

## 计算机体系结构 Exercise 02

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### 3.5. Solution:

The truth table is as follows.

$A$	$B$	$C$	$OUT$
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

### 3.7. Solution:

When  $A = 1, B = 0$ , the high voltage is directly linked to low voltage and the voltage of  $OUT$  is ambiguous. The ambiguity of  $OUT$  under this circumstance is a fatal error.

**3.21. Solution:** There are  $2^{14}=16384$  units. Thus, memory contains 16384 Bytes = 32768 Nibbles.

### 3.24. Solution:

a)  $X$  acts as a control that decides whether  $B$  or  $C$  is the operand. When  $X = 0$ , the result is  $A + B$ ; when  $X = 1$ , the result is  $A + C$ .

b) Let  $C = \bar{B}$ . Also, let  $X$  be the carry-in of the 0-th adder (shown in the following figure).

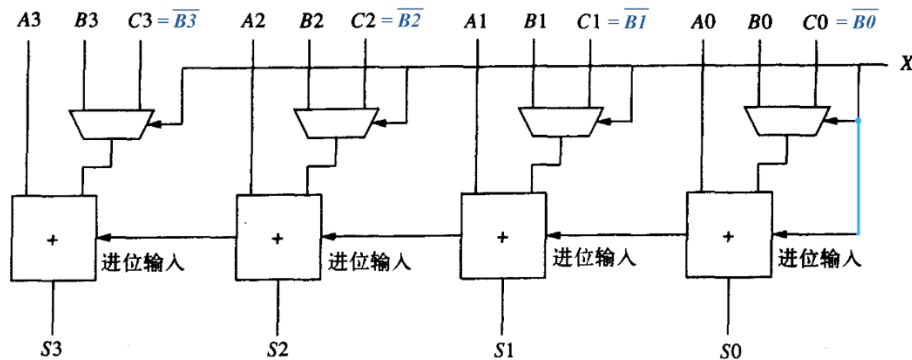
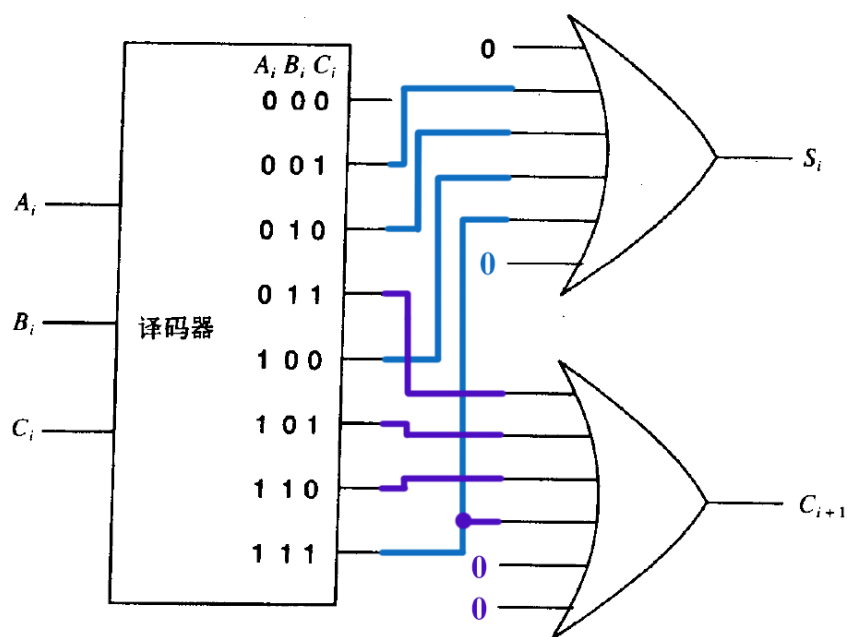


图3-39 结构框图

When  $X = 0$ , the circuit in the figure above acts as an adder.

When  $X = 1$ , the circuit acts as subtracter.

3.26. Solution:

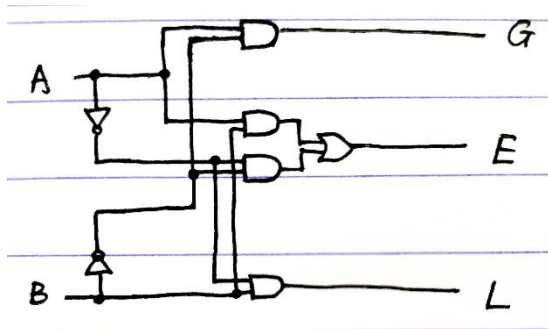


3.30. Solution:

a) The truth table is as follows.

