

A large, stylized 'X' logo composed of two overlapping chevron shapes, rendered in a dark blue-grey color, serves as a background for the text.


Unstructured Recursion

Recursion

```
fac :: Int -> Int
fac n = if n == 0
        then 1
        else fac (n - 1)
```

Recursion

```
fac :: Int -> Int
fac n = if n == 0
        then 1
        else fac (n - 1)
```



Explicit Recursion

`fac = fix fac'`

Explicit Recursion

non-recursive
definition

`fac = fix fac'`

Explicit Recursion

non-recursive
definition

`fac = fix fac'`

generic fixpoint
combinator

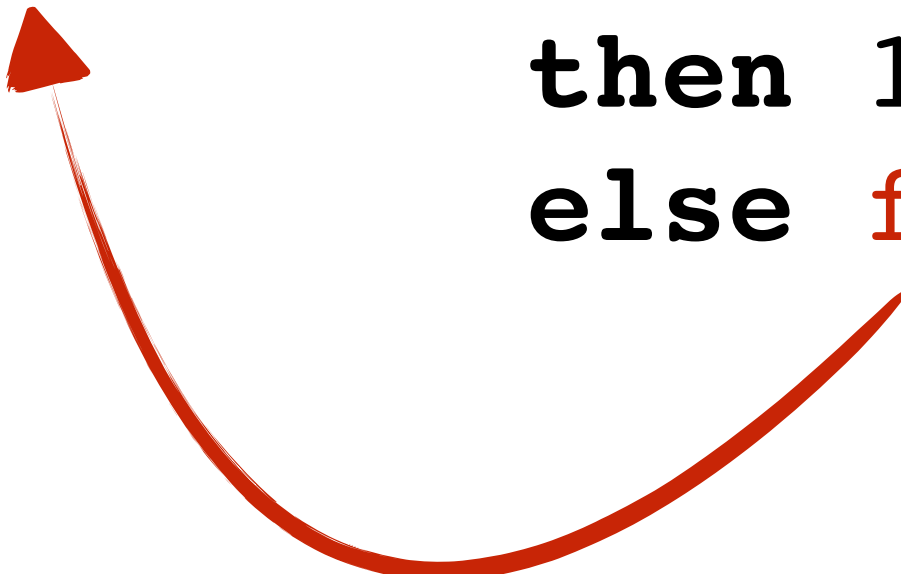
Non-Recursive Version

```
fac' f n = if n == 0  
           then 1  
           else f (n - 1)
```

1 level of the
recursion

Non-Recursive Version

```
fac' f n = if n == 0  
          then 1  
          else f (n - 1)
```

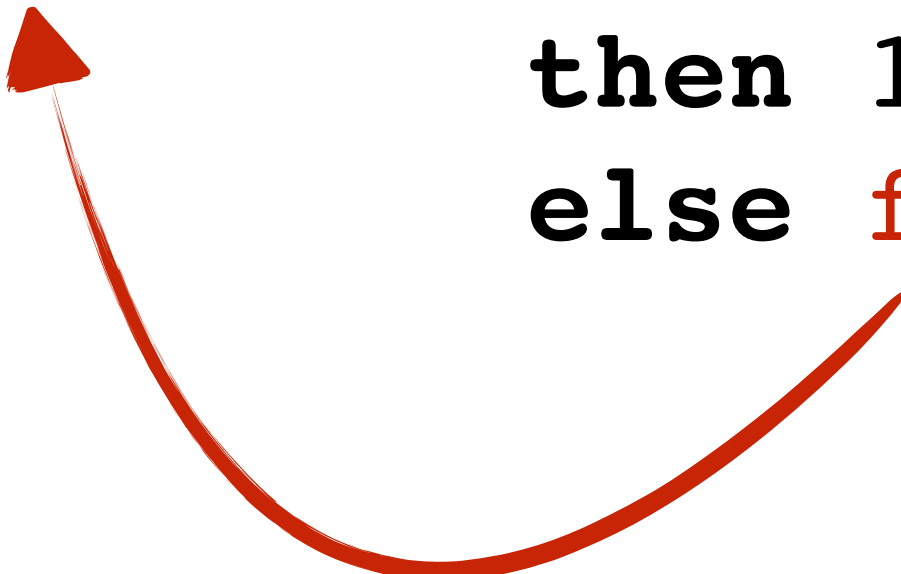


1 level of the
recursion

parameter for
the other levels

Non-Recursive Version

```
fac' :: (Int -> Int) -> (Int -> Int)
fac' f n = if n == 0
           then 1
           else f (n - 1)
```



1 level of the
recursion

parameter for
the other levels

Fixpoint Combinator

`fix :: (a -> a) -> a`

Fixpoint Combinator

expression with
hole

`fix :: (a -> a) -> a`

Fixpoint Combinator

expression with
hole

hole plugged
with itself

`fix :: (a -> a) -> a`

Fixpoint Combinator

expression with
hole

hole plugged
with itself

`fix :: (a -> a) -> a`

`fix f = f (fix f)`

Fixpoint Combinator

expression with
hole

hole plugged
with itself

```
fix :: (a -> a) -> a
fix f = f (fix f)
      = f (f (fix f))
```

Fixpoint Combinator

expression with
hole

hole plugged
with itself

`fix :: (a -> a) -> a`

`fix f = f (fix f)`

`= f (f (fix f))`

`= f (f (f (fix f)))`

Fixpoint Combinator

expression with
hole

hole plugged
with itself

$\text{fix} :: (a \rightarrow a) \rightarrow a$

$\text{fix } f = f (\text{fix } f)$
 $= f (f (\text{fix } f))$
 $= f (f (f (\text{fix } f)))$
 $= f (f (f (f \dots)))$

f all the way
down

General Recursion

```
fix :: (a -> a) -> a  
fix f = f (fix f)
```

General Recursion

`fix :: (a -> a) -> a`
`fix f = f (fix f)`

makes language
Turing complete

General Recursion

`fix :: (a -> a) -> a`
`fix f = f (fix f)`

can express all
recursion

makes language
Turing complete

General Recursion Considered Harmful



General Recursion Considered Harmful

`loop = fix id`



General Recursion Considered Harmful



```
loop = fix id
```

```
diverge = fix go
```

where

```
go f n = f (n+1)
```


General Recursion Considered Harmful



```
loop = fix id
```

Easy to go over
the egde

```
diverge = fix go
```

where

```
go f n = f (n+1)
```

Recursion Schemes



A handrail to hold on to

```
data List
  = Nil
  | Cons Int List
```

A handrail to hold on to

```
data List
  = Nil
  | Cons Int List
```

No irregular
syntax

A handrail to hold on to

```
data List
  = Nil
  | Cons Int List
```

No irregular
syntax

No type
parameter

A handrail to hold on to

Keeping things
simple

```
data List
  = Nil
  | Cons Int List
```

No irregular
syntax

No type
parameter

Recursion follows Recursion

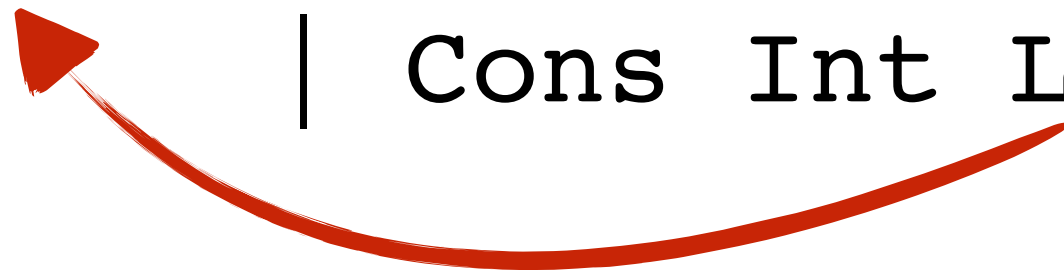
`sum Nil = 0`

`sum (Cons x xs) = x + sum xs`



`data List = Nil`

`| Cons Int List`



Sliding down the railing

`sum Nil = 0`

`sum (Cons x xs) = x + sum xs`

`prod Nil = 1`

`prod (Cons x xs) = x * prod xs`

Sliding down the railing

`sum Nil` = 0

`sum (Cons x xs)` = `x + sum xs`

`prod Nil` = 1

`prod (Cons x xs)` = `x * prod xs`

Sliding down the railing

`sum Nil` = 0

`sum (Cons x xs)` = `x` + `sum xs`

`prod Nil` = 1

`prod (Cons x xs)` = `x` * `prod xs`

Fold Recursion Scheme

`sum Nil = 0`

`sum (Cons x xs) = x + sum xs`

`prod Nil = 1`

`prod (Cons x xs) = x * prod xs`

 **Abstract**

`fold n c Nil = n`

`fold n c (Cons x xs) = c x (fold n c xs)`

Fold Recursion Scheme

`sum Nil` = 0

`sum (Cons x xs)` = `x + sum xs`

`prod Nil` = 1

`prod (Cons x xs)` = `x * prod xs`

 **Abstract**

`fold n c Nil` = n

`fold n c (Cons x xs)` = `c x (fold n c xs)`

Fold Recursion Scheme

`sum Nil` = 0

`sum (Cons x xs)` = `x` + `sum xs`

`prod Nil` = 1

`prod (Cons x xs)` = `x` * `prod xs`



Abstract

`fold n c Nil` = `n`

`fold n c (Cons x xs)` = `c` `x` (`fold n c xs`)

Fold Recursion Scheme

```
sum = fold 0 (+)
```

```
prod = fold 1 (*)
```

```
fold n c Nil = n
```

```
fold n c (Cons x xs) = c x (fold n c xs)
```

Structured Recursion

```
fold n c Nil
```

```
  = n
```

```
fold n c (Cons x xs)
```

```
  = c x (fold n c xs)
```

Structured Recursion

```
fold n c Nil
  = n
fold n c (Cons x xs)
  = c x (fold n c xs)
```

Recursion follows
list structure

Structured Recursion

```
fold n c Nil
  = n
fold n c (Cons x xs)
  = c x (fold n c xs)
```

Recursion follows
list structure

Recursive call
over recursive list

Structured Recursion

```
fold n c Nil
    = n
fold n c (Cons x xs)
    = c x (fold n c xs)
```

Recursion follows
list structure

Recursive call
over recursive list

Terminates
(for finite list)



Recursive Types

Explicit Recursion

`fac = fix fac'`

Explicit Recursion

non-recursive
definition

`fac = fix fac'`

Explicit Recursion

non-recursive
definition

`fac = fix fac'`

generic fixpoint
combinator

Explicit Recursion

fac = fix fac'



Explicit Recursion

fac = **fix** fac'



List <~> **Fix** List'

Explicit Recursion

`fac = fix fac'`



`List <~> Fix List'`

generic fixpoint
combinator

Explicit Recursion

`fac = fix fac'`



`List <~> Fix List'`

non-recursive
definition

generic fixpoint
combinator

Fixpoint Combinator

```
fix :: (a -> a) -> a  
fix f = f (fix f)
```

Fixpoint Combinator

`fix :: (a -> a) -> a`

`fix f = f (fix f)`



Fixpoint Combinator

```
fix :: (a -> a) -> a  
fix f = f (fix f)
```



```
-- Fix :: ( * -> * ) -> *  
type Fix f = f (Fix f)
```

Fixpoint Combinator

```
fix :: (a -> a) -> a  
fix f = f (fix f)
```



Recursive Type
Synonyms not
supported

```
-- Fix :: ( * -> * ) -> *  
type Fix f = f (Fix f)
```

Fixpoint Combinator

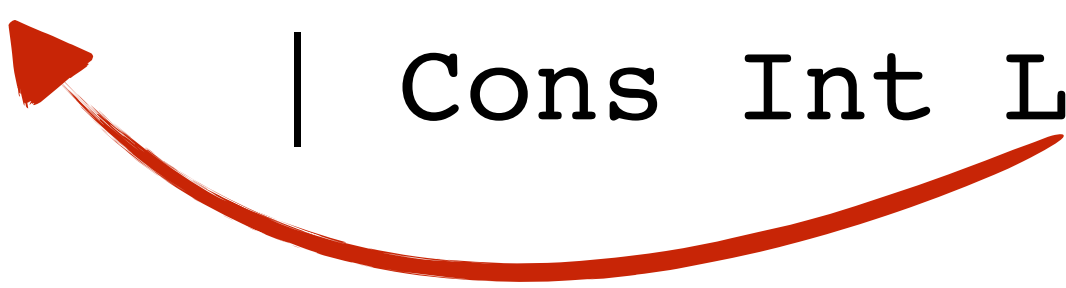
```
fix :: (a -> a) -> a  
fix f = f (fix f)
```



```
-- Fix :: ( * -> * ) -> *  
data Fix f = In (f (Fix f))
```

