Finding Communities

Social Networks Analysis and Graph Algorithms

Prof. Carlos Castillo — https://chato.cl/teach

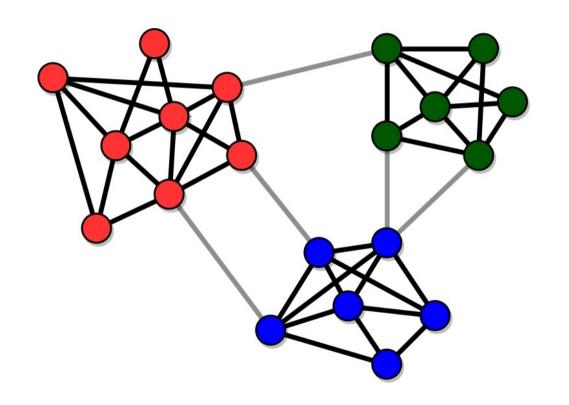


Sources

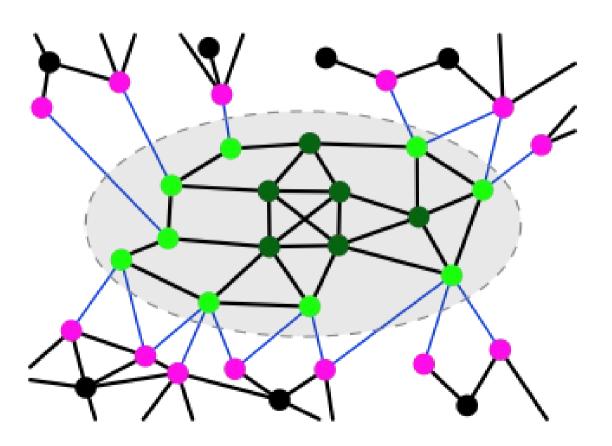
- A. L. Barabási (2016). Network Science Chapter 09
- D. Easly and J. Kleinberg (2010). Networks, Crowds, and Markets
 Chapter 03
- F. Menczer, S. Fortunato, C. A. Davis (2020). A First Course in Network Science Chapter 06
- URLs cited in the footer of slides

Defining communities

Example with clear community structure



Communities are connected and dense



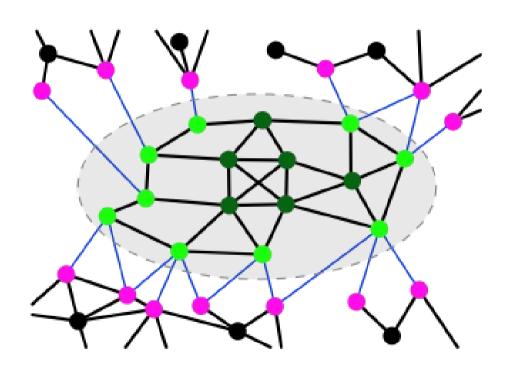
Given a community C

Internal degree $k^{int}(C)$ considers only nodes inside the community

External degree $k^{ext}(C)$ considers only nodes outside the community

$$k_i = k_i^{\text{int}}(C) + k_i^{\text{ext}}(C)$$

Strong community

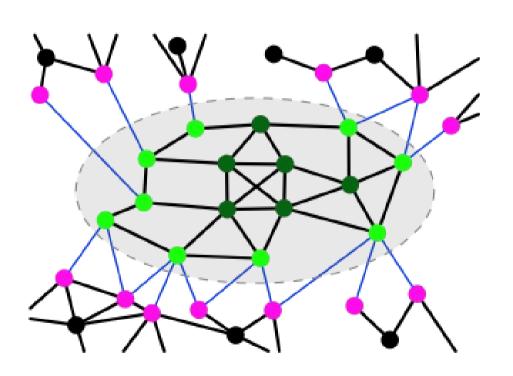


A community C is **strong** if **every node** *i* within the community satisfies:

$$k_i^{\text{int}}(C) > k_i^{\text{ext}}(C)$$

- Is the community of green nodes (dark green and light green) a strong community?
- What is the difference between dark green and light green nodes?

