

# Finding Communities

## Social Networks Analysis and Graph Algorithms

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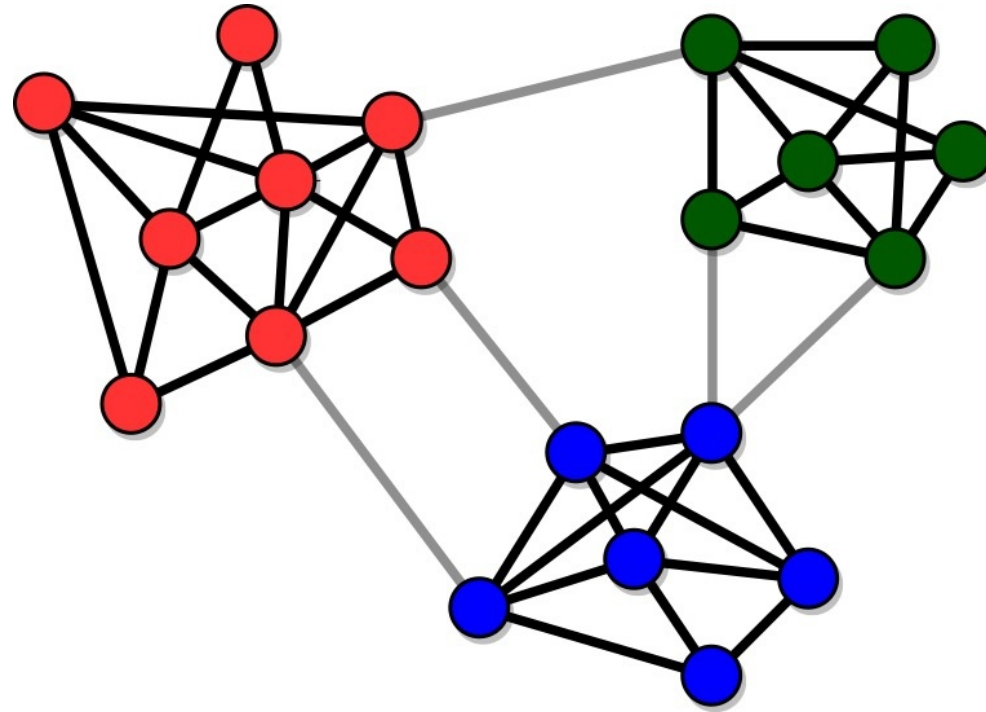


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*Barcelona*

# Sources

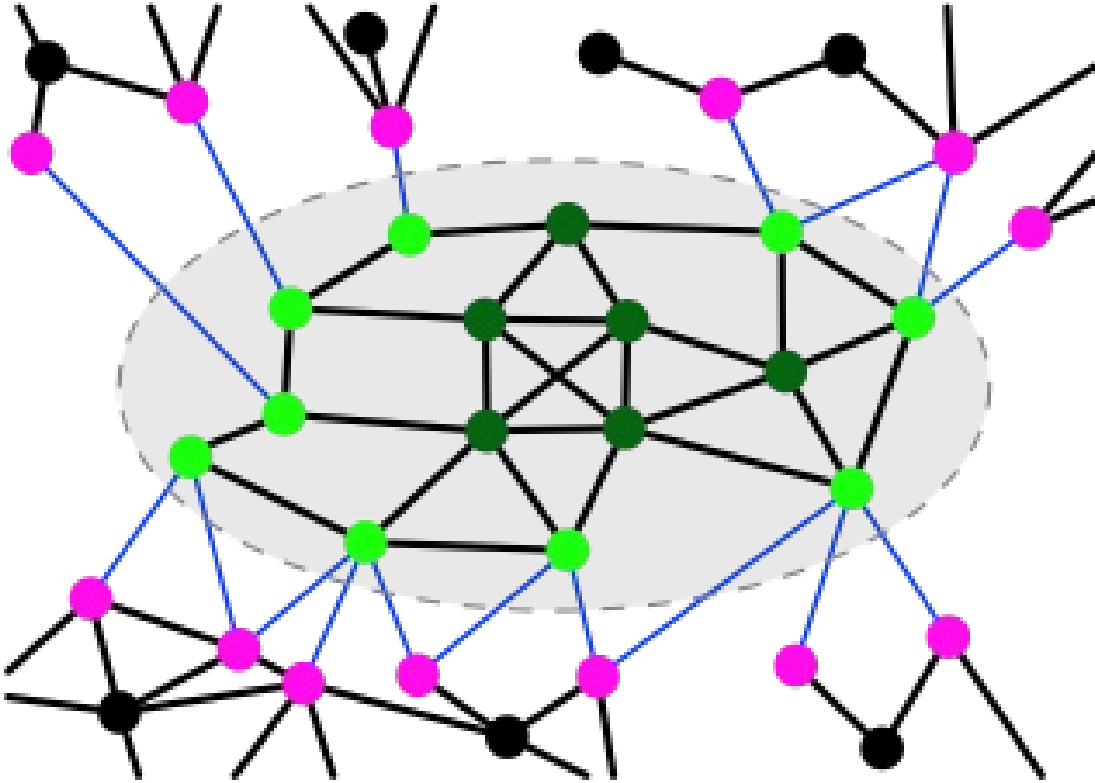
- A. L. Barabási (2016). Network Science – [Chapter 09](#)
- D. Easley and J. Kleinberg (2010). Networks, Crowds, and Markets – [Chapter 03](#)
- F. Menczer, S. Fortunato, C. A. Davis (2020). A First Course in Network Science – Chapter 06
- URLs cited in the footer of slides

# Example with clear community structure



# Characterizing one community

# Communities are **connected** and **dense**



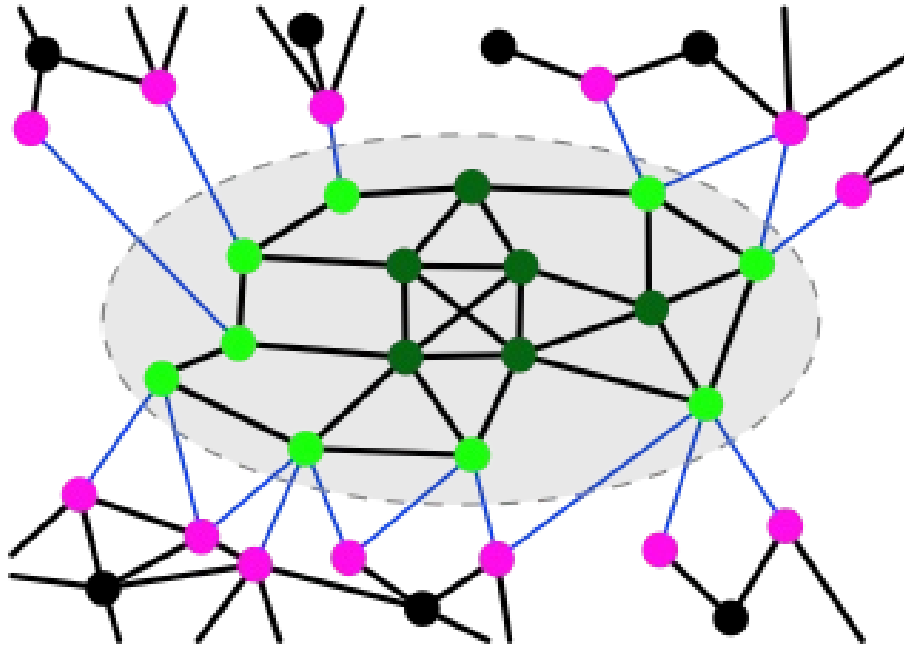
Given a community  $C$

**Internal degree**  $k^{int}(C)$  considers only nodes inside the community

**External degree**  $k^{ext}(C)$  considers only nodes outside the community

$$k_i = k_i^{int}(C) + k_i^{ext}(C)$$

# Strong community

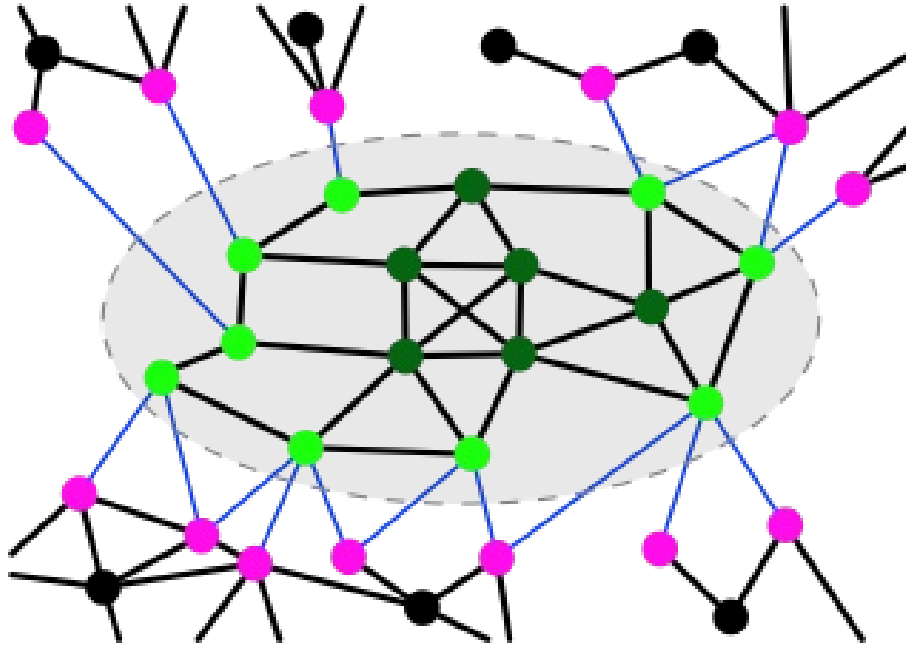


A community  $C$  is **strong** if **every** **node**  $i$  within the community satisfies:

$$k_i^{\text{int}}(C) > k_i^{\text{ext}}(C)$$

- Is the community of green nodes (dark green and light green) a strong community?
- What is the difference between dark green and light green nodes?

