

A Formal Verification of Reversible Primitive Permutations

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1. The definition

1.1 Reversible computing

Reversible computing is a model of computation in which every process can be run backwards. Simply put, in a reversible setting any program takes inputs and gives outputs (like usual), but can also go the other way around: provided the output it can reconstruct the input. In a mathematical sense, every function is expected to be invertible.

Why do we care about such a thing?

Firstly, having a programming language in which every function (or even a subset of functions) is reversible could lead to interesting and practical applications.

But we can also imagine reversible computers, in which the underlying architecture is inherently reversible: Toffoli gates provides a way to do so. The opposite of reversibility is loss of information, which (for thermodynamic reasons) leads to loss of energy and heat dissipation. This means that a non-reversible gate dissipates energy each time information is discarded, while in principle a reversible computer wouldn't.

Lastly, reversible computing is directly related to quantum computing, as each operation in a quantum computer must be reversible.

1.2 Reversible Primitive Permutations

In the article I decided to formalize, the authors focus on providing a functional model of reversible computation. They develop an inductively defined set of functions, called **Reversible Primitive Permutations** or **RPP**, which are expressive enough to represent all Primitive Recursive Functions (we talk about what this means in section ?). Here is the definition that we will use:

Definition 1 (Reversible Primitive Permutations). The class of **Reversible Primitive Permutations** or RPP is the smallest subset of functions $\mathbb{Z}^n \rightarrow \mathbb{Z}^n$ satisfying the following conditions:

- The n -ary **identity** $\text{Id}_n(x_1, \dots, x_n) = (x_1, \dots, x_n)$ belongs to RPP.

$$\begin{array}{ccc} x_1 & & x_1 \\ \vdots & \text{Id}_n & \vdots \\ x_n & & x_n \end{array}$$

The meaning of these diagrams should be fairly obvious: if the values on the left of a function are provided as inputs to that function, we get the values on the right as outputs.

- The **sign-change** $\text{Ne}(x) = -x$ belongs to RPP.

$$x \quad \text{Ne} \quad -x$$

- The **successor function** $\text{Su}(x) = x + 1$ belongs to RPP.

$$x \quad \text{Su} \quad x + 1$$

- The **predecessor function** $\text{Pr}(x) = x - 1$ belongs to RPP.

$$x \quad \text{Pr} \quad x - 1$$

- The **swap** $\text{Sw}(x, y) = (y, x)$ belongs to RPP.

$$\begin{array}{ccc} x & & y \\ y & \text{Sw} & x \end{array}$$

- If $f : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ and $g : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ belongs to RPP, then the **series composition** $(f \circ g) : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ belongs to RPP and is such that:

$$(f \circ g)(x_1, \dots, x_n) = g(f(x_1, \dots, x_n)) = (g \circ f)(x_1, \dots, x_n).$$

We remark that $f \circ g$ means that f is applied first, and then g , in opposition to the standard functional composition (denoted by \circ).

$$\begin{array}{ccc} x_1 & & z_1 \\ \vdots & f \circ g & \vdots \\ x_n & & z_n \end{array} = \begin{array}{ccc} x_1 & & y_1 \\ \vdots & f & \vdots \\ x_n & & y_n \end{array} \begin{array}{ccc} y_1 & & z_1 \\ \vdots & g & \vdots \\ y_n & & z_n \end{array}$$

- If $f : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ and $g : \mathbb{Z}^m \rightarrow \mathbb{Z}^m$ belongs to RPP, then the **parallel composition** $(f \parallel g) : \mathbb{Z}^{n+m} \rightarrow \mathbb{Z}^{n+m}$ belongs to RPP and is such that:

$$(f \parallel g)(x_1, \dots, x_n, y_1, \dots, y_m) = (f(x_1, \dots, x_n), g(y_1, \dots, y_m)).$$

$$\begin{array}{ccc} x_1 & & w_1 \\ \vdots & & \vdots \\ x_n & & w_n \\ y_1 & & z_1 \\ \vdots & & \vdots \\ y_m & & z_m \end{array} \begin{array}{ccc} & f \parallel g & \\ & & \end{array} = \begin{array}{ccc} x_1 & & w_1 \\ \vdots & & \vdots \\ x_n & & w_n \\ y_1 & & z_1 \\ \vdots & & \vdots \\ y_m & & z_m \end{array} \begin{array}{ccc} & f & \\ & & \\ & g & \end{array}$$

- If $f : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ belongs to RPP, then then **finite iteration** $\text{lt}[f] : \mathbb{Z}^{n+1} \rightarrow \mathbb{Z}^{n+1}$ belongs to RPP and is such that:

$$\text{lt}[f](x, x_1, \dots, x_n) = (x, \overbrace{(f \circ \dots \circ f)}^{\downarrow x \text{ times}}(x_1, \dots, x_n))$$

where $\downarrow (\cdot) : \mathbb{Z} \rightarrow \mathbb{N}$ is defined as

$$\downarrow x = \begin{cases} x, & \text{if } x \geq 0 \\ 0, & \text{if } x < 0 \end{cases}.$$

This means that the function f is applied $\downarrow x$ times to (x_1, \dots, x_n) .

