

## PROJECT KEY POINTS

### SUMMARY & OPERATION DESCRIPTION

#### a) **PARAMETERS** (N,R,W,T)

- N = replication
- R = read QUORUM
- W = write QUORUM
- T = timeout

#### b) **ITEM** (fields):

- Key
- Value-String

#### b) **NODE**:

##### b.1) FIELDS:

- Key
- List of items (key, value)
- List of node (key)

##### b.2) METHODS:

###### – **JOIN**:

###### 1) CONTACT A NODE

- get nodes information from command line (java Node join remote\_ip remote\_port)
- ping the node to check if it is present or not
- send a request for getting list of nodes

###### 2) GET THE RESPONSE

- store the list of nodes

###### 3) REQUEST ITEMS & GETTING ITEMS

- ask to the next node (in the ring) to send the list of items for which the sender is responsible for
- receive the list and store it

###### 4) ANNOUNCING MSG & DELETE

- sender broadcasts a "announcing" msg to all other nodes
- all the other nodes have to remove items (they are not responsible for anymore)

###### – **RECOVERY**:

###### 1) RECOVERY NODE

- from a node terminal send a request to recover to a specific node (java Node recover remote\_ip remote\_port)

\* (——-Note: I have some doubts here——-)

###### 2) REQUEST LIST OF NODES

- node (recovered) sends a request to the node specified from command line in order to get the list of nodes (now in the network)

###### 3) UPDATE LISTS

- if node recovers and nothing happened (no join, no leaves, no items added/deleted) lists remain the same
- if node recovers and someone joined with a smaller key, remove the item/s acquired by the node with smaller key and insert it/them in the list of the node recovered
- if node recovers and someone joined with a higher key, remove from the list of the recovered node the item for which it is not more responsible.

##### b.3) LOCAL STORAGE:

- FIELDS
  - 1) Key
  - 2) Value
  - 3) Version

c) **CLIENT:**

c.1) FIELDS:

- Key

c.2) METHODS:

- **READ:**

1) SEND READ MSG

- client sends a read msg to one of the node (we call this node CORDINATOR) of the network

2) CHECK CORDINATOR

- if CORDINATOR is not present, node sends back a "not present" msg
- if CORDINATOR is present, node sends the read request to the N clockwise nodes and waits.

3) SEND RESPONSE

- if quorum is reached ( $R$  less than  $N$ ), node sends back to the client the item with the highest version number
- if quorum is not reached after a timeout  $T$ , node sends back to the client a msg to inform it.

- **WRITE:**

1) SEND WRITE MSG

- client sends a write msg to one of the node (we call it CORDINATOR) of the network

2) CHECK CORDINATOR

- if CORDINATOR is not present, node sends back a "not present" msg
- if CORDINATOR is present, node sends the write request to the N clockwise nodes and waits.

3) UPDATE/RESPONSE

- if quorum is reached ( $\max(R, W)$ ), node sends back to the client a success msg and to the other  $N$  nodes the update version of the item ( $\text{version} := \text{version} + 1$ )
- if quorum is not reached after a timeout  $T$ , node sends back to the client a msg to inform it.

## TECHNICAL PART

- key = (16-bit unsigned integers)
- Value = (16-bit unsigned integers)
- List items/values = HashSet
- Local storage = File

For point A we can think to store them in a file..... Remember to check  $R+W$  greater than  $N$

## DOUBTS :

How can we implement crash mechanism? — a crash emulation is present in the 2PC lab exercise

When contacting a node -¿ suppose to node a priori the ip/port of contacted node? (SOLVED)