Reverse engineering Golang binaries with Ghidra





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Introduction

Go (also called Golang) is an open source programming language, designed by Google in 2007 and made available to the public in 2012. During the years it gained popularity among developers and as usually, it is not only used for good purposes, but attracted attention of malware developers as well. The fact that Go supports cross compiling to run binaries on various operating systems makes it a tempting choice for malware developers. The possibility to compile the same code for all major platforms (Windows, Linux, MacOS) makes the attackers life much easier, they don't have to develop and maintain different codebase for each target environment.

Some special features of the Go programming language give a hard time for reverse engineers when investigating Go binaries. Reverse engineering tools (e.g. disassemblers) can do a great job in analyzing binaries that are written in more popular languages (e.g. C, C++, .NET), but Go gives new challenges that makes the analysis more cumbersome.

Go binaries are usually statically linked, which means that all the necessary libraries are included in the compiled binary. This results in large binaries. On one hand this makes malware distribution more difficult for the attackers, but on the other hand some of the security products also have issues with handling such large files. It can help avoiding detection. The other advantage of statically linked binaries for the attackers is that the malware can run on the target systems without dependency issues.

As we see a continuous growth in the number of malware written in Go and we expect more families to emerge, we decided to dive deeper into the Go programming language and enhance our toolset to be more effective in investigating Go malware.

Below we will discuss two of the difficulties that reverse engineers have to face during Go binary analysis and we will show our solutions for those.

Ghidra is an open source reverse engineering tool developed by the National Security Agency, which we frequently use for static malware analysis. It is possible to create custom scripts and plugins for Ghidra to provide specific functionalities that are needed by researchers. We used this feature of Ghidra and created custom scripts to aid our Go binary analysis.

The topics discussed in this article were presented at the <u>Hacktivity2020</u> online conference. The slides and other materials are available in our <u>Github repository</u>.

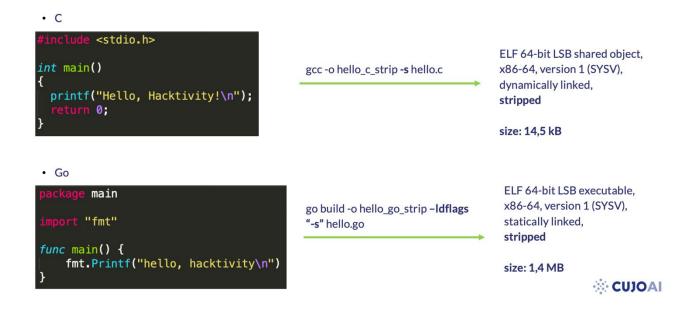
Lost function names

The first issue is not specific to Go binaries, but stripped binaries in general. Compiled executable files can contain debug symbols which make debugging and analysis easier. When reverse engineering a program that was compiled with debugging information included, analysts can see not only memory



addresses but also the names of the routines and variables. However, in order to reduce the size and make reverse engineering more difficult, malware authors usually compile the files without this information, creating so-called stripped binaries. In this case analysts cannot rely on the function names to help them finding their way around the code. For statically linked Go binaries, where all the necessary libraries are included, this can significantly slow down the analysis.

To illustrate this issue, we used simple "Hello Hacktivity" examples written in C^[1] and Go^[2] for comparison and compiled them to stripped binaries. Note the size difference of the two executables.



Ghidra function window lists all the defined functions within the binaries. In the non-stripped versions function names are nicely visible and provide a great help for reverse engineers.



∄ Functions - 18 items			■ △ ×
Label	Location	Function Sign	Function Size
_init	00101000	int _init(E	27
cxa_finalize	00101040	thunk undef	11
puts	00101050	thunk int p	11
_start	00101060	undefined	47
deregister_tm_clones	00101090	undefined d	34
register_tm_clones	001010c0	undefined r	51
do_global_dtors_aux	00101100	undefined	54
frame_dummy	00101140	thunk undef	9
main	00101149	undefined m	27
libc_csu_init	00101170	undefined	101
libc_csu_fini	001011e0	undefined	5
_fini	001011e8	undefined	13
_ITM_deregisterTMCloneTable	00105000	thunk undef	1
puts	00105008	thunk int p	1
libc_start_main	00105010	thunk undef	1
gmon_start	00105018	thunk undef	1
_ITM_registerTMCloneTable	00105020	thunk undef	1
cxa_finalize	00105028	thunk undef	1
Filter:			② ≑ ·

Figure 1 - hello_c^[3] function list

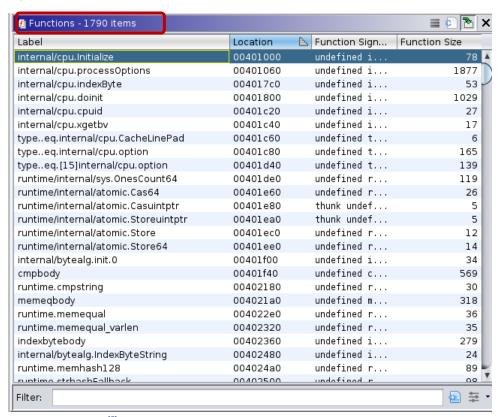


Figure 2 - hello_go^[5] function list



For stripped binaries the function lists look the following:

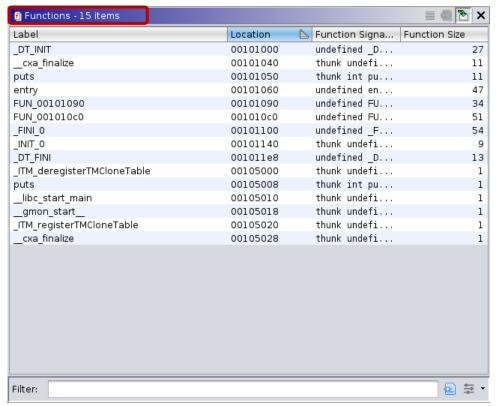


Figure 3 - hello_c_strip^[4] function list



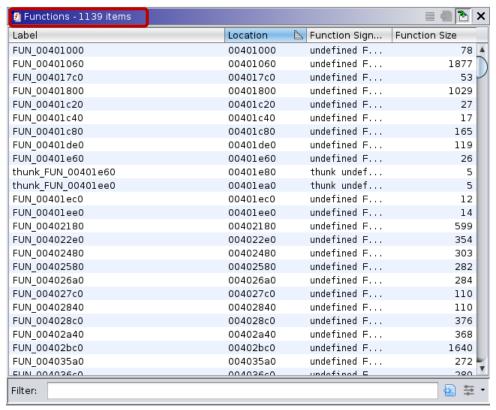


Figure 4 - hello_go_strip^[6] function list

These examples nicely show that even a simple hello world Go binary is huge with more than a thousand functions and in the stripped version reverse engineers cannot rely on the function names to aid their analysis.

Note: Due to stripping not only the function names disappeared, but instead of 1790 defined functions only 1139 were recognized by Ghidra.

We were interested if there is any way to recover the function names within stripped binaries. First by a simple string search it can be checked if the function names are still available within the binaries. For the C example we were looking for the function name "main", while in the Go example it is "main.main".

```
pad0rka in hacktivity2020 % strings hello_c | grep -o ".\{0,10\}main.\{0,10\}" ibc_start_main ibc_start_main@@GLIBC_2. main  
Figure 5 - hello_c^{[3]} strings - "main" was found pad0rka in hacktivity2020 % strings hello_c_strip | grep -o ".\{0,10\}main.\{0,10\}" ibc_start_main  
Figure 6 - hello_c strip^{[4]} strings - "main" was not found
```



```
pad0rka in hacktivity2020 % strings hello_go | grep -o ".\{0,10\}main.\{0,10\}"
 hasmain
 edruntime.main not on m0
 p stateremaining pointe
 out of domainpanic whil
 e space remainingreflect
 routines (main called ru
 runtime.main
 runtime.main.func1
 runtime.main.func2
main.main
 main..inittask
 runtime.main_init_done
 runtime.mainStarted
 runtime.mainPC
 runtime.main
 runtime.main.func1
 runtime.main.func2
main.main
```

Figure 7 - hello_go^[5] strings – "main.main" was found

```
pad0rka in hacktivity2020 % strings hello_go_strip | grep -o ".\{0,10\}main.\{0,10\}" hasmain edruntime.main not on m0 p stateremaining pointe out of domainpanic whil e space remainingreflect routines (main called ru runtime.main runtime.main.func1 runtime.main.func2 main.main
```

Figure 8 - hello_go_strip^[6] strings - "main.main" was found

While in the stripped C binary^[4] the function name cannot be found with the strings utility, in the Go version^[6] "main.main" is still available. This discovery gave us some hope that in stripped Go binaries function name recovery might be possible.

Loading the binary^[6] to Ghidra and searching for the "main.main" string will show the exact location. As it can be seen on the image below the function name string is located within the .gopclntab section.



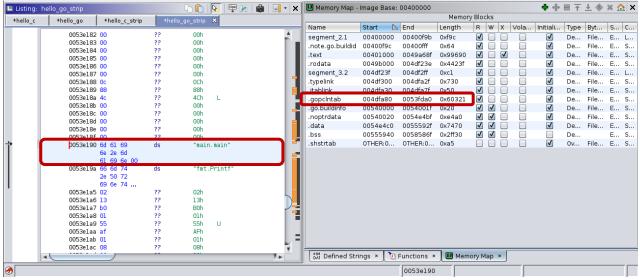
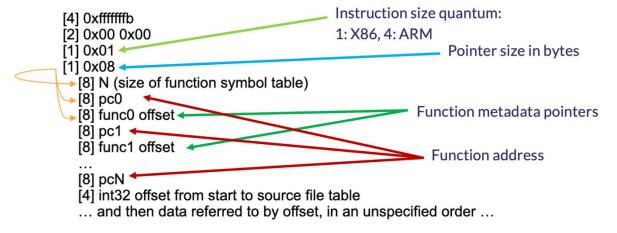


Figure 9 - hello_go_strip[6] main.main string in Ghidra

The pcIntab structure is available since Go 1.2 and nicely <u>documented</u>. The structure starts with a magic value followed by information about the architecture. Then the function symbol table holds information about the functions within the binary. The address of the entry point of each function is followed by a function metadata table.



The function metadata table, among other important information, stores an offset to the function name.



```
struct
                Func
{
         uintptr
                          entry; // start pc
                                // name (offset to C string)
        int32 name;
                                // size of arguments passed to function
         int32 args;
         int32 frame;
                                // size of function frame, including saved caller PC
         int32
                        pcsp;
                                                 // pcsp table (offset to pcvalue table)
         int32
                                           // pcfile table (offset to pcvalue table)
                        pcfile;
                                                   // pcln table (offset to pcvalue table)
         int32
                        pcln;
         int32
                        nfuncdata;
                                               // number of entries in funcdata list
         int32
                        npcdata;
                                            // number of entries in pcdata list
};
```

Using this information, it is possible to recover the function names. Our team created a <u>script</u> (go_func.py) for Ghidra to recover function names in stripped Go ELF files by executing the following steps:

- Locate pclntab structure
- Extract function addresses
- Find function name offsets

After executing our script not only the function names will be restored, but the previously unrecognized functions will be defined as well.

1-1-1	1 11	в Е
Label	Location	Func Func
fmt.(*pp).badVerb	00492f40	und 1649
fmt.(*pp).fmtBool	004935c0	und 111
fmt.(*pp).fmt0x64	00493640	und 149
fmt.(*pp).fmtInteger	004936e0	und 820
fmt.(*pp).fmtFloat	00493a20	und 408
fmt.(*pp).fmtComplex	00493bc0	und 583
fmt.(*pp).fmtString	00493e20	und 457
fmt.(*pp).fmtBytes	00494000	und 2303
fmt.(*pp).fmtPointer	00494900	und 1358
fmt.(*pp).catchPanic	00494e60	und 1534
fmt.(*pp).handleMethods	00495460	und 1748
fmt.(*pp).printArg	00495b40	und 2348
fmt.(*pp).print∨alue	004964a0	und 9767
fmt.intFromArg	00498b00	und 529
fmt.parseArgNumber	00498d20	und 293
fmt.(*pp).argNumber	00498e60	und 278
fmt.(*pp).badArgNum	00498f80	und 367
fmt.(*pp).missingArg	00499100	und 367
fmt.(*pp).doPrintf	00499280	und 4490
fmt.globfuncl	0049a420	und 84
fmt.init	0049a480	und 197
typeeq.fmt.fmt	0049a560	und 172
main.main	0049a620	und 112

Figure 10 - hello_go_strip^[6] function list after executing go_func.py



To see a real-world example let's look at an eCh0raix ransomware sample [9]:

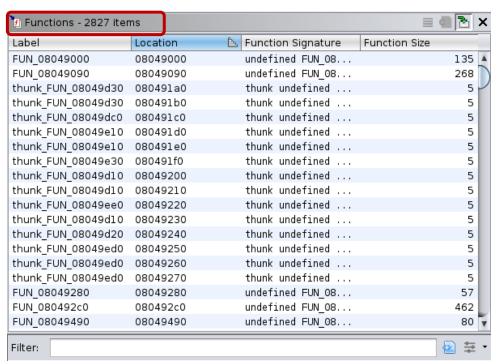


Figure 11 - eCh0raix[9] function list

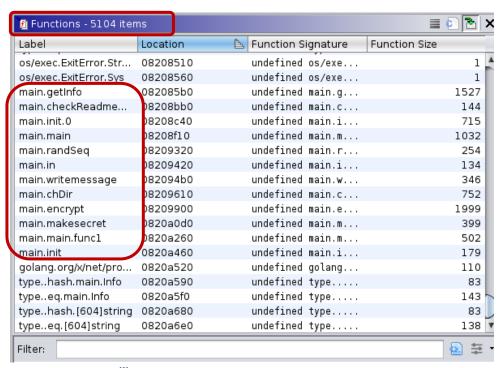


Figure 12 - eCh0raix[9] function list after executing go_func.py



This example clearly shows how much help this simple function name recovery script can be during reverse engineering. Only by looking at the function names analysts can assume that they are dealing with a ransomware.

Note: In Windows Go binaries there is no specific section for the pcIntab structure, rather researchers need explicitly search for the fields of this structure (e.g. magic value, possible field values). For MacOS _gopcIntab section is available, similarly to .gopcIntab in Linux binaries.

Challenges

If a function name string is not defined by Ghidra, then the function name recovery script will fail to rename that specific function, since it cannot find the function name string at the given location. To overcome this issue our script always checks if a defined data type is located at the function name address and if not, then before renaming a function it tries to define a string data type at the given address.

In the below example the function name string "log.New" is not defined in an eCh0raix ransomware sample [9], so the corresponding function cannot be renamed without string creation first.