$$\begin{array}{c} x \\ x_1 \\ \vdots \\ x_n \end{array} \begin{array}{|c|} \hline \text{lt}[f] \\ \hline \end{array} \begin{array}{c} x \\ y_1 \\ \vdots \\ y_n \end{array} = \begin{array}{c} x \\ x_1 \\ \vdots \\ x_n \end{array} \underbrace{\begin{array}{|c|} \hline f \\ \hline \end{array} \dots \begin{array}{|c|} \hline f \\ \hline \end{array}}_{\downarrow x \text{ times}} \begin{array}{c} x \\ y_1 \\ \vdots \\ y_n \end{array}$$

- If $f, g, h : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ belongs to RPP, then the **selection** $\text{lf}[f, g, h] : \mathbb{Z}^{n+1} \rightarrow \mathbb{Z}^{n+1}$ belongs to RPP and is such that:

$$\text{lf}[f, g, h](x, x_1, \dots, x_n) = \begin{cases} (x, f(x_1, \dots, x_n)), & \text{if } x > 0 \\ (x, g(x_1, \dots, x_n)), & \text{if } x = 0 \\ (x, h(x_1, \dots, x_n)), & \text{if } x < 0 \end{cases}.$$

We remark that the argument x which determines which among f , g and h must be used cannot be among the arguments of f , g and h , as that would break reversibility.

$$\begin{array}{c} x \\ x_1 \\ \vdots \\ x_n \end{array} \begin{array}{|c|} \hline \text{lf}[f, g, h] \\ \hline \end{array} \begin{array}{c} x \\ y_1 \\ \vdots \\ y_n \end{array} = \begin{cases} \begin{array}{c} f(x_1, \dots, x_n) \\ g(x_1, \dots, x_n) \\ h(x_1, \dots, x_n) \end{array} & \begin{array}{l} \text{if } x > 0 \\ \text{if } x = 0 \\ \text{if } x < 0 \end{array} \end{cases}$$

Remark 1. If we have two functions of different arity, for example $f : \mathbb{Z}^3 \rightarrow \mathbb{Z}^3$ and $g : \mathbb{Z}^5 \rightarrow \mathbb{Z}^5$, then we will still write $f \circ g$ to mean the function with arity $\max(3, 5)$ given by $(f \parallel \text{Id}_2) \circ g$. In general, the arity of the "smaller" function can be enlarged by a suitable parallel composition with the identity. The same goes for the arguments of the selection $\text{lf}[f, g, h]$.

$$\begin{array}{c} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{array} \begin{array}{|c|} \hline f \circ g \\ \hline \end{array} \begin{array}{c} z_1 \\ z_2 \\ z_3 \\ z_4 \\ z_5 \end{array} = \begin{array}{c} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{array} \begin{array}{|c|} \hline f \\ \hline \end{array} \begin{array}{c} y_1 \\ y_2 \\ y_3 \\ y_4 \\ y_5 \end{array} \begin{array}{|c|} \hline g \\ \hline \end{array} \begin{array}{c} z_1 \\ z_2 \\ z_3 \\ z_4 \\ z_5 \end{array} = \begin{array}{c} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{array} \begin{array}{|c|} \hline f \\ \hline \end{array} \begin{array}{c} y_1 \\ y_2 \\ y_3 \\ y_4 \\ y_5 \end{array} \begin{array}{|c|} \hline \text{Id}_2 \\ \hline \end{array} \begin{array}{c} z_1 \\ z_2 \\ z_3 \\ z_4 \\ z_5 \end{array}$$

1.3 Some examples

In order to get accustomed to this definition, let's see some examples.

Increment and decrement Let's try to imagine what addition should look like in RPP. Of course, addition is usually thought of as a function which takes two inputs and yields their sum: something like $\text{add}(x, y) = x + y$. But notice that this operation is not reversible: given only the output (the value $x + y$) it is impossible to obtain the original values (x, y) . As we will see, every function in RPP is reversible, so it's impossible to define addition in this way.

Instead, we can define a function inc in RPP which, given $n \in \mathbb{N}$ and $x \in \mathbb{Z}$, yields

$$\begin{array}{ccc} n & \boxed{\text{inc}} & n \\ x & & x + n \end{array}$$

If n is negative the output is just (n, x) . The fact that the above diagram is only valid for $n \in \mathbb{N}$ might bother some of you; we'll explain later why it is so, and how we can also make it work for $n \in \mathbb{Z}$.

For now let's focus on the output: we don't just have $x + n$ but also n , and indeed, given both n and $x + n$ we can reconstruct n (obviously) and x (by $(x + n) - n$). As a matter of fact, the following function dec also belongs to RPP:

$$\begin{array}{ccc} n & \boxed{\text{dec}} & n \\ x & & x - n \end{array}$$

and if we try to compose inc and dec we get this remarkable result:

$$\begin{array}{ccccc} n & \boxed{\text{inc}} & n & \boxed{\text{dec}} & n \\ x & & x + n & & x \end{array}$$

and similarly for $\text{dec} \circ \text{inc}$. So indeed dec is the inverse of inc , and we can write $\text{dec} = \text{inc}^{-1}$.

But we haven't said how to actually define inc . Well, just like this:

$$\text{inc} = \text{It}[\text{Su}]$$

This means that we apply the successor function Su to the value x , for $\downarrow n$ times. If $n \in \mathbb{N}$ then $\downarrow n = n$, so we effectively add n to the value x . If instead n is negative then $\downarrow n = 0$ and nothing changes.

Can you guess how dec is defined?

In a very similar manner, using the predecessor function:

$$\text{dec} = \text{It}[\text{Pr}]$$

and as we will shortly see, finding the inverse is not something that we have to do by hand.

Multiplication and square