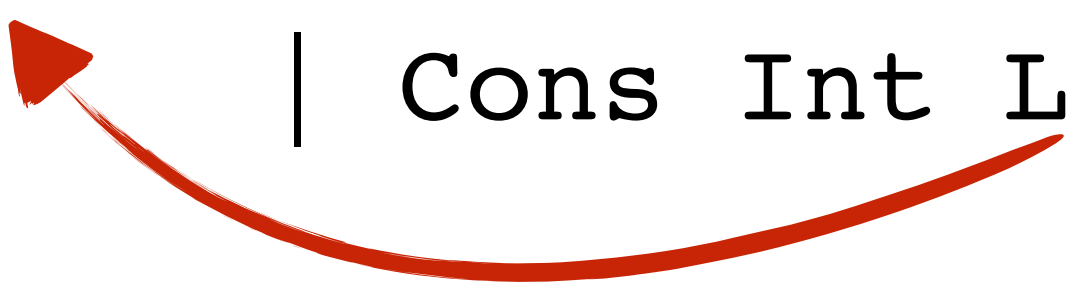
Non-Recursive Version

```
-- List :: *  
data List  = Nil  
           | Cons Int List
```



Non-Recursive Version

```
-- List :: *  
data List = Nil  
          | Cons Int List
```



```
-- List' :: * -> *  
data List' r = Nil'  
             | Cons' Int r
```

Isomorphism

```
iso :: List <~> Fix List'
iso = Iso t f where
    t Nil = In Nil'
    t (Cons x xs) = In (Cons' x (t xs))

    f (In Nil') = Nil
    f (In (Cons' x xs)) = Cons x (f xs)
```




Generic Fold

What fold does

```
Cons 1 (Cons 2 (Cons 3 Nil))
```

```
fold 0 (+)
```

What fold does

Cons 1 (Cons 2 (Cons 3 Nil))

(+) 1 ((+) 2 (+) 3 0))

fold 0 (+)

What fold does

Cons 1 (Cons 2 (Cons 3 Nil))



(+ 1 ((+ 2 ((+ 3 0))))

fold 0 (+)

What fold does

Cons 1 (Cons 2 (Cons 3 Nil))



(+ 1 ((+ 2 ((+ 3 0))))

fold 0 (+)

The Type of Fold

`fold`

`:: a -> (Int -> a -> a) -> List -> a`

`fold n c Nil`

`= n`

`fold n c (Cons x xs)`

`= c x (fold n c xs)`

The Type of Fold

`fold`

`:: a -> (Int -> a -> a) -> List -> a`

`fold n c Nil`

`= n`

`fold n c (Cons x xs)`

`= c x (fold n c xs)`

The List Algebra

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$
 $\quad \quad \quad \langle \sim \rangle$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} * a \rightarrow a) \rightarrow \dots$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} * a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} * a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$((\text{One} + \text{Int} * a) \rightarrow a) \rightarrow \dots$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} * a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$((\text{One} + \text{Int} * a) \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

The List Algebra

$a \rightarrow (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$a * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} \rightarrow a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{One} \rightarrow a) * (\text{Int} * a \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$((\text{One} + \text{Int} * a) \rightarrow a) \rightarrow \dots$

$\langle \sim \rangle$

$(\text{List}' a \rightarrow a) \rightarrow \dots$

Towards Generic Fold

fold

$:: a \rightarrow (Int \rightarrow a \rightarrow a) \rightarrow List \rightarrow a$

fold n c Nil

= n

fold n c (Cons x xs)

= c x (fold n c xs)

Towards Generic Fold

```
fold
  :: (List' a -> a) -> List -> a
fold alg Nil
  = alg Nil'
fold alg (Cons x xs)
  = alg (Cons' x (fold alg xs))
```

Towards Generic Fold

fold

$:: (\text{List}' \ a \rightarrow a) \rightarrow \text{Fix List}' \rightarrow a$

fold alg (**In Nil'**)

= alg Nil'

fold alg (**In (Cons' x xs)**)

= alg (Cons' x (fold alg xs))

Functor

```
class Functor f where
```

```
  fmap :: (r -> s) -> (f r -> f s)
```

```
instance Functor List' where
```

```
  fmap f Nil' = Nil'
```

```
  fmap f (Cons' x r) = Cons' x (f r)
```

Towards Generic Fold

```
fold
  :: (List' a -> a) -> Fix List' -> a
fold alg (In s)
  = alg (fmap (fold alg) s)
```


Generic Fold

```
fold
  :: Functor f
  => (f a -> a) -> Fix f -> a
fold alg (In s)
  = alg (fmap (fold alg) s)
```

Example

```
data Exp = Lit Int | Add Exp Exp
```

```
data Exp' r = Lit' Int | Add' r r
```

```
instance Functor Exp' where
```

```
  fmap f (Lit' n)
```

```
    = Lit' n
```

```
  fmap f (Add' e1 e2)
```

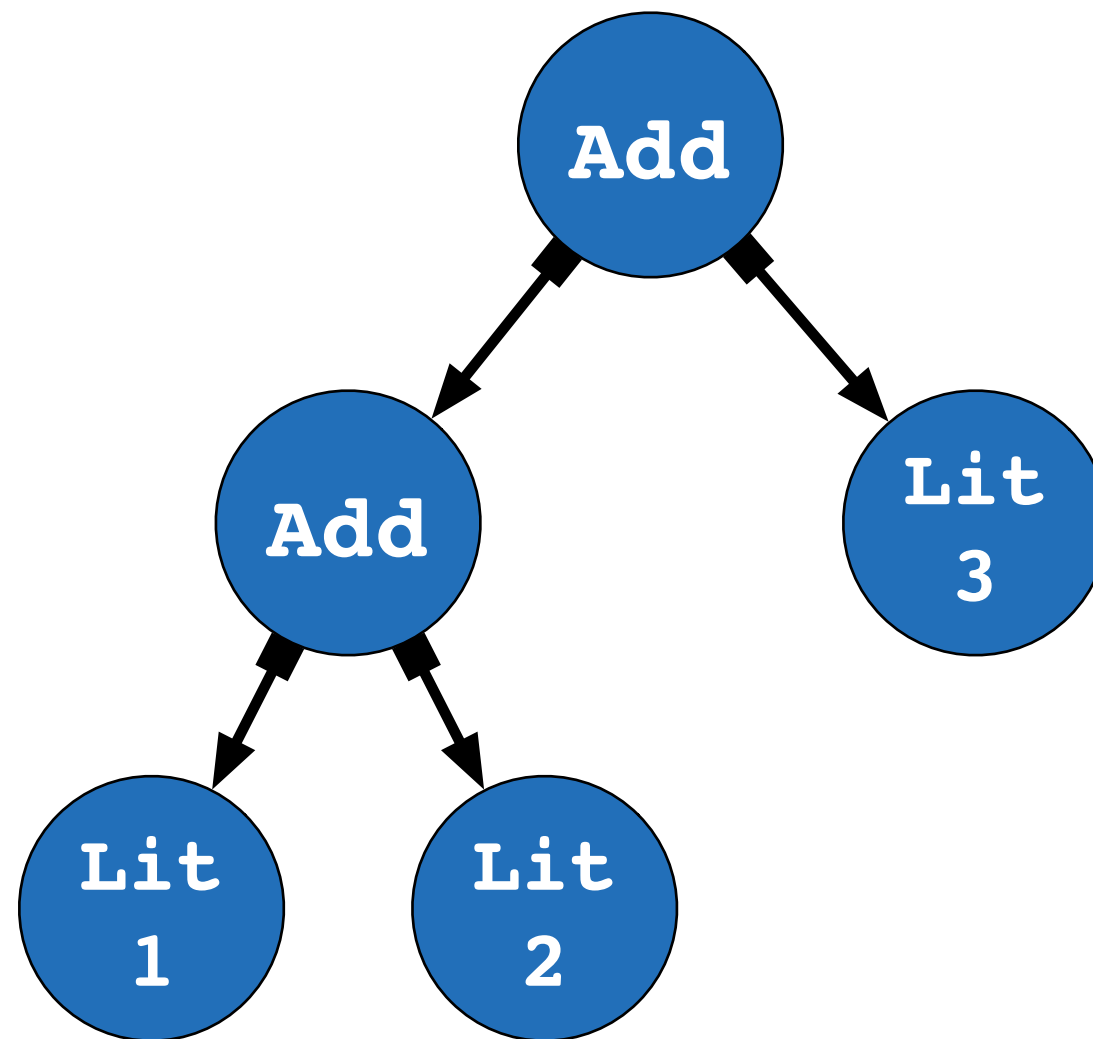
```
    = Add' (f e1) (f e2)
```

Example

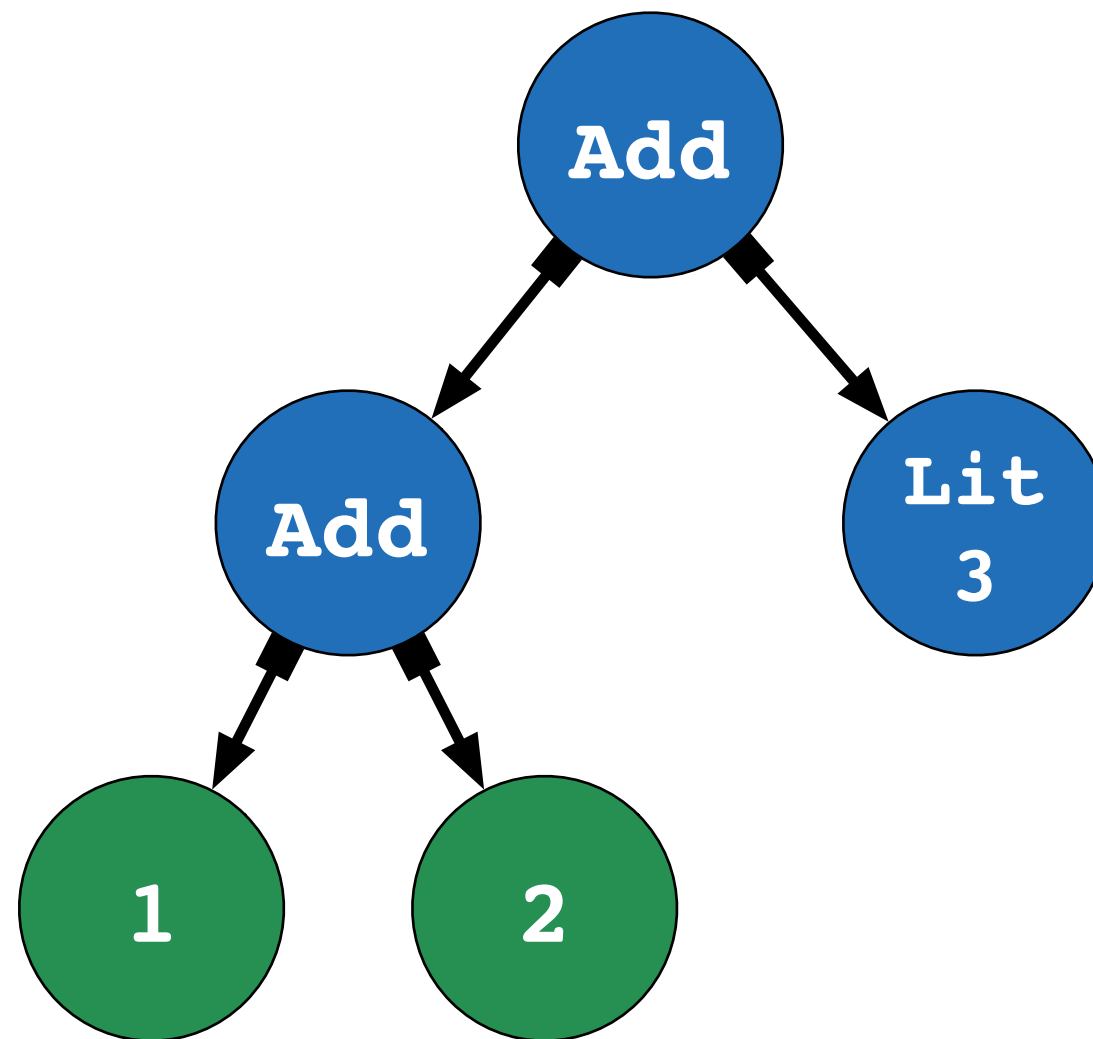
```
eval :: Fix Exp' -> Int  
eval = fold alg where
```

```
    alg :: Exp' Int -> Int  
    alg (Lit' n)      = n  
    alg (Add' x y)    = x + y
```

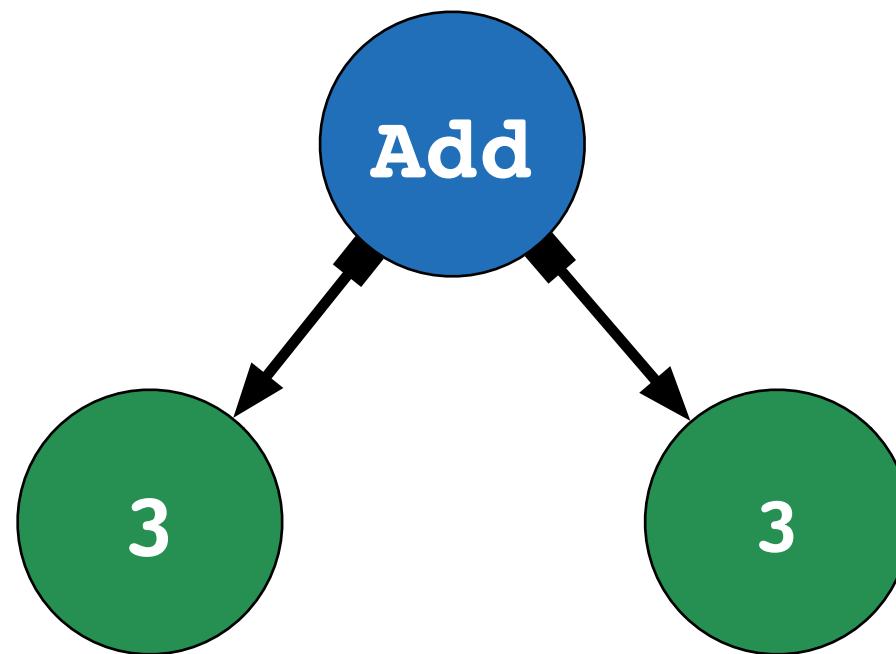
Example



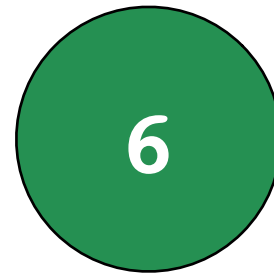
Example



Example



Example



A large, stylized 'X' logo composed of two overlapping, slightly offset rectangular shapes, one in a lighter shade of purple and one in a darker shade, creating a 3D effect. The logo is centered in the background.

Summary

Fold

★ Unstructured Recursion

- ✦ powerful
- ✦ dangerous

★ Structured Recursion with Fold

- ✦ safe
- ✦ often expressive enough
- ✦ datatype generic concept

More to Learn

- ★ Fold fusion

- ◆ parallel fold fusion

- ◆ fold / build fusion

- ◆ deforestation

- ★ Other recursion schemes

Parallel Fold

```
average :: List Float -> Float
average l =
    sum l
    /
    length l
```

Parallel Fold

```
average :: List Float -> Float
average l =
    fold 0 (+) l
    /
    fold 0 (\_ y -> 1 + y) l
```

Parallel Fold Fusion

```
average :: List Float -> Float
average l =
    uncurry (/) (fold nil cons l)
where
    nil          = (0,0)
    cons x (s,l) = (x + s, 1 + l)
```



Jasper Van der Jeugt

Enabling Fusion

2012
-
2013

GHC
compiler

pipeline of
recursive functions



Jasper Van der Jeugt

Enabling Fusion

2012
-
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GHC
compiler

compiler
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Jasper Van der Jeugt

Enabling Fusion

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-
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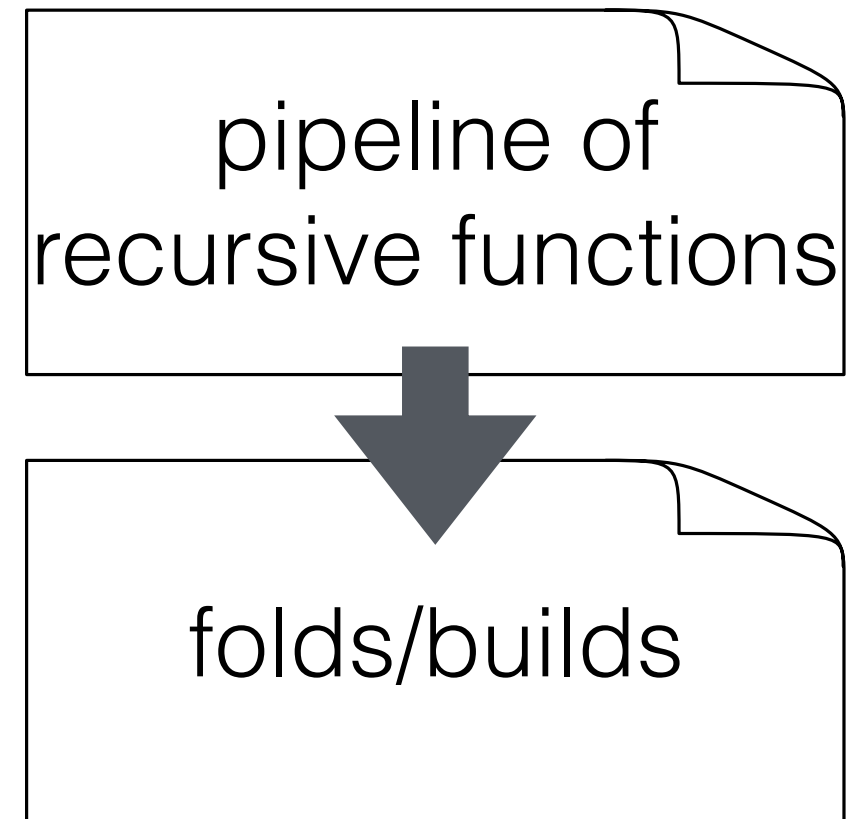
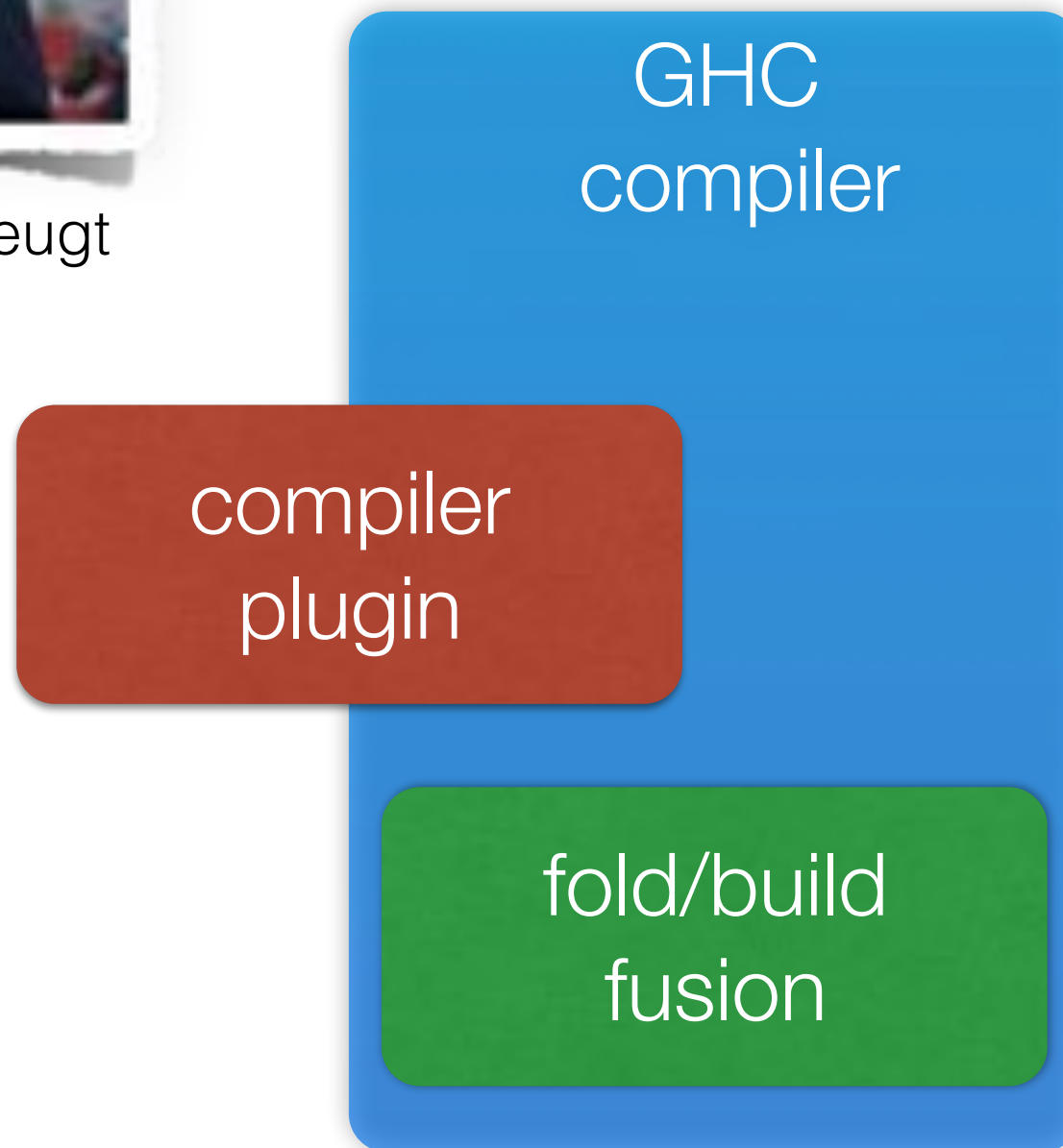
folds/builds



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Enabling Fusion

2012
-
2013

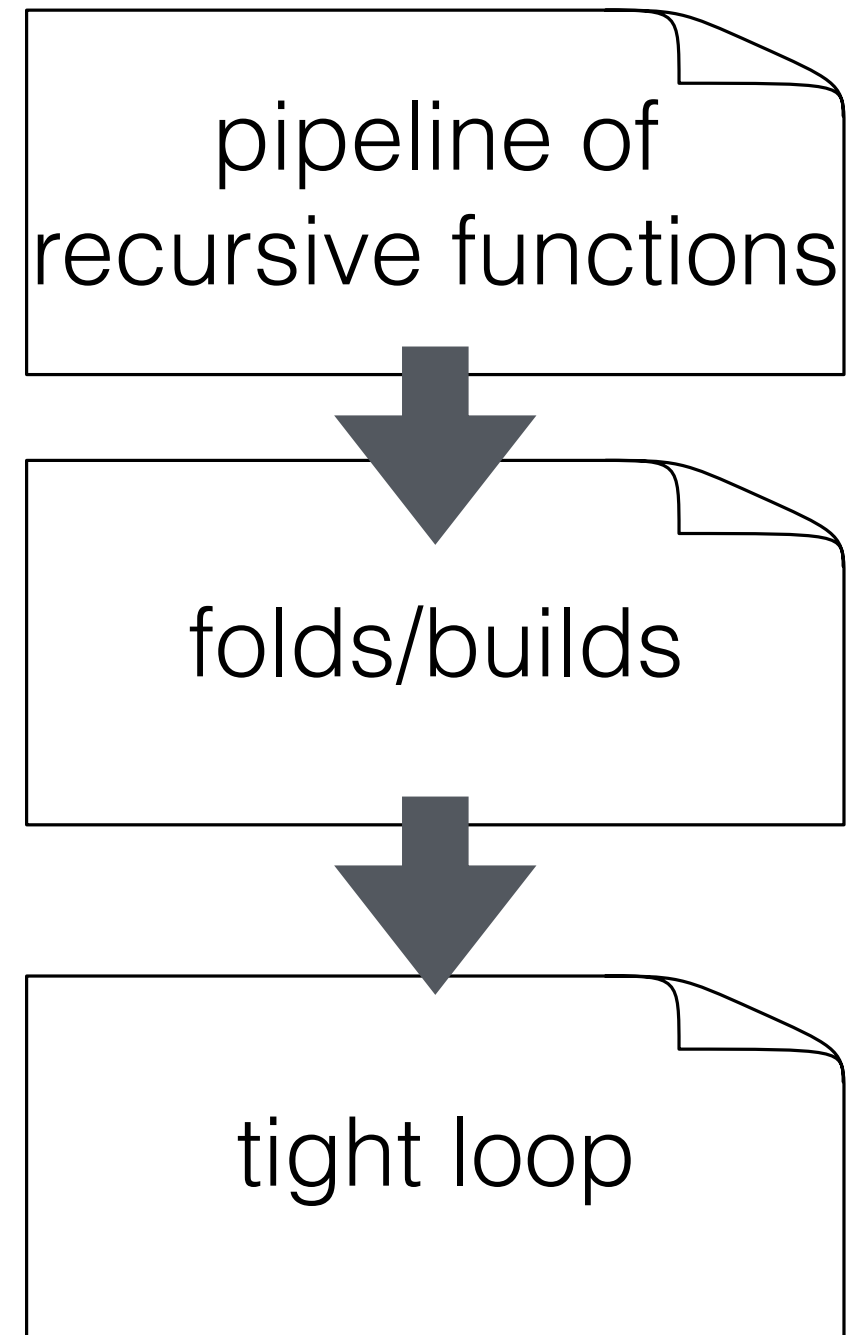
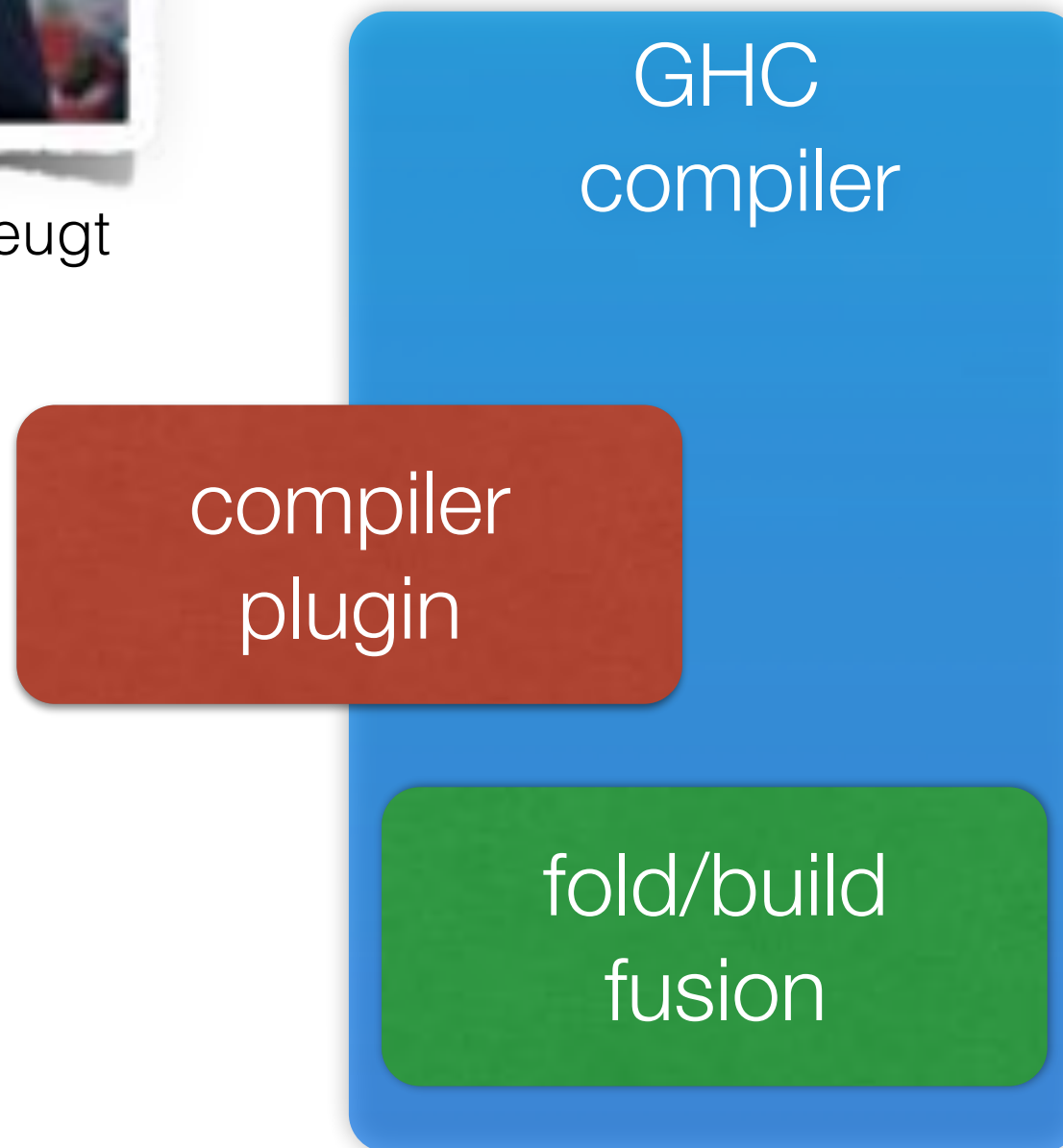




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2012
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2013

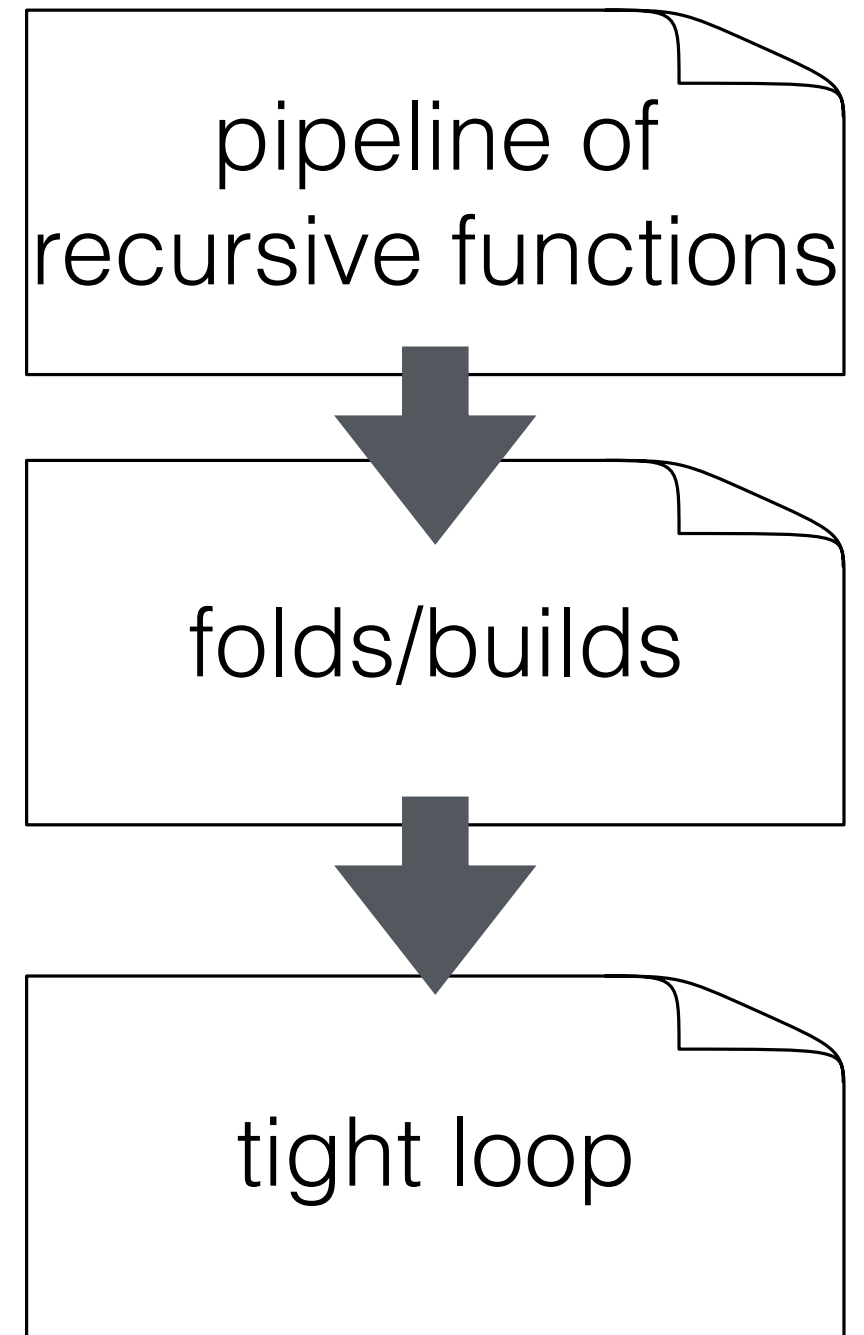
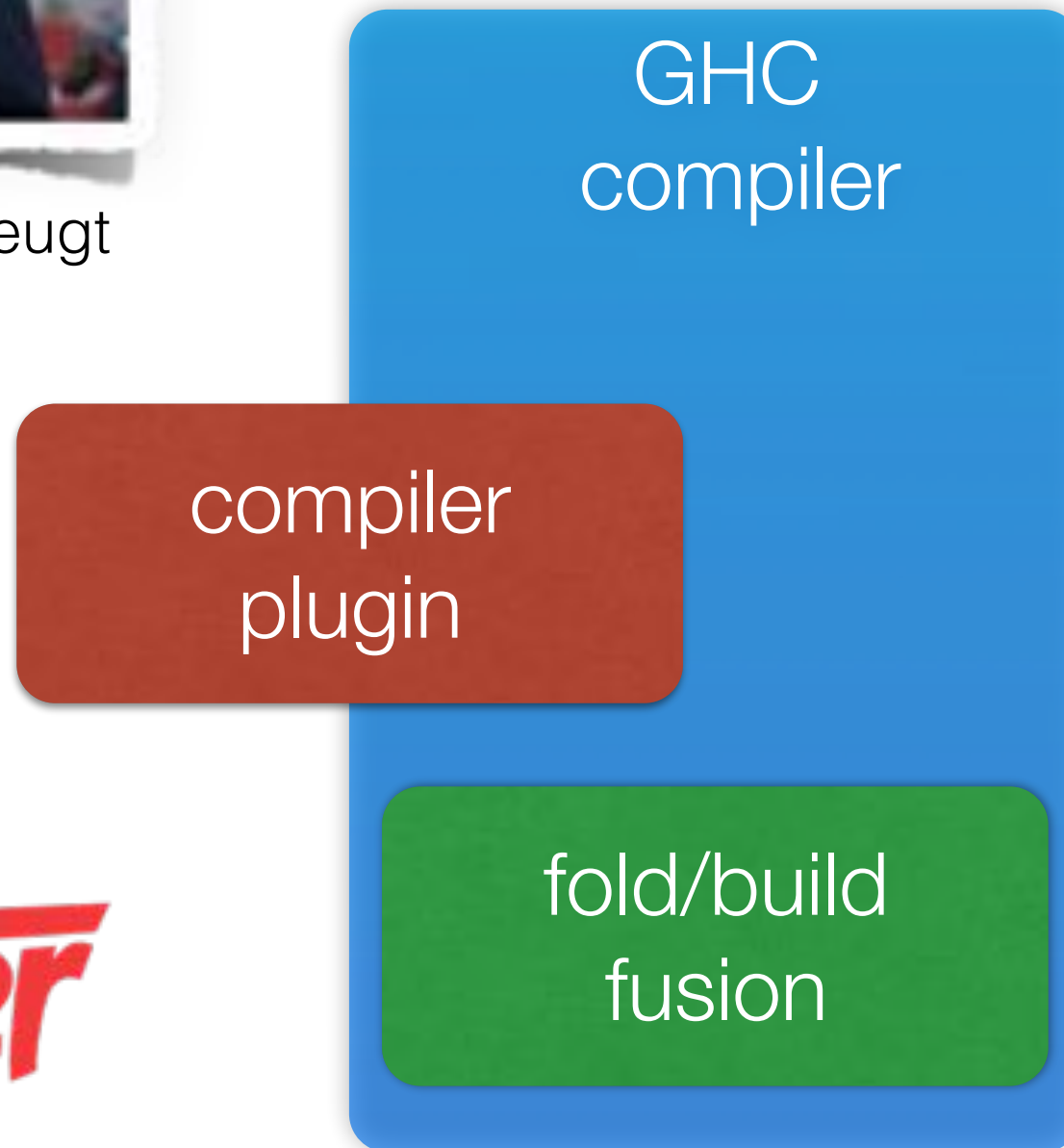




Jasper Van der Jeugt

Enabling Fusion

2012
-
2013



Fast Monads for Free

2015

-

2016



You

GHC
compiler

stack of monads

Fast Monads for Free

2015

-

2016



You

GHC
compiler

compiler
plugin

stack of monads

Fast Monads for Free

2015
-
2016

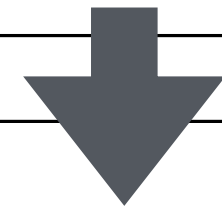


You

GHC
compiler

compiler
plugin

stack of monads



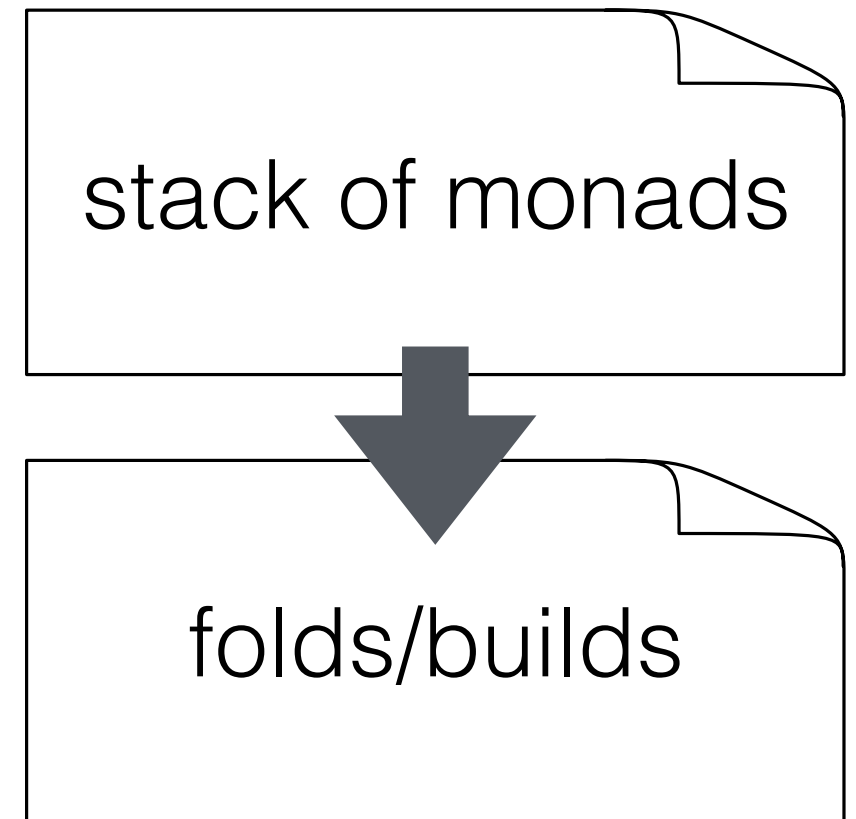
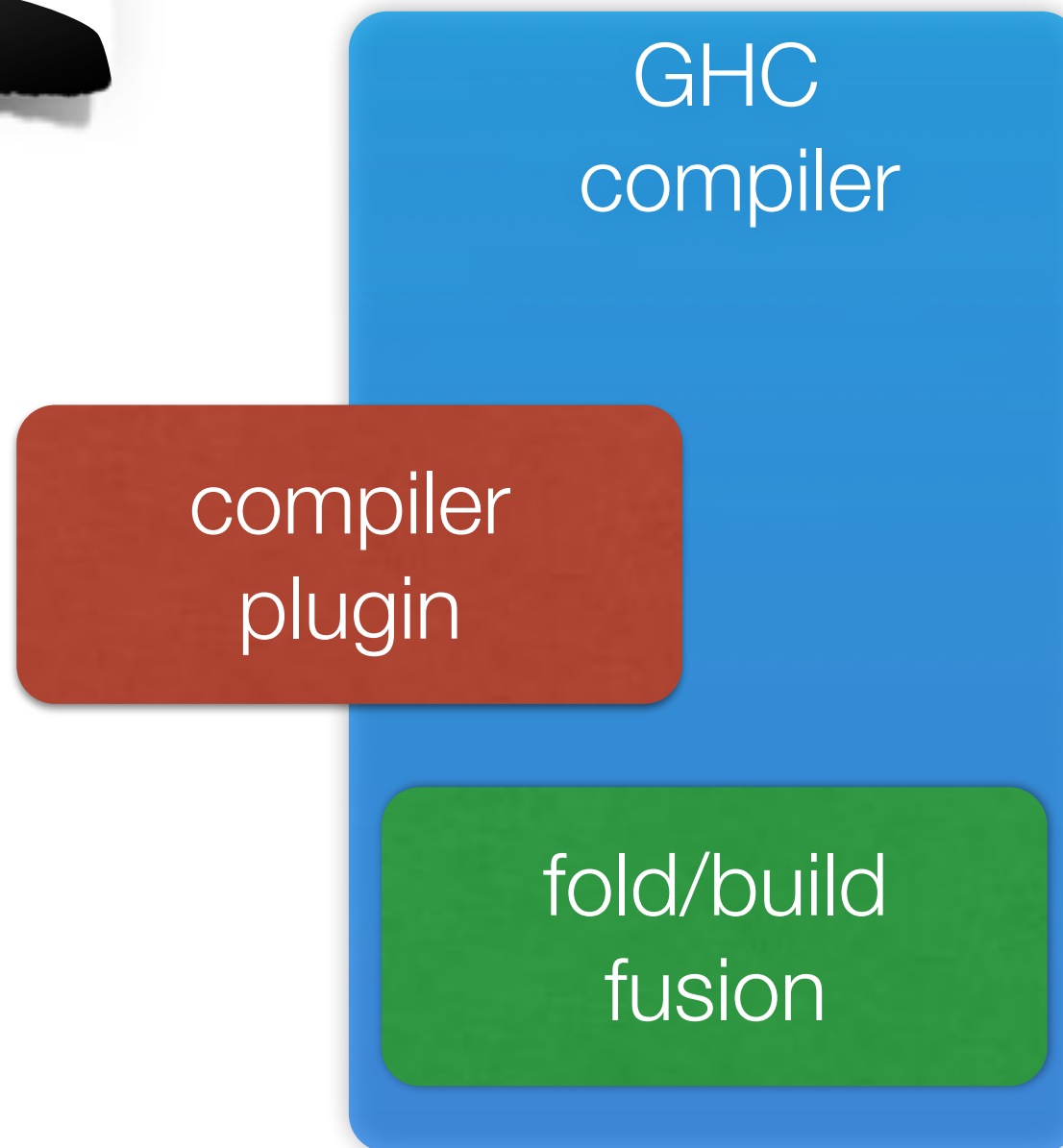
folds/builds

Fast Monads for Free

2015
-
2016



You

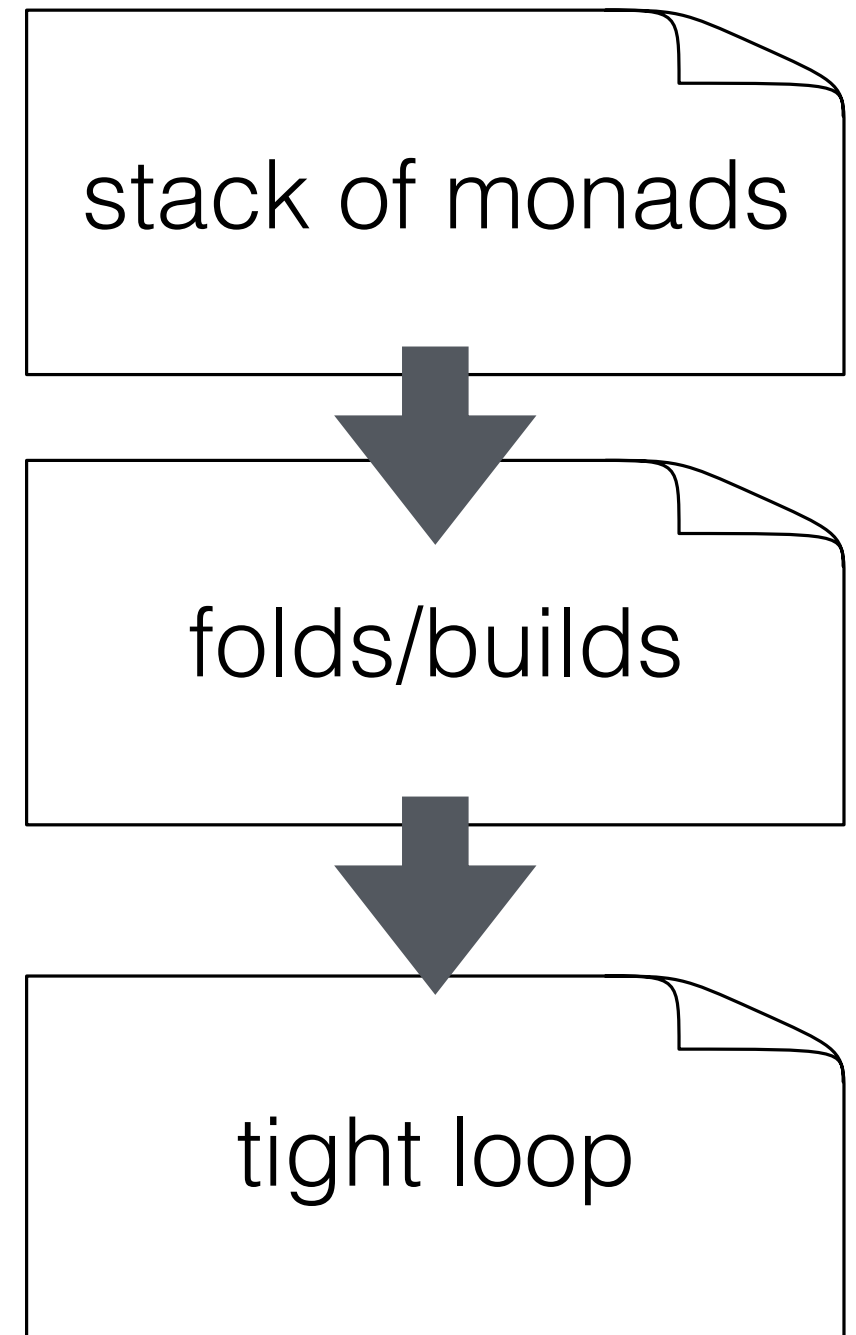
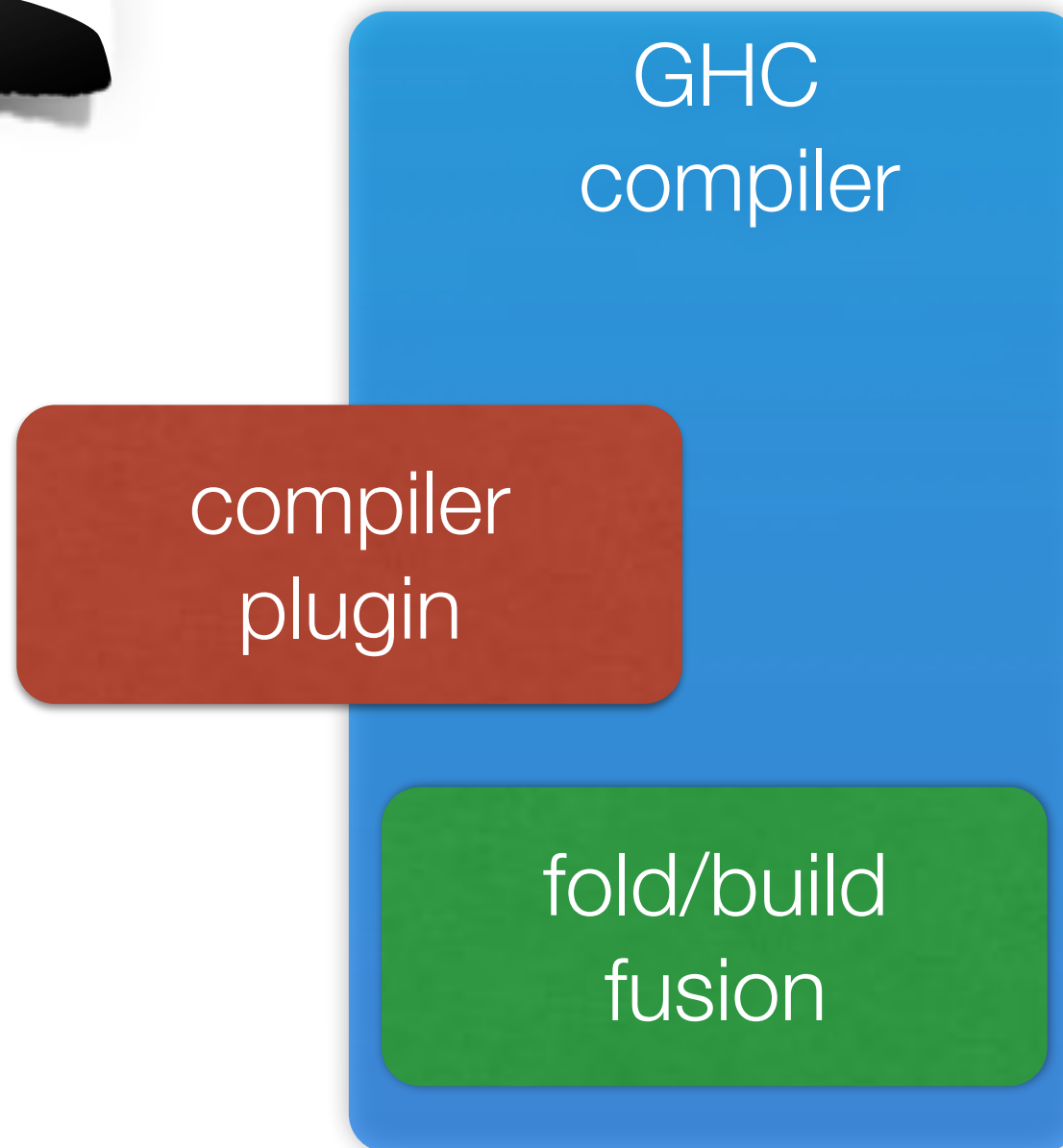


