

## ESO207 Programming Assignment 3

Arpit Kumar, Roll no.: 200189

Kumar Arpit, Roll no.: 200532

### 1. Pseudo Code

```
fun Bipartite(g)                                //'g' is the input graph
{
    color[V]                                    //Creating an array to store the
    color of a vertex which is either 0 or 1
    for i from 0 to V-1                        //Initializing the color of each
    vertex to -1
        color[i]=-1
        if(modifiedDFS(g,color,0,0)==true)      //Calls the subroutine modifiedDFS
    from vertex 0 and the color to be assigned to it as 0
            V1[V], V2[V]                      //Creating 2 arrays to store the
    partitions
            for i from 0 to g->V-1            //Looping through the array to add a
    vertex to either V1 or V2
                if(color[i]==0)
                    V1.append(i)
                else
                    V2.append(i)
            return (V1,V2)
    else
        return false                          //Returning false if g is not
    bipartite
}

fun modifiedDFS(g,color[],vertex,c)             //A modified DFS routine. Note that
    color kind of plays the role of mark[] as well.
{
    color[vertex]=c
    temp=g->adjLists[vertex]                   //Creating a copy of the pointer at
    Linked list of the vertex 'vertex' in the graph g
    while(temp!=NULL)
        if(color[temp->vertex]==-1)           //If -1 then the node has not been
    visited
            if(dfs(g,color,temp->vertex,(c+1)%2)==0) //Call the function recursively with
    the other color
                return false                  //If the result of call was zero then
    it wasn't possible to color the graph with two colors. Thus the sequence of return 0 is
    passed down the recursion and finally to Bipartite.
            else if(color[temp->vertex]==c)
                return false                  //If color same as its parent then 0
    is returned instantly.
            temp=temp->next                    //Proceeds to the next neighboring
    vertex of temp
```

```

    return true //If it comes out of the while loop
without returning zero then the subgraph can be successfully colored using two colors and
true is returned.
}

//Structures used:
struct node
{
    long long int vertex;
    struct node* next;
};
struct Graph
{
    long long int V;
    struct node** adjLists;
};

```

## 2. Uniqueness:

In the given case, graph  $G$  is connected. Thus, in this case, the partition  $V_1$  and  $V_2$  is unique (assuming the order  $V_1$  and  $V_2$  does not matter.)

This can be seen from the method used to solve the given problem i.e., the two coloring itself. Assuming we know that a bipartite connected graph can always be colored using two colors say  $A$  and  $B$  then if we choose any vertex in  $V$  for partition  $V_1$  and assigns it the color say  $A$  then all its neighbors will be assigned color  $B$  and their neighbors  $A$  and so on. Since the graph is connected, all vertices will be assigned a color. Since the coloring is unique or complementary for different choices of the first vertex, the partition is unique.

However, when a graph is unconnected it has **always** has more than one way of partitioning it into  $V_1$  and  $V_2$  if it is bipartite. This can be seen as follow:

Suppose the graph has  $n$  components, let  $a_1$  and  $a_2$  be any two of them.

Since  $G$  is bipartite,  $a_1$  and  $a_2$  themselves are bipartite as well.

Say  $x_1$  and  $x_2$  are bipartite partitions of  $a_1$  and  $y_1$  and  $y_2$  of  $a_2$ . Let  $z_1$  and  $z_2$  be two bipartitions of the remaining  $n-2$  components (again we know that the union of all those components must be bipartite as well assuming a third component otherwise they both are empty.)

We then partition  $G$  as  $x_1 \cup y_1 \cup z_1$  and  $x_2 \cup y_2 \cup z_2$  once and  $x_1 \cup y_2 \cup z_1$  and  $x_2 \cup y_1 \cup z_2$  again.

Note that these partitions must be different as  $G$  was not connected and thus had more than one component.

Hence, we prove by contradiction that a non-connected bipartite graph can always be partitioned in more than one way.

## 3. Implementation:

[On Hackerrank]