083aa0e4	6c	??	6Ch	ι
083aa0e5	6f	??	6Fh	0
083aa0e6	67	??	67h	g
083aa0e7	2e	??	2Eh	Ţ
083aa0e8	4e	??	4Eh	Ν
083aa0e9	65	??	65h	е
083aa0ea	77	??	77h	W
083aa0eb	00	??	00h	

Figure 13 - eCh0raix^[9] log.New function name undefined



```
**************************
                                                         FUNCTION
                          *********************
                         undefined FUN 08184fa0(undefined4 param 1, undefined4 pa...
                     AL:1 <RETURN:
Stack[0x4]:4 param_1
Stack[0x8]:4 param_2
      undefined
      undefined4
                                                                                                  XREF[1]: 08184fc7(R)
      undefined4
                                                                                                  XREF[2]: 08184fd8(R),
                                                                                                                  0818501d(R)
     undefined4
                     Stack[Oxc]:4 param 3
                                                                                                  XREF[2]:
                                                                                                              08184ff0(R),
                                                                                                                 0818500b (R)

        undefined4
        Stack[0x10]:4 param_4

        undefined4
        Stack[0x14]:4 param_5

        undefined4
        Stack[0x18]:4 param_6

        undefined4
        Stack[-0x4]:4 local_4

        undefined4
        Stack[-0x8]:4 local_8

        FUN_08184fa0

                                                                                                  XREF[1]: 08184fdf(R)
                                                                                                  XREF[1]: 08184ff7(R)
                                                                                                  XREF[1]: 08184ffe(W)
                                                                                                  XREF[1]: 08184fc3(R)
                                                                                                 XREF[1]:
                                                                                                                 08184fbb(*)
                                                                                       XREF[2]: 0818502f(c),
                                                                                                       log.init:08186012(c)
08184fa0 65 8b 0d
                              MOV
                                             ECX, dword ptr GS: [0x0]
          00 00 00 00
08184fa7 8b 89 fc
                                             ECX, dword ptr [ECX + Oxfffffffc]
           ff ff ff
```

Figure 14 – eCh0raix^[9] log.New function couldn't be renamed

The following lines in our script are responsible to solve this challenge:

```
func_name = getDataAt(name_address)

#Try to define function name string.
if func_name is None:
    try:
        func_name = createAsciiString(name_address)
    except:
        print "ERROR: No name"
        continue
```

Figure 15 - go_func.py

Unrecognized strings

The second issue that our scripts are solving is related to strings within Go binaries. Let's turn back to the "Hello Hacktivity" examples and take a look at the defined strings within Ghidra.

In the C binary^[3] 70 strings are defined, among which "Hello, Hacktivity!" can be found. Meanwhile the Go binary^[5] includes 6540 strings but searching for "hacktivity" gives no result. Having such a high number of strings already gives a hard time for reverse engineers to find the relevant ones, but in this case, the string that we would expect to find, is not even recognized by Ghidra.



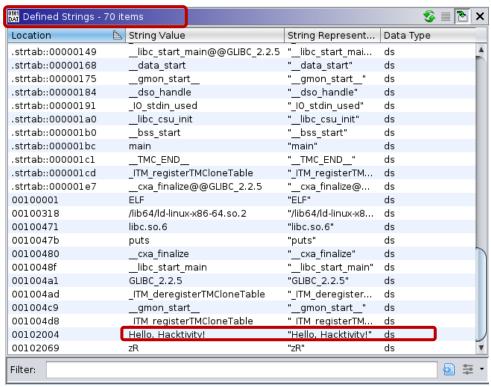


Figure 16 - hello_c^[3] defined strings with "Hello, Hacktivity!"

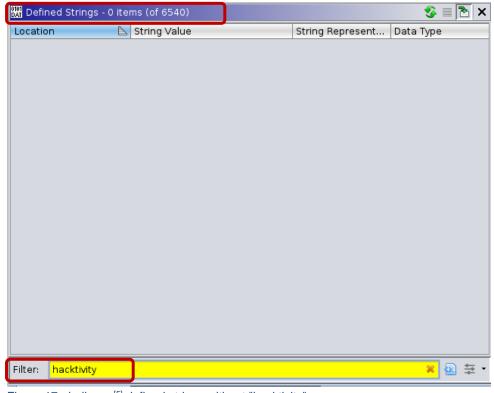


Figure 17 - hello_go^[5] defined strings without "hacktivity"



To understand the problem here, the first step is to understand what a string in Go is. Unlike in C-like languages, where strings are sequence of characters terminated with a null character, in Go strings are sequences of bytes with a fixed length. Strings are Go specific structures, built up by a pointer to the location of the string and an integer, which is the length of the string.

```
type stringStruct struct {
          str unsafe.Pointer
          len int
}
```

These strings are stored within Go binaries as a large string blob, which consists of the concatenation of the strings without null character between them. So, while searching for "Hacktivity" using strings and grep gives the expected result in C, in case of Go a huge string blob is returned containing somewhere "hacktivity".

```
pad0rka in hacktivity2020 % strings hello_c | grep Hacktivity Hello, Hacktivity!
```

Figure 18 - hello_c^[3] string search for "Hacktivity"

padOrka in hacktivity2020 % strings hello_go | grep hacktivity

object is remotepacer: H_m_prev=reflect mismatchremote I/O errorruntime: g: g=runtime: addr = runtime:

base = runtime: gp: gp=runtime: head = runtime: nelems=schedule: in cgosigaction failedtime: bad [0-9]*

workbuf is empty initialHeapLive= spinningthreads=, s.searchAddr = 0123456789ABCDEFX0123456789abcdefx119

2092895507812559604644775390625: missing method GC assist markingOld_North_ArabianOld_South_ArabianOther

_ID_ContinueSIGBUS: bus errorSIGCONT: continueSIGINT: interruptSentence_TerminalUnified_Ideographbad Tin

ySizeClassdebugPtrmask.lockentersyscallblockexec format errorfutexwakeup addr=g already scannedglobalAll

oc.mutexlocked m0 woke upmark - bad statusmarkBits overflowno data availablenotetsleepg on g0permission
deniedreflect.Value.Intreflect.Value.Lenreflect: New(nil)reflect: call of runtime/internal/runtime: leve
l = runtime: nameOff runtime: next_gc=runtime: pointer runtime: summary[runtime: textOff runtime: typeOff
f scanobject n == 0select (no cases)stack: frame={sp:swept cached spanthread exhaustionunknown caller pc

wait for GC cyclewrong medium type but memory size because dotdotdot to non-Go memory , locked to thre
ad298023223876953125Caucasian_AlbanianRFS specific errorRegional_IndicatorVariation_Selectorbad lfnode a
ddressbad manualFreeListconnection refusedfaketimeState.lockfile name too longforEachP: not donegarbage
collection(hello, hacktivity)

Figure 19 - hello_go^[5] string search for "hacktivity"

Since the definition of strings is different and as a result referencing them within the assembly code is also differ from the usual C-like solutions, Ghidra has a hard time to define the strings within Go binaries.

The string structure can be allocated in many different ways, it can be created statically or dynamically during runtime, it varies over architecture and even within one architecture multiple solutions are possible. Our team created two scripts to help Ghidra identifying strings.



Dynamically allocated string structures

In the first case string structures are created runtime. A sequence of assembly instructions is responsible for setting up the structure before a string operation. Thanks to the different instruction sets it varies over architectures. In the next few paragraphs we will go through a couple of use cases and show the instruction sequences that our script (find_dynamic_strings.py) is looking for.

x86

First let's start with the "Hello Hacktivity" example [5].

