

COURSEWORK SUBMISSION FORM

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Module Title Introduction to Algorithms		Module Convenor Aidin JALILZADEH
Coursework Title CW_20411905_N086		Module Code CELEN086

Compulsory

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Question 1

(a)

Algorithm: addN(x,n)

Requires: a number **x** and a positive integer **n**

Return: $n \times x$

```
1. if n == 1 then
2.   return x
3. else
4.   return x + addN(x,n-1)
5. endif
```

(b)

Algorithm: power(x,n)

Requires: a positive integer **x** and a positive integer **n**

Return: the value of x^n

```
1. if n == 0 then
2.   return 1
3. else
4.   return addN(x,power(x,n-1))
5. endif
```

(c)

*trace **power(x,n)** for **power(4,3)***

```
power(4,3) n!=0
    return addN(4,power(4,2))
power(4,2) n!=0
    return addN(4,power(4,1))
power(4,1) n!=0
    return addN(4,power(4,0))
power(4,0) n==0
    return 1
```

addN(4,16)=64

addN(4,4)=16

addN(4,1)=4



Question 2

(a)

Algorithm: `concat(L1,L2)`

Requires: 2 lists L1 L2

Return: a new list with L1 attached to the head of L2

```
1. if isEmpty(L1) then
2.   return L2
3. else
4.   return cons(head(L1),concat(tail(L1),L2))
5. endif
```

(b)

Algorithm: `subset(L1,L2)`

Requires: 2 lists L1 L2

Return: TRUE or FALSE

```
1. if length(L1) > length(L2) then
2.   return False
3. elseif isEmpty(L1) then
4.   return True
5. elseif linSearch(head(L1),L2) == false then
6.   return False
7. else
8.   return subset(tail(L1),L2)
9. endif
```

Algorithm: `linSearch(x,list)`

Requires: a positive integer x and a list

Return: TRUE or FALSE

```
1. if x == head(list) then
2.   return True
3. elseif isEmpty(list) then
4.   return False
5. else
6.   return linSearch(x,tail(list))
7. endif
```

(c)

trace subset(L1,L2) for subset([1,5,2],[4,5,1,0,2,9])

```
subset([1,5,2],[4,5,1,0,2,9])
    length(L1) > length(L2)??      NO!
    isEmpty(L1)??                  NO!
    linSearch(1,[4,5,1,0,2,9])??    TRUE!
    return subset([5,2],[4,5,1,0,2,9])
subset([5,2],[4,5,1,0,2,9])
    length(L1) > length(L2)??      NO!
    isEmpty(L1)??                  NO!
    linSearch(5,[4,5,1,0,2,9])??    TRUE!
    return subset([2],[4,5,1,0,2,9])
subset([2],[4,5,1,0,2,9])
    length(L1) > length(L2)??      NO!
    isEmpty(L1)??                  NO!
    linSearch(2,[4,5,1,0,2,9])??    TRUE!
    return subset([], [4,5,1,0,2,9])
subset([], [4,5,1,0,2,9])
    length(L1) > length(L2)??      NO!
    isEmpty(L1)??                  YES!
    return TRUE
```

Question 3

(a)

Algorithm: level(x,binT)

Requires: a binary tree and a node value x

Return: the generation level to which the node x belongs

```
1. if x == root(binT) then
2.     return 0
3. elseif searchBT(x,left(binT)) then
4.     return 1 + level(x,left(binT))
5. else
6.     return 1 + level(x,right(binT))
7. endif
```

Algorithm: searchBT(x,binT)

Requires: a binary tree and a node value x

Return: true or false

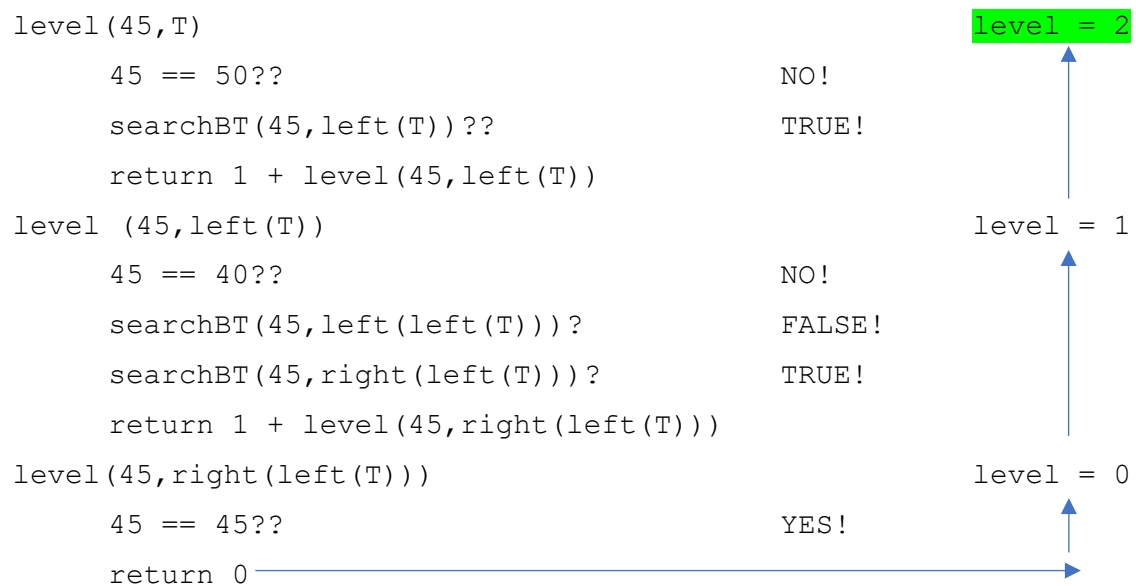
```

1. if isLeaf(binT) then
2.   return false
3. elseif x == root(binT) then
4.   return true
5. elseif isLeaf(left(binT)) && isLeaf(Right(binT)) then
6.   return false
7. else
8.   return searchBT(x,left(binT)) || searchBT(x,right(binT))
9. endif

