (salary) from instructor group by dept_name;



Views Defined Using Other Views

- create view physics_fall_2009 as
 select course.course_id, sec_id, building, room_number
 from course, section
 where course.course_id = section.course_id
 and course.dept_name = 'Physics'
 and section.semester = 'Fall'
 and section.year = '2009';
- create view physics_fall_2009_watson as select course_id, room_number from physics_fall_2009 where building= 'Watson';



View Expansion

Expand use of a view in a query/another view

```
create view physics_fall_2009_watson as
(select course_id, room_number
from (select course.course_id, building, room_number
    from course, section
    where course.course_id = section.course_id
        and course.dept_name = 'Physics'
        and section.semester = 'Fall'
        and section.year = '2009')
where building= 'Watson';
```



Views Defined Using Other Views

- One view may be used in the expression defining another view,
- A view relation v_1 is said to *depend directly* on a view relation v_2 if v_2 is used in the expression defining v_1
- A view relation v_1 is said to depend on view relation v_2 if either v_1 depends directly to v_2 or there is a path of dependencies from v_1 to v_2
- A view relation *v* is said to be *recursive* if it depends on itself.



View Expansion

- A way to define the meaning of views defined in terms of other views.
- Let view v_1 be defined by an expression e_1 that may itself contain uses of view relations.
- View expansion of an expression repeats the following replacement step:

repeat

Find any view relation v_i in e_1 Replace the view relation v_i by the expression defining v_i until no more view relations are present in e_1

As long as the view definitions are not recursive, this loop will terminate.



Update of a View

Add a new tuple to faculty view which we defined earlier insert into faculty values ('30765', 'Green', 'Music');

This insertion must be represented by the insertion of the tuple ('30765', 'Green', 'Music', null)

into the instructor relation.



Some Updates cannot be Translated Uniquely

- create view instructor_info as select ID, name, building from instructor, department where instructor.dept_name= department.dept_name;
- insert into instructor info values ('69987', 'White', 'Taylor');
 - which department, if multiple departments in Taylor?
 - what if no department is in Taylor?
- Most SQL implementations allow updates only on simple views
 - The from clause has only one database relation.
 - The select clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
 - Any attribute not listed in the select clause can be set to null
 - The query does not have a group by or having clause.



And Some Not at All

- create view history_instructors as
 select *
 from instructor
 where dept_name= 'History';
- Insert ('25566', 'Brown', 'Biology', 100000) into history_instructors



Transactions

- Unit of work
- Atomic transaction
 - either fully executed or rolled back as if it never occurred
- Isolation from concurrent transactions
- Transactions begin implicitly
 - Ended by commit work or rollback work
- But default on most databases: each SQL statement commits automatically
 - Can turn off auto commit for a session (e.g. using API)
 - In SQL:1999, can use: begin atomic end



Integrity Constraints

- Integrity constraints guard against accidental damage to the database, by ensuring that authorized changes to the database do not result in a loss of data consistency.
 - A checking account must have a balance greater than \$10,000.00.
 - A salary of a bank employee must be at least \$4.00 an hour.
 - A customer must have a (non-null) phone number.



Constraints on a Single Relation

- not null
- primary key
- unique
- **check** (P), where P is a predicate



Not Null and Unique Constraints

not null

Declare name and budget to be not null

```
name varchar(20) not null budget numeric(12,2) not null
```

- unique $(A_1, A_2, ..., A_m)$
 - The unique specification states that the attributes A1, A2, ... Am form a candidate key.
 - Candidate keys are permitted to be null (in contrast to primary keys).



The check clause

check (P) where P is a predicate Example: ensure that semester is one of fall, winter, spring or summer: create table section (course_id varchar (8), sec_id varchar (8), semester varchar (6), year numeric (4,0), building varchar (15), room_number varchar (7), time slot id varchar (4), primary key (course_id, sec_id, semester, year), **check** (semester in ('Fall', 'Winter', 'Spring', 'Summer'))



Referential Integrity

- Ensures that a value that appears in one relation for a given set of attributes also appears for a certain set of attributes in another relation.
 - Example: If "Biology" is a department name appearing in one of the tuples in the *instructor* relation, then there exists a tuple in the *department* relation for "Biology".
- Let A be a set of attributes. Let R and S be two relations that contain attributes A and where A is the primary key of S. A is said to be a **foreign key** of R if for any values of A appearing in R these values also appear in S.



Cascading Actions in Referential Integrity

```
create table course (
   course_id char(5) primary key,
             varchar(20),
   title
   dept_name varchar(20) references department
create table course (
   dept_name varchar(20),
  foreign key (dept_name) references department
         on delete cascade
         on update cascade,
```

alternative actions to cascade: set null, set default



Integrity Constraint Violation During Transactions

■ E.g.,

```
create table person (
    ID char(10),
    name char(40),
    mother char(10),
    father char(10),
    primary key ID,
    foreign key father references person,
    foreign key mother references person)
```

- How to insert a tuple?
- What if mother or father is declared not null?
 - constraint father_ref foreign key father references person,
 constraint mother_ref foreign key mother references person)
 - set constraints father_ref, mother_ref deferred



Complex Check Clauses

- check (time_slot_id in (select time_slot_id from time_slot))
 - why not use a foreign key here?
- Every section has at least one instructor teaching the section.
 - how to write this?
- Unfortunately: subquery in check clause not supported by pretty much any database
 - Alternative: triggers (later)
- create assertion <assertion-name> check cpredicate>;
 - Also not supported by anyone



Built-in Data Types in SQL

- date: Dates, containing a (4 digit) year, month and date
 - Example: date '2005-7-27'
- **time:** Time of day, in hours, minutes and seconds.
 - Example: time '09:00:30'time '09:00:30.75'
- **timestamp**: date plus time of day
 - Example: timestamp '2005-7-27 09:00:30.75'
- interval: period of time
 - Example: interval '1' day
 - Subtracting a date/time/timestamp value from another gives an interval value
 - Interval values can be added to date/time/timestamp values



Other Features

- create table student (ID varchar (5), name varchar (20) not null, dept_name varchar (20), tot_cred numeric (3,0) default 0, primary key (ID))
- create index studentID index on student(ID)
- Large objects
 - book review clob(10KB)
 - image blob(10MB)
 - movie blob(2GB)



User-Defined Types

create type construct in SQL creates user-defined type

create type Dollars as numeric (12,2) final

 create table department (dept_name varchar (20), building varchar (15), budget Dollars);



Domains

create domain construct in SQL-92 creates user-defined domain types

create domain person_name char(20) not null

- Types and domains are similar. Domains can have constraints, such as **not null**, specified on them.
- create domain degree_level varchar(10) constraint degree_level_test check (value in ('Bachelors', 'Masters', 'Doctorate'));



Large-Object Types

- Large objects (photos, videos, CAD files, etc.) are stored as a large object.
 - blob: binary large object -- object is a large collection of uninterpreted binary data (whose interpretation is left to an application outside of the database system)
 - clob: character large object -- object is a large collection of character data
 - When a query returns a large object, a pointer is returned rather than the large object itself.



Authorization

Forms of authorization on parts of the database:

- Read allows reading, but not modification of data.
- Insert allows insertion of new data, but not modification of existing data.
- Update allows modification, but not deletion of data.
- Delete allows deletion of data.

Forms of authorization to modify the database schema

- Index allows creation and deletion of indices.
- Resources allows creation of new relations.
- Alteration allows addition or deletion of attributes in a relation.
- Drop allows deletion of relations.



Authorization Specification in SQL

- The grant statement is used to confer authorization grant <pri>privilege list>
 - on <relation name or view name> to <user list>
- <user list> is:
 - a user-id
 - public, which allows all valid users the privilege granted
 - A role (more on this later)
- Granting a privilege on a view does not imply granting any privileges on the underlying relations.
- The grantor of the privilege must already hold the privilege on the specified item (or be the database administrator).



Privileges in SQL

- select: allows read access to relation, or the ability to query using the view
 - Example: grant users U_1 , U_2 , and U_3 **select** authorization on the *branch* relation:

grant select on instructor to U_1 , U_2 , U_3

- insert: the ability to insert tuples.
- update: the ability to update using the SQL update statement.
- delete: the ability to delete tuples.
- **all privileges**: used as a short form for all the allowable privileges.



Revoking Authorization in SQL

The revoke statement is used to revoke authorization.

revoke <privilege list>
on <relation name or view name> from <user list>

Example:

revoke select on branch from U_1 , U_2 , U_3

- <pri><pri><ipri>ilege-list may be all to revoke all privileges the revokee may hold.
- If <revokee-list> includes **public**, all users lose the privilege except those granted it explicitly.
- If the same privilege was granted twice to the same user by different grantees, the user may retain the privilege after the revocation.
- All privileges that depend on the privilege being revoked are also revoked.



