

# Discrete Mathematics Exercise 15

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1. **Solution:** a), b), d), f) are trees.

2. **Proof:**

We prove there are no simple circuits in  $G$  by contradiction.

Assume there is a simple circuit in  $G$ .

First, since  $G$  is connected, there are at least  $(n - 1)$  edges in  $G$ . (because the addition of an edge will decrease at most one connected component and if there are no edges in a graph with  $n$  vertices, there are  $n$  connected components.)

Thus, if you remove any edge from  $G$ , the new graph has only  $(n - 2)$  edges and is definitely not connected.

If there is a simple circuit in  $G$ , if you remove an edge involved in the simple circuit, the new graph is still connected. **Contradiction.**

Thus, there are no simple circuits in  $G$ , i.e.  $G$  is a connected graph with no simple circuits.

In other words,  $G$  is a tree.

**QED**

3. **Proof:**

We prove  $G$  is connected by induction.

First, we prove that if  $G$  has no simple circuits, there are at most  $(n - 1)$  edges in  $G$  and when there are  $(n - 1)$  edges in  $G$  without simple circuits,  $G$  is connected.

**BASE STEP.**  $n = 1$ . Obvious.  $n = 2$ . Obvious.

**INDUCTION STEP.**

If  $n = k$ , there are at most  $(k - 1)$  edges in  $G$  if  $G$  has no simple circuits and when there are  $(n - 1)$  edges in  $G$  without simple circuits,  $G$  is connected.

When  $n = k + 1$ , we can at most add an edge between the new vertex  $v$  and a vertex  $u$  in  $G$ . If we add two edges, the second edge is either from a vertex in  $G$  to another in  $G$  or from  $v$  to a vertex  $x$  (not  $u$ ) in  $G$ .

In the former case, there is already a simple path between the two vertices. Thus, the addition of the new edge generates a circuit in the new graph.

In the latter case, there is a simple path from  $u$  to  $x$ . Thus, the addition of the first edge ensures there is a simple path from  $v$  to  $x$ . The addition of the second edge (from  $v$  to  $x$ ) generates a circuit in the new graph.

Thus, when  $n = k + 1$ , there are at most  $(n - 1)$  edges in  $G$  if  $G$  has no simple circuits and when there are  $(n - 1)$  edges in  $G$  without simple circuits,  $G$  is connected.

Thus,  $G$  is connected, i.e.  $G$  is a connected graph with no simple circuits.

In other words,  $G$  is a tree.

**QED**