Fast Monads for Free

2015
-
2016



You

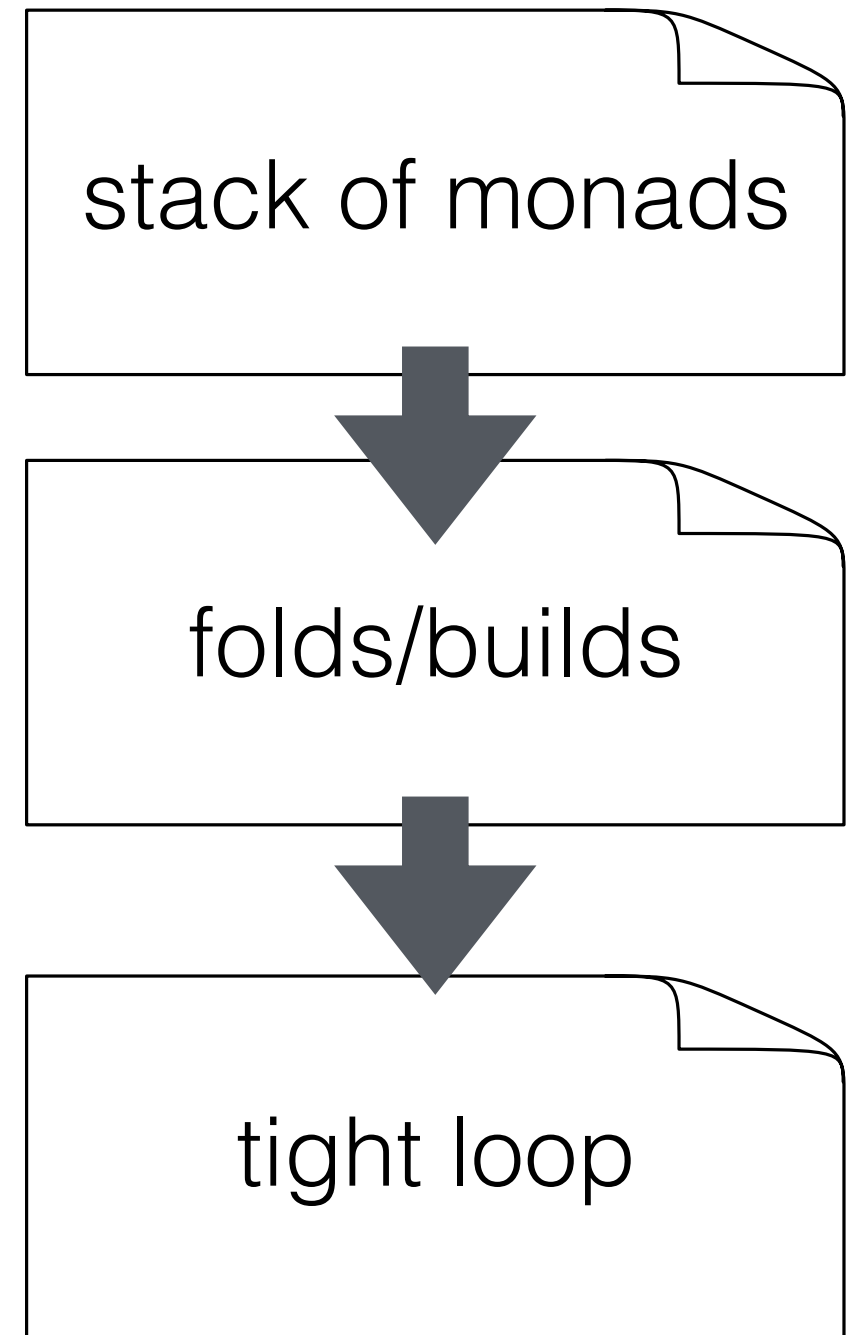
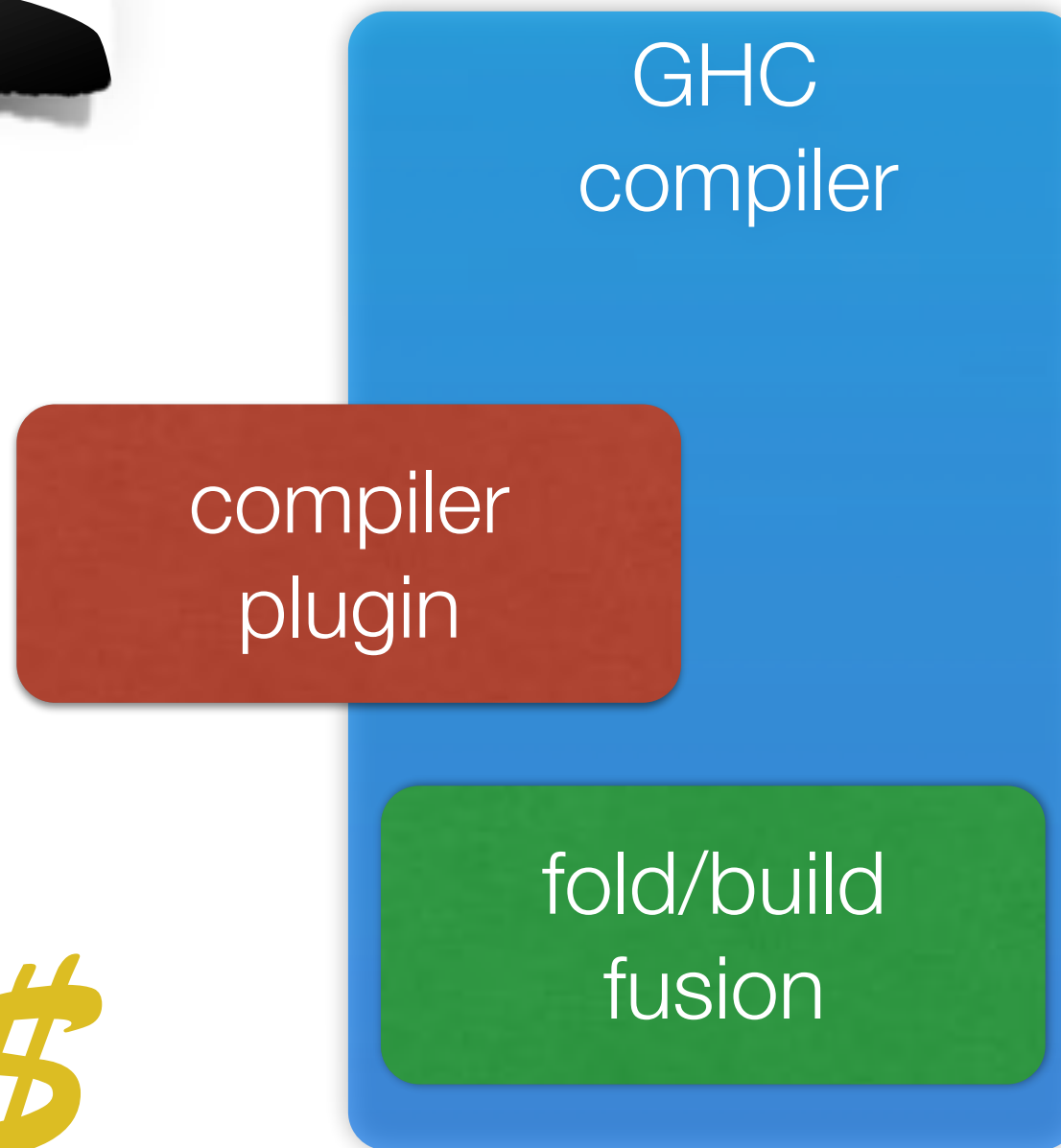