# Weak community



A community  $C$  is **weak** if **on aggregate** nodes satisfy:

$$\sum_{i \in C} k_i^{\text{int}}(C) > \sum_{i \in C} k_i^{\text{ext}}(C)$$

- All communities satisfying the strong property satisfy the weak one

# Exercise

A community  $C$  is **strong** if, for all nodes  $i$  within the community:

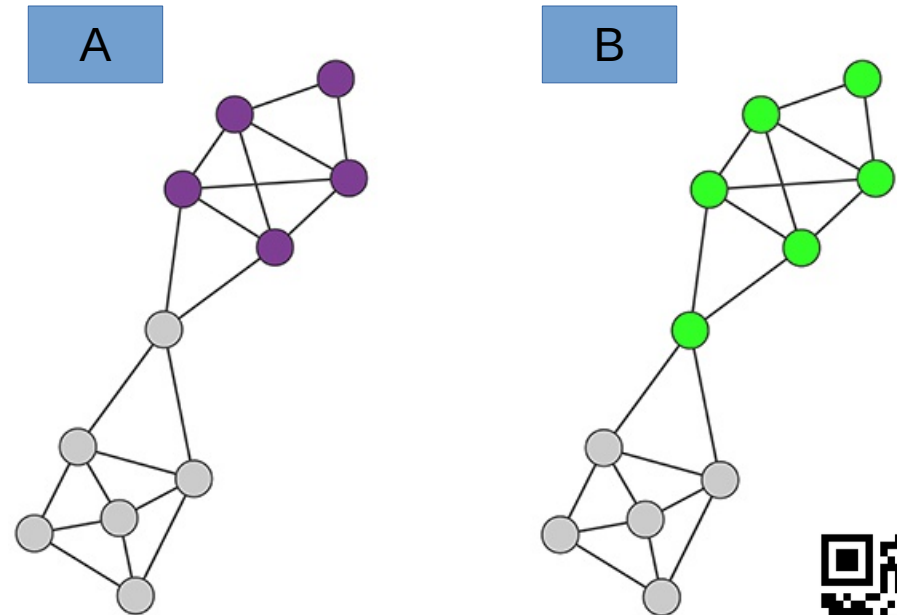
$$k_i^{\text{int}}(C) > k_i^{\text{ext}}(C)$$

A community  $C$  is **weak** if:

$$\sum_{i \in C} k_i^{\text{int}}(C) > \sum_{i \in C} k_i^{\text{ext}}(C)$$

Is **community A** strong, weak, both?

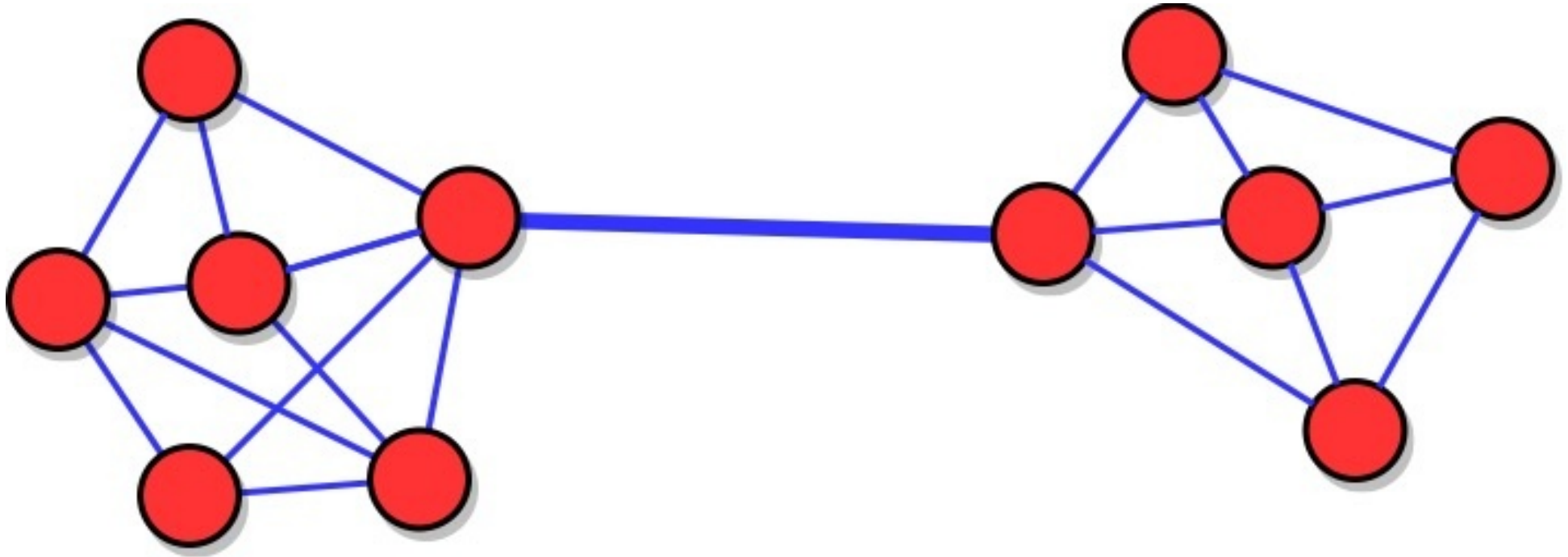
Is **community B** strong, weak, both?



