#### **KU LEUVEN**

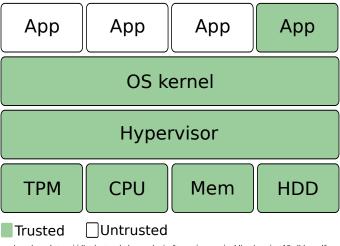


# A Security Analysis of Interrupts in Embedded Enclaved Execution

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Distrinet, KU Leuven
June 28, 2019

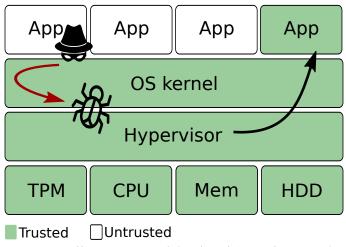


#### Conventional Software Isolation



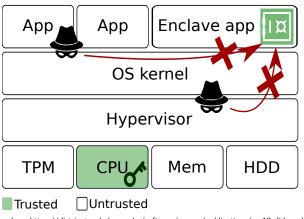
 $based\ on:\ https://distrinet.cs.kuleuven.be/software/sancus/publications/ccs18-slides.pdf$ 

## TCB: Security Risk



 $based\ on:\ https://distrinet.cs.kuleuven.be/software/sancus/publications/ccs18-slides.pdf$ 

#### **Enclaved Execution**



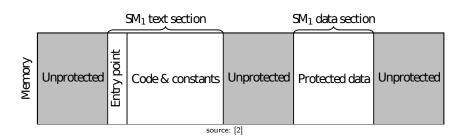
based on: https://distrinet.cs.kuleuven.be/software/sancus/publications/ccs18-slides.pdf

Key feature: Isolation and Attestation

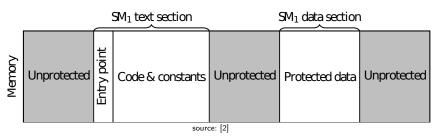
Embedded PMA

- Embedded PMA
- Hardware TCB only

- Embedded PMA
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- Enclave layout



- Embedded PMA
- Hardware TCB only
- Enclave layout



• Fully abstract compilation [1]

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2

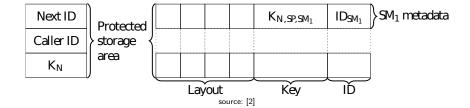
### Sancus: Isolation

Program counter based access control policy

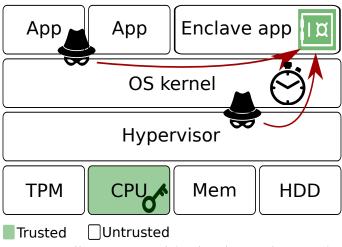
$From \backslash To$	Entry	Text	Data	Unprotected
Entry	r-x	r-x	rw-	rwx
Text	r-x	r-x	rw-	rwx
Other	X			rwx

source: [2]

### Sancus: Attestation



#### Side-Channel Attacks



 $based\ on:\ https://distrinet.cs.kuleuven.be/software/sancus/publications/ccs18-slides.pdf$ 

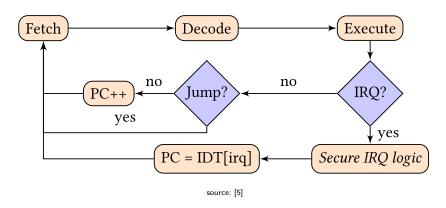
### Side-Channels inside Intel



Abusing microarchitectural optimizations

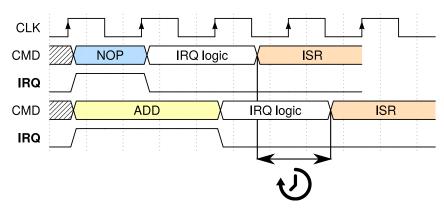
 $\rightarrow$  Not present in embedded

## Nemesis [5]: Fetch-Decode-Execute Cycle



Interrupts only served on instruction retirement

## Nemesis [5]: Interrupt Latency

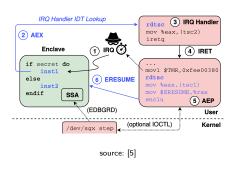


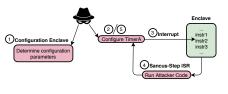
source: https://distrinet.cs.kuleuven.be/software/sancus/publications/ccs18-slides.pdf

