

J-NVM: Off-Heap Persistent Objects in Java

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new persistent medium (in-between SSD and DRAM)

Durable

resists reboots, power loss

High-density

smallest DIMM = 128 GB

Byte addressable

persistent memory abstraction

High-performance

low latency (seq. read/write ~ 160/90ns) high bandwidth (up to 8.10GB/s, 2nd gen)



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Direct byte-addressability of durable data

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1- Dramatic **throughput** and **latency** improvement for persistent data applications

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Direct byte-addressability of durable data Durable resists reboots, power loss **High-density** 1- Dramatic throughput and latency smallest DII improvement for persistent data applications Byte addressable persistent r 2- Simpler code bases with single data representation and no file I/Os High-performal low latency (seq. read/write ~ rou/90/15) high bandwidth (up to 8.10GB/s, 2nd gen)

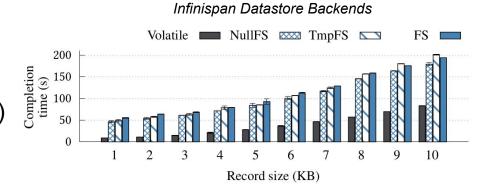
Why Java?

Many data stores & processing frameworks

Spark, Hadoop, Kafka, Flink, Cassandra, HBase, Elasticsearch, etc.

Lack of efficient interfaces

- FS/ext4-dax
 - almost as slow as tmpfs
 - dual representation (consistency)
 - cost of marshalling
- PCJ (JNI+PMDK)
 - slower than FS on YCSB benchmark



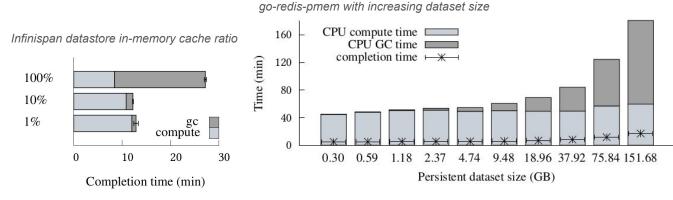
Varying record size in YCSB-F.

Problematic: Java-native NVMM interface

Prior works: internal design

= [Espresso, AutoPersist, go-pmem]

Managed Persistent Objects



Fixed dataset size - 80GB on heap for 100% cache

Varying cache ratio (YCSB-F)

Increasing dataset (YCSB-F, go-pmem)

Garbage collectors can not scale to large persistent datasets

Features

managed persistent objects

orthogonal persistence (pnew, @persistentRoot)

heavily-modified runtime

failure-atomic blocks

non-scalable

Prior works: internal design

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Managed Persistent Objects

go-redis-pmem with increasing dataset size

In [go-pmem]: "as the applications become complicated it becomes increasingly difficult to keep track of exactly which variables and pointers are in persistent memory".



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code instrumentation = made whole JVM 51% slower in [Autopersist]

10ry".

4.74 9.48 18.96 37.92 75.84 151.68

ersistent dataset size (GB)

Completion time (min)

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Overview

J-NVM = Off-Heap Persistent Objects

Challenges

single data representation

programming model

direct access to NVMM

durability abstraction

scalability (large persistent dataset)

Features

off-heap persistent objects

class-centric model (code generator + PDT library)

sun.misc.Unsafe

failure-atomic blocks + fine-grained

see evaluation

Outline

Introduction

- NVMM
- why Java?
- prior works
- overview

System Design

- programming model
 - persistent objects
 - code generator
- JPFA
- JPDT

Evaluation

- YCSB benchmark
- recovery

Conclusion

Overview

J-NVM = Off-Heap Persistent Objects

Key idea

each persistent object is decoupled into

- a persistent data structure: unmanaged, allocated off-heap (NVMM)
- a proxy: managed, allocated on-heap (DRAM)