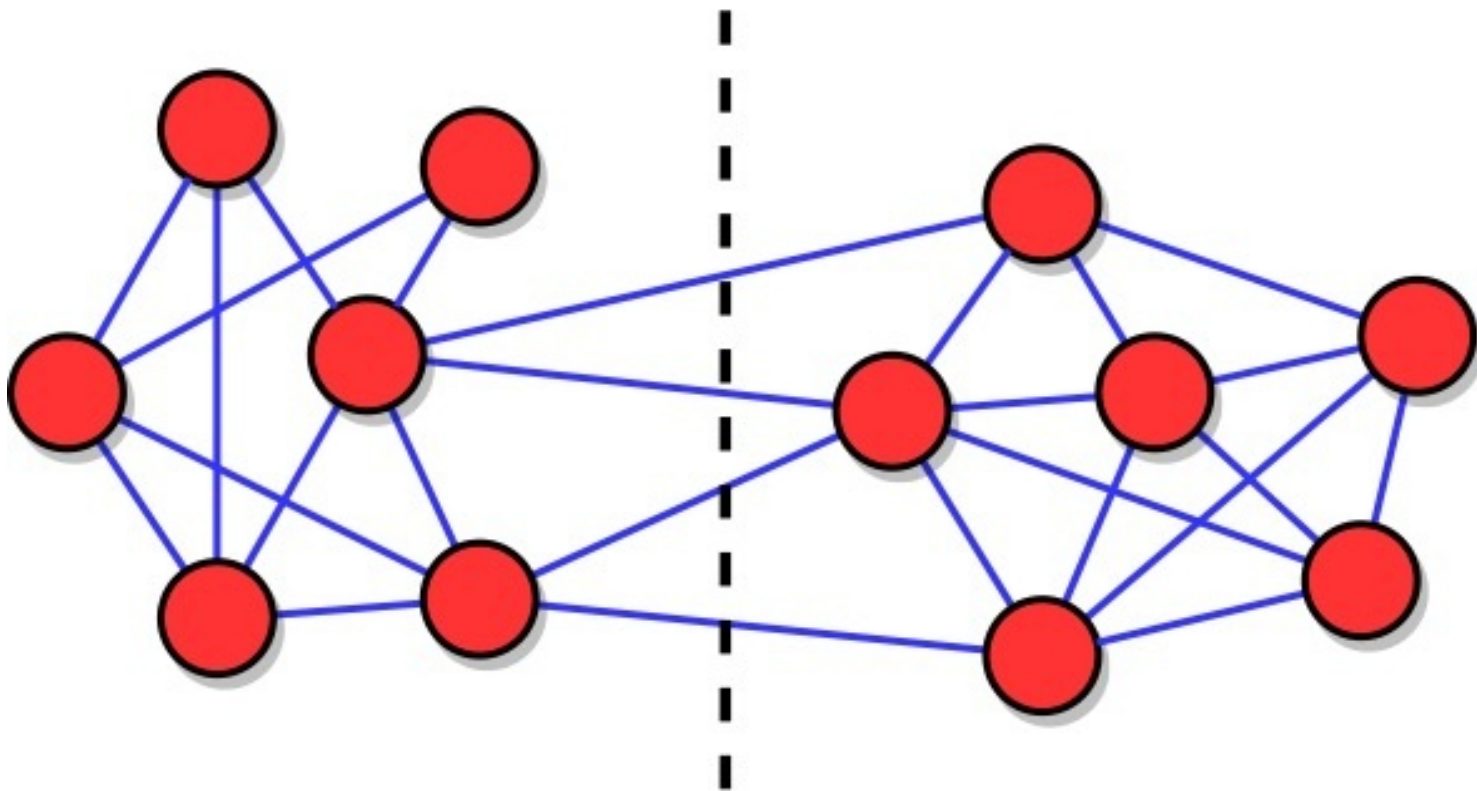


# Finding two communities: network bisection

# A graph that is easy to bisect



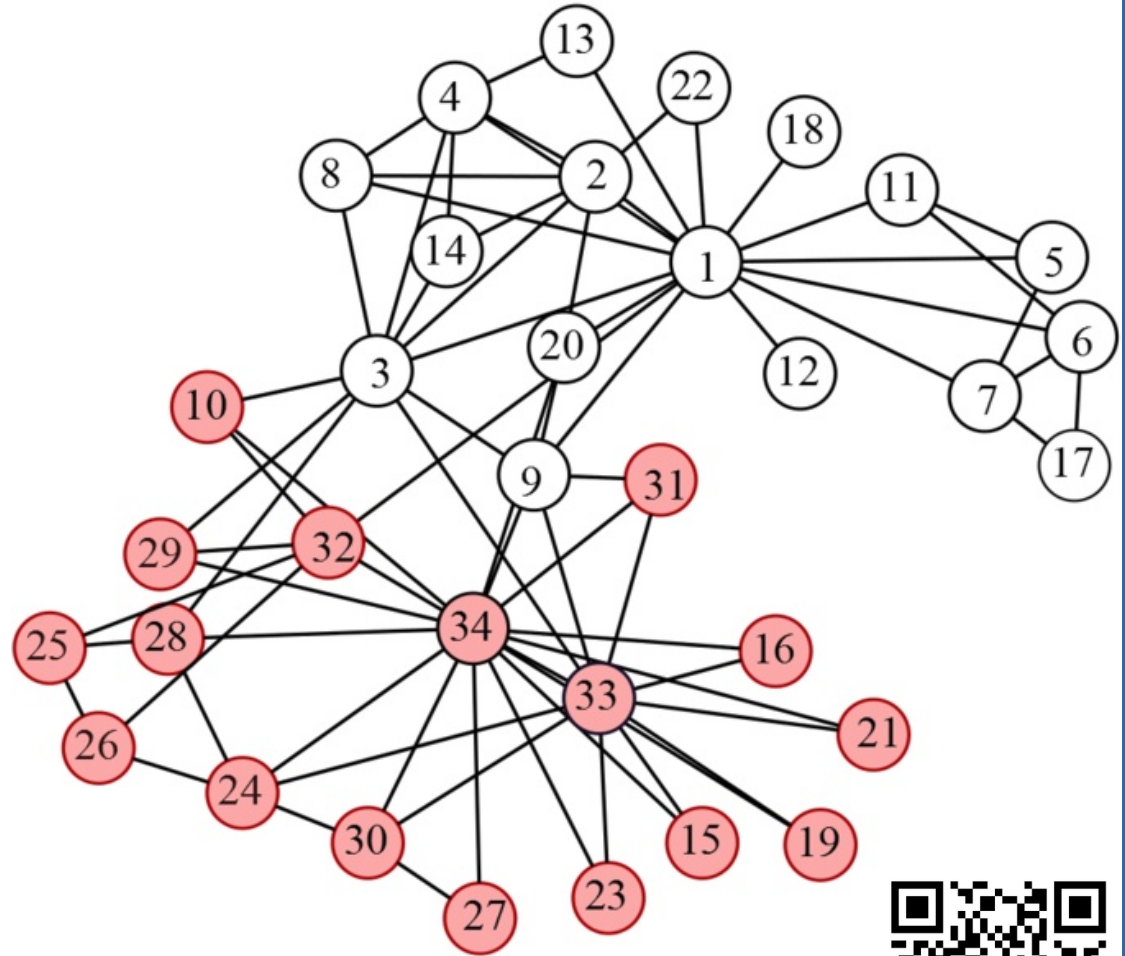
# Graph bisection: finding a minimal “cut”