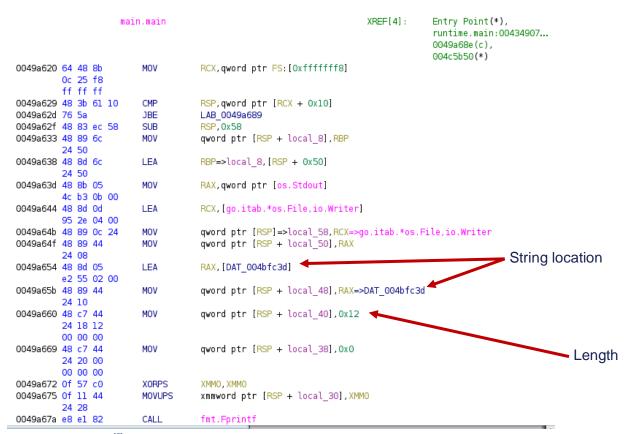


Figure 20 - hello_go^[5] dynamic allocation of string structure



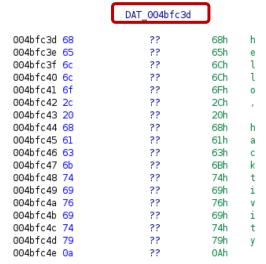


Figure 21 - hello_go^[5] undefined "hello, hacktivity" string

After executing the script, the code looks the following:

```
0049a654 48 8d 05
                                     RAX, [s_hello,_hacktivity_004bfc3d]
                         LEA
         e2 55 02 00
0049a65b 48 89 44
                         MOV
                                     qword ptr [RSP + local_48], RAX=>s_hello, _hacktivity_004bfc3d
         24 10
0049a660 48 c7 44
                         MOV
                                     qword ptr [RSP + local_40],0x12
         24 18 12
         00 00 00
0049a669 48 c7 44
                         MOV
                                     qword ptr [RSP + local_38],0x0
         24 20 00
         00 00 00
0049a672 Of 57 c0
                         XORPS
                                     XMMO, XMMO
0049a675 Of 11 44
                         MOVUPS
                                     xmmword ptr [RSP + local_30],XMMO
         24 28
0049a67a e8 e1 82
                         CALL
                                     fmt.Fprintf
         ff ff
```

Figure 22 - hello_go^[5] dynamic allocation of string structure after executing find_dynamic_strings.py

The string is defined:

Figure 23 - hello_go^[5] defined "hello hacktivity" string

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And "hacktivity" can be found in the defined strings view in Ghidra:

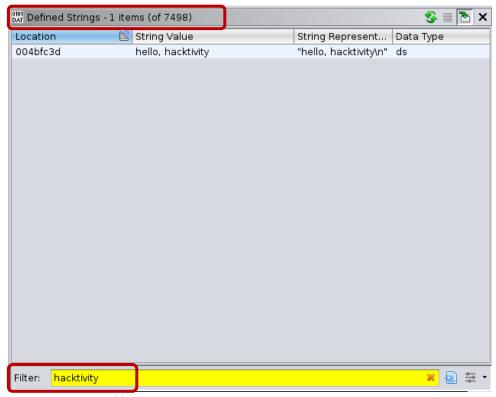


Figure 24 - hello_go^[5] defined strings with "hacktivity"

The script is looking for the following instruction sequences in case of 32-bit and 64-bit x86 binaries:

Figure 25 - eCh0raix^[9] dynamic allocation of string structure

```
#x86_64
#LEA REG, [STRING_ADDRESS]
#MOV [RSP + ..], REG
#MOV [RSP + ..], STRING_SIZE
```



Figure 26 - hello_go^[5] dynamic allocation of string structure

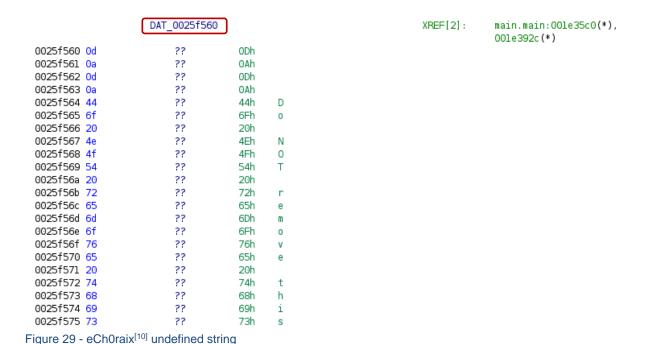
ARM

For the 32-bit ARM architecture an eCh0raix ransomware sample^[10] will be used to illustrate the string recovery.

```
Pointer to address
                                          r2, [PTR_DAT_001e392c]
                                                                                   containing the
 001e35bc 68 23 9f e5
                             ldr
                                          r2=>DAT_0025f560,[sp,#local_90]
 001e35c0 10 20 8d e5
                             str
                                                                                   string location
                                          r2,#0x44
 001e35c4 44 20 a0 e3
                             mov
 001e35c8 14 20 8d e5
                             str
                                          r2,[sp,#losal_8c]
                                                                            String location
 001e35cc 18 00 8d e5
                             str
                                          r0,[sp,#local_88]
 001e35d0 lc 10 8d e5
                                          rl,[sp,#local 84]
                             str
 001e35d4 44 cc f9 eb
                             bl
                                          runtime.concatstring3
                                                                            Length
Figure 27 - eCh0raix<sup>[10]</sup> dynamic allocation of string structure
```

Figure 28 - eCh0raix^[10] pointer to string address





After executing the script, the code looks the following:

```
001e35bc 68 23 9f e5
                                    r2, [PTR_s_Do_NOT_remove_this_file_and_NOT_001e392c]
001e35c0 10 20 8d e5
                                    r2=>s Do NOT remove this file and NOT 0025f560, [sp,#local 90]
                        str
00le35c4 44 20 a0 e3
                        mov
                                    r2,#0x44
001e35c8 14 20 8d e5
                        str
                                    r2,[sp,#local_8c]
001e35cc 18 00 8d e5
                        str
                                    r0,[sp,#local_88]
001e35d0 1c 10 8d e5
                        str
                                    rl,[sp,#local_84]
001e35d4 44 cc f9 eb
                        bl
                                    runtime.concatstring3
```

Figure 30 - eCh0raix[10] dynamic allocation of string structure after executing find_dynamic_strings.py

The pointer is renamed, and the string is defined:

```
PTR_s_Do_NOT_remove_this_file_and_NOT_00le392c XREF[1]: main.main:00le35bc(R) 00le392c 60 f5 25 00 addr s_Do_NOT_remove_this_file_and_NOT_0025f560
```

Figure 31 - eCh0raix[10] pointer to string address after executing find_dynamic_strings.py

Figure 32 – eCh0raix^[10] defined string after executing find_dynamic_strings.py

The script is looking for the following instruction sequence in case of 32-bit ARM binaries:



```
#ARM, 32-bit
#LDR REG, [STRING_ADDRESS_POINTER]
#STR REG, [SP, ..]
#MOV REG, STRING_SIZE
#STR REG, [SP, ..]
```

