

# FSPL – Assignment 3

冯诗伟 161220039

## 1 Problem 2.1

*Syntax* :  $\langle comm \rangle ::= \langle var \rangle, \langle var \rangle := \langle intexp \rangle, \langle intexp \rangle$

*Semantic* :  $\llbracket v_1, v_2 := e_1, e_2 \rrbracket_{comm} \sigma = [\sigma | v_1 : \llbracket e_1 \rrbracket_{intexp} \sigma, v_2 : \llbracket e_2 \rrbracket_{intexp} \sigma]$

## 2 Problem 2.2

2.1 a

*Syntax* :  $\langle comm \rangle ::= \mathbf{repeat} \ \langle comm \rangle \ \mathbf{until} \ \langle boolexp \rangle$

*Semantic* :  $\llbracket \mathbf{repeat} \ c \ \mathbf{until} \ b \rrbracket_{comm} \sigma = Y_{\Sigma \rightarrow \Sigma_{\perp}} F$

where  $F \ f \ \sigma = (\mathbf{if} \ \neg b \ \mathbf{then} \ f \ \mathbf{else} \ \mathbf{skip})_{\perp} (\llbracket c \rrbracket_{comm} \sigma)$  ? ? ?

2.2 b

*Syntactic sugar* :  $\mathbf{repeat} \ c \ \mathbf{until} \ b \stackrel{def}{=} c ; \mathbf{while} \ \neg b \ \mathbf{do} \ c$

2.3 c

For all  $\sigma \in \Sigma$ , according to the definition in (a),

-25

$$\begin{aligned} \llbracket \mathbf{repeat} \ c \ \mathbf{until} \ b \rrbracket_{comm} \sigma &= F \llbracket \mathbf{repeat} \ c \ \mathbf{until} \ b \rrbracket_{comm} \sigma \\ &= (\mathbf{if} \ \neg b \ \mathbf{then} \ (\mathbf{repeat} \ c \ \mathbf{until} \ b) \ \mathbf{else} \ \mathbf{skip})_{\perp} (\llbracket c \rrbracket_{comm} \sigma) \\ &= \mathbf{if} \ \neg (\llbracket b \rrbracket_{boolexp} \perp (\llbracket c \rrbracket_{comm} \sigma)) \ \mathbf{then} \ (\mathbf{repeat} \ c \ \mathbf{until} \ b) \ \perp (\llbracket c \rrbracket_{comm} \sigma) \ \mathbf{else} \ (\llbracket c \rrbracket_{comm} \sigma) \end{aligned}$$

According to the definition in (b),

$$\begin{aligned} \llbracket \mathbf{repeat} \ c \ \mathbf{until} \ b \rrbracket_{comm} \sigma &= \llbracket c ; \mathbf{while} \ \neg b \ \mathbf{do} \ c \rrbracket_{comm} \sigma \\ &= \llbracket \mathbf{while} \ \neg b \ \mathbf{do} \ c \rrbracket_{comm} \perp (\llbracket c \rrbracket_{comm} \sigma) \\ &= \mathbf{if} \ \neg (\llbracket b \rrbracket_{boolexp} \perp (\llbracket c \rrbracket_{comm} \sigma)) \ \mathbf{then} \ \llbracket c ; \mathbf{while} \ \neg b \ \mathbf{do} \ c \rrbracket_{comm} \perp (\llbracket c \rrbracket_{comm} \sigma) \ \mathbf{else} \ (\llbracket c \rrbracket_{comm} \sigma) \\ &= \mathbf{if} \ \neg (\llbracket b \rrbracket_{boolexp} \perp (\llbracket c \rrbracket_{comm} \sigma)) \ \mathbf{then} \ (\mathbf{repeat} \ c \ \mathbf{until} \ b) \ \perp (\llbracket c \rrbracket_{comm} \sigma) \ \mathbf{else} \ (\llbracket c \rrbracket_{comm} \sigma) \end{aligned}$$

Therefore, the two definitions are equivalent.

### 3 Problem 2.3

We can divide  $\sigma x$  into three conditions and discuss them respectively.

-5

(1) When  $\sigma x$  is even and  $\sigma x < 0$ ,  $\sigma x - 2 < 0$ .

