

```
In [ ]: import math
import numpy as np
import matplotlib.pyplot as plt
```

```
In [ ]: def plot_theta(f, g, cp, cpp, n_0, min_x, max_x):
    xs = np.arange(min_x, max_x, 1, dtype='int64')
    fn_ys = f(xs)
    gn_ysp = cp*g(xs)
    gn_yspp = cpp*g(xs)
    plt.axvline(x = n_0, color = 'k', linestyle='--')
    plt.scatter(x=xs, y=fn_ys, marker='x')
    plt.scatter(x=xs, y=gn_ysp, marker='.')
    plt.scatter(x=xs, y=gn_yspp, marker='.')
    plt.xlabel('n')
    plt.ylabel('$\lambda(n)$')
    plt.legend(['n_0 = '+str(n_0), 'f(n)', 'c\g(n)', 'c\\'g(n)'])

def plot_oh_incr(f, g, c, n_0, min_x, max_x, incr):
    xs = np.arange(min_x, max_x, incr, dtype='int64')
    fn_ys = f(xs)
    gn_ys = c*g(xs)
    plt.axvline(x = n_0, color = 'k', linestyle='--')
    plt.scatter(x=xs, y=fn_ys, marker='x')
    plt.scatter(x=xs, y=gn_ys, marker='.')
    plt.xlabel('n')
    plt.ylabel('$\lambda(n)$')
    plt.legend(['n_0 = '+str(n_0), 'f(n)', 'cg(n)'])

def plot_oh(f, g, c, n_0, min_x, max_x):
    plot_oh_incr(f, g, c, n_0, min_x, max_x, 1)
```

## Big-Oh Definition

Given positive functions  $f(n)$  and  $g(n)$ , we can say that  $f(n)$  is  $O(g(n))$  if and only if there exists positive constants  $c$  and  $n_0$  such that:

$$f(n) \leq c \cdot g(n), \forall n \geq n_0$$

## Big-Oh Using Rules

**Drop smaller terms rule:**

If  $f(n) = (1 + h(n))$  and  $h(n) \rightarrow 0$  as  $n \rightarrow \infty$ , then  $f(n)$  is  $O(1)$

**Q1. Prove that  $n^3 + 2n^2$  is  $O(n^3)$  using the multiplication and drop smaller terms rules**

**Q2. Prove that  $n^3 + 2n^2 \log(n)$  is  $O(n^3)$  using the multiplication and drop smaller terms rules**

## Big-Omega and Big-Theta Definitions

**Big-Omega:** Given positive functions  $f(n)$  and  $g(n)$ , we can say that  $f(n)$  is  $\Omega(g(n))$  if and only if there exists strictly positive constants  $c$  and  $n_0$  such that:

$$f(n) \geq c \cdot g(n), \forall n \geq n_0$$

$\Omega$  expresses that a function  $f(n)$  grows at least as fast as  $g(n)$ .

**Big-Theta:** Given positive functions  $f(n)$  and  $g(n)$ , we can say that  $f(n)$  is  $\Theta(g(n))$  if and only if there exists positive constants  $c'$ ,  $c''$  and  $n_0$  such that:

$$f(n) \leq c' \cdot g(n), f(n) \geq c'' \cdot g(n), \forall n \geq n_0$$

$\Theta$  expresses that a function  $f(n)$  grows exactly as fast as  $g(n)$ .

**Q3. Prove that  $2n + 1$  is  $\Omega(3n)$  and hence  $2n + 1$  is  $\Theta(3n)$**

**Plot  $\Omega$**

```
In [ ]: # c = ???
# n_0 = ???
# f = Lambda n: ???
# g = Lambda n: ???

plot_oh(f, g, c, n_0, 0, 50)
```

**Plot  $\Theta$**

```
In [ ]: # cp = ???
# cpp = ???
# n_0 = ???
# f = Lambda n: ???
# g = Lambda n: ???

plot_theta(f, g, cp, cpp, n_0, 0, 50)
```

## Big-Omega/Big-Theta Questions: Basics

**Q4. Prove that 5 is  $\Omega(1)$ , and hence 5 is  $\Theta(1)$**

**Q5. Prove that 4 is  $\Omega(2)$ , and hence 4 is  $\Theta(2)$**

**Q6. Prove that  $2n + 1$  is  $\Omega(n)$ , and hence  $2n + 1$  is  $\Theta(n)$**

## Big-Omega/Big-Theta Questions: Medium Difficulty

Q7. Prove that  $n^2$  is  $\Omega(2n^2)$

Q8. Prove that  $n^2 - 3$  is  $\Omega(n^2)$

Q9. Prove that  $n^2 - 5n$  is  $\Omega(n^2)$ , and hence is  $\Theta(n^2)$

Q10. Prove that  $n^2 + 1$  is  $\Omega(n^2)$

## Little-Oh Definition

Given positive functions  $f(n)$  and  $g(n)$ , we can say that  $f(n)$  is  $o(g(n))$  if **for all positive real constants**  $c > 0$  there exists  $n_0$  such that:

$$f(n) < c \cdot g(n), \forall n \geq n_0$$

### Important:

- $c \in \mathbb{R}_*^+$
- $\mathbb{R}_*^+ = \{x \in \mathbb{R} | x > 0\}$

Q11. Prove or disprove that 5 is  $o(1)$

Q12. Prove or disprove that 5 is  $o(n)$

## Little-Oh Questions

Q13. Prove or disprove that  $n$  is  $o(n^2)$

Q14. Prove or disprove that 1 is  $o(\log n)$

Q15. Prove or disprove that  $\log n$  is  $o(1)$

## Additional Practice Questions (more challenging)

Q16. Prove or disprove that 1 is  $\Omega(n)$

Q17. Prove or disprove that  $n$  is  $\Omega(1)$

Q18. Prove or disprove that  $n^2$  is  $\Omega(n)$

Q19. Prove or disprove that  $n$  is  $o(n \log n)$

Q20. Given that  $f(n) = n^2$  if  $n$  is even, and  $f(n) = n$  if  $n$  is odd. From the definitions find the  $O$  and  $\Omega$  behaviours of  $f(n)$ .

Warning: be careful to find a single  $c$  that works for all  $n$ , not separate  $c$  for even and odd  $n$

Q21. Prove or disprove that  $n \log n$  is  $o(n^2)$