For the 64-bit ARM architecture a Kaiji sample^[12] will be used to illustrate the string recovery. Here two instruction sequences are used, that only differ in one instruction.

```
LAB 0020b59c
                                                                       XREF[2]:
                                                                                    0020b814(j), 0020b988(j)
0020b59c 00 04 00 b0
                                     x0,0x28c000
0020b5a0 00 c4 lc 91
                         add
                                     x0, x0, #0x731
0020b5a4 e0 07 00 f9
                                     x0=>DAT 0028c73l [sn.
                         str
                                                           #local 681

    String location

0020b5a8 e0 07 7e b2
                                     x0,xzr,#0xc
                         orr
0020b5ac e0 0b 00 f9
                         str
                                     x0,[sp, #local 60]
0020b5b0 e4 d3 ff 97
                                     ddos.PathExists
                                                                         Length
                         bl
0020b5b4 e0 63 40 39
                         ldrb
                                     w0,[sp, #local 58]
0020b5b8 60 05 00 b5
                                     x0, LAB_0020b664
                         cbnz
                    LAB 0020b5bc
                                                                       XREF[2]:
                                                                                    0020b680(j), 0020b7f4(j)
0020b5bc 00 04 00 f0
                         adrp
                                     x0,0x28e000
0020b5c0 00 84 28 91
                         add
                                     x0,x0,#0xa21
0020b5c4 e0 07 00 f9
                         str
                                     x0=>DAT_0028ea21,[sp, #local_68]
0020b5c8 80 02 80 d2
                                     x0,#0x14
                         mov
                                                                         String location
0020b5cc e0 0b 00 f9
                         str
                                     x0,[sp, #local 6
0020b5d0 dc d3 ff 97
                         bl
                                     ddos.PathExists
                                                                       Length
0020b5d4 e0 63 40 39
                         ldrb
                                     w0,[sp, #local_58]
                                     x0, LAB_0020b5e8
0020b5d8 80 00 00 b5
                         cbnz
```

Figure 33 – Kaiji^[12] dynamic allocation of string structure

After executing the script, the code looks the following:

```
LAB 0020b59c
                                                                       XREF[2]:
                                                                                    0020b814(j), 0020b988(j)
                                     x0,0x28c000
0020b59c 00 04 00 b0
                         adrp
0020b5a0 00 c4 lc 91
                         add
                                     x0, x0, #0x731
0020b5a4 e0 07 00 f9
                         str
                                     x0=>s_/etc/init.d/_0028c731,[sp, #local_68]
0020b5a8 e0 07 7e b2
                         orr
                                     x0,xzr,#0xc
0020b5ac e0 0b 00 f9
                                     x0,[sp, #local 60]
                         str
0020b5b0 e4 d3 ff 97
                                     ddos.PathExists
0020b5b4 e0 63 40 39
                         ldrb
                                     w0,[sp, #local 58]
0020b5b8 60 05 00 b5
                                     x0,LAB 0020b664
                         cbnz
                     LAB_0020b5bc
                                                                       XREF[2]:
                                                                                    0020b680(j), 0020b7f4(j)
0020b5bc 00 04 00 f0
                                     x0,0x28e000
                                     x0, x0, \#0xa21
0020b5c0 00 84 28 91
                         add
0020b5c4 e0 07 00 f9
                         str
                                     x0=>s /etc/systemd/system/ 0028ea21,[sp, #local 68]
0020b5c8 80 02 80 d2
                         mov
                                     x0,[sp, #local_60]
0020b5cc e0 0b 00 f9
                         str
0020b5d0 dc d3 ff 97
                         bl
                                     ddos.PathExists
0020b5d4 e0 63 40 39
                         ldrb
                                     w0,[sp, #local 58]
0020b5d8 80 00 00 b5
                         cbnz
                                     x0,LAB_0020b5e8
```

Figure 34 – Kaiji^[12] dynamic allocation of string structure after executing find_dynamic_strings.py



The strings are defined:

Figure 35 – Kaiji^[12] defined strings after executing find_dynamic_strings.py

The script is looking for the following instruction sequences in case of 64-bit ARM binaries:

```
#ARM, 64-bit - version 1
#ADRP REG, [STRING_ADDRESS_START]
#ADD REG, REG, INT
#STR REG, [SP, ..]
#ORR REG, REG, STRING_SIZE
#STR REG, [SP, ..]

#ARM, 64-bit - version 2
#ADRP REG, [STRING_ADDRESS_START]
#ADD REG, REG, INT
#STR REG, [SP, ..]
#MOV REG, STRING_SIZE
#STR REG, [SP, ..]
```

As the above examples show, after executing the script dynamically allocated string structures can be recovered. It gives a great help to reverse engineers to read the assembly code or look for interesting strings within the defined string window in Ghidra.

Challenges

The biggest drawback of this approach is that for each architecture and even for different solutions within the same architecture a new branch has to be added to the script. Also, it is very easy to evade these predefined instruction sets. In the example below in a Kaiji 64-bit ARM malware sample^[12] the length of the string is moved to a register earlier than our script would expect, therefore this string will be missed.





```
001fd734 21 01 80 d2
                                        param 2,#0x9
                           mov
 001fd738 el 4b 00 f9
                                        param 2,[sp, #local c0]
                           str
 001fd73c 62 04 00 d0
                           adrp
                                        param 3,0x28b000
 001fd740 42 fc 2f 91
                                        param_3=>DAT_0028bbff,param_3,#0xbff
                            add
 001fd744 e2 4f 00 f9
                                        param 3=>DAT 0028bbff_[sp, #local b8]
                            str
 001fd748 el 53 00 f9
                                        param_2,[sp, #local_b0]
                           str
Figure 36 – Kaiji[12] dynamic allocation of string structure in an unusual way
                                                                           String location
```

```
DAT_0028bbff
                                                                     XREF[6]:
                                                                                 ddos.sshgo:001fd740(*),
                                                                                 ddos.sshgo:001fd744(*),
                                                                                 ddos.sshqo:001fd788(*),
                                                                                 ddos.sshgo:001fd7a4(*),
                                                                                 ddos.sshgo:001fd7c0(*),
                                                                                 ddos.sshgo:001fd7dc(*)
0028bbff 6c
                         ??
                                    6Ch
0028bc00 69
                         ??
                                    69h
                                           i
                        ??
0028bc01 6e
                                    6Eh
                                           n
0028bc02 75
                                    75h
                                           u
0028bc03 78
                                           Х
                        ??
                                    5Fh
0028bc04 5f
                                           a
0028bc05 61
                        ??
                                    61h
                        ??
0028bc06 72
                                    72h
                                           r
0028bc07 6d
                        ??
                                    6Dh
```

Figure 37 - Kaiji^[12] undefined string

Statically allocated string structures

In the next case our <u>script</u> (find_static_strings.py) looks for string structures that are statically allocated, meaning the string pointer is followed by the string length within the data section of the code.

To illustrate this let's look at the x86 eCh0raix ransomware sample [9].



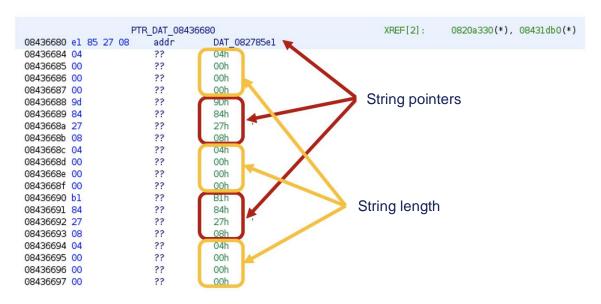


Figure 38 - eCh0raix^[9] static allocation of string structures

On the image above string pointers are followed by string length values, however Ghidra couldn't recognize the addresses, neither the integer data types, except for the first pointer, which is directly referenced from the code.

```
0820a30f 8b 44 24 20
                                      EAX, dword ptr [ESP + 0x20]
                          MOV
                                      dword ptr [ESP], EAX
0820a313 89 04 24
                          MOV
0820a316 8b 44 24 1c
                          MOV
                                      EAX, dword ptr [ESP + 0xlc]
0820a31a 89 44 24 04
                          MOV
                                      dword ptr [ESP + 0x4], EAX
0820a3le 8b 05 b0
                          MOV
                                      EAX, dword ptr [PTR_PTR_DAT_08431db0]
         1d 43 08
0820a324 8b 0d b4
                          MOV
                                      ECX, dword ptr [DAT_08431db4]
         1d 43 08
0820a32a 8b 15 b8
                          MOV
                                      EDX, dword ptr [DAT 08431db8]
         1d 43 08
                                      dword ptr [ESP + 0x8], EAX=>PTR DAT 08436680
                          MOV
0820a330 89 44 24 08
                                      dword ptr [ESP + 0xc], ECX
0820a334 89 4c 24 0c
                          MOV
                          MOV
                                      dword ptr [ESP + 0x10], EDX
0820a338 89 54 24 10
0820a33c e8 df f0
                          CALL
                                      FUN_08209420
         ff ff
Figure 39 – eCh0raix<sup>[9]</sup> pointer
```

Following the string addresses, the undefined strings can be found.