Weak community



A community C is **weak** if **on aggregate** nodes satisfy:

$$\sum_{i \in C} k_i^{\text{int}}(C) > \sum_{i \in C} k_i^{\text{ext}}(C)$$

 All communities satisfying the strong property satisfy the weak one

Exercise

A community C is **strong** if, for all nodes *i* within the community:

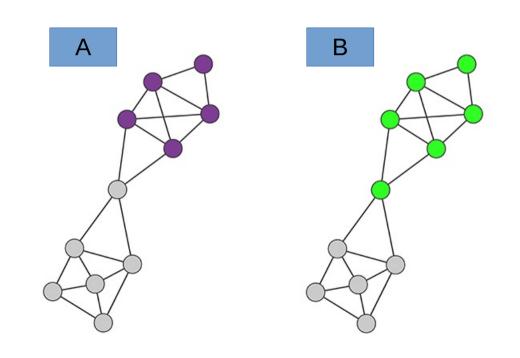
$$k_i^{\text{int}}(C) > k_i^{\text{ext}}(C)$$

A community C is **weak** if:

$$\sum_{i \in C} k_i^{\text{int}}(C) > \sum_{i \in C} k_i^{\text{ext}}(C)$$

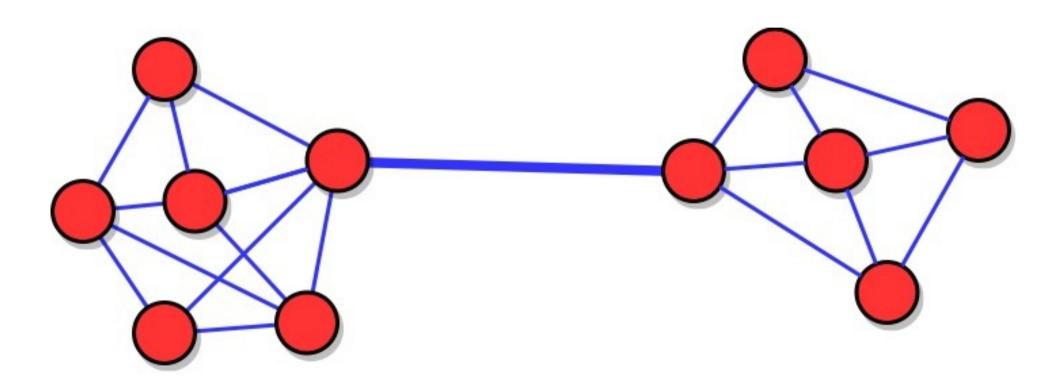
Is community A strong, weak, both?

Is community B strong, weak, both?

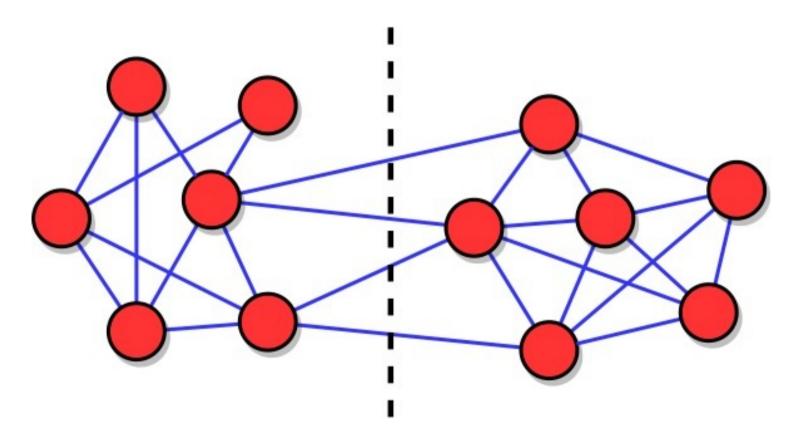


Network bisection

A graph that is easy to bisect



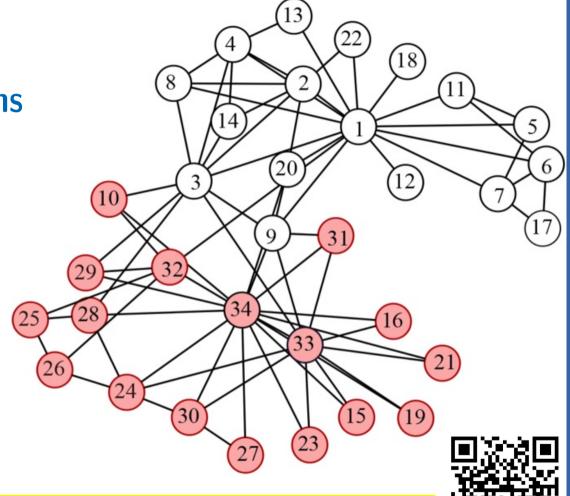
Graph bisection: finding a minimal "cut"



Simple exercise

Cut size under two partitions

- What is the size of the white-red cut?
- If node 9 goes to the red component, what is the size of the white-red cut?



Pin board: https://upfbarcelona.padlet.org/chato/4qz0k8ro0zquen1

A divisive method

Hierarchical graph partitioning

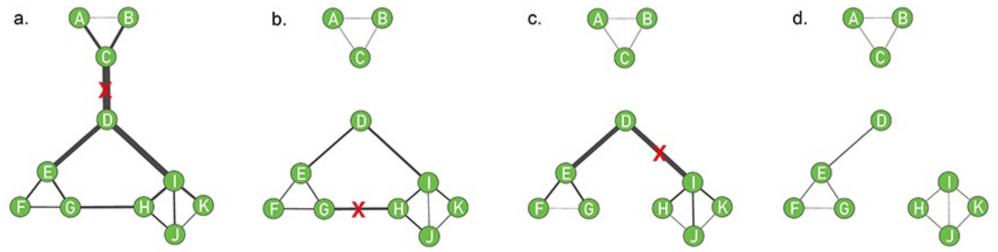
Until there are edges in the graph

Find an edge e that bridges two communities

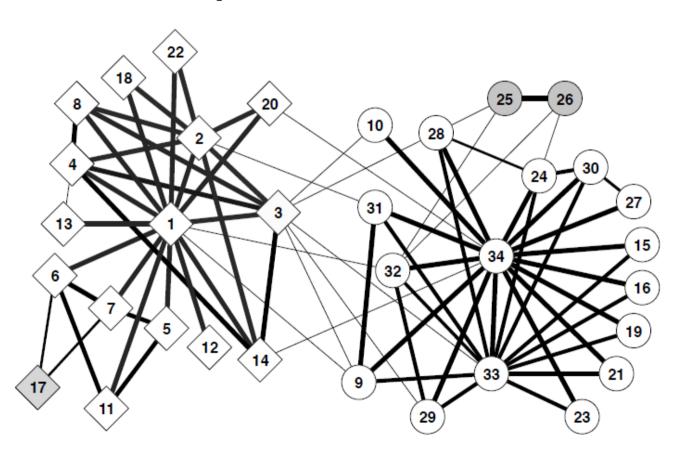
Remove edge e

The Girvan-Newman algorithm

- Repeat:
 - Compute edge betweenness
 - Remove edge with larger betweenness



Example: Karate Club



https://slidetodoc.com/online-social-networks-and-media-community-detection-1/

Modularity

Measuring a partition in a graph

- Modularity (or one of its variants) is a popular method to determine how good a partition is on a graph
- It compares the observed number of internal links in each partition, against the expected number of internal links if those internal links had been placed at random

Modularity of a partition

$$Q = \frac{1}{L} \sum_{C} \left(L_C - \frac{k_C^2}{4L} \right)$$

- L = number of links in the network
- L_c = number of internal links in community C
- k_C = sum of degree of nodes in C

Modularity of a partition (cont.)

