## Mathematics of Cryptography: Algebraic Structure

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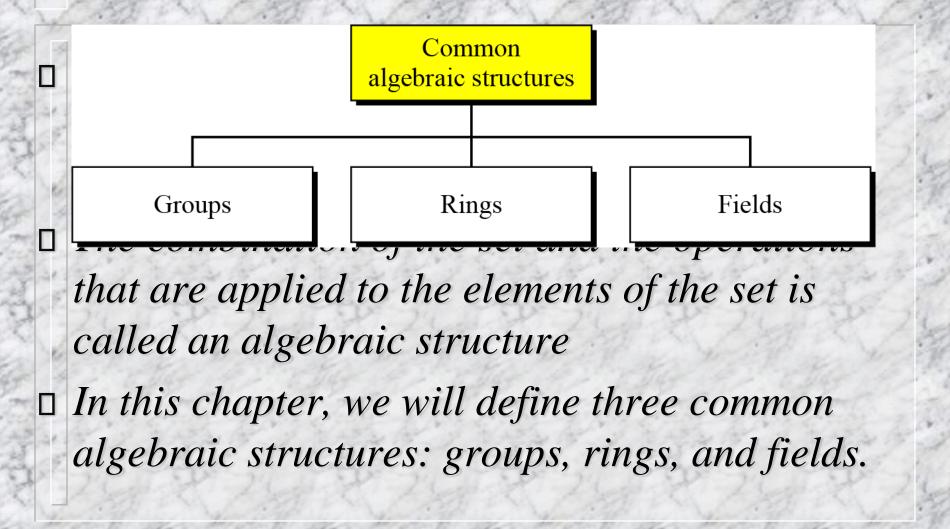
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## Objectives

- □ Review the concept of algebraic structures
- Define and give some examples of groups
- Define and give some examples of rings
- Define and give some examples of fields
- Emphasize the finite fields of type  $GF(2^n)$  that make it possible to perform operations such as addition, subtraction, multiplication, and division on n-bit words in modern block ciphers

## **ALGEBRAIC STRUCTURES**



## Groups

1. Closure

**Properties** 

- 1. Closure
- 2. Associativity
- 3. Commutativity (See note)
- □ A group 4. Existence of identity
  - operatio 5. Existence of inverse
  - axioms)

{a, b, c, ...} Set



Note:

The third property needs

to be satisfied only for a

commutative group.

#### Closure:

Associativit,.

Group,

Existence of identity: for all  $a \in G$ , there is  $e \in G$  such that  $a \cdot e = e \cdot a = a$ , e is the identity element

Existence of inverse: for all  $a \in G$ , there is  $b \in G$  such that  $a \cdot b = b \cdot a = e$ , b is inverse of a and vice versa

## Groups (contd...)

- A group (G, •) is called commutative or abelian group if the operator '•' satisfies the commutative property
  - Commutative property: for all  $a, b \in G$ ,  $a \cdot b = b \cdot a$

Although a group involves a single operation, the properties imposed on the operation allow the use of a pair of operations as long as they are inverses of each other.

## Example: groups

The set of residue integers with the addition operator,

$$G = < Z_n, +>,$$

is a commutative group. We can perform addition and subtraction on the elements of this set without moving out of the set.

The set  $Z_n^*$  with the multiplication operator,  $G = \langle Z_n^*, \times \rangle$ , is also an abelian group.

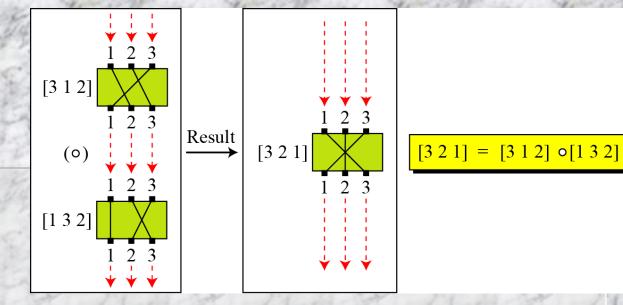
Let us define a set  $G = \langle \{a, b, c, d\}, \bullet \rangle$  and the operation as shown in

