

Exercise 7: Financial Forecasting

```
Program    Debug    Debug Result    Java (11) [v] [x]
1- import java.util.*;
2- class SimpleForecast {
3-     public static void main(String[] args) {
4-         Scanner s=new Scanner(System.in);
5-         int n=sc.nextInt();
6-         int[] revenue=new int[n];
7-
8-         for (int i=0;i<n;i++) {
9-             revenue[i]=sc.nextInt();
10-        }
11-
12-        System.out.println("Monthly Revenue:");
13-        for (int r:revenue) {
14-            System.out.print("₹"+r+" ");
15-        }
16-        System.out.println();
17-
18-        int sum=0;
19-        for (int i=0;i<revenue.length;i++) {
20-            sum+=revenue[i];
21-        }
22-
23-        int average=sum/revenue.length;
24-        System.out.println("Forecasted Revenue for Next Month: ₹"+average);
25-    }
26- }
27- }
```

OUTPUT:

```
5
12000 15000 20202 23023 25202
```

Compiler Message

Compilation successful

Custom Testcase

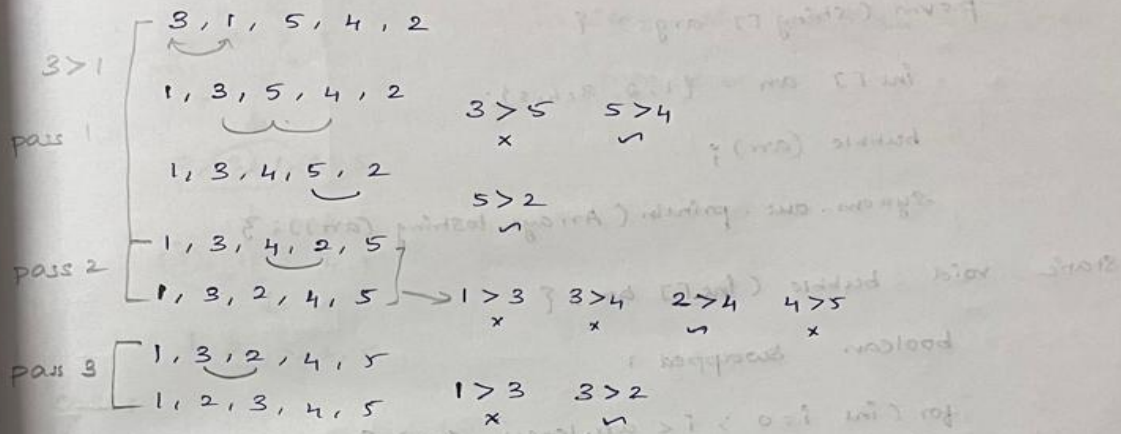
Output

```
Monthly Revenue:
₹12000 ₹15000 ₹20202 ₹23023 ₹25202
Forecasted Revenue for Next Month: ₹19085
```

NOTES:

Sorting :

Bubble Sort : \Rightarrow also known as Sinking Sort, exchange sort



Why Bubble Sort?

With the 1st pass through the array, the largest element came to the end.

With the 2nd pass the second largest element will be the last.

How it works \Rightarrow Refer notes book.

Space Complexity : $O(1)$ // Constant

No extra space is required. (i.e) copying the

This space complexity is also known as inplace Sorting algorithm.

Best case \rightarrow Array is Sorted $O(N)$

Worst case \rightarrow Sorting descending order to ascending order.

$O(N^2)$

Code // Bubble Sort

```
public class Main {  
    psvm (String[] args) {  
        int[] arr = {1, 2, 3, 4, 5};  
        bubble (arr);  
        System.out.println (Arrays.toString (arr));  
    }  
  
    static void bubble (int[] arr) {  
        boolean swapped;  
        for (int i = 0; i < arr.length; i++) {  
            swapped = false;  
            // for each step max value will be in the last  
            for (int j = 1; j < arr.length - i; j++) {  
                if (arr[j] < arr[j-1]) {  
                    // swap  
                    int temp = arr[j];  
                    arr[j] = arr[j-1];  
                    arr[j-1] = temp;  
                    swapped = true; }  
            }  
            if (!swapped) {  
                break;  
            }  
        }  
    }  
}
```

I/p:

1, 2, 3, 4, 5

O/p

[1, 2, 3, 4, 5]

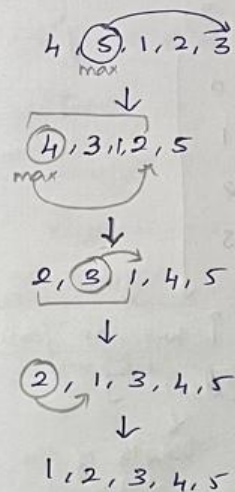
I/p:

5, 3, 2, 1, 4

O/p

[1, 2, 3, 4, 5]

Selection Sort : \Rightarrow Refer Notes too



i) Find the max Element

ii) Swap the last index with that max element

Code:

```
public class Main { psvm (String[] args) {
    int[] arr = { 3, 1, 5, 4, 2 };
    selection (arr);  $\rightarrow$  This method stores the sorted array
                        that is the original array is
    sop (Arrays.toString (arr)); sorted itself. So, we should not
    static void selection (int[] arr) { store it in any of the arr
        for (int i = 0; i < arr.length; i++) {
            int last = arr.length - i - 1;
            int maxIndex = getMaxIndex (arr, start: 0, last);
            swap (arr, maxIndex, last); } }

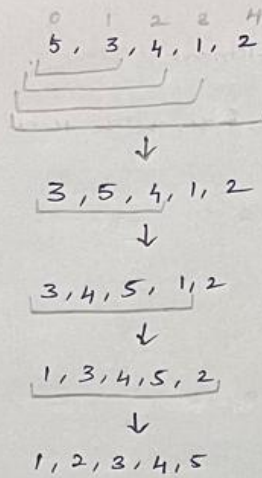
    static void swap (int[] arr, int first, int second) {
        int temp = arr [1st]; arr [1st] = arr [2nd]; arr [2nd] = temp;

    static int getMaxIndex (int[] arr, int start, int end) {
        int max = start;
        for (int i = start; i <= end; i++) {
            if (arr [max] < arr [i]) {
                max = i; }
            return max; }
    }
```

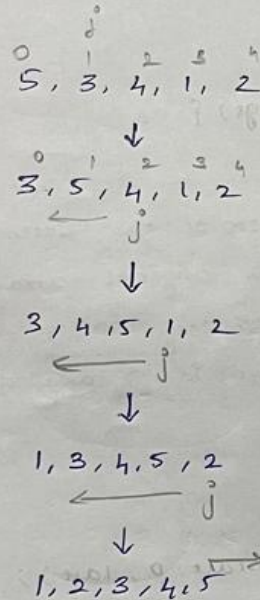
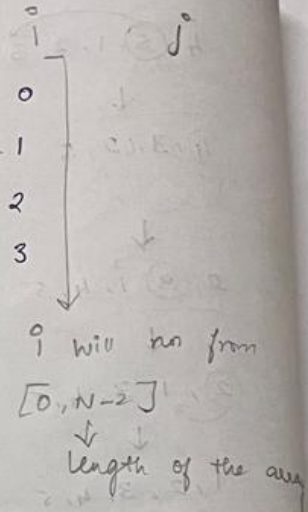
I/p: { 3, 1, 5, 4, 2 }

O/p: [1, 2, 3, 4, 5]

Insertion Sort :



Sort array ← pass No ← 0
 till index 1
 Sort array ← 2
 till index 2
 till index 3
 till index 4



$i \leq N-2$ $j > 0$

0 1 2
 1 2

When $(j > j-1)$ the L.H.S is sorted.

Complexity : Worst case : $O(N^2)$

Best Case : $O(N)$ ⇒ Array is already sorted

Why we use insertion sort?

→ Steps gets reduced compared to Binary Search

→ Stable

→ Used for smaller values of N

→ Used when Array is partially sorted

Code :: Insertion Sort

```
public static void main (String args[]) {  
    int[] arr = {5, 3, 4, 1, 2};  
    insertion (arr);  
    System.out.println (Arrays.toString(arr));  
}  
static void insertion (int[] arr) {  
    for (int i = 0; i < arr.length - 1; i++) {  
        for (int j = i + 1; j > 0; j--) {  
            if (arr[j] < arr[j - 1]) {  
                swap (arr, j, j - 1);  
            }  
            else {  
                break; }  
        }  
    }  
}  
static void swap (int[] arr, int first, int second) {  
    int temp = arr[first];  
    arr[first] = arr[second];  
    arr[second] = temp; }  
}
```

I/p:

5, 3, 4, 1, 2

O/p:

[1, 2, 3, 4, 5]

In Insertion Sort the number of swaps is reduced as compared to the bubble sort

Insertion Sort is a stable Sorting Algorithm.

Used for the smaller values of n:

works good when array is partially sorted.

Cyclic Sort: \Rightarrow In all other sorting the worst case time complexity is $O(N^2)$.

When given numbers from range 1 to $N \Rightarrow$ use Cyclic Sort.

Example: 3, 5, 2, 1, 4 \Rightarrow 1 to N then Cyclic Sort
1, 2, 3, 4, 5 \Rightarrow Sorted Array & Bubble Sort.

3, 5, 2, 1, 4

After Sorting:

0	1	2	3	4	\rightarrow index = value - 1
1	2	3	4	5	

 $4 = 5 - 1$

0	1	2	3	4
3	5	2	1	4

Swap coz $3 - 1 = 2$ index

0	1	2	3	4
2	5	3	1	4

Swap coz $2 - 1 = 1^{\text{st}}$ index

0	1	2	3	4
5	2	3	1	4

Swap coz $5 - 1 = 4^{\text{th}}$ index

0	1	2	3	4
4	2	3	1	5

Swap coz $4 - 1 = 3^{\text{rd}}$ index

1, 2, 3, 4, 5 \Rightarrow Finally Sorted.

0	1	2	3	4
---	---	---	---	---

\hookrightarrow index = value - 1

$4 = 5 - 1$

$4 = 4$

Complexity: $O(N)$

\hookrightarrow Worst Case.

Code // Refer Notes :

```
public class CyclicSort {  
    public static void main (String[] args) {  
        int[] arr = {3, 5, 2, 1, 4};  
        sort (arr);  
        System.out.println (Array.toString (arr));  
    }  
    static void sort (int[] arr) {  
        int i = 0;  
        while (i < arr.length) {  
            int correct = arr[i] - 1;  
            if (arr[i] != arr[correct]) {  
                swap (arr, i, correct);  
            }  
            else {  
                i++;  
            }  
        }  
    }  
    static void swap (int[] arr, int first, int second) {  
        int temp = arr[first];  
        arr[first] = arr[second];  
        arr[second] = temp;  
    }  
}
```

I/p :

3, 5, 2, 1, 4

↓

1 to N

O/p :

[1, 2, 3, 4, 5]
0 1 2 3 4

↳ Index = Val

4 = 5 - 1

4 = 4