Programming model - persistent objects

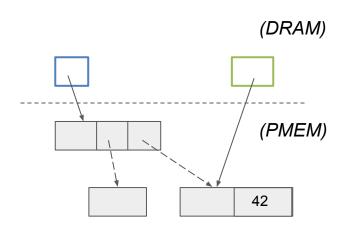
Persistent object is

- a persistent data structure
 - holds object fields
- a proxy
 - holds object methods
 - implement PObject interface
 - intermediate access to pers. data structure
 - instantiated lazily (low GC pressure)

Alive when reachable (from persistent root)

Class-centric model

- safe references thanks to the type system



```
Map root = JNVM.root();
Simple s = root.get("Simple");
s.setX(42);
```

Constructor

- allocate NVMM
- attach persistent data structure

Re-Constructor

- re-attach proxy
- re-build soft state via resurrect()

- explicit JNVM.free() to reclaim NVMM
- detach proxy
- ready to be GCed

```
Simple s = new Simple(42);

(DRAM)

-----
(PMEM)
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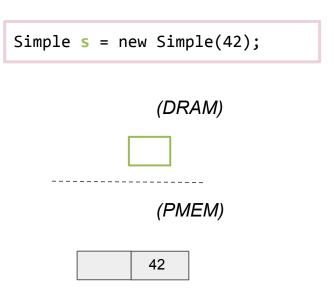
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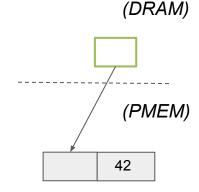
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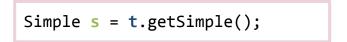
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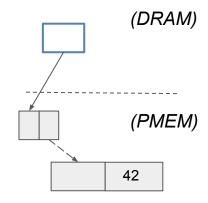
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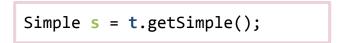
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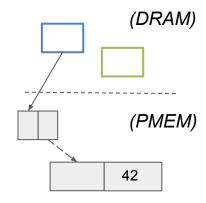
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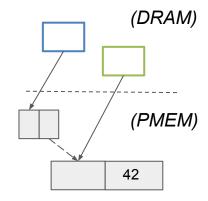
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Simple s = t.getSimple();



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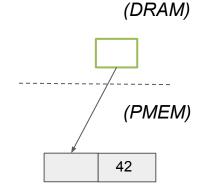
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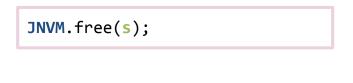
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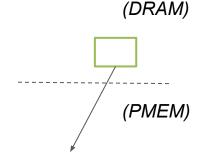
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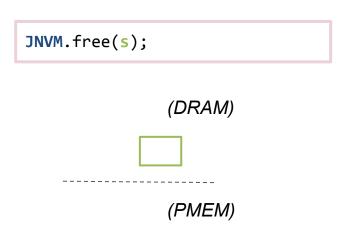
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```
JNVM.free(s);

(DRAM)

(PMEM)
```

Overview

J-NVM = Off-Heap Persistent Objects

Tooling

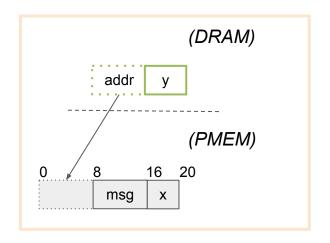
- built-in off-heap memory management for NVMM
- code generator: automatic decoupling for POJOs
- J-PFA: automatic failure-atomic code
- J-PDT: data types + collections for persistent memory
- low-level API (for experts)
- recovery-time GC

```
@Persistent(fa="non-private")
class Simple {
  PString msg;
  int x;
  transient int y;
  Simple(int x) {
    this.x = x;
    this.msg = new PString("Hello, NVMM!");
  void inc() { x++; }
```

- class-wide off-heap layout
- generate constructor, re-constructor
- replace (non-transient) field accesses
- wrap non-private methods

```
// transformed code
class Simple implements PObject {
  transient int y;
  long addr; // persistent data structure
  Simple(int x) {
    JNVM.faStart();
    this.addr = JNVM.alloc(getClass(), size());
    setX(x);
    setMsg(new PