

Discrete Mathematics Exercise 8

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1. Solution:

For any $x \in \mathbb{R}$, $x - x = 0 \in \mathbb{Q}$. Thus, $\forall x \in \mathbb{R} (x\mathcal{R}x)$, i.e. \mathcal{R} is reflexive.

For any $x, y \in \mathbb{R}$, $x - y \in \mathbb{Q} \rightarrow y - x = -(x - y) \in \mathbb{Q}$.

Thus, $\forall x, y \in \mathbb{R} (x\mathcal{R}y \rightarrow y\mathcal{R}x)$, i.e. \mathcal{R} is symmetric.

Exists $x = 2, y = 3 \in \mathbb{R}$ s.t. $x\mathcal{R}y \wedge y\mathcal{R}x$ but $x \neq y$. Thus, \mathcal{R} is not antisymmetric.

For any $x, y, z \in \mathbb{R}$, $x - y \in \mathbb{Q} \wedge y - z \in \mathbb{Q} \rightarrow x - z = (x - y) + (y - z) \in \mathbb{Q}$.

Thus, $\forall x, y, z \in \mathbb{R} (x\mathcal{R}y \wedge y\mathcal{R}z \rightarrow x\mathcal{R}z)$, i.e. \mathcal{R} is transitive.

In conclusion, \mathcal{R} is reflexive, symmetric and transitive, but is not antisymmetric.

2. b) Solution:

Exists $x = 0 \in \mathbb{Z}$ s.t. $x^2 = 0 \not\geq 1$, i.e. $\neg x\mathcal{R}x$. Thus, \mathcal{R} is not reflexive.

For any $x, y \in \mathbb{Z}$, $xy \geq 1 \rightarrow yx \geq 1$. Thus, $\forall x, y \in \mathbb{Z} (x\mathcal{R}y \rightarrow y\mathcal{R}x)$, i.e. \mathcal{R} is symmetric.

Exists $x = 1, y = 2 \in \mathbb{Z}$ s.t. $x\mathcal{R}y \wedge y\mathcal{R}x$ but $x \neq y$. Thus, \mathcal{R} is not antisymmetric.

Exists $x = 0.5, y = 4, z = 0.3 \in \mathbb{Z}$ s.t. $x\mathcal{R}y \wedge y\mathcal{R}z$ but $\neg x\mathcal{R}z$. Thus, \mathcal{R} is not transitive.

In conclusion, \mathcal{R} is symmetric but is not reflexive or antisymmetric or transitive.

f) Solution:

For any $x \in \mathbb{Z}$, x and x are for sure both negative or both nonnegative. Thus, $\forall x \in \mathbb{Z} (x\mathcal{R}x)$, i.e. \mathcal{R} is reflexive.

For any $x, y \in \mathbb{Z}$, that x and y are both negative or both nonnegative implies y and x are both negative or both nonnegative. Thus, $\forall x, y \in \mathbb{Z} (x\mathcal{R}y \rightarrow y\mathcal{R}x)$, i.e. \mathcal{R} is symmetric.

Exists $\forall x = 1, y = 2 \in \mathbb{Z}$ s.t. $x\mathcal{R}y \wedge y\mathcal{R}x$ but $x \neq y$. Thus, \mathcal{R} is not antisymmetric.

For any $x, y, z \in \mathbb{R}$, that x and y are both negative or both nonnegative and that y and z are both negative or both nonnegative implies x and z are both negative or both nonnegative.

Thus, $\forall x, y, z \in \mathbb{Z} (x\mathcal{R}y \wedge y\mathcal{R}z \rightarrow x\mathcal{R}z)$, i.e. \mathcal{R} is transitive.

In conclusion, \mathcal{R} is reflexive, symmetric and transitive but is not antisymmetric.

3. Solution:

a) $\mathcal{R}_1 \circ \mathcal{R}_1 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$.

For any $(a, b) \in \mathcal{R}_1 \circ \mathcal{R}_1$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_1, (c, b) \in \mathcal{R}_1$ i.e. $(a > c) \wedge (c > b)$, i.e. $a > b$.

Thus, $\mathcal{R}_1 \circ \mathcal{R}_1 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$.

b) $\mathcal{R}_1 \circ \mathcal{R}_2 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$

For any $(a, b) \in \mathcal{R}_1 \circ \mathcal{R}_2$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_2, (c, b) \in \mathcal{R}_1$ i.e. $(a \geq c) \wedge (c > b)$, i.e. $a > b$.

Thus, $\mathcal{R}_1 \circ \mathcal{R}_2 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$.

c) $\mathcal{R}_1 \circ \mathcal{R}_3 = \mathbb{R}^2$

For any $(a, b) \in \mathcal{R}_1 \circ \mathcal{R}_3$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_3, (c, b) \in \mathcal{R}_1$ i.e. $(a < c) \wedge (c > b)$.

Thus, $\mathcal{R}_1 \circ \mathcal{R}_3 = \mathbb{R}^2$.

e) $\mathcal{R}_1 \circ \mathcal{R}_5 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$

For any $(a, b) \in \mathcal{R}_1 \circ \mathcal{R}_5$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_5, (c, b) \in \mathcal{R}_1$ i.e. $(a = c) \wedge (c > b)$, i.e. $a > b$.

Thus, $\mathcal{R}_1 \circ \mathcal{R}_5 = \{(a, b) \in \mathbb{R}^2 \mid a > b\}$.

f) $\mathcal{R}_1 \circ \mathcal{R}_6 = \mathbb{R}^2$

For any $(a, b) \in \mathcal{R}_1 \circ \mathcal{R}_6$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_6, (c, b) \in \mathcal{R}_1$ i.e. $(a \neq c) \wedge (c > b)$.

Thus, $\mathcal{R}_1 \circ \mathcal{R}_6 = \mathbb{R}^2$.

g) $\mathcal{R}_2 \circ \mathcal{R}_3 = \mathbb{R}^2$

For any $(a, b) \in \mathcal{R}_2 \circ \mathcal{R}_3$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_3, (c, b) \in \mathcal{R}_2$ i.e. $(a < c) \wedge (c \geq b)$.

Thus, $\mathcal{R}_2 \circ \mathcal{R}_3 = \mathbb{R}^2$.

h) $\mathcal{R}_3 \circ \mathcal{R}_3 = \{(a, b) \in \mathbb{R}^2 \mid a < b\}$

For any $(a, b) \in \mathcal{R}_3 \circ \mathcal{R}_3$, $\exists c \in \mathbb{R} (a, c) \in \mathcal{R}_3, (c, b) \in \mathcal{R}_3$ i.e. $(a < c) \wedge (c < b)$, i.e. $a < b$.

Thus, $\mathcal{R}_3 \circ \mathcal{R}_3 = \{(a, b) \in \mathbb{R}^2 \mid a < b\}$.

4. Proof:

To prove the composition operator is associative, we only need to prove that

$$\forall x \forall y ((x, y) \in (\mathcal{R}_3 \circ \mathcal{R}_2) \circ \mathcal{R}_1 \leftrightarrow (x, y) \in \mathcal{R}_3 \circ (\mathcal{R}_2 \circ \mathcal{R}_1)).$$

$$(a, b) \in (\mathcal{R}_3 \circ \mathcal{R}_2) \circ \mathcal{R}_1 \text{ iff. } \exists c ((a, c) \in \mathcal{R}_1 \wedge (c, b) \in \mathcal{R}_3 \circ \mathcal{R}_2)$$

$$\text{iff. } \exists c \exists d ((a, c) \in \mathcal{R}_1 \wedge (c, d) \in \mathcal{R}_2 \wedge (d, b) \in \mathcal{R}_3)$$

$$\text{iff. } \exists d ((a, d) \in \mathcal{R}_2 \circ \mathcal{R}_1 \wedge (d, b) \in \mathcal{R}_3)$$

$$\text{iff. } (a, b) \in \mathcal{R}_3 \circ (\mathcal{R}_2 \circ \mathcal{R}_1)$$

Thus, the composition operator is associative over relations.

QED

5. Proof:

To prove that *a relation \mathcal{R} on a set \mathbb{A} is transitive iff. $\mathcal{R} \circ \mathcal{R} \subseteq \mathcal{R}$* , we just need to prove that *a relation \mathcal{R} on a set \mathbb{A} is not transitive iff. $\mathcal{R} \circ \mathcal{R} \not\subseteq \mathcal{R}$* .

a relation \mathcal{R} on a set \mathbb{A} is not transitive

$$\text{iff. } \exists x, y, z \in \mathbb{A} \neg(x\mathcal{R}y \wedge y\mathcal{R}z \rightarrow x\mathcal{R}z)$$

$$\text{iff. } \exists x, y, z \in \mathbb{A} \neg(\neg(x\mathcal{R}y \wedge y\mathcal{R}z) \vee x\mathcal{R}z)$$

$$\text{iff. } \exists x, y, z \in \mathbb{A} (x\mathcal{R}y \wedge y\mathcal{R}z) \wedge \neg x\mathcal{R}z$$

$$\text{iff. } \exists x, z \in \mathbb{A} (\exists y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (y, z) \in \mathcal{R} \wedge (x, z) \notin \mathcal{R}))$$

$$\text{iff. } \exists x, z \in \mathbb{A} ((x, z) \in \mathcal{R} \circ \mathcal{R} \wedge (x, z) \notin \mathcal{R})$$

$$\text{iff. } \mathcal{R} \circ \mathcal{R} \not\subseteq \mathcal{R}$$

QED

6. Proof:

Since a relation \mathcal{R} on a set \mathbb{A} is antisymmetric iff. $\forall x, y \in \mathbb{A} (x\mathcal{R}y \wedge y\mathcal{R}x \rightarrow x = y)$ and $\mathcal{R} \cap \mathcal{R}^{-1} \subseteq I_{\mathbb{A}}$ iff. $\forall x, y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (x, y) \in \mathcal{R}^{-1} \rightarrow (x, y) \in I_{\mathbb{A}})$, we just need to prove $\forall x, y \in \mathbb{A} (x\mathcal{R}y \wedge y\mathcal{R}x \rightarrow x = y)$ iff. $\forall x, y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (x, y) \in \mathcal{R}^{-1} \rightarrow (x, y) \in I_{\mathbb{A}})$.

$$\begin{aligned} & \forall x, y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (x, y) \in \mathcal{R}^{-1} \rightarrow (x, y) \in I_{\mathbb{A}}) \\ \text{iff. } & \forall x, y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (x, y) \in \mathcal{R}^{-1} \rightarrow x = y) \\ \text{iff. } & \forall x, y \in \mathbb{A} ((x, y) \in \mathcal{R} \wedge (y, x) \in \mathcal{R} \rightarrow x = y) \\ \text{iff. } & \forall x, y \in \mathbb{A} (x\mathcal{R}y \wedge y\mathcal{R}x \rightarrow x = y) \end{aligned}$$

QED

7. Disproof:

$\mathcal{R}_1 \cup \mathcal{R}_2$ might not be an equivalence relation on \mathbb{A} .

For example:

$\mathbb{A} = \mathbb{N}$, $\mathcal{R}_1 = \{(a, b) \in \mathbb{N} \times \mathbb{N} \mid a \text{ and } b \text{ are congruent modulo } 7\}$,

$\mathcal{R}_2 = \{(a, b) \in \mathbb{N} \times \mathbb{N} \mid a \text{ and } b \text{ are congruent modulo } 2\}$.

Since $9\mathcal{R}_1 2$ and $2\mathcal{R}_2 4$, $(9, 2), (2, 4) \in \mathcal{R}_1 \cup \mathcal{R}_2$. However, $(9, 4) \notin \mathcal{R}_1 \cup \mathcal{R}_2$.

Therefore, $\mathcal{R}_1 \cup \mathcal{R}_2$ is not transitive, and is thereby not an equivalence relation on \mathbb{A} .

Thus, when \mathcal{R}_1 and \mathcal{R}_2 are equivalence relations on \mathbb{A} , $\mathcal{R}_1 \cup \mathcal{R}_2$ might not be an equivalence relation on \mathbb{A} . ***QED***

8. Proof:

Since \mathcal{R}_1 and \mathcal{R}_2 are equivalence relations,

$$\begin{aligned} (\forall a \in \mathbb{A} ((a, a) \in \mathcal{R}_1)) \wedge (\forall a \in \mathbb{A} ((a, a) \in \mathcal{R}_2)) & \Rightarrow \forall a \in \mathbb{A} ((a, a) \in \mathcal{R}_1 \wedge (a, a) \in \mathcal{R}_2) \\ & \Rightarrow \forall a \in \mathbb{A} ((a, a) \in \mathcal{R}_1 \cap \mathcal{R}_2). \end{aligned}$$

Thus, $\mathcal{R}_1 \cap \mathcal{R}_2$ is reflexive.

Since \mathcal{R}_1 and \mathcal{R}_2 are equivalence relations,

$$\begin{aligned} & (\forall a, b \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \rightarrow (b, a) \in \mathcal{R}_1)) \wedge (\forall a, b \in \mathbb{A} ((a, b) \in \mathcal{R}_2 \rightarrow (b, a) \in \mathcal{R}_2)) \\ \Rightarrow & \forall a, b \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \rightarrow (b, a) \in \mathcal{R}_1) \wedge ((a, b) \in \mathcal{R}_2 \rightarrow (b, a) \in \mathcal{R}_2) \\ \Rightarrow & \forall a, b \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \wedge (a, b) \in \mathcal{R}_2 \rightarrow ((b, a) \in \mathcal{R}_1 \wedge (b, a) \in \mathcal{R}_2)) \\ \Rightarrow & \forall a, b \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \cap \mathcal{R}_2 \rightarrow (b, a) \in \mathcal{R}_1 \cap \mathcal{R}_2). \end{aligned}$$

Thus, $\mathcal{R}_1 \cap \mathcal{R}_2$ is symmetric.

Since \mathcal{R}_1 and \mathcal{R}_2 are equivalence relations,

$$\begin{aligned} & (\forall a, b, c \in \mathbb{A} (a\mathcal{R}_1 b \wedge b\mathcal{R}_1 c \rightarrow a\mathcal{R}_1 c)) \wedge (\forall a, b, c \in \mathbb{A} (a\mathcal{R}_2 b \wedge b\mathcal{R}_2 c \rightarrow a\mathcal{R}_2 c)) \\ \Rightarrow & \forall a, b, c \in \mathbb{A} ((a\mathcal{R}_1 b \wedge b\mathcal{R}_1 c \rightarrow a\mathcal{R}_1 c) \wedge (a\mathcal{R}_2 b \wedge b\mathcal{R}_2 c \rightarrow a\mathcal{R}_2 c)) \\ \Rightarrow & \forall a, b, c \in \mathbb{A} ((a\mathcal{R}_1 b \wedge b\mathcal{R}_1 c \wedge a\mathcal{R}_2 b \wedge b\mathcal{R}_2 c) \rightarrow (a\mathcal{R}_1 c \wedge a\mathcal{R}_2 c)) \\ \Rightarrow & \forall a, b, c \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \wedge (a, b) \in \mathcal{R}_2 \wedge (b, c) \in \mathcal{R}_1 \wedge (b, c) \in \mathcal{R}_2) \rightarrow ((a, c) \in \mathcal{R}_1 \wedge (a, c) \in \mathcal{R}_2) \\ \Rightarrow & \forall a, b, c \in \mathbb{A} ((a, b) \in \mathcal{R}_1 \cap \mathcal{R}_2 \wedge (b, c) \in \mathcal{R}_1 \cap \mathcal{R}_2) \rightarrow ((a, c) \in \mathcal{R}_1 \cap \mathcal{R}_2) \end{aligned}$$

Thus, $\mathcal{R}_1 \cap \mathcal{R}_2$ is transitive.

Since $\mathcal{R}_1 \cap \mathcal{R}_2$ is reflexive, symmetric and transitive, $\mathcal{R}_1 \cap \mathcal{R}_2$ is an equivalence relation on \mathbb{A} .

QED