So it will never terminates. So when  $\sigma x < 0$ ,  $\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} \sigma = \perp$ .

(2) When  $\sigma x$  is odd,  $\sigma x - 2$  is always add. But 0 is an even number, so  $x$  will never equal 0. So it will never terminates. So when  $\sigma x < 0$ ,  $\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} \sigma = \perp$ .

(3) When  $\sigma x$  is even and  $\sigma x \geq 0$ ,  $x$  must be in the form of  $2k$  where  $k \geq 0$ . Let's do the induction.

When  $\sigma x = 0$ , it will terminate directly and go to the state  $[\sigma|x : 0]$ .

Suppose that when  $\sigma x = 2k$ , it will terminate in the state of  $[\sigma|x : 0]$ .

$$\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} [\sigma|x : 2k + 2] \rightarrow^* \llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} [\sigma|x : 2k]$$

According the induction hypothesis,  $\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} [\sigma|x : 2k]$  will terminate in the state of  $[\sigma|x : 0]$ . So we can get that  $\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} [\sigma|x : 2k + 2]$  will terminate in the state of  $[\sigma|x : 0]$ .

Now we come to the conclusion that for all  $k \geq 0$ ,  $\llbracket \text{while } x \neq 0 \text{ do } x := x - 2 \rrbracket_{comm} [\sigma|x : 2k]$  will terminate in the state of  $[\sigma|x : 0]$ .

In conclusion, by combining(1)(2)(3) together we can prove the equation in Problem2.3.

### 4 Problem 2.4

$$F f \sigma = \text{if } \llbracket b \rrbracket_{boolexp} \sigma \text{ then } f_{\perp}(\llbracket c \rrbracket_{comm} \sigma) \text{ else } \sigma$$

First we have to prove that  $F$  is monotone. Consider two functions  $f$  and  $h$ , where  $f \sqsubseteq h$ .

For all  $\sigma \in \Sigma$ , if  $\llbracket b \rrbracket_{boolexp} \sigma = \text{false}$ , or if  $\llbracket b \rrbracket_{boolexp} \sigma = \text{true}$  and  $(\llbracket c \rrbracket_{comm} \sigma) = \perp$ , it is obvious that  $F f \sigma = F h \sigma$ , which means  $F f \sigma \sqsubseteq F h \sigma$ . If  $\llbracket b \rrbracket_{boolexp} \sigma = \text{true}$  and  $(\llbracket c \rrbracket_{comm} \sigma) \neq \perp$ ,  $F f \sigma = f_{\perp}(\llbracket c \rrbracket_{comm} \sigma)$ ,  $F h \sigma = h_{\perp}(\llbracket c \rrbracket_{comm} \sigma)$ . Because  $f \sqsubseteq h$ , we have that for all  $x$ ,  $f x \sqsubseteq h x$ . So  $f_{\perp}(\llbracket c \rrbracket_{comm} \sigma) \sqsubseteq h_{\perp}(\llbracket c \rrbracket_{comm} \sigma)$  which means  $F f \sigma \sqsubseteq F h \sigma$ .

Next we have to prove that  $F$  is continous. Consider an interesting chain of functions  $f_0 \sqsubseteq f_1 \sqsubseteq f_2 \sqsubseteq \dots$  and  $g = \sqcup_{n=0}^{\infty} f_n$ .

For all  $\sigma \in \Sigma$ , if  $\llbracket b \rrbracket_{boolexp} \sigma = \text{false}$ , or if  $\llbracket b \rrbracket_{boolexp} \sigma = \text{true}$  and  $(\llbracket c \rrbracket_{comm} \sigma) = \perp$ , it is obvious that  $F(\sqcup_{n=0}^{\infty} f_n) \sigma = (\sqcup_{n=0}^{\infty} F f_n) \sigma$ , which means  $F(\sqcup_{n=0}^{\infty} f_n) \sigma \sqsubseteq \sqcup_{n=0}^{\infty} (F f_n) \sigma$ .