following table

•	а	b	c	d
а	а	b	С	d
b	b	С	d	а
c	С	d	а	b
d	d	а	b	С

This is an abelian group

# Permutation group



0	[1 2 3]	[1 3 2]	[2 1 3]	[2 3 1]	[3 1 2]	[3 2 1]
[1 2 3]	[1 2 3]	[1 3 2]	[2 1 3]	[2 3 1]	[3 1 2]	[3 2 1]
[1 3 2]	[1 3 2]	[1 2 3]	[2 3 1]	[2 1 3]	[3 2 1]	[3 1 2]
[2 1 3]	[2 1 3]	[3 1 2]	[1 2 3]	[3 2 1]	[1 3 2]	[2 3 1]
[2 3 1]	[2 3 1]	[3 2 1]	[1 3 2]	[3 1 2]	[1 2 3]	[2 1 3]
[3 1 2]	[3 1 2]	[2 1 3]	[3 2 1]	[1 2 3]	[2 3 1]	[1 3 2]
[3 2 1]	[3 2 1]	[2 3 1]	[3 1 2]	[1 3 2]	[2 1 3]	[1 2 3]

## Permutation group

- □ set of permutations with the composition operation is a group
  - This implies that using two permutations one after another cannot strengthen the security of a cipher
    - □ because we can always find a permutation that can do the same job because of the closure property

## Groups (contd...)

☐ Finite Group: a group with finite elements; otherwise, infinite group

□ Order of a Group: |G|, number of elements if finite; otherwise, infinite

☐ Subgroups: A subset H of G is a subgroup of G if H is a group under the operation of G

## subgroups

- □ If  $a, b \in G, H \rightarrow c = a \cdot b \in G, H$
- $\Box$  e  $\in$  G, H
- □ If  $a \in G$ ,  $H \rightarrow If a^{-1} \in G$ , H
- $\square$  ({e}, •) is subgroup of G, H
- □ G is a subgroup of itself

## Example: subgroup

Is the group  $H = \langle Z_{10}, + \rangle$  a subgroup of the group  $G = \langle Z_{12}, + \rangle$ ?

The answer is no. Although H is a subset of G, the operations defined for these two groups are different. The operation in H is addition modulo 10; the operation in G is addition modulo 12.

## Cyclic Subgroups

If a subgroup of a group can be generated using the power of an element, the subgroup is called the cyclic subgroup

$$a^n \to a \bullet a \bullet \dots \bullet a \quad (n \text{ times})$$

## Example: cyclic subgroup

Four cyclic subgroups can be made from the group  $G = \langle Z_6, + \rangle$ . They are  $H_1 = \langle \{0\}, + \rangle$ ,  $H_2 = \langle \{0, 2, 4\}, + \rangle$ ,  $H_3 = \langle \{0, 3\}, + \rangle$ , and  $H_4 = G$ .

## Example: cyclic subgroup

Three cyclic subgroups can be made from the group  $G = \langle Z_{10}^*, \times \rangle$ . G has only four elements: 1, 3, 7, and 9. The cyclic subgroups are  $H_1 = \langle \{1\}, \times \rangle$ ,  $H_2 = \langle \{1, 9\}, \times \rangle$ , and  $H_3 = G$ .

$$1^0 \mod 10 = 1$$

$$3^0 \mod 10 = 1$$
  
 $3^1 \mod 10 = 3$ 

$$3^2 \mod 10 = 9$$

$$3^3 \mod 10 = 7$$

$$7^0 \mod 10 = 1$$

$$7^1 \mod 10 = 7$$

$$7^2 \mod 10 = 9$$

$$7^3 \mod 10 = 3$$

$$9^0 \mod 10 = 1$$

$$9^1 \mod 10 = 9$$

## Cyclic Groups

A cyclic group is a group that is its own cyclic subgroup.

$$\{e, g, g^2, \dots, g^{n-1}\}\$$
, where  $g^n = e$ 

## Cyclic Groups

Three cyclic subgroups can be made from the group  $G = \langle Z_{10}^*, \times \rangle$ . G has only four elements: 1, 3, 7, and 9. The cyclic subgroups are  $H_1 = \langle \{1\}, \times \rangle$ ,  $H_2 = \langle \{1, 9\}, \times \rangle$ , and  $H_3 = G$ .