```

(b)

*trace **level(x,binT)** for **level(45,T)***



(c)

Algorithm: `level(x, BST)`

Requires: a BST and a node value `x`

Return: the generation level to which the node `x` belongs

```

1. if isLeaf(BST) then
2.   return 0
3. elseif x == root(BST) then
4.   return 0
5. elseif x > root(BST) then
6.   return 1 + level(x, right(BST))
7. else
8.   return 1 + level(x, left(BST))
9. endif

```

*trace **level(x, BST)** for **level(45, T)***

<code>level(45, T)</code>		level=2
<code>isLeaf(T)??</code>	NO!	↑
<code>45 == 50??</code>	NO!	
<code>45 > 50??</code>	NO!	
<code>return 1 + level(45, left(T))</code>		
<code>level(45, left(T))</code>		level=1
<code>isLeaf(left(T))</code>	NO!	↑
<code>45 == 40??</code>	NO!	
<code>45 > 40??</code>	YES!	
<code>return 1 + level(45, right(left(T)))</code>		
<code>level(45, right(left(T)))</code>		level=0
<code>isLeaf(right(left(T)))??</code>	NO!	↑
<code>45 == 45??</code>	YES!	
<code>return 0</code>		→

level(x, BST) is faster than **level(x, T)**

level(x, BST) is $O(\lg n)$

level(x, T) is $O(n)$

Question 4

(a)

Algorithm: partition(L)

Requires: a list

Return: LP, PP and RP

```
1. if isEmpty(L) then
2.   return [], [], []
3. else
4.   return pHelper(head(L), tail(L), [], [])
5. endif
```

Algorithm: pHelper(x, list, L1, L2)

Requires: a positive integer x and 3 lists

Return: 3 new lists

```
1. if isEmpty(list) then
2.   return L1, cons(x, nil), L2
3. elseif x < head(list) then
4.   return pHelper(x, tail(list), L1, cons(head(list), L2))
5. else
6.   return pHelper(x, tail(list), cons(head(list), L1), L2)
7. endif
```

(b)

Algorithm: quicksort(list)

Requires: an unsorted list

Return: a sorted list of numbers

```
1. if isEmpty(list) || isEmpty(tail(list)) then
2.   return list
3. else
4.   let (L1, L2, L3) = partition(list)
5.   let S1 = quicksort(L1)
6.   let S2 = quicksort(L3)
7.   return concat(S1, concat(L2, S2))
8. endif
```

Algorithm: concat(L1,L2)

Requires: 2 lists L1 L2

Return: a new list

```

1. if isEmpty(L1) then
2.   return L2
3. else
4.   return cons(head(L1),concat(tail(L1),L2))
5. endif

```

(c)