# Simple exercise

## Cut size under bisection

- What is the size of the white-red cut?
- If node 9 goes to the red component, what is the size of the white-red cut?



Pin board: <https://upfbarcelona.padlet.org/chato/4qz0k8ro0zquen1>

# Finding multiple communities: a divisive method

# Hierarchical graph partitioning

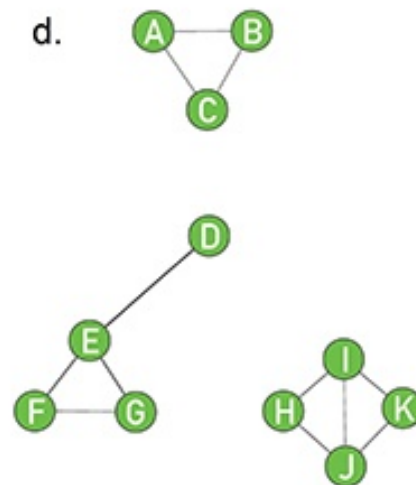
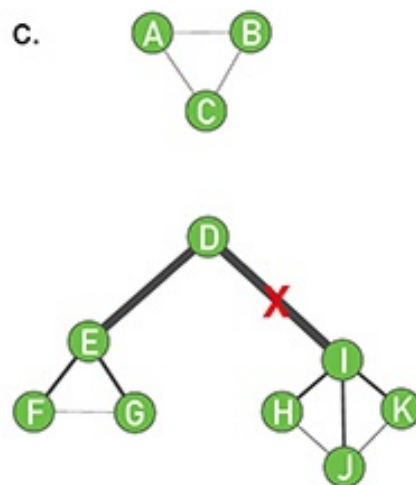
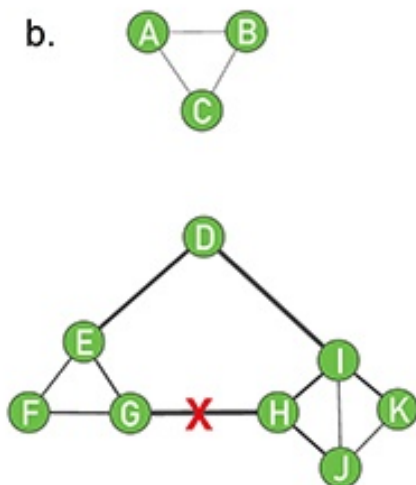
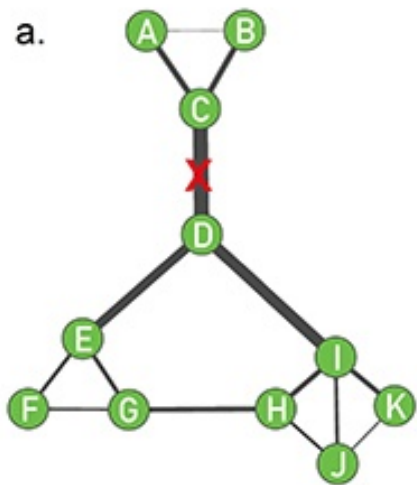
Until there are edges in the graph