- a. The group  $G = \langle Z_6, + \rangle$  is a cyclic group with two generators, g = 1 and g = 5.
- b. The group  $G = \langle Z_{10}^*, \times \rangle$  is a cyclic group with two generators, g = 3 and g = 7.

## Subgroup property

Assume that G is a group, and H is a subgroup of G. If the order of G and H are |G| and |H|, respectively, then, based on this theorem, |H| divides |G|.

Order of an Element

The order of an element a, ord(a), is the smallest number n such that  $a^n = e$ 

In other words, ord(a) is the order of the cyclic group generated by a

## Example: ord(a)

a. In the group  $G = \langle Z_6, + \rangle$ , the orders of the elements are:

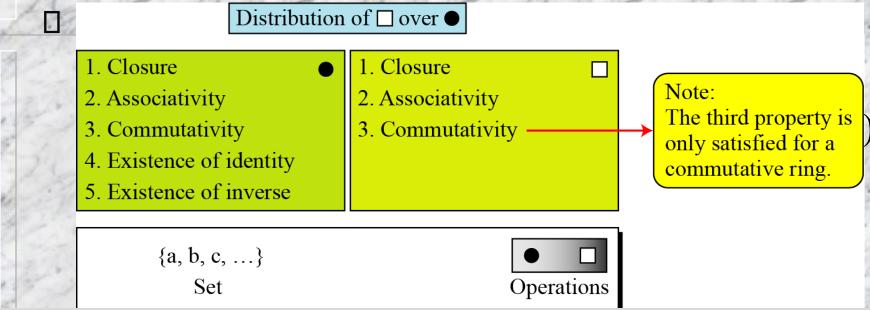
ord(0) = 1, ord(1) = 6, ord(2) = 3, ord(3) = 2, ord(4) = 3,

ord(5) = 6.

b. In the group  $G = \langle Z_{10}^*, \times \rangle$ , the orders of the elements are:

$$ord(1) = 1$$
,  $ord(3) = 4$ ,  $ord(7) = 4$ ,  $ord(9) = 2$ .





The set Z with two operations, addition and multiplication, is a commutative ring. We show it by  $R = \langle Z, +, \times \rangle$ . Addition satisfies all of the five properties; multiplication satisfies only three properties.

If '\*' is commutative, R is commutative ring

### Field

- $\square$  A field, denoted by  $F = \langle \{...\}, +, *\rangle$  is



- C 1. Closure
  - 2. Associativity
  - 3. Commutativity
  - 4. Existence of identity
  - 5. Existence of inverse

1. Closure

Field

- 2. Associativity
- 3. Commutativity
- 4. Existence of identity
- 5. Existence of inverse

Note:

The identity element of the first operation has no inverse with respect to the second operation.

•

Operations

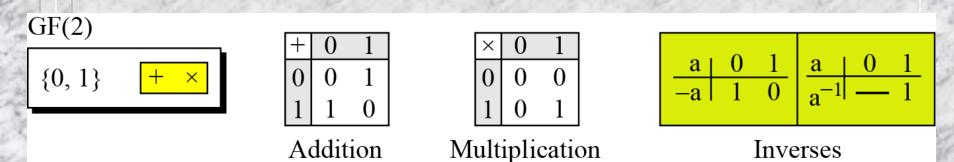
### Finite field

- ☐ Finite field is important in cryptography
- A field with finite number is called finite field
- Galois showed that for a field to be finite, the number of elements should be  $p^n$ , where p is a prime and n is a positive integer

## Galois field GF(p)

When n = 1, we have GF(p) field. This field can be the set  $Z_p$ ,  $\{0, 1, ..., p - 1\}$ , with two arithmetic operations

A very common field in this category is GF(2) with the set {0, 1} and two operations, addition and multiplication, as shown in Figure 4.6.



## Galois field GF(p) (contd...)

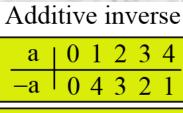
We can define GF(5) on the set  $Z_5$  (5 is a prime) with addition and multiplication operators as shown below

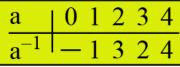
GF(5)  $\{0, 1, 2, 3, 4\} + \times$ 

+	0	1	2	3	4
0	0	1	2 3	3	4
2	2	- 4	4	0	0
3	3	4	0	1	2
4	4	0	1	2	3

×	0	1	2	3	4
0	0	0	0 2 4 1 3	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

Multiplication





Multiplicative inverse

## Summary

Algebraic Structure	Supported Typical Operations	Supported Typical Sets of Integers
Group	$(+ -) \text{ or } (\times \div)$	$\mathbf{Z}_n$ or $\mathbf{Z}_n^*$
Ring	(+ −) and (×)	Z
Field	$(+ -)$ and $(\times \div)$	$\mathbf{Z}_{p}$

## $GF(2^n)$

- ☐ In cryptography, we often need to use four operations (addition, subtraction, multiplication, and division)
- □ In other words, we need to use fields
- □ We can work in GF(p) where p is the largest number less than  $2^n$ 
  - But, numbers between p and  $2^n$  -1 cannot be handled
- $\square$  In  $GF(2^n)$ , we have a set of  $2^n$  elements
  - The elements in this set are n-bit words

## GF(2<sup>n</sup>) (contd...)

Let us define a  $GF(2^2)$  field in which the set has four 2-bit words:  $\{00, 01, 10, 11\}$ . We can redefine addition and multiplication for this field in such a way that all properties of these operations are satisfied

Addition						Mu	ltip	lica	tion
$\bigoplus$	00	01	10	11	$\otimes$	00	01	10	11
00	00	01	10	11	00	00	00	00	00
01	01	00	11	10	01	00	01	10	11
10	10	11	00	01	10	00	10	11	01
11	11	10	01	00	11	00	11	01	10
	Ida	enti	<b>1</b> v·	00	'	Id	enti	tv•	01

00, 01,..., 11 cannot be considered as integer from 0 to 3

Addition and multiplication are defined in terms of polynomial

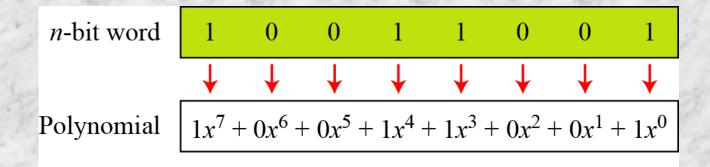
## Polynomials

 $\square$  A polynomial of degree n-1 is an expression like

$$f(x) = a_{n-1}x^{n-1} + a_{n-2}x^{n-2} + \dots + a_1x^1 + a_0x^0$$

where  $a_i$  is called coefficient of the  $i^{th}$  term.

8-bit word 10011001 represents as



First simplification 
$$1x^7 + 1x^4 + 1x^3 + 1x^0$$

Second simplification

$$x^7 + x^4 + x^3 + 1$$

## Polynomials (contd...)

 $\square$  To find the 8-bit word related to the polynomial  $x^5$ 

$$+ x^2 + x$$

- we first supply the omitted terms
  - Since n = 8, it means the polynomial is of degree 7
  - The expanded polynomial is

$$0x^7 + 0x^6 + 1x^5 + 0x^4 + 0x^3 + 1x^2 + 1x^1 + 0x^0$$

Related 8-bit word is 00100110

## Operations on polynomials

- Any operation on polynomial involves two operations:
  - operation on coefficients and
  - operations on two polynomials
- □ Operations on coefficients (0/1) use GF(2)
- □ For operations on polynomials need GF(2<sup>n</sup>)

## Modulus respect to polynomial

8	Degree	Irreducible Polynomials	
1	1	(x+1),(x)	
Contract of the	2	$(x^2 + x + 1)$	
9	3	$(x^3 + x^2 + 1), (x^3 + x + 1)$	
1	4	$(x^4 + x^3 + x^2 + x + 1), (x^4 + x^3 + 1), (x^4 + x + 1)$	
	5	$(x^5 + x^2 + 1), (x^5 + x^3 + x^2 + x + 1), (x^5 + x^4 + x^3 + x + 1),$ $(x^5 + x^4 + x^3 + x^2 + 1), (x^5 + x^4 + x^2 + x + 1)$	

- We need a polynomial of degree n, respect to which we have to take remainder
- The modulus polynomial takes as prime polynomial
- Prime polynomial is irreducible, i.e., no polynomial can divides it

## Addition operation

- □ Addition (or subtraction) over GF(2)
- $\Box$   $(x^5 + x^2 + x) \oplus (x^3 + x^2 + 1)$  in GF(28), the symbol
  - to show that we mean polynomial addition

$$0x^{7} + 0x^{6} + 1x^{5} + 0x^{4} + 0x^{3} + 1x^{2} + 1x^{1} + 0x^{0} \oplus 0x^{7} + 0x^{6} + 0x^{5} + 0x^{4} + 1x^{3} + 1x^{2} + 0x^{1} + 1x^{0}$$

 $0x^{7} + 0x^{6} + 1x^{5} + 0x^{4} + 1x^{3} + 0x^{2} + 1x^{1} + 1x^{0} \rightarrow x^{5} + x^{3} + x + 1$ 

Additive identity: zero polynomial

Additive inverse: polynomial itself

## Multiplication

- 1. The coefficient multiplication is done in GF(2).
- 2. The multiplying  $x^i$  by  $x^j$  results in  $x^{i+j}$ .
- 3. The multiplication may create terms with degree more than n-1, which means the result needs to be reduced using a modulus polynomial.

## Multiplication: $e^{x^8 + x^4 + x^3 + x + 1}$

 $x^4 + 1$ 

$$x^{12} + x^7 + x^2$$
$$x^{12} + x^8 + x^7 + x^5 + x^4$$

$$x^{8} + x^{5} + x^{4} + x^{2}$$
$$x^{8} + x^{4} + x^{3} + x + 1$$

Find the result of  $(x^5 + x^2 + x^3)$ irreducible polynomial ( $x^8$ represent multiplication of tv

Remainder 
$$x^5 + x^3 + x^2 + x + 1$$

Solution

$$P_{1} \otimes P_{2} = x^{5}(x^{7} + x^{4} + x^{3} + x^{2} + x) + x^{2}(x^{7} + x^{4} + x^{3} + x^{2} + x) + x(x^{7} + x^{4} + x^{3} + x^{2} + x)$$

$$P_{1} \otimes P_{2} = x^{12} + x^{9} + x^{8} + x^{7} + x^{6} + x^{9} + x^{6} + x^{5} + x^{4} + x^{3} + x^{8} + x^{5} + x^{4} + x^{3} + x^{2}$$

$$P_{1} \otimes P_{2} = (x^{12} + x^{7} + x^{2}) \mod (x^{8} + x^{4} + x^{3} + x + 1) = x^{5} + x^{3} + x^{2} + x + 1$$

To find the final result, divide the polynomial of degree 12 by the polynomial of degree 8 (the modulus) and keep only the remainder

## Multiplication

- □ Multiplicative identity: 1 i.e., 00000...0001
- Multiplicative inverse: extended Euclidean algorithm on the given polynomial and modulus polynomial

In GF (2<sup>4</sup>), find the inverse of  $(x^2 + 1)$  modulo  $(x^4 + x + 1)$ .

q	$r_{I}$	$r_2$	r	$t_I$	$t_2$	t
$(x^2 + 1)$	$(x^4 + x + 1)$	$(x^2 + 1)$	(x)	(0)	(1)	$(x^2 + 1)$
(x)	$(x^2 + 1)$	(x)	(1)	(1)	$(x^2 + 1)$	$(x^3 + x + 1)$
(x)	(x)	(1)	(0)	$(x^2 + 1)$	$(x^3 + x + 1)$	(0)
	(1)	(0)		$(x^3 + x + 1)$	(0)	

## Multiplicative inverse

In GF(28), find the inverse of (x5) modulo  $(x^8 + x^4 + x^3 + x + 1)$ .

q	$r_I$	$r_2$	r	$t_1$	$t_2$	t
$(x^3)$	$(x^8 + x^4 + x^3 + x^3)$	$(x+1) \qquad (x^5)$	$(x^4 + x^3 + x + 1)$	(0)	(1)	$(x^3)$
(x+1)	$(x^5)$ $(x^4)$	$+x^3+x+1)$	$(x^3 + x^2 + 1)$	(1)	$(x^3)$	$(x^4 + x^3 + 1)$
(x)	$(x^4 + x^3 + x + 1)$	$(x^3 + x^2 + 1)$	(1)	$(x^3)$	$(x^4 + x^3 + 1)$	$(x^5 + x^4 + x^3 + x)$
$(x^3 + x^2 + 1)$	$(x^3 + x^2 + 1)$	(1)	(0)	$(x^4 + x^3 + 1)$	$(x^5 + x^4 + x^3 + x)$	(0)
	(1)	(0)		$(x^5 + x^4 + x^3)$	(0)	