	DAT_082785e3	L	
082785el <mark>2e</mark>	??	2Eh	
082785e2 64	??	64h	d
082785e3 61	??	61h	а
082785e4 74	??	74h	t
082785e5 2e	??	2Eh	
082785e6 64	??	64h	d
082785e7 62	??	62h	b
082785e8 30	??	30h	0
082785e9 2e	??	2Eh	
082785ea 64	??	64h	d
082785eb 62	??	62h	b
082785ec 61	??	61h	а
082785ed 2e	??	2Eh	
082785ee 64	??	64h	d
082785ef 62	??	62h	b
082785f0 66	??	66h	f
082785fl 2e	??	2Eh	
082785f2 64	??	64h	d
082785f3 62	??	62h	b
082785f4 6d	??	6Dh	m

Figure 40 - eCh0raix[9] undefined strings

After executing the script, the string addresses will be defined, along with the string length values and the strings themselves.

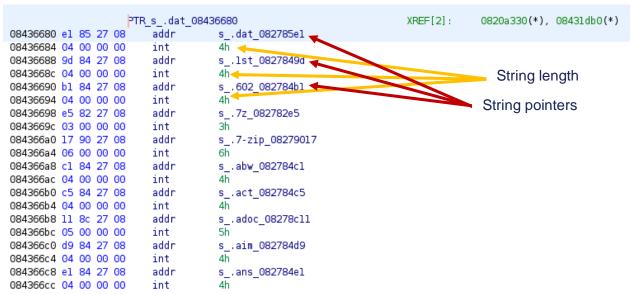


Figure 41 - eCh0raix^[9] static allocation of string structures after executing find_static_strings.py



```
s_.dat_082785el
                                                                XREF[2]:
                                                                             08436680(*), 084378b0(*)
082785el 2e 64 61 74
                     ds
                   s_.db0_082785e5
                                                                 XREF[1]:
                                                                             08437248(*)
                                 ".db0"
082785e5 2e 64 62 30
                     ds
                                                                 XREF[1]:
                                                                             08437250(*)
                   s_.dba_082785e9
                                 ".dba"
082785e9 2e 64 62 61 ds
                  s_.dbf_082785ed
                                                                 XREF[1]:
                                                                             08437258(*)
082785ed 2e 64 62 66
                     ds
                   s_.dbm_082785fl
                                                                XREF[1]:
                                                                             08436ed8(*)
082785f1 2e 64 62 6d ds
                   s_.dbx_082785f5
                                                                 XREF[1]:
                                                                             08437260(*)
082785f5 2e 64 62 78 ds
                                 ".dbx"
                  s_.dcr_082785f9
                                                                             084369d8(*), 08437268(*)
                                                                 XREF[2]:
082785f9 2e 64 63 72 ds
                  s_.der_082785fd
                                                                XREF[2]:
                                                                             08436d70(*), 08437270(*)
082785fd 2e 64 65 72
                     ds
```

Figure 42 – eCh0raix^[9] defined strings after executing find_static_strings.py

Challenges

To eliminate false positives, we limit the string length, search only for printable characters and only in data sections of the binaries. Obviously, as a result of these limitations strings can be easily missed. If you use the script feel free to experiment with it, change the values and find the best settings for your analysis. The following lines in the code are responsible for the length and character set limitations:

```
#Look for strings with printable characters only to eliminate FPs.
def isPrintable(s, l):
    for i in range(l):
        if getByte(s) not in range(32,126):
            return False
        s = s.add(1)
    return True
```

Figure 43 - find_static_strings.py

```
length = getInt(length_address)
#Set the possible length to eliminate FPs.
if length not in range(1,100):
    continue
```

Figure 44 - find_static_strings.py



Further challenges in string recovery

It can happen that Ghidra auto analysis falsely identify certain data types. When it happens, our script will fail to create the correct data at that specific location. To overcome this issue, first the incorrect data type has to be removed, then the new one can be created.

As an example, let's take a look at the eCh0riax ransomware^[9] with statically allocated string structures.

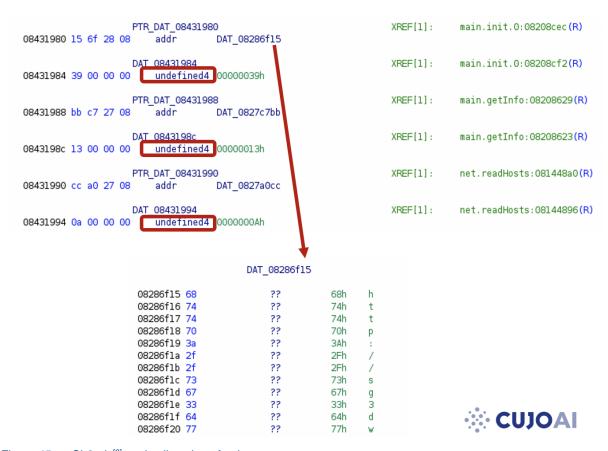


Figure 45 – eCh0raix^[9] static allocation of string structures

Here the addresses are correctly identified, however the string length values, that supposed to be integer data types, are falsely defined as undefined4 values.

The following lines in our script are responsible for removing the incorrect data types:

```
if getDataAt(length_address) is not None:
    data_type = getDataAt(length_address).getDataType()
    #Remove undefined data to be able to create int.
    #Keep an eye on other predefined data types.
    if data_type.getName() in ["undefined4", "undefined8"]:
        removeData(getDataAt(length_address))
```

Figure 46 - find_static_strings.py



After executing the script all the data types are correctly identified and the strings are defined.



Figure 47 - eCh0raix[9] static allocation of string structures after executing find_static_strings.py

Another issue is coming from the fact that in Go binaries strings are stored concatenated, in a large string blob. In certain cases, Ghidra define these blobs as one string. These can be identified by the high number of offcut references. Offcut references are references to certain parts of the defined string, not the address where the string starts, rather somewhere inside the string.