If  $\llbracket b \rrbracket_{boolexp} \sigma = \text{true}$  and  $(\llbracket c \rrbracket_{comm} \sigma) \neq \perp$ ,

$$F(\sqcup_{n=0}^{\infty} f_n) \sigma = F g \sigma = g(\llbracket c \rrbracket_{comm} \sigma)$$

$$\sqcup_{n=0}^{\infty} (F f_n) \sigma = \sqcup_{n=0}^{\infty} f_n(\llbracket c \rrbracket_{comm} \sigma) \sqsubseteq g(\llbracket c \rrbracket_{comm} \sigma)$$

$$\therefore F(\sqcup_{n=0}^{\infty} f_n) \sigma \sqsubseteq \sqcup_{n=0}^{\infty} (F f_n) \sigma$$

In conclusion,  $F$  in the semantic equation for the *while* command is continous.

## 5 Problem 2.5

Define a series of commands:  $w_0, w_1, w_2, \dots$

$$w_0 \stackrel{\text{def}}{=} \text{while true do skip}$$

$$w_{i+1} \stackrel{\text{def}}{=} \text{if } b \text{ then } (c; w_i) \text{ else skip}$$

Define the function  $F$ :

$$F f \sigma = \text{if } \llbracket b \rrbracket_{\text{boolexp}} \sigma \text{ then } f_{\perp}(\llbracket c \rrbracket_{\text{comm}} \sigma) \text{ else } \sigma$$

$$\text{where } F \llbracket \text{while } b \text{ do } c \rrbracket_{\text{comm}} = \llbracket \text{while } b \text{ do } c \rrbracket_{\text{comm}}$$

For all  $\sigma \in \Sigma$ , consider the left side of the equation:

$$\begin{aligned} \llbracket \text{while } b \text{ do } c \rrbracket_{\text{comm}} \sigma &= \sqcup_{n=0}^{\infty} \llbracket w_n \rrbracket_{\text{comm}} \sigma \\ &= \sqcup_{n=0}^{\infty} F^n \perp \sigma \end{aligned}$$

Consider the right side of the equation:

$$\llbracket \text{while } b \text{ do } (c; \text{if } b \text{ then } c \text{ else skip}) \rrbracket_{\text{comm}} \sigma = \llbracket \text{while } b \text{ do } (c; w_1) \rrbracket_{\text{comm}} \sigma$$

Suppose  $\text{while } b \text{ do } (c; w_1)$  terminates after  $\text{while}$  testing  $b$  exactly  $n(> 0)$  times. It indicates that  $\text{while } b \text{ do } (c; w_1)$  terminates after testing  $b$   $2n-2$  or  $2n-1$  times because both  $\text{while}$  and  $w_1$  test  $b$  and it can terminate at both places. Therefore we can operate on the equation above:

$$\begin{aligned} \llbracket \text{while } b \text{ do } (c; \text{if } b \text{ then } c \text{ else skip}) \rrbracket_{\text{comm}} \sigma &= \llbracket \text{while } b \text{ do } (c; w_1) \rrbracket_{\text{comm}} \sigma \\ &= \sqcup_{n=1}^{\infty} (F^{2n-2} \perp \sigma \sqcup F^{2n-1} \perp \sigma) \\ &= \sqcup_{n=0}^{\infty} F^n \perp \sigma \end{aligned}$$

$$\therefore \llbracket \text{while } b \text{ do } c \rrbracket_{\text{comm}} = \llbracket \text{while } b \text{ do } (c; \text{if } b \text{ then } c \text{ else skip}) \rrbracket_{\text{comm}}$$

## 6 Problem 2.9

$$\begin{aligned} \text{for } v := e_0 \text{ to } e_1 \text{ do } c &\stackrel{\text{def}}{=} \text{newvar } w := e_1 \text{ in newvar } v := e_0 \text{ in} \\ &(\text{while } v < w \text{ do } (c; v := v + 1)); \text{if } v = w \text{ then } c \text{ else skip} \end{aligned}$$

-1

## 7 Problem 2.10

If  $c$  doesn't contain occurrence of *dotwice*,

-8

$$\text{dotwice } c = c; c$$

If  $c$  contains occurrences of *dotwice*,  $c = \text{dotwice } d$

$$\text{dotwice } c = \text{dotwice } (\text{dotwice } d) = \text{dotwice } (d; d) = d; d; d; d$$

The right side will eventually contain no occurrence of *dotwice*. If the  $d$  mentioned above also contains the occurrence of *dotwice*, the length of final result will just explode in an exponential way.