### Attacker Frameworks

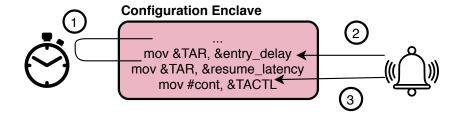
SGX-Step: Increase granularity and ease of use

Sancus-Step: Ease of use

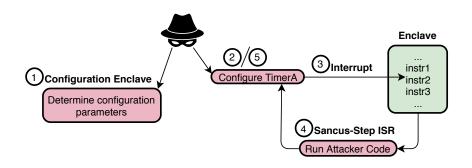




## Sancus-Step Configuration



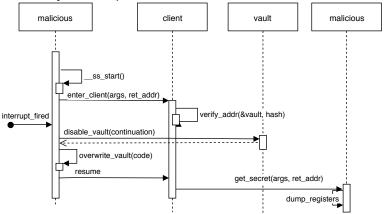
## Sancus-Step Overview



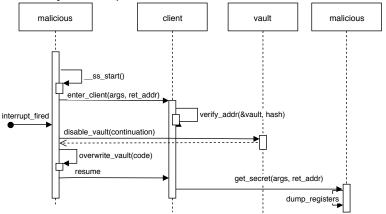
• Making enclaves interruptible  $\rightarrow$  well known issues [2, 3, 4]

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- ... In theory, not in practice

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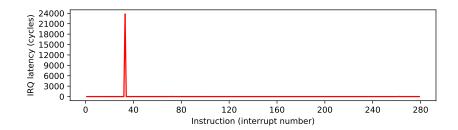
- Making enclaves interruptible  $\rightarrow$  well known issues [2, 3, 4]
- ... In theory, not in practice



Our ad-hoc implementation reveals vulnerabilities



## Verifcation TOCTOU: Automatic Exploit



# Side-Channel Analysis in Compiler Generated Code

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- Stubs:
  - Multiplication stub
  - Other arithmetic stubs
  - Entry and Exit stub

# Side-Channel Analysis in Compiler Generated Code

- Stubs:
  - Multiplication stub
  - Other arithmetic stubs
  - Entry and Exit stub
- Both start-to-end and Nemesis

## Reversing the Multiplication Stub

```
__sm_mulhi3:
                       unsigned int mul(unsigned int a,
            r15, r13
                                         unsigned int b)
    mov
    clr
            r15
1: tst r14
                            unsigned int rv = 0;
                            while (a!=0)
    jz
          3f
    clrc
            r13
                                if (b & 1)
    rrc
    jnc
            2f
                                    rv += a;
    add
            r14, r15
                               b >>= 1;
2: rla
            r14
                                a <<= 1:
    tst
            r13
                                if (b == 0)
    jnz
            1b
                                    break;
3:
    ret
                            return rv;
                       }
```

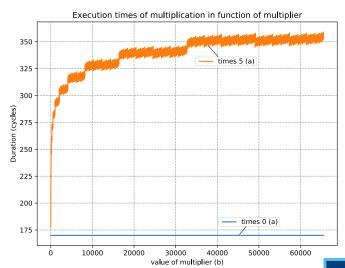
### Intuition of the Multiplication Stub

```
unsigned int mul(unsigned int a,
                 unsigned int b)
                          10(a)
x 101(b)
    unsigned int rv = 0;
    while (a!=0)
        if (b & 1)
                                  10
            rv += a;
        b >>= 1;
        a <<= 1:
        if (b == 0)
                             1010
            break;
    return rv;
}
```

## Case Study: SM\_mul

```
#include <sancus/sm support.h>
DECLARE SM(SM mul, 0x1234);
int SM DATA(SM mul) a = -63;
int SM DATA(SM mul) b = 6;
int SM ENTRY(SM mul) sm multiply(void)
    volatile int c = a * b;
    return 0;
```

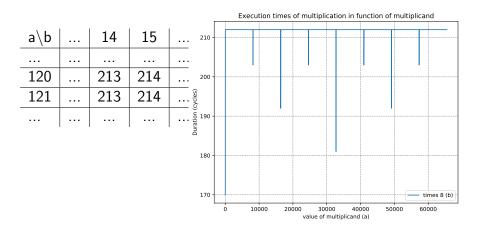
## Start-to-End Side-Channel in Multiplication



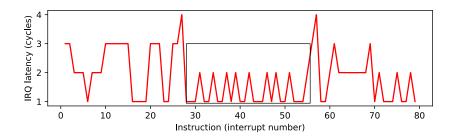
# Start-to-End Side-Channel in Multiplication