A	B	G	E	L
0	0	0	1	0
0	1	0	0	1
1	0	1	0	0
1	1	0	1	0

b)



c)

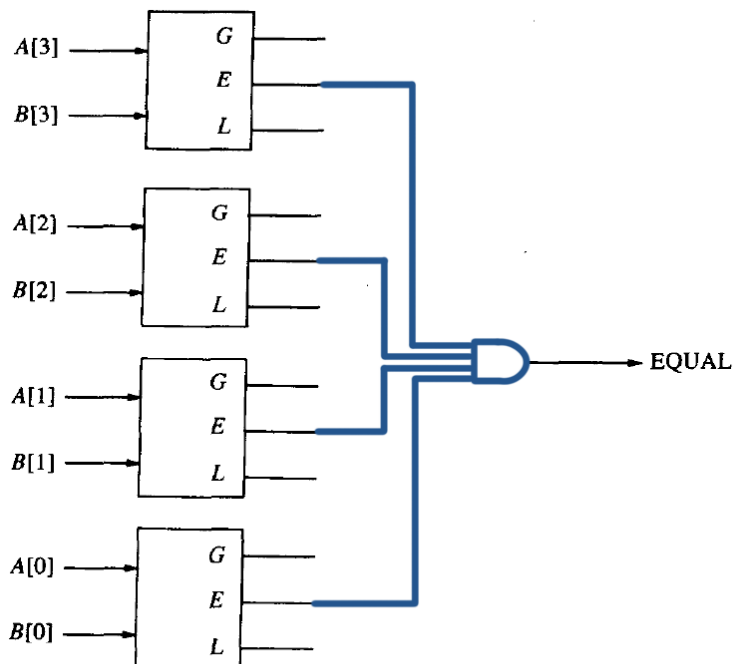


图3-41 习题3.30的结构框图

4.5 Solution:

a) 地址 3 存放的数值是 0000 0000 0000 0000; 地址 6 存放的数值是 1111 1110 1101 0011。

b) (1) 地址 0 对应的补码数:  $(0001\ 1110\ 0100\ 0011)_2 = \underline{7747}$ ;

地址 1 对应的补码数:  $(1111\ 0000\ 0010\ 0101)_2 = \underline{-4059}$

(2) 地址 4 对应的数值为  $(0000\ 0000\ 0110\ 0101)_2 = (101)_{10}$ 。ASCII 码 101 对应的字符为 e。

(3) 地址 6 和 7 对应的二进制数值为 1 111 1110 1 101 0011 0000 0110 1101 1001, IEEE 标准下指数部分为 +126, 尾数为  $(0.10100110000011011011001)_2 = (0.6486464739)_{10}$ , 故对应的浮点数为  $-1.6486464739 \times 10^{126}$ 。

(4) 地址 0 对应的无符号数: 7747, 地址 1 对应的无符号数为 61477。

c) 地址 0 对应的数值为 0001 111 001 000 011，对应的 LC3 标准下指令为  $R7 \leftarrow R1 + R3$ 。

d) 地址 5 保存的数值为 0110，即指向的单元地址是 6。

该地址中存放的数值为 1111 1110 1101 0011。

#### 4.7 Solution:

Since there exist 60 opcodes, we need at least 6 bits to represent these opcodes (since  $2^6 = 64 > 60$ ).

Since there exist 32 registers, we need at least 5 bits and 5 bits to represent SR and DR respectively (since  $2^5 = 32$ ).

Thus, in a 32-bit instruction, there are at most  $32 - 6 - 5 - 5 = 16$  bits to represent immediate number.

Since the immediate number is represented as 2's complement,  $-2^{16} \leq IMM \leq 2^{16} - 1$ , i.e. the range of the immediate number is  $-2^{16} \sim 2^{16} - 1$ .

#### 4.13 Solution:

1) For Intel x86 instruction "ADD [eax], edx":

FETCH – 1, DECODE – 1; EVALUATE ADDRESS – 1; FETCH OPERAND – 100; EXECUTE – 1; STORE RESULT – 100. Thus, it takes 204 periods in total.

2) For LC-3 instruction "ADD R6,R2,R6":

FETCH – 1, DECODE – 1; FETCH OPERAND (from register files) – 1; EXECUTE – 1; STORE RESULT (into register files) – 1. Thus, it takes 5 periods in total.