Find an edge  $e$  that bridges two communities

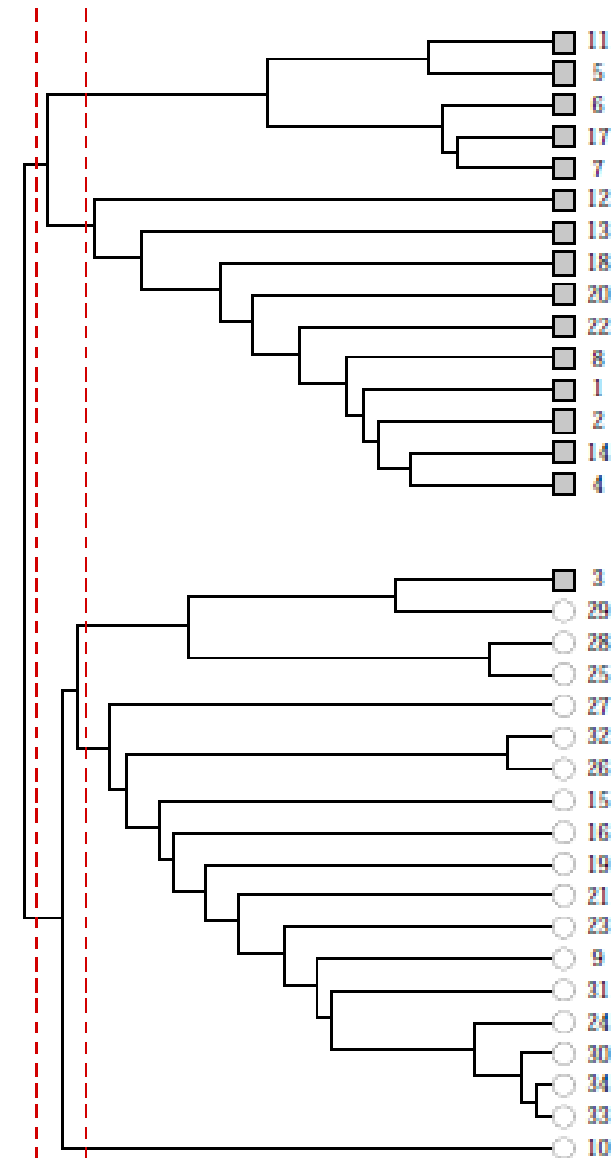
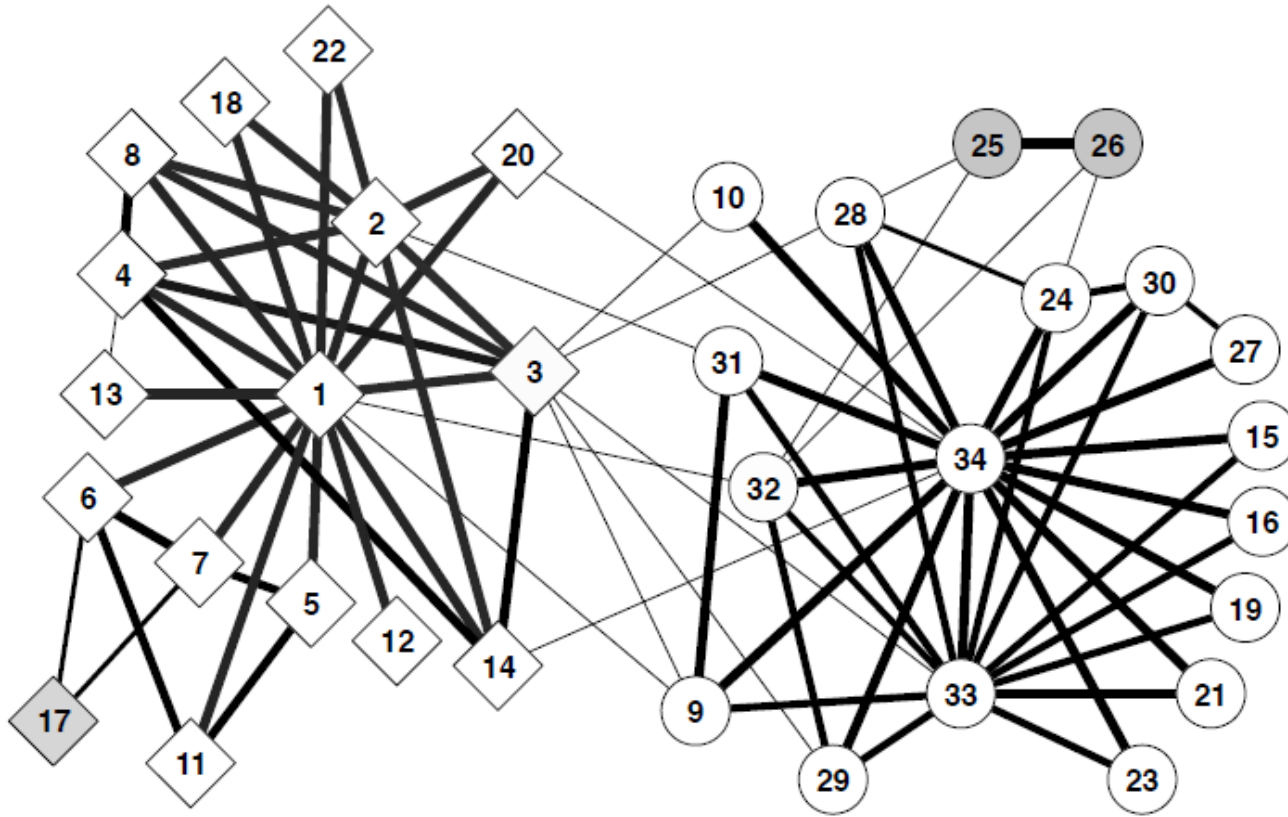
Remove edge  $e$

# The Girvan-Newman algorithm

- Repeat:
  - Compute edge betweenness
  - Remove edge with larger betweenness



# Example: Karate Club





# Quantifying multiple communities: modularity

# Measuring a partition in a graph

- **Modularity** (or one of its variants) is a popular method to determine how good a partition is on a graph
- It compares the **observed number of internal links** in each partition, against the **expected number of internal links** if those internal links had been placed at random

# Modularity of a partition

$$Q = \frac{1}{L} \sum_C \left( L_C - \frac{k_C^2}{4L} \right)$$

- $L$  = number of links in the network
- $L_C$  = number of internal links in community  $C$
- $k_C$  = sum of degree of nodes in  $C$