Roles

- **create role** instructor,
- Privileges can be granted to roles:
 - grant select on takes to instructor,
- Roles can be granted to users, as well as to other roles
 - create role student
 - grant instructor to Amit;
 - create role dean;
 - grant instructor to dean;
 - grant dean to Satoshi;



Authorization on Views

- create view geo_instructor as
 (select *
 from instructor
 where dept_name = 'Geology');
- **grant select on** *geo_instructor* **to** *staff*
- Suppose that a staff member issues
 - select * from geo_instructor,
- What if
 - staff does not have permissions on instructor?
 - creator of view did not have some permissions on instructor?



Other Authorization Features

- references privilege to create foreign key
 - grant reference (dept_name) on department to Mariano;
 - why is this required?
- transfer of privileges
 - grant select on department to Amit with grant option;
 - revoke select on department from Amit, Satoshi cascade;
 - revoke select on department from Amit, Satoshi restrict;
- Etc. read Section 4.6 for more details we have omitted here.



End of Chapter 4

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ID	name	dept_name	tot_cred
00128	Zhang	Comp. Sci.	102
12345	Shankar	Comp. Sci.	32
19991	Brandt	History	80
23121	Chavez	Finance	110
44553	Peltier	Physics	56
45678	Levy	Physics	46
54321	Williams	Comp. Sci.	54
55739	Sanchez	Music	38
70557	Snow	Physics	0
76543	Brown	Comp. Sci.	58
76653	Aoi	Elec. Eng.	60
98765	Bourikas	Elec. Eng.	98
98988	Tanaka	Biology	120



ID	course_id	sec_id	semester	year	grade
00128	CS-101	1	Fall	2009	A
00128	CS-347	1	Fall	2009	A-
12345	CS-101	1	Fall	2009	C
12345	CS-190	2	Spring	2009	Α
12345	CS-315	1	Spring	2010	Α
12345	CS-347	1	Fall	2009	Α
19991	HIS-351	1	Spring	2010	В
23121	FIN-201	1	Spring	2010	C+
44553	PHY-101	1	Fall	2009	B-
45678	CS-101	1	Fall	2009	F
45678	CS-101	1	Spring	2010	B+
45678	CS-319	1	Spring	2010	В
54321	CS-101	1	Fall	2009	A-
54321	CS-190	2	Spring	2009	B+
55739	MU-199	1	Spring	2010	A-
76543	CS-101	1	Fall	2009	Α
76543	CS-319	2	Spring	2010	Α
76653	EE-181	1	Spring	2009	C
98765	CS-101	1	Fall	2009	C-
98765	CS-315	1	Spring	2010	В
98988	BIO-101	1	Summer	2009	Α
98988	BIO-301	1	Summer	2010	null



ID	name	dept_name	tot_cred	course_id	sec_id	semester	year	grade
00128	Zhang	Comp. Sci.	102	CS-101	1	Fall	2009	A
00128	Zhang	Comp. Sci.	102	CS-347	1	Fall	2009	A-
12345	Shankar	Comp. Sci.	32	CS-101	1	Fall	2009	C
12345	Shankar	Comp. Sci.	32	CS-190	2	Spring	2009	A
12345	Shankar	History	32	CS-315	1	Spring	2010	A
12345	Shankar	Finance	32	CS-347	1	Fall	2009	A
19991	Brandt	Music	80	HIS-351	1	Spring	2010	В
23121	Chavez	Physics	110	FIN-201	1	Spring	2010	C+
44553	Peltier	Physics	56	PHY-101	1	Fall	2009	B-
45678	Levy	Physics	46	CS-101	1	Fall	2009	F
45678	Levy	Physics	46	CS-101	1	Spring	2010	B+
45678	Levy	Physics	46	CS-319	1	Spring	2010	В
54321	Williams	Comp. Sci.	54	CS-101	1	Fall	2009	A-
54321	Williams	Comp. Sci.	54	CS-190	2	Spring	2009	B+
55739	Sanchez	Music	38	MU-199	1	Spring	2010	A-
76543	Brown	Comp. Sci.	58	CS-101	1	Fall	2009	A
76543	Brown	Comp. Sci.	58	CS-319	2	Spring	2010	A
76653	Aoi	Elec. Eng.	60	EE-181	1	Spring	2009	C
98765	Bourikas	Elec. Eng.	98	CS-101	1	Fall	2009	C-
98765	Bourikas	Elec. Eng.	98	CS-315	1	Spring	2010	В
98988	Tanaka	Biology	120	BIO-101	1	Summer	2009	A
98988	Tanaka	Biology	120	BIO-301	1	Summer	2010	null



