

Exercise Sheet 3

Discrete Mathematics, 2020.9.25

1. a) Prove that $p \rightarrow (q \rightarrow p)$ is a tautology.
b) Prove that $(p \rightarrow q \rightarrow r) \rightarrow (p \rightarrow q) \rightarrow (p \rightarrow r)$ is a tautology.
c) Prove that $p \rightarrow q \rightarrow r \equiv (p \wedge q) \rightarrow r$.
Here, \rightarrow is right associative, i.e. $\phi \rightarrow \psi \rightarrow \chi$ means $\phi \rightarrow (\psi \rightarrow \chi)$.
2. (P35, Ex.22-24, [R])
a) Show that $(p \rightarrow q) \wedge (p \rightarrow r)$ and $p \rightarrow (q \wedge r)$ are logically equivalent.
b) Show that $(p \rightarrow r) \wedge (q \rightarrow r)$ and $(p \vee q) \rightarrow r$ are logically equivalent.
c) Show that $(p \rightarrow q) \vee (p \rightarrow r)$ and $p \rightarrow (q \vee r)$ are logically equivalent.
3. (P35, Ex.31, [R]) Show that $(p \rightarrow q) \rightarrow r$ and $p \rightarrow (q \rightarrow r)$ are not logically equivalent.
4. (P35, Ex.32, [R]) Show that $(p \wedge q) \rightarrow r$ and $(p \rightarrow r) \wedge (q \rightarrow r)$ are not logically equivalent.
5. Consider the compound proposition $\phi = p \rightarrow (q \oplus r)$ where p, q, r are propositional variables.
(a) Find a compound proposition ψ in CNF such that $\phi \equiv \psi$.
(b) Use the algorithm that we learned in class to construct a compound proposition ψ in CNF such that ϕ is satisfiable if and only if ψ is satisfiable.
6. (P36, Ex.52, [R]) In this exercise we will show that $\{|\}$ is a functionally complete collection of logical operators. (**Note:** $p|q$ means p NAND q . The proposition p NAND q is true when either p or q , or both, are false; and it is false when both p and q are true. The operators $|$ is called the Sheffer stroke after H. M. Sheffer)
a) Show that $\phi|\phi$ is logically equivalent to $\neg\phi$.
b) Show that $(\phi|\psi)|(\phi|\psi)$ is logically equivalent to $\phi \wedge \psi$.
c) Show that $\{|\}$ is a functionally complete collection of logical operators based on the results above.
7. Consider potential process of determining whether
$$(\neg p_1 \vee p_2) \wedge (\neg p_1 \vee p_3 \vee p_5) \wedge (\neg p_2 \vee p_4) \wedge (\neg p_3 \vee \neg p_4) \wedge (p_1 \vee p_5 \vee \neg p_2) \wedge (p_2 \vee p_3) \wedge (p_2 \vee \neg p_3) \wedge (p_6 \vee \neg p_5)$$
is SAT or UNSAT.
a) Calculus UnitPro(\mathcal{J}_1) where $\mathcal{J}_1 = [p_1 \mapsto \mathbf{T}, p_3 \mapsto \mathbf{F}]$.
b) Calculus UnitPro(\mathcal{J}_2) where $\mathcal{J}_2 = [p_3 \mapsto \mathbf{F}]$.