*Trace **partition(L)** for [10,20,40,50,45]*

```

partition([10,20,40,50,45])
  isEmpty([10,20,40,50,45])??          NO!
  return pHelper(10,[20,40,50,45],[],[20])

pHelper(10,[20,40,50,45],[],[20])
  isEmpty([20,40,50,45])??          NO!
  10 < 20??                          YES!
  return pHelper(10,[40,50,45],[],[20])

pHelper(10,[40,50,45],[],[20])
  isEmpty([40,50,45])??          NO!
  10 < 40??                          YES!
  return pHelper(10,[50,45],[],[40,20])

pHelper(10,[50,45],[],[40,20])
  isEmpty([50,45])??          NO!
  10 < 50??                          YES!
  return pHelper(10,[45],[],[50,40,20])

pHelper(10,[45],[],[50,40,20])
  isEmpty([45])??          NO!
  10 < 45??                          YES!
  return pHelper(10,[],[],[45,50,40,20])

pHelper(10,[],[],[45,50,40,20])
  isEmpty([])??          YES!
  return [], [10], [45,50,40,20]

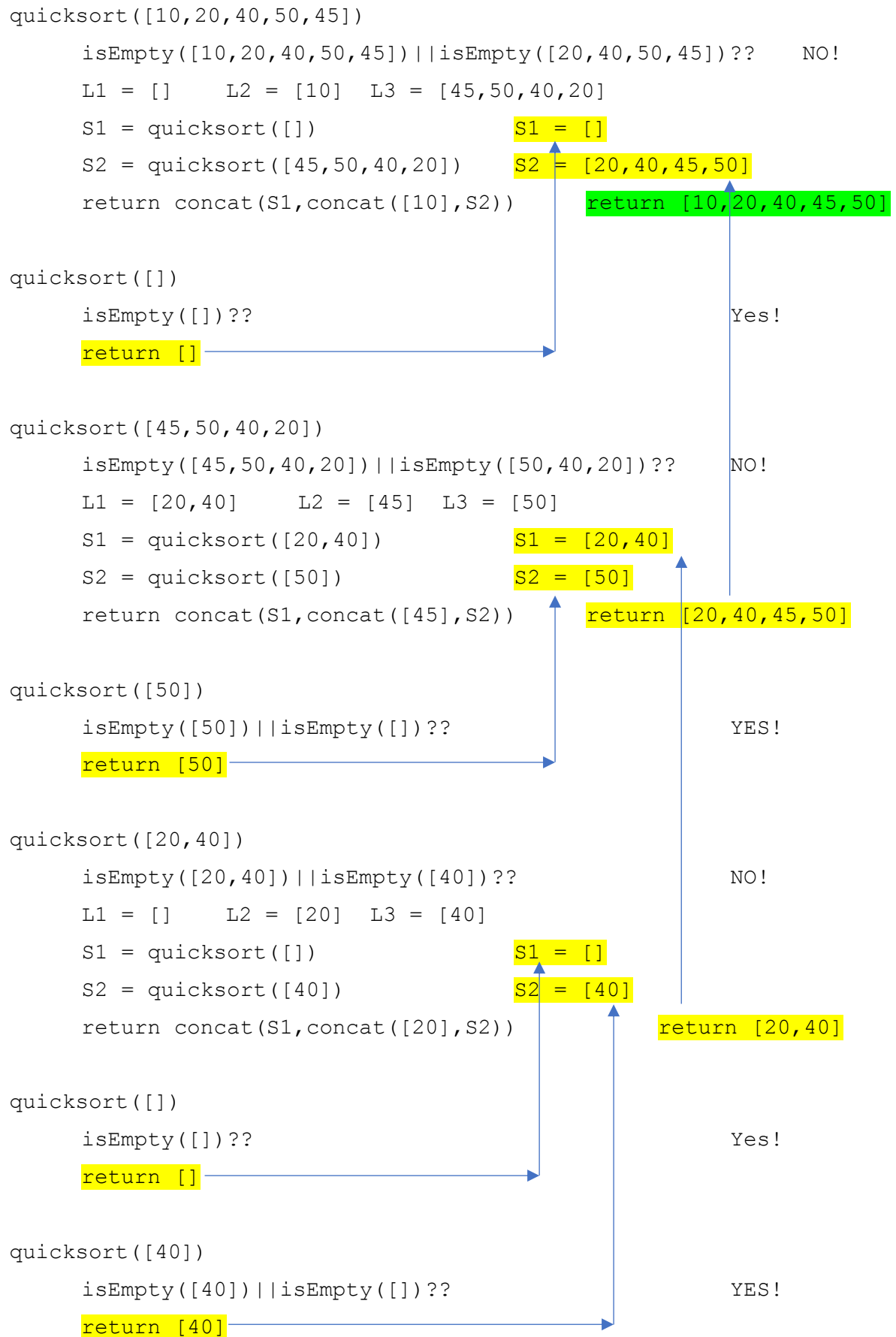
```