The example below is from an ARM Kaiji sample^[12].

```
s_runtime:_panic_before_malloc_hea_002978ff
                                                                                                        runtime.casgstatus:00043ef4(*).
                          s_ru| "*-+*-+####@@@@!!!!first path segment in URL cannot contain colonln -s /etc/rc.d/init.d/linux_kill
                          s_rul /etc/rc.d/rcmath/big: mismatched montgomery number lengthsmemory reservation exceeds address space
                          s_sl:|limitpanicwrap: unexpected string after type name: reflect.Value.Slice: slice index out of
                          S_SSI boundsreflect: nil type passed to Type.ConvertibleToreleased less than one physical page of S_SyI memoryruntime: debugCallV1 called by unknown caller runtime: failed to create new OS thread (have
                                runtime: name offset base pointer out of rangeruntime: panic before malloc heap
                                initialized\nruntime: text offset base pointer out of rangeruntime: type offset base pointer out of
                                rangeslice bounds out of range [:%x] with length %yssh: unmarshal error for field %s of type
                                %s%sstopTheWorld: not stopped (status != _Pgcstop)sysGrow bounds not aligned to pallocChunkBytestls:
                                failed to parse certificate from server: tls: received new session ticket from a clienttls: server
                                chose an unconfigured cipher suitetls: server did not echo the legacy session IDx509: failed to parse rfc822Name constraint %qx509: failed to unmarshal elliptic curve pointx509: invalid elliptic
                                curve private key valueP has cached GC work at end of mark terminationattempting to link in too many
                                shared librariesbufio: reader returned negative count from Readchacha20poly1305: message
                                authentication failedcurve25519: global Basepoint value was modifiedexplicit string type given to
                                non-string memberfirst record does not look like a TLS handshakeslice bounds out of range [::%x]
                                with length %ytls: incorrect renegotiation extension contentstls: internal error: pskBinders length
                          S_Chi mismatchtls: server selected TLS 1.3 in a renegotiationtls: server sent two HelloRetryRequest 
S_Cu messagesx509: internal error: IP SAN %x failed to parsebufio: writer returned negative count from
                          S_ex| Writecrypto/rsa: key size too small for PSS signaturefailed to parse certificate #%d in the chain:
002976f3 2a 2d 2b
                                %wparsing/packing of this type isn't available yetruntime: cannot map pages i...
           2a 2d 2b
           23 23 23 ...
```

Figure 48 – Kaiji^[12] falsely defined string in Ghidra



```
s runtime: panic before malloc hea 002978ff
                                                                                    runtime.casgstatus:00043ef4(*),
                     s_runtime:_text_offset_base_pointe_0029792d
                                                                                    runtime.doInit:0004eefc(*).
                     s_runtime:_type_offset_base_pointe_0029795b
                                                                                    runtime.sigpanic:00055da4(*),
                     s_slice_bounds_out_of_range_[:%x]_w_00297989
                                                                                    runtime.sigpanic:00055de4(*),
                     s_ssh:_unmarshal_error_for_field_%_002979b7
                                                                                    runtime.sigpanic:00055f24(*),
                                                                                   runtime.sigpanic:00055f64(*),
                     s_sysGrow_bounds_not_aligned_to_pa_00297al3
                     s_tls:_failed_to_parse_certificate_00297a4l
                                                                                    runtime.getStackMap:0005a7d4(*),
                     s led to parse certificate from se 00297a49
                                                                                    runtime.morestackc:0005a834(*),
                                                                                    runtime.resolveNameOff:00065blc(...
                     s_tls:_received_new_session_ticket_00297a6f
                     s_tls:_server_chose_an_unconfigure_00297a9d
                     s_tls:_server_did_not_echo_the_leg_00297acb
                    s_x509:_failed_to_parse_rfc822Name_00297af9
s_x509:_failed_to_unmarshal_ellipt_00297b27
                     s_x509:_invalid_elliptic_curve_pri_00297b55
                     s_P_has_cached_GC_work_at_end_of_m_00297b83
                     s_attempting_to_link_in_too_many_s_00297bb2
                     s_bufio:_reader_returned_negative_c_00297bel
                     s_chacha20poly1305:_message_authen_00297c10
                     s curve25519:_global_Basepoint_val_00297c3f
                     s_explicit_string_type_given_to_no_00297c6e
002976f3 2a 2d 2b
                                     "*-+*-+###@@@@!!!!first path segment in URL cannot contain colonln -s /etc/rc.d.
         2a 2d 2b
         23 23 23 ...
```

Figure 49 – Kaiji^[12] offcut references of a falsely defined string

To find falsely defined strings, one can use the defined strings window of Ghidra and sort the strings by offcut reference count. Large strings with numerous offcut references can be undefined manually before executing the string recovery scripts, so the scripts can successfully create the correct string data types.

Defined Strings -	10814 items			ॐ ≣ 🔁 :
Location	String Value	Data Type	Byte Count	Offcut Reference Count 🕒
0022073d	certificateAuthorities	ds	23	1
00220ecl	ReplaceAllLiteralString	ds	24	1
00220ef5	responseMessageReceived	ds	24	1
00220f29	verifyServerCertificate	ds	24	1
00221561	hashForClientCertificate	ds	25	1
00221ele	asn1:"explicit,tag:1"	ds	22	1
00221e53	handlePostHandshakeMessage	ds	27	1
00222552	secureRenegotiationSupported	ds	30	1
00222ebd	asn1:"optional,tag:2"	ds	23	1
00290069	ckunpa	ds	6	1
002903f7	queuefinalizer during GC	ds	24	1
00330cff	runtime.dropg	ds	14	1
00460248	END	ds	12	1
00460258	BEGIN	ds	16	1
0029bb9c	0001020304050607080910111	ds	969	2
002e9100	expand 32-byte k	ds	20	3
002e91a0	expand 32-byte k	ds	20	3
00293a08	3552713678800500929355621	ds	170	4
0028b3b3	= is not_mcount= minutes nallo	ds	225	23
002976f3	*-+*-+###@@@@!!!!first pat	ds	4517	95

Figure 50 - Kaiji^[12] defined strings

At last we will show an issue in Ghidra decompiler view. Once a string is successfully defined by either manually or by one of our scripts, it will be nicely visible in the listing view of Ghidra, giving a great help to reverse engineers when reading the assembly code. However, the decompiler view in Ghidra cannot



handle fixed length strings correctly and regardless of the length of the string it will display everything until it finds a null character. Luckily this issue will be solved for the next release of Ghidra (9.2). This issue is illustrated below using the eCh0raix sample^[9].

```
main.checkReadmeExists
                                                                       XREE[2]:
                                                                                    08208c3b(c).
                                                                                    main.init.0:08208cda(c)
08208bb0 65 8b 0d
                                     ECX, dword ptr GS: [0x0]
                         MOV
         00 00 00 00
08208bb7 8b 89 fc
                         MOV
                                     ECX, dword ptr [ECX + Oxfffffffc]
         ff ff ff
08208bbd 3b 61 08
                         CMP
                                     ESP, dword ptr [ECX + 0x8]
08208bc0 76 74
                         JBE
                                     LAB 08208c36
08208bc2 83 ec 1c
                         SUB
                                     ESP, 0x1c
08208bc5 c7 04 24
                         MOV
                                     dword ptr [ESP]=>local_lc,0x0
         00 00 00 00
08208bcc 8b 44 24 20
                         MOV
                                     EAX, dword ptr [ESP + param_1]
08208bd0 89 44 24 04
                                     dword ptr [ESP + local 18], EAX
08208bd4 8b 44 24 24
                                     EAX, dword ptr [ESP + param_2]
                         MOV
08208bd8 89 44 24 08
                                     dword ptr [ESP + local_14], EAX
                         MOV
08208bdc 8d 05 0e
                         LEA
                                     EAX, [s_/README_FOR_DECRYPT.txt_0827de0e]
         de 27 08
08208be2 89 44 24 0c
                         MOV
                                     dword ptr [ESP + local_10], EAX=>s_/README_FOR_DECRYPT.txt_0827de0e
08208be6 c7 44 24
                         MOV
                                     dword ptr [ESP + local_c],0x17
         10 17 00
         00 00
08208bee e8 dd c1
                         CALL
                                     runtime.concatstring2
```