# Modularity of a partition (cont.)

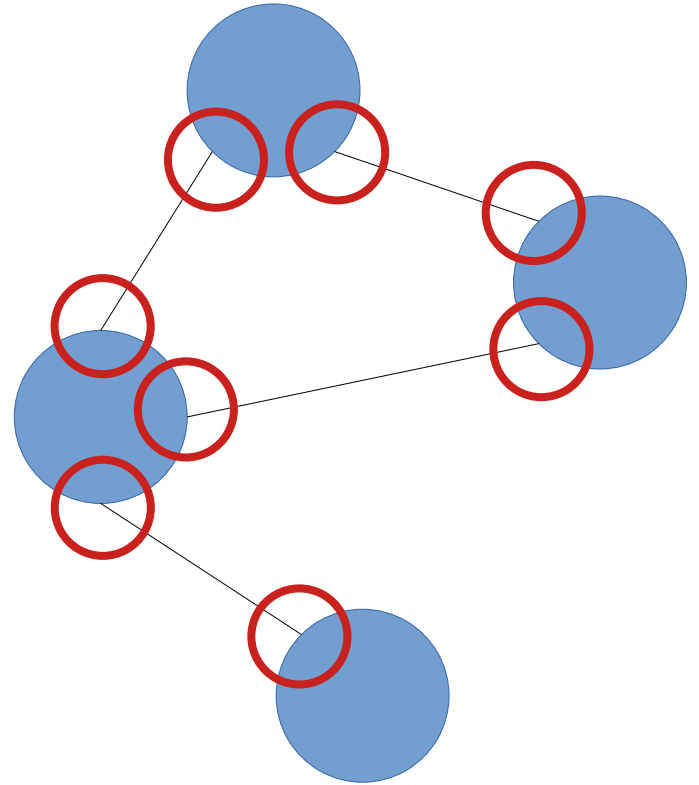
$$Q = \frac{1}{L} \sum_C \left( L_C - \frac{k_C^2}{4L} \right) \longrightarrow$$

Expression in parenthesis is the difference between observed and expected internal links in community  $C$

- $L$  = number of links in the network
- $L_C$  = number of internal links in community  $C$
- $k_C$  = sum of degree of nodes in  $C$
- $k_C^2/4L$  = **expected** number of internal links in community  $C$

# Where does $k_c^2/4L$ comes from?

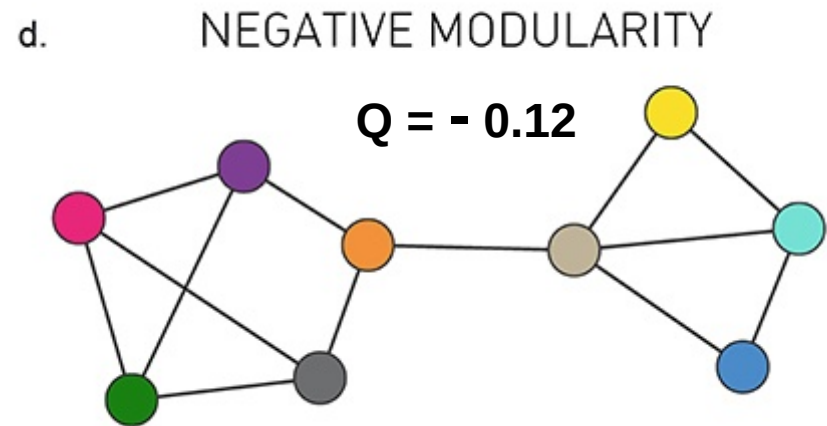
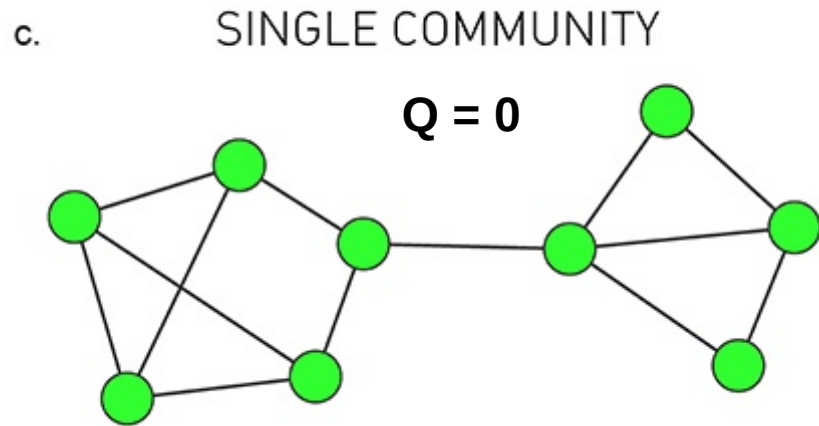
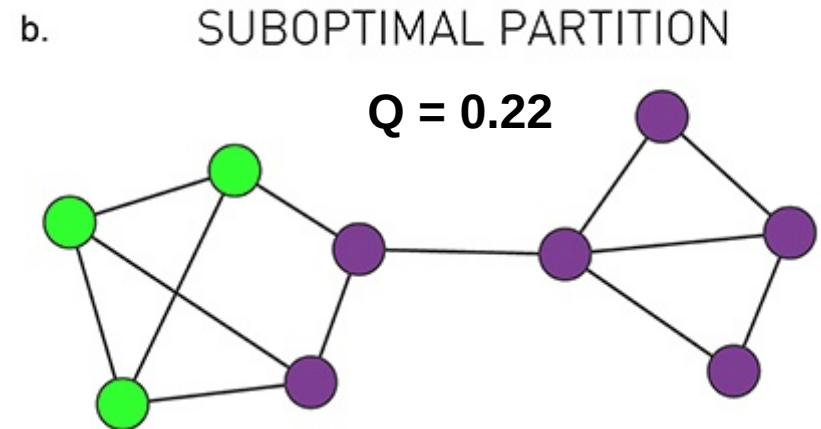
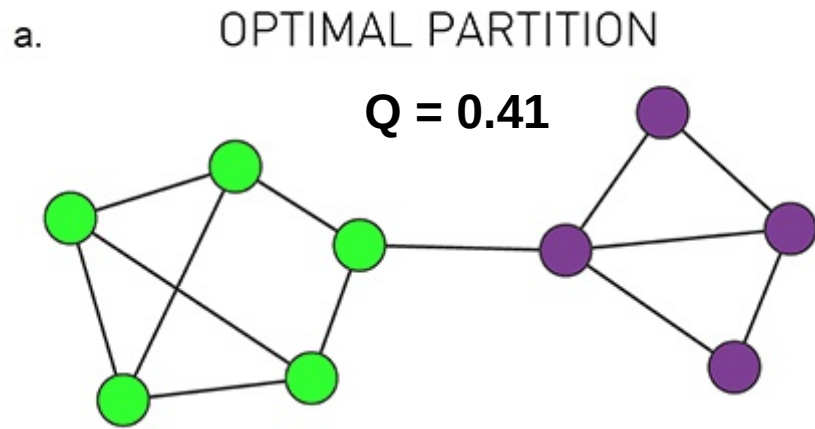
- A **link “stub”** is a connection between a link and a node
- There are  $2L$  stubs in a network
- There are as many stubs as the sum of the degree of nodes