ID	name	dept_name	tot_cred	course_id	sec_id	semester	year	grade
00128	Zhang	Comp. Sci.	102	CS-101	1	Fall	2009	A
00128	Zhang	Comp. Sci.	102	CS-347	1	Fall	2009	A-
12345	Shankar	Comp. Sci.	32	CS-101	1	Fall	2009	C
12345	Shankar	Comp. Sci.	32	CS-190	2	Spring	2009	A
12345	Shankar	History	32	CS-315	1	Spring	2010	A
12345	Shankar	Finance	32	CS-347	1	Fall	2009	Α
19991	Brandt	Music	80	HIS-351	1	Spring	2010	В
23121	Chavez	Physics	110	FIN-201	1	Spring	2010	C+
44553	Peltier	Physics	56	PHY-101	1	Fall	2009	B-
45678	Levy	Physics	46	CS-101	1	Fall	2009	F
45678	Levy	Physics	46	CS-101	1	Spring	2010	B+
45678	Levy	Physics	46	CS-319	1	Spring	2010	В
54321	Williams	Comp. Sci.	54	CS-101	1	Fall	2009	A-
54321	Williams	Comp. Sci.	54	CS-190	2	Spring	2009	B+
55739	Sanchez	Music	38	MU-199	1	Spring	2010	A-
70557	Snow	Physics	0	null	null	null	null	null
76543	Brown	Comp. Sci.	58	CS-101	1	Fall	2009	Α
76543	Brown	Comp. Sci.	58	CS-319	2	Spring	2010	Α
76653	Aoi	Elec. Eng.	60	EE-181	1	Spring	2009	C
98765	Bourikas	Elec. Eng.	98	CS-101	1	Fall	2009	C-
98765	Bourikas	Elec. Eng.	98	CS-315	1	Spring	2010	В
98988	Tanaka	Biology	120	BIO-101	1	Summer	2009	Α
98988	Tanaka	Biology	120	BIO-301	1	Summer	2010	null



ID	course_id	sec_id	semester	year	grade	пате	dept_name	tot_cred
00128	CS-101	1	Fall	2009	A	Zhang	Comp. Sci.	102
00128	CS-347	1	Fall	2009	A-	Zhang	Comp. Sci.	102
12345	CS-101	1	Fall	2009	C	Shankar	Comp. Sci.	32
12345	CS-190	2	Spring	2009	Α	Shankar	Comp. Sci.	32
12345	CS-315	1	Spring	2010	Α	Shankar	History	32
12345	CS-347	1	Fall	2009	Α	Shankar	Finance	32
19991	HIS-351	1	Spring	2010	В	Brandt	Music	80
23121	FIN-201	1	Spring	2010	C+	Chavez	Physics	110
44553	PHY-101	1	Fall	2009	B-	Peltier	Physics	56
45678	CS-101	1	Fall	2009	F	Levy	Physics	46
45678	CS-101	1	Spring	2010	B+	Levy	Physics	46
45678	CS-319	1	Spring	2010	В	Levy	Physics	46
54321	CS-101	1	Fall	2009	A-	Williams	Comp. Sci.	54
54321	CS-190	2	Spring	2009	B+	Williams	Comp. Sci.	54
55739	MU-199	1	Spring	2010	A-	Sanchez	Music	38
70557	null	null	null	null	null	Snow	Physics	0
76543	CS-101	1	Fall	2009	Α	Brown	Comp. Sci.	58
76543	CS-319	2	Spring	2010	Α	Brown	Comp. Sci.	58
76653	EE-181	1	Spring	2009	C	Aoi	Elec. Eng.	60
98765	CS-101	1	Fall	2009	C-	Bourikas	Elec. Eng.	98
98765	CS-315	1	Spring	2010	В	Bourikas	Elec. Eng.	98
98988	BIO-101	1	Summer	2009	Α	Tanaka	Biology	120
98988	BIO-301	1	Summer	2010	null	Tanaka	Biology	120



ID	name	dept_name	salary
10101	Srinivasan	Comp. Sci.	65000
12121	Wu	Finance	90000
15151	Mozart	Music	40000
22222	Einstein	Physics	95000
32343	El Said	History	60000
33456	Gold	Physics	87000
45565	Katz	Comp. Sci.	75000
58583	Califieri	History	62000
76543	Singh	Finance	80000
76766	Crick	Biology	72000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
69987	White	null	null

instructor

dept_name	building	budget
Biology	Watson	90000
Comp. Sci.	Taylor	100000
Elec. Eng.	Taylor	85000
Finance	Painter	120000
History	Painter	50000
Music	Packard	80000
Physics	Watson	70000
null	Painter	null

department



Join types

inner join left outer join right outer join full outer join

Join conditions

natural on < predicate> **using** $(A_1, A_2, ..., A_n)$