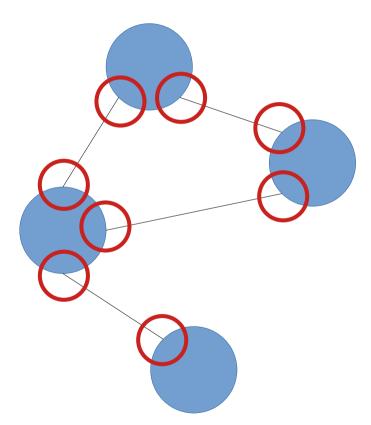
$$Q = \frac{1}{L} \sum_{C} \left(L_C - \frac{k_C^2}{4L} \right) \longrightarrow$$

Expression in parenthesis is the difference between observed and expected internal links in community *C*

- L = number of links in the network
- L_c = number of internal links in community C
- k_C = sum of degree of nodes in C
- $k^2_C/4L =$ expected number of internal links in community C

Where does $k_c^2/4L$ comes from?

- A link "stub" is a connection between a link and a node
- There are 2L stubs in a network
- There are as many stubs as the sum of the degree of nodes



Modularity formula explained $Q = \frac{1}{L} \sum \left(L_C - \frac{k_C^2}{4L} \right)$

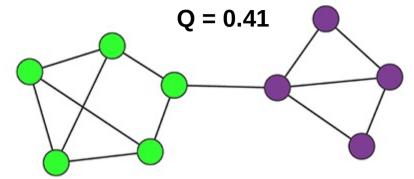
$$Q = rac{1}{L} \sum_{C} \left(L_C - rac{k_C^2}{4L}
ight)$$

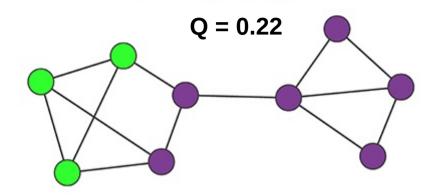
- There are L_C internal links in C
- Total number of stubs in nodes in C is k_c
- Total number of stubs in the network is 2L
- Probability of chosing two stubs in C: $(k_c/L)^2 = k_c^2/4L^2$
- The expected number of links joining two stubs in C is $L(k_c^2/4L^2) = k_c^2/4L$
- The observed number is L_c

OPTIMAL PARTITION a.

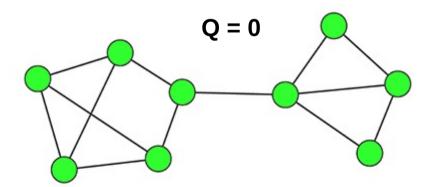


SUBOPTIMAL PARTITION

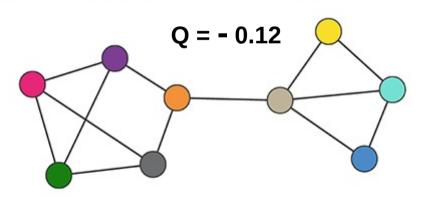




SINGLE COMMUNITY C.



NEGATIVE MODULARITY d.



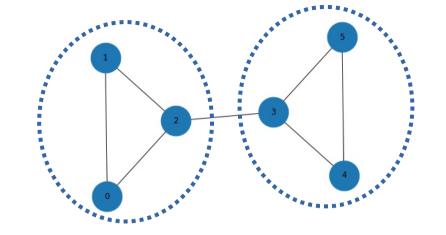


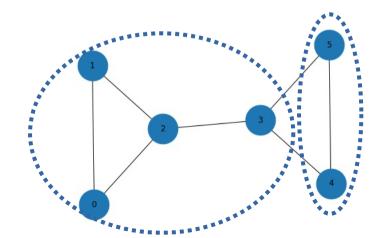
b.

Exercise

- What is the modularity of the partition {0, 1, 2}, {3, 4, 5}?
- What is the modularity of the partition {0, 1, 2, 3}, {4, 5}?

$$Q = \frac{1}{L} \sum_{C} \left(L_C - \frac{k_C^2}{4L} \right)$$





Summary

Things to remember

- Strong and weak community
- The concept of "cut" in graph bisection
- Girvan-Newman's algorithm
- Modularity

Practice on your own

- Check the modularity computations in the example on the slide marked ★: (a) optional partitioning into two communities, (b) suboptimal partitioning into two communities, (c) all the nodes in a single community, (d) one community per node
- You can check your answers with networkx.algorithms.community.modularity