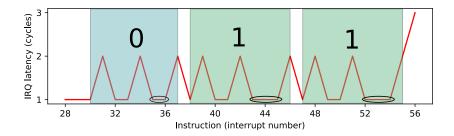
## Start-to-End Side-Channel in Multiplication



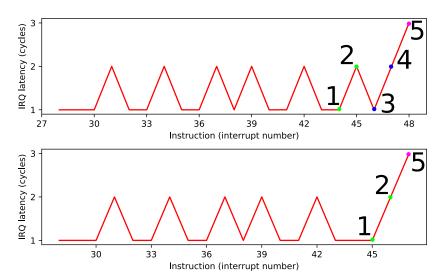
# Nemesis Side-Channel in Multiplication



## Leaking the Multiplier



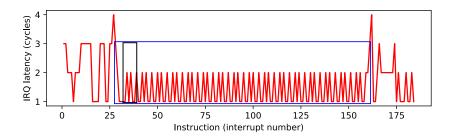
## Determining the Stop Condition



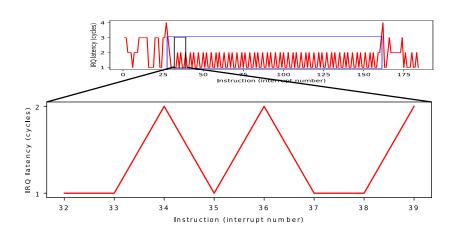
## Hardening the Multiplication Stub

```
__sm_mulhi3:
                        unsigned int mul(unsigned int a,
            #16, r12
                                          unsigned int b)
    mov
            r15, r13 {
    mov
    clr r15
                            unsigned int rv = 0;
1: clrc
                            for (int i = 0; i < 16; i++)
            r13
    rrc
                                if (b & 1)
    jnc
            2f
    add
            r14, r15
                                     rv += a:
            4f
                                else
    jmp
2:
                                     asm("nop");
    nop
            4f
                                b >>= 1:
    jmp
4:
    rla
            r14
                                a <<= 1;
    sub
            #1, r12
    jnz
            1b
                            return rv;
                        }
3:
    ret
```

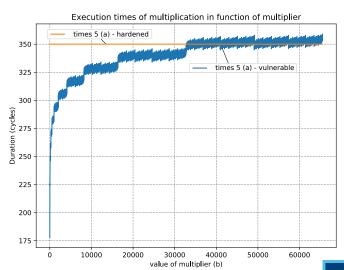
# **Closing Nemesis**



# **Closing Nemesis**



## Performance Comparison



## Future and Ongoing Work

- Analyse other stubs
  - Side-channel leakage
  - Atomicity issues
- Defense
  - Shortterm  $\rightarrow$  software
  - Longterm  $\rightarrow$  hardware (ongoing work [6])

Thank you! Questions?

### References I

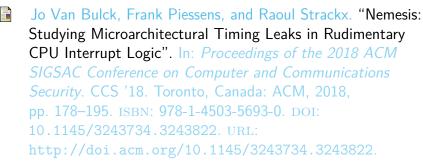
- Pieter Agten et al. "Secure Compilation to Modern Processors". In: 2012 IEEE 25th Computer Security Foundations Symposium. IEEE, June 2012. DOI: 10.1109/csf.2012.12. URL: https://doi.org/10.1109/csf.2012.12.
- Job Noorman et al. "Sancus 2.0: A Low-Cost Security Architecture for IoT Devices". In: *ACM Transactions on Privacy and Security (TOPS)* 20.3 (Sept. 2017), 7:1–7:33.

### References II

Raoul Strackx and Frank Piessens. "Fides: Selectively Hardening Software Application Components Against Kernel-level or Process-level Malware". In: Proceedings of the 2012 ACM Conference on Computer and Communications Security. CCS '12. Raleigh, North Carolina, USA: ACM, 2012, pp. 2–13. ISBN: 978-1-4503-1651-4. DOI: 10.1145/2382196.2382200. URL: http://doi.acm.org/10.1145/2382196.2382200.

Raoul Strackx, Frank Piessens, and Bart Preneel. "Efficient Isolation of Trusted Subsystems in Embedded Systems". In: vol. 50. Sept. 2010, pp. 344–361. DOI: 10.1007/978-3-642-16161-2 20.

### References III



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