Trace **partition(L)** for **[10,9,12,8,15]**

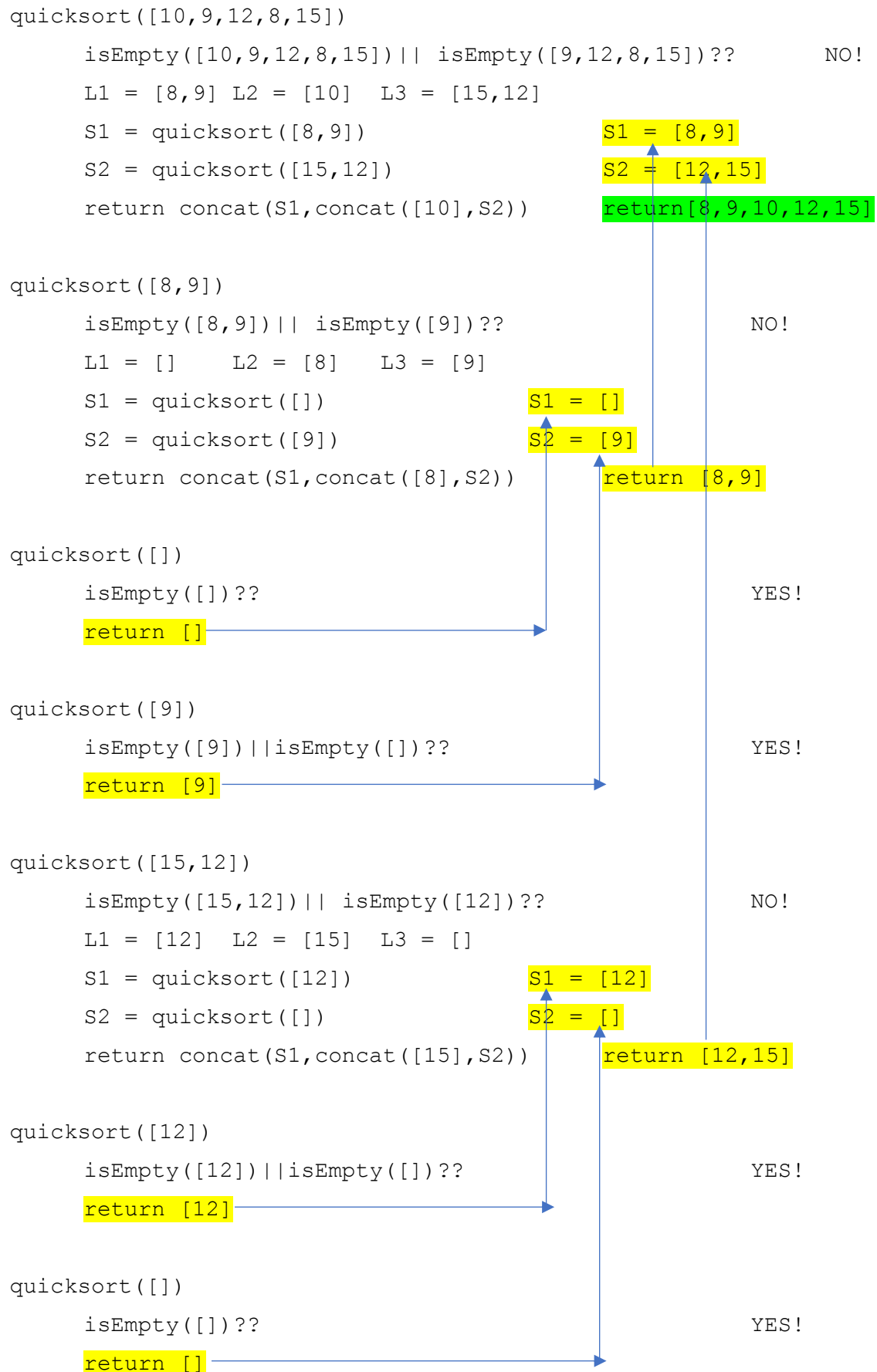
```
partition([10,9,12,8,15])
    isEmpty([10,9,12,8,15])??          NO!
    return pHelper(10,[9,12,8,15],[],[])

pHelper(10,[9,12,8,15],[],[])
    isEmpty([9,12,8,15])??          NO!
    10 < 9??                        NO!
    return pHelper(10,[12,8,15],[9],[])
pHelper(10,[12,8,15],[9],[])
    isEmpty([12,8,15])??          NO!
    10 < 12??                      YES!
    return pHelper(10,[8,15],[9],[12])
pHelper(10,[8,15],[9],[12])
    isEmpty([8,15])??          NO!
    10 < 8??                      NO!
    return pHelper(10,[15],[8,9],[12])
pHelper(10,[15],[8,9],[12])
    isEmpty([15])??          NO!
    10 < 15??                  YES!
    return pHelper(10,[],[8,9],[15,12])
pHelper(10,[],[8,9],[15,12])
    isEmpty([])??          YES!
    return [8,9],[10],[15,12]
```

Trace **quicksort(list)** for **[10,20,40,50,45]**



Trace **quicksort(list)** for **[10,9,12,8,15]**



(d)

As for List1, the head of List1 is the smallest number in List1.
This leads to longer length of S2.
And longer length of S2 leads to more operations for partition,
thus leading to more operations for L1.

As for List2, the head of List2 is the middle number in List2.
This leads to the same length of S1 and S2.
The length of S1(S2) is nearly half of the length of List2.
And shorter length of S1 and S2 leads to less operations for
partition, thus leading to less operations for L2.