## Algorithm for multiplication

Find the result of multiplying  $P_1 = (x^5 + x^2 + x)$  by  $P_2 = (x^7 + x^4 + x^3 + x^2 + x)$  in GF(28) with irreducible polynomial  $(x^8 + x^4 + x^3 + x + 1)$  $x^5P_2 + x^2P_2 + xP_2 \mod IRP$ 

Multiply  $P_2$  by  $x, x^2, x^3, ...$ 

Powers	Operation	New Result	Reduction		
$x^0 \otimes P_2$		$x^7 + x^4 + x^3 + x^2 + x$	No		
$x^1 \otimes P_2$	$x \otimes (x^7 + x^4 + x^3 + x^2 + x)$	$x^5 + x^2 + x + 1$	Yes		
$x^2 \otimes P_2$	$\boldsymbol{x} \otimes (x^5 + x^2 + x + 1)$	$x^6 + x^3 + x^2 + x$	No		
$x^3 \otimes P_2$	$\boldsymbol{x} \otimes (x^6 + x^3 + x^2 + x)$	$x^7 + x^4 + x^3 + x^2$	No		
$x^4 \otimes P_2$	$\boldsymbol{x} \otimes (x^7 + x^4 + x^3 + x^2)$	$x^5 + x + 1$	Yes		
$x^5 \otimes P_2$	$\boldsymbol{x} \otimes (x^5 + x + 1)$	$x^6 + x^2 + x$	No		
$\mathbf{P_1} \times \mathbf{P_2} = (x^6 + x^2 + x) + (x^6 + x^3 + x^2 + x) + (x^5 + x^2 + x + 1) = x^5 + x^3 + x^2 + x + 1$					

- □ Multiplication by x can be achieved by one bit left shift of  $P_2$
- □ Need to be reduced after multiplication if degree greater than n-1
  - i.e., previously degree was n-1 (leading bit was 1)
  - Reduction, after multiplication result is XOR-ed with IRP

## Algorithm

- ☐ if leading bit of previous result '0'
  - One left shift
  - if leading bit of previous result '1'
    - One left shift

- P1 = 00100110,
- P2 = 100111110,
- modulus = 100011011
- XOR the result with least n-1 bits of IRP
  - □ (note: IRP is with degree n

Powers	Shift-Left Operation	Exclusive-Or			
$x^0 \otimes P_2$		10011110			
$x^1 \otimes P_2$	00111100	$(00111100) \oplus 00011011 = 00100111$			
$x^2 \otimes P_2$	01001110	01001110			
$x^3 \otimes P_2$	10011100	10011100			
$x^4 \otimes P_2$	00111000	$(00111000) \oplus 00011011 = 00100011$			
$x^5 \otimes P_2$	01000110	<u>01000110</u>			
$P_1 \otimes P_2 = (00100111) \oplus (01001110) \oplus (01000110) = 00101111$					

### Note

- □ A filed GF(2<sup>n</sup>) may have more than one irreducible polynomials
- ☐ In addition no role of irreducible polynomial where as the result of multiplication highly depends on irreducible polynomial (like mod p)

## Using generator

- $\square$  It is easier to define the elements of  $GF(2^n)$  using a generator
- □ Generator is 'g' where f(g)=0 for an irreducible polynomial of  $GF(2^n)$
- $\square$  Using the generator, the elements of  $GF(2^n)$  are

$$\{0, g, g, g^2, ..., g^N\}$$
, where  $N = 2^n - 2$ 

## Elements of GF(2<sup>4</sup>): $x^4 + x + 1$

## Inverse

□ Additive inverse: the element itself

□ Multiplicative inverse: for  $g^i$ , it is  $g^{-i}$  where  $-i \equiv k \mod 2^n - 1$ 

## **Operations**

## Operations (contd...)