Figure 51 - eCh0raix[9] defined string in listing view

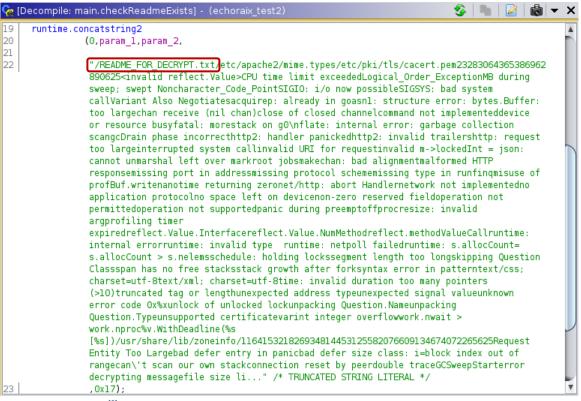


Figure 52 - eCh0raix^[9] defined string in decompile view



Future work

In this article we proposed solutions for two issues within Go binaries to help reverse engineers when they are using Ghidra to statically analyze malware written in Go. In the first topic we discussed how to recover function names in stripped Go binaries. Then we proposed multiple solutions for defining strings within Ghidra. The scripts that we created and files we used for the examples in this article are publicly available, the links can be found below.

There are even more possibilities to aid Go reverse engineering, the two topics that we discussed here are just the beginning. As a next step we are planning to dive deeper into Go function call conventions and type system.

In Go binaries arguments and return values are passed to functions using the stack, rather than registers. Currently Ghidra has a hard time to correctly detect these. Helping Ghidra to support Go's calling convention will help reverse engineers to understand the purpose of the analyzed functions.

The other interesting topic is types within Go binaries. Just like it was possible to extract function names from the investigated files, Go binaries also store information about the used types. Recovering these types can be a great help during reverse engineering. In the example below we recovered the main. Info structure in an eCh0raix ransomware sample^[9]. This structure tells us what information the malware is expecting from the C2 server.

```
main.info struct
                                                                      XREF[3]:
                                                                                   main.getInfo:082085fc(*),
                                                                                   main.getInfo:08208602(*),
                                                                                   08225100(*)
0824bd20 10 00 00 00
0824bd24 0c 00 00 00
                         ddw
                                     Ch
0824bd28 15 e7 c0 27
                         ddw
                                     27C0E715h
0824bd2c 07
                         db
0824bd2d 04
                         db
                                     4h
0824bd2e 04
                         db
                                     4h
0824bd2f 19
                         db
0824bd30 28 c8 20 08
                         addr
                                     PTR PTR type..hash.main.Info 0820c828
0824bd34 fc a0 2b 08
                                     DAT 082ba0fc
                         addr
0824bd38 20 75 00 00
                         ddw
                                     7520h
0824bd3c e0 a0 01 00
                                     1A0E0h
                         ddw
0824bd40 00 00 00 00
                         ddw
0824bd44 60 bd 24 08
                         addr
                                     PTR_rsapublickey_structfield_0824bd60
0824bd48 02 00 00 00
                         ddw
0824bd4c 02 00 00 00
                         ddw
                                     2h
0824bd50 5c 0d 00 00
                         ddw
                                     D5Ch
0824bd54 00 00
                         dw
                                     0h
0824bd56 00 00
                         dw
                                     0h
0824bd58 28 00 00 00
                         ddw
                                     28h
0824bd5c 00 00 00 00
```

Figure 53 - eCh0raix[9] main.info structure



```
PTR_rsapublickey_structfield_0824bd60
                                                                       XREF[1]: 0824bd44(*)
  0824bd60 60 aa 22 08
                           addr
                                      rsapublickey_structfield
  0824bd64 a0 a7 23 08
                           addr
                                      string_type
  0824bd68 00 00 00 00
                          ddw
                                      0h
  0824bd6c 18 cf 21 08
                           addr
                                      readme_structfield
  0824bd70 a0 a7 23 08
                           addr
                                      string_type
  0824bd74 10 00 00 00
                           ddw
                                      10h
Figure 54 - eCh0raix[9] main.info fields
```

```
type main.Info struct{
    RsaPublicKey string
    Readme string
}
```

Figure 55 - eCh0raix^[9] main.info structure

As these examples illustrated there are still a lot of interesting areas to discover within Go binaries from reverse engineering point of view. So, stay tuned for our next write-up.



GitHub repository with scripts and additional materials

- https://github.com/getCUJO/ThreatIntel/tree/master/Scripts/Ghidra
- https://github.com/getCUJO/ThreatIntel/tree/master/Research materials/Golang reversing

Files used during the research

	File name	SHA-256
[1]	hello.c	ab84ee5bcc6507d870fdbb6597bed13f858bbe322dc566522723fd8669a6d073
[2]	hello.go	2f6f6b83179a239c5ed63cccf5082d0336b9a86ed93dcf0e03634c8e1ba8389b
[3]	hello_c	efe3a095cea591fe9f36b6dd8f67bd8e043c92678f479582f61aabf5428e4fc4
[4]	hello_c_strip	95bca2d8795243af30c3c00922240d85385ee2c6e161d242ec37fa986b423726
[5]	hello_go	4d18f9824fe6c1ce28f93af6d12bdb290633905a34678009505d216bf744ecb3
[6]	hello_go_strip	45a338dfddf59b3fd229ddd5822bc44e0d4a036f570b7eaa8a32958222af2be2
[7]	hello_go.exe	5ab9ab9ca2abf03199516285b4fc81e2884342211bf0b88b7684f87e61538c4d
[8]	hello_go_strip.exe	ca487812de31a5b74b3e43f399cb58d6bd6d8c422a4009788f22ed4bd4fd936c
[9]	eCh0raix - x86	154dea7cace3d58c0ceccb5a3b8d7e0347674a0e76daffa9fa53578c036d9357
[10]	eCh0raix - ARM	3d7ebe73319a3435293838296fbb86c2e920fd0ccc9169285cc2c4d7fa3f120d
[11]	Kaiji - x86_64	f4a64ab3ffc0b4a94fd07a55565f24915b7a1aaec58454df5e47d8f8a2eec22a
[12]	Kaiji - ARM	3e68118ad46b9eb64063b259fca5f6682c5c2cb18fd9a4e7d97969226b2e6fb4



References and further reading

- https://rednaga.io/2016/09/21/reversing go binaries like a pro/
- https://2016.zeronights.ru/wp-content/uploads/2016/12/GO Zaytsev.pdf
- https://carvesystems.com/news/reverse-engineering-go-binaries-using-radare-2-and-python/
- https://www.pnfsoftware.com/blog/analyzing-golang-executables/
- https://github.com/strazzere/golang-loader-assist/blob/master/Bsides-GO-Forth-And-Reverse.pdf
- https://github.com/radareorg/r2con2020/blob/master/day2/r2 Gophers-AnalysisOfGoBinariesWithRadare2.pdf

Solutions by other researchers for various tools

IDA Pro

- https://github.com/sibears/IDAGolangHelper
- https://github.com/strazzere/golang loader assist

radare2 / Cutter

- https://github.com/f0rki/r2-go-helpers
- https://github.com/JacobPimental/r2-gohelper/blob/master/golang helper.py
- https://github.com/CarveSvstems/gostringsr2

Binary Ninja

• https://github.com/f0rki/bn-goloader

Ghidra

- https://github.com/felberj/gotools
- https://github.com/ghidraninja/ghidra-scripts/blob/master/golang-renamer.py