# Modularity formula explained

$$Q = \frac{1}{L} \sum_C \left( L_C - \frac{k_C^2}{4L} \right)$$

- There are  $L_C$  internal links in  $C$
- Total number of stubs in nodes in  $C$  is  $k_C$
- Total number of stubs in the network is  $2L$
- Probability of choosing two stubs in  $C$ :  $(k_C/2L)^2 = k_C^2/4L^2$
- The **expected number** of links joining two stubs in  $C$  is  $L(k_C^2/4L^2) = k_C^2/4L$
- The **observed number** is  $L_{caa}$
- $Q$  has a range:  $Q \in [-1, +1]$

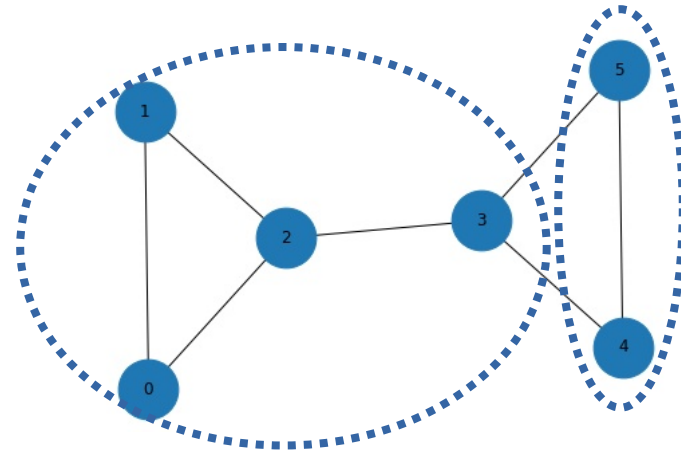
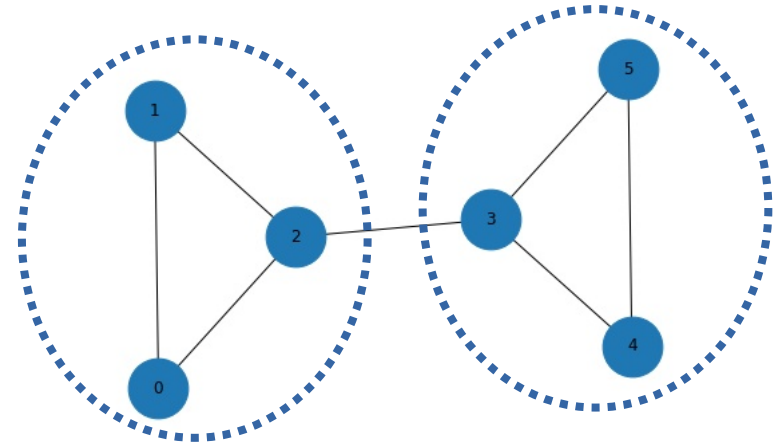


$$Q = \frac{1}{L} \sum_C \left( L_C - \frac{k_C^2}{4L} \right)$$

# Exercise

- What is the modularity of the partition  $\{0, 1, 2\}, \{3, 4, 5\}$ ?
- What is the modularity of the partition  $\{0, 1, 2, 3\}, \{4, 5\}$ ?

$$Q = \frac{1}{L} \sum_C \left( L_C - \frac{k_C^2}{4L} \right)$$





# Summary

# Things to remember

- Strong and weak community
- The concept of “cut” in graph bisection
- Girvan-Newman’s algorithm
- Modularity

# Practice on your own

- Check the modularity computations in the example on the slide marked ★ : (a) optimal partitioning into two communities, (b) suboptimal partitioning into two communities, (c) all the nodes in a single community, (d) one community per node
- You can check your answers with [networkx.algorithms.community.modularity](https://networkx.org/documentation/stable/reference/algorithms/community/modularity.html)