Fast Monads for Free

2015
-
2016



You





catamorphism

Cata-morphism



catamorphism

Cata-morphism

Greek:
down



catamorphism

Cata-morphism

Greek:
down

Category theory:
structure-preserving function

slide from
Nicolas Wu



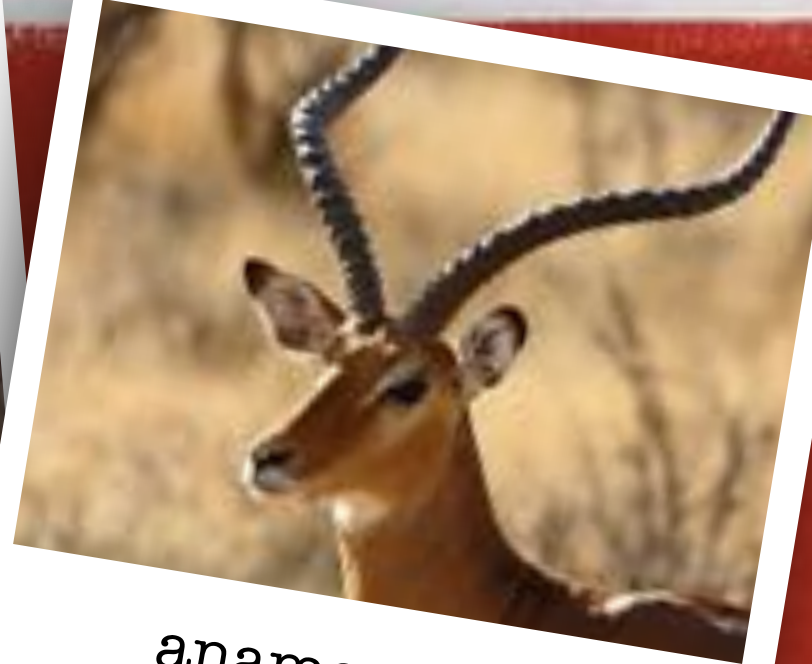
catamorphism

ZOO OF MORPHISMS

slide from
Nicolas Wu



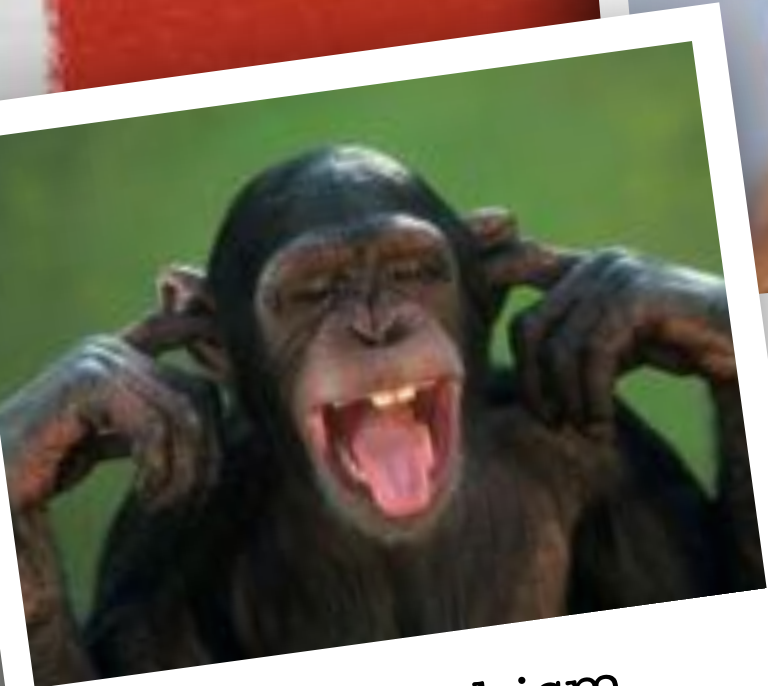
catamorphism



anamorphism

ZOO OF MORPHISMS

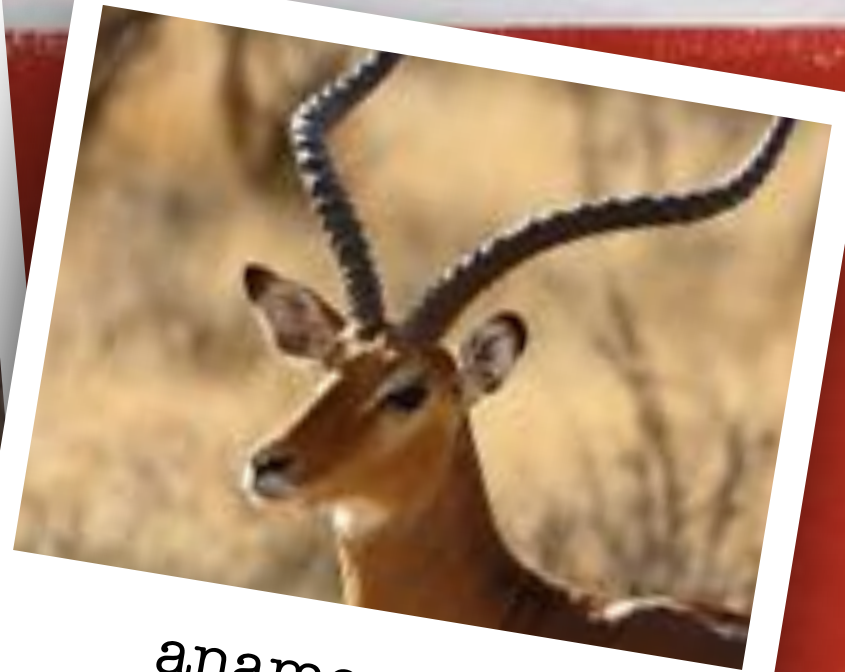
slide from
Nicolas Wu



apomorphism



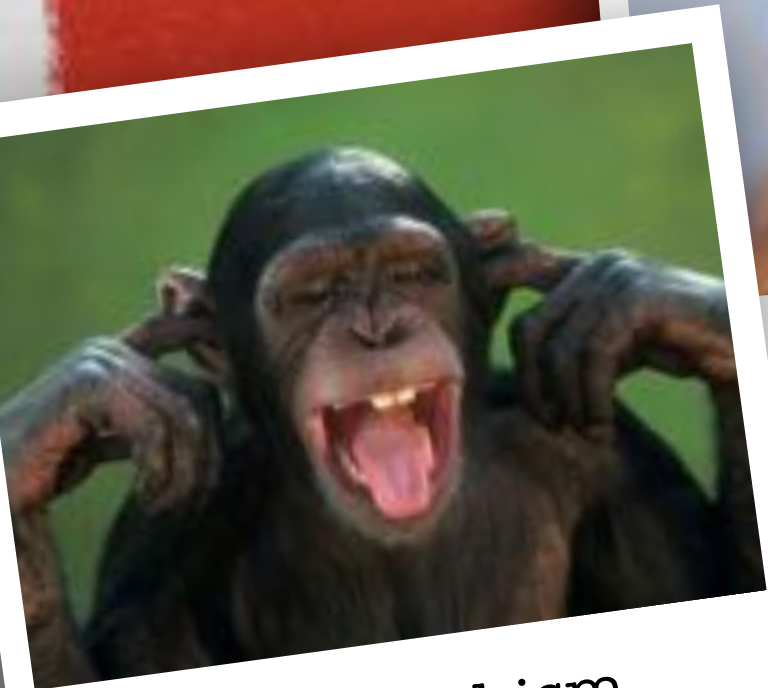
catamorphism



anamorphism

OF MORPHISMS

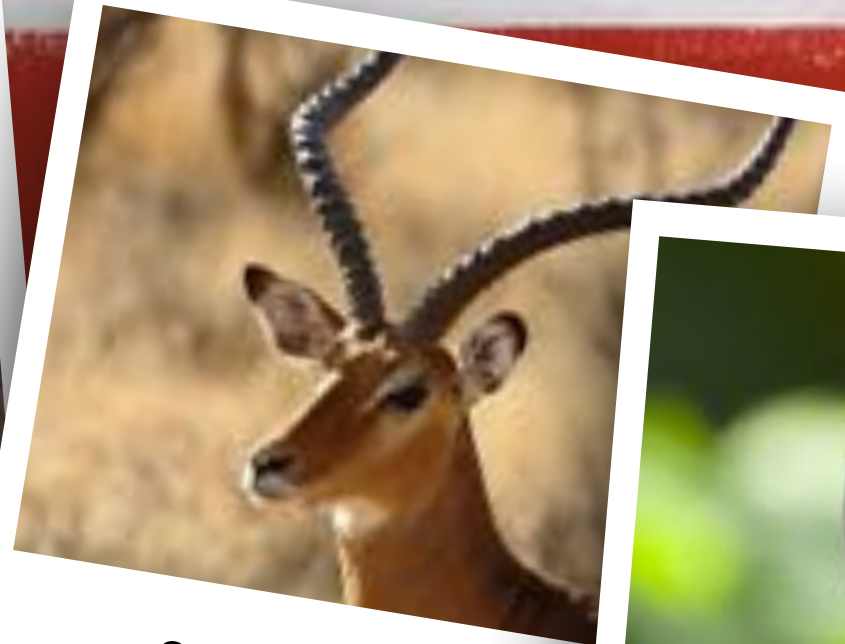
slide from
Nicolas Wu



apomorphism



catamorphism



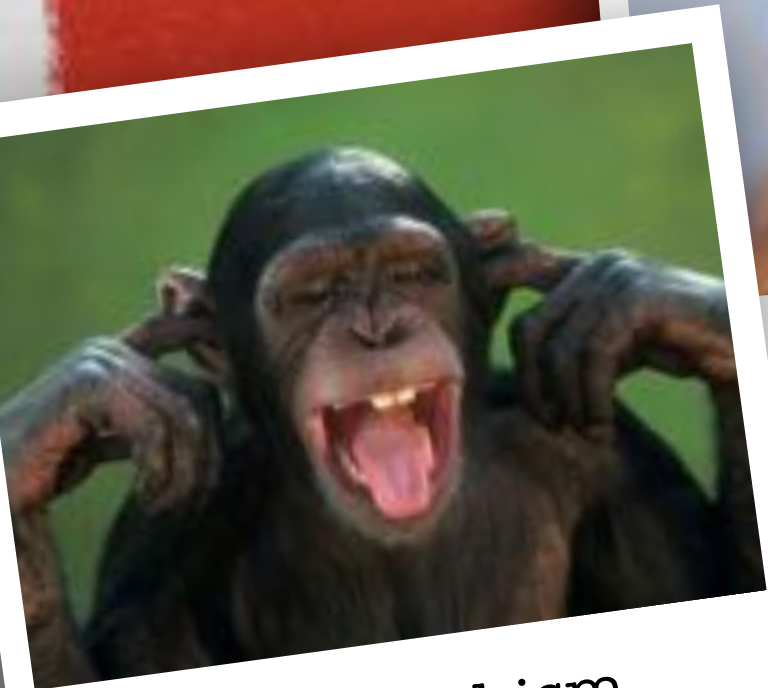
anamorphism



paramorphism

OF MORPHISMS

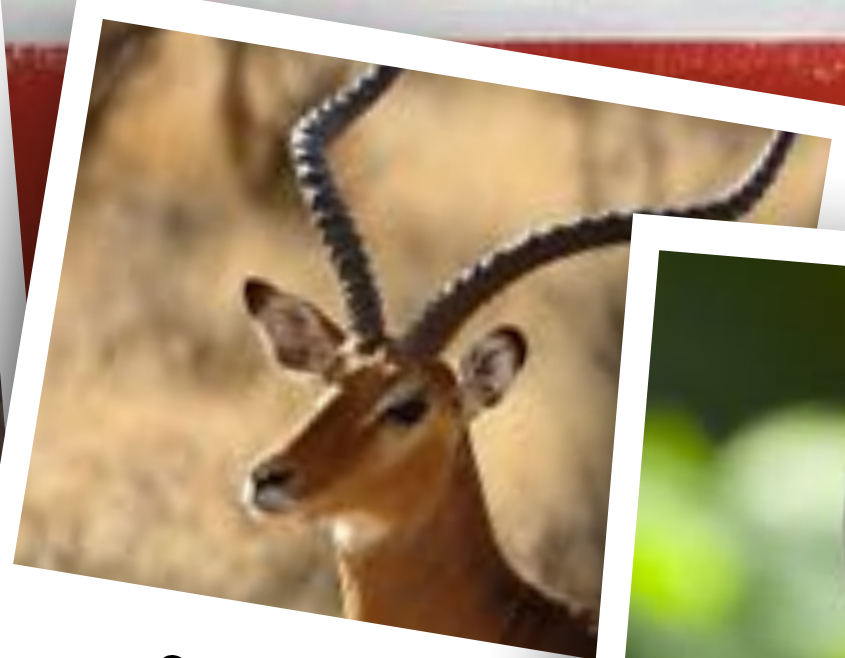
slide from
Nicolas Wu



apomorphism



catamorphism



anamorphism



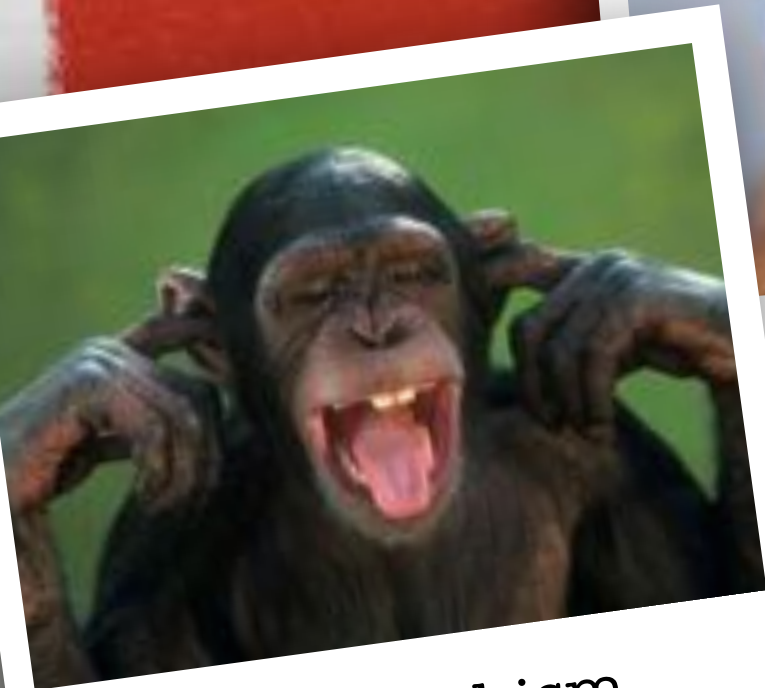
paramorphism

D OF MORPHISMS



mutumorphism

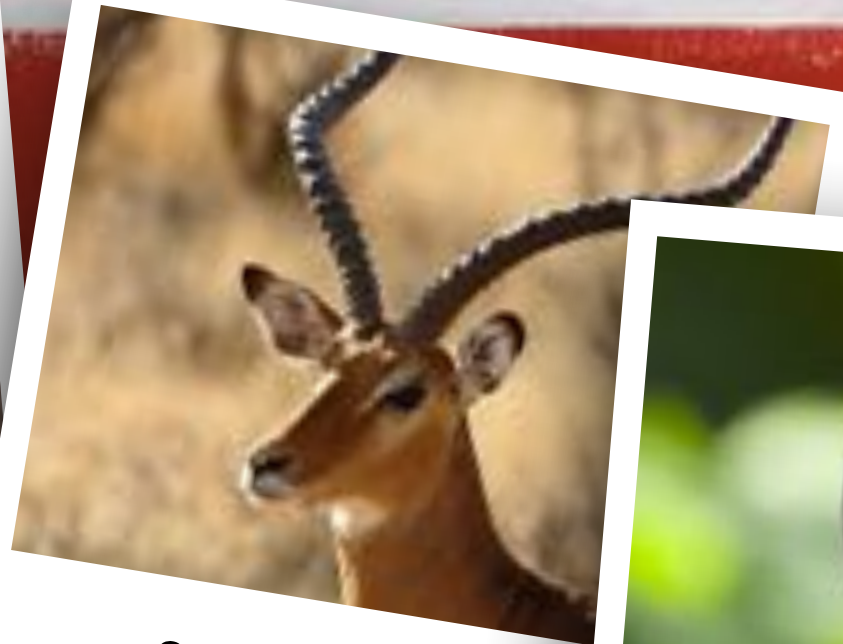
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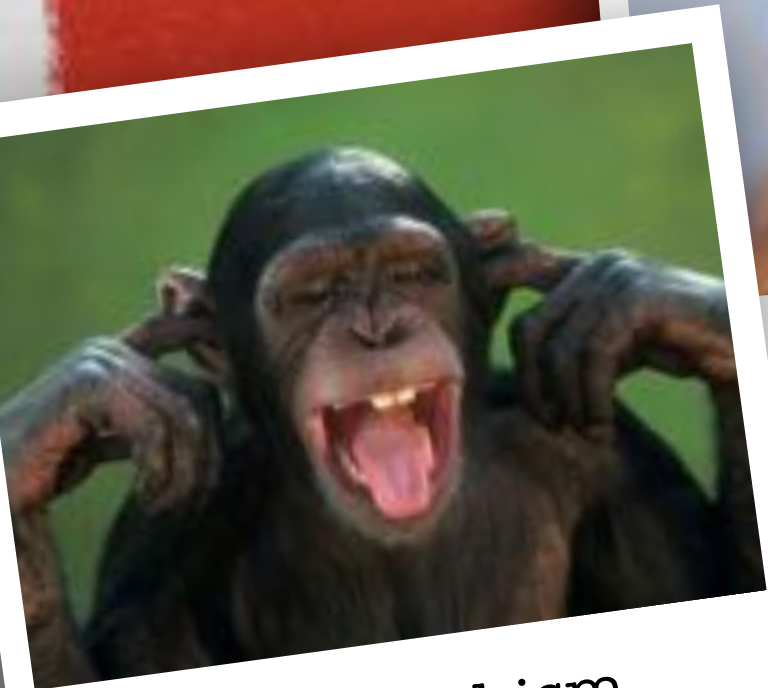


mutumorphism



zygomorphism

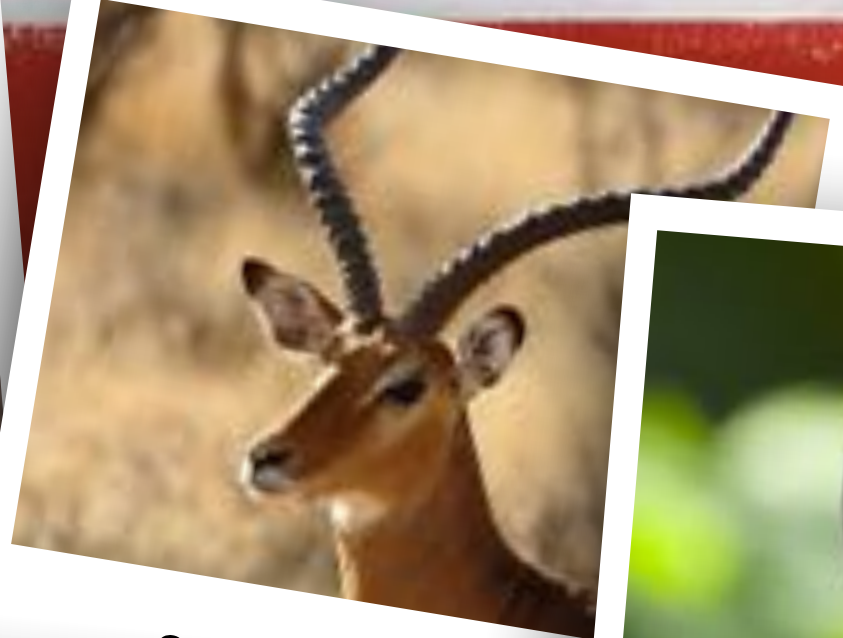
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apomorphism



cataractism



anamorphism



paramorphism



histomorphism



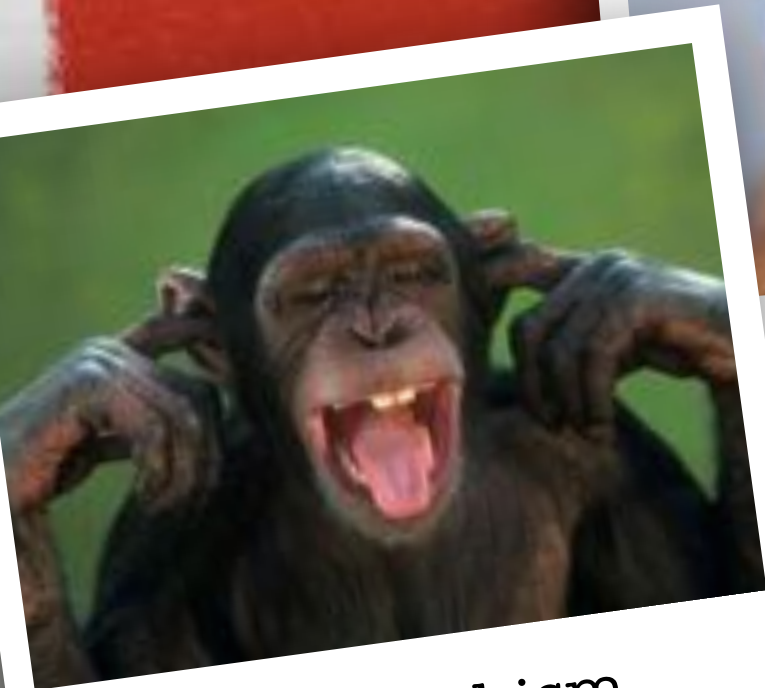
mutumorphism



zygomorphism

HISMS

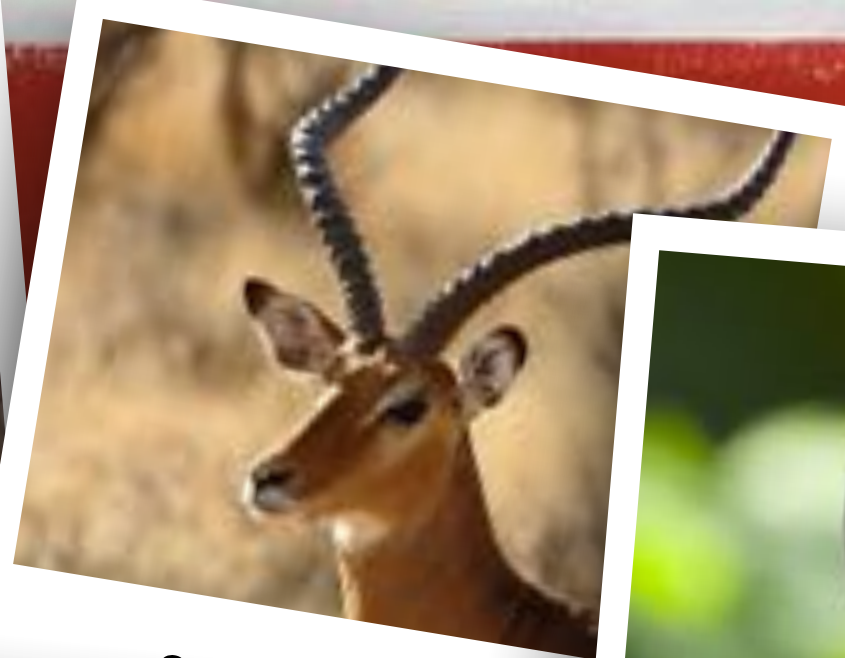
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paramorphism



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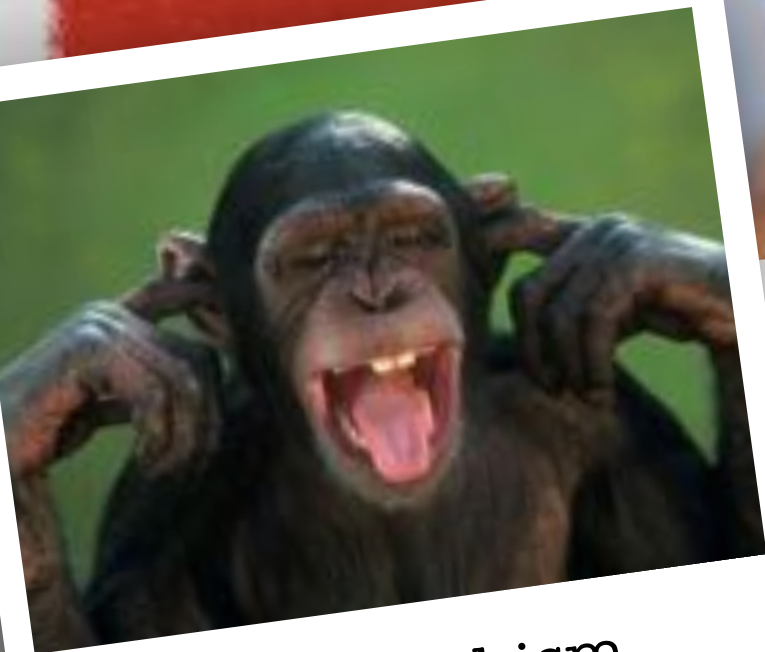
zygomorphism



dynamorphism

HISMS

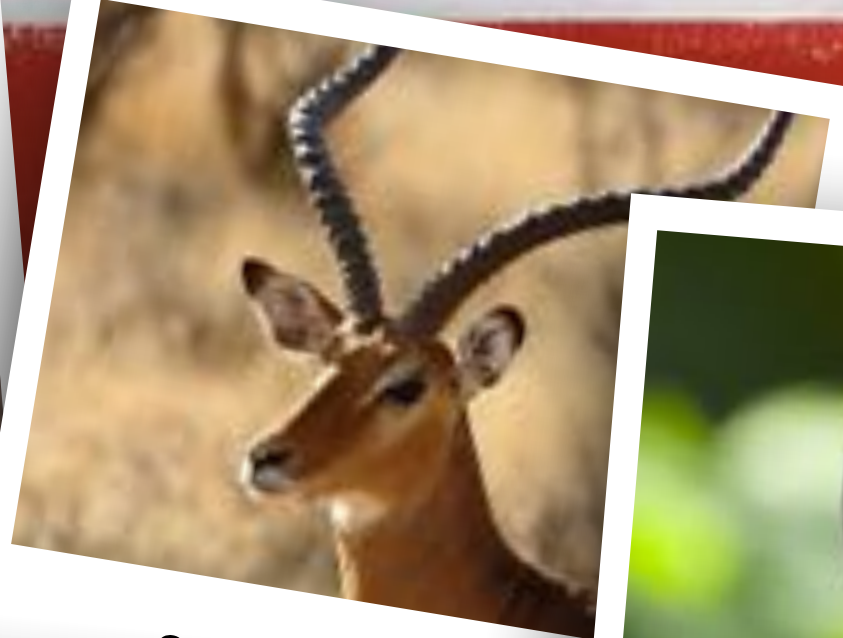
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anamorphism



paramorphism



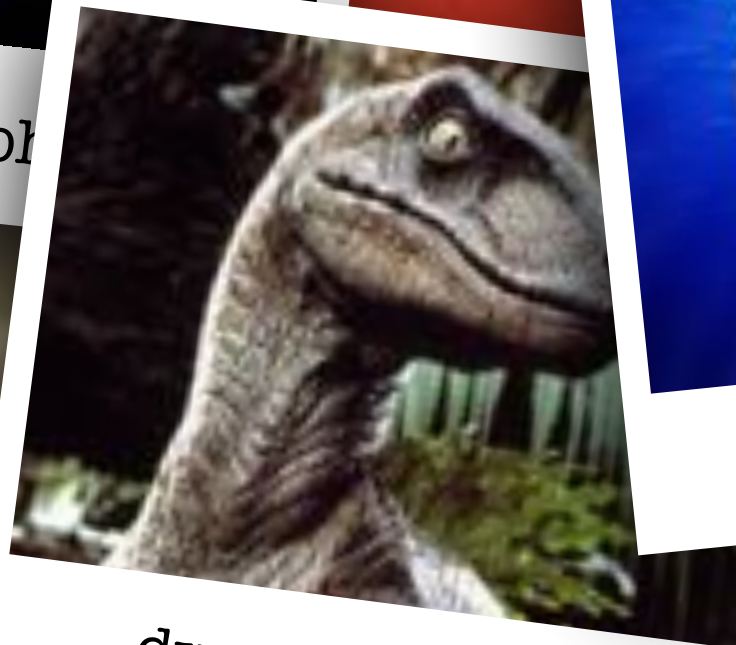
histomorph



mutumorphism



zygomorphism



dynamorphism

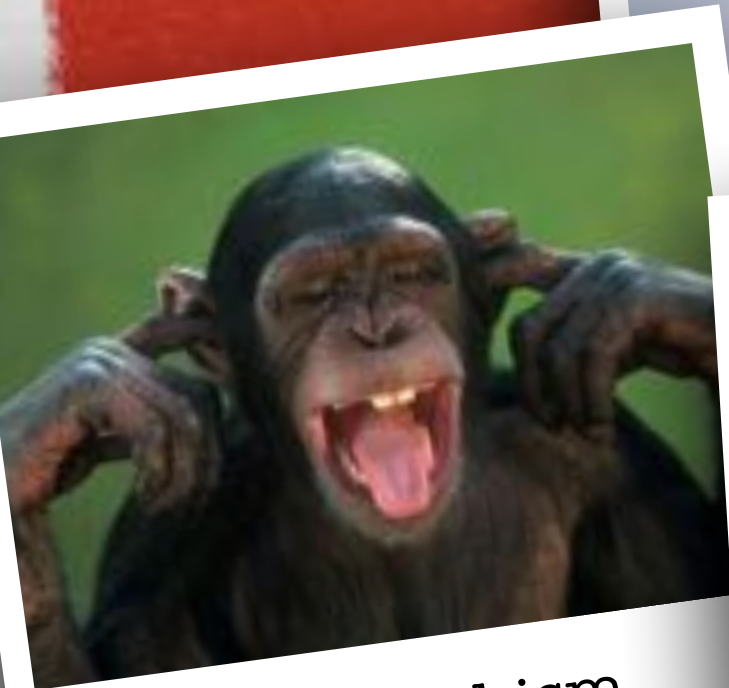


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paramorphism



apomorphism



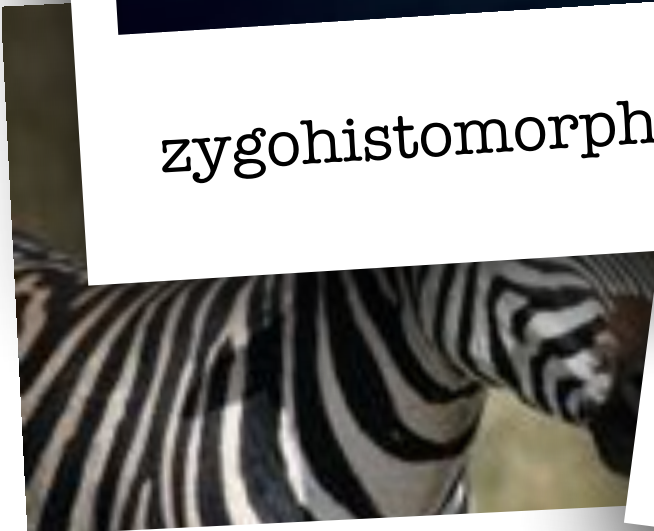
zygohistomorphic prepromorphism



futumorphism



mutumorphism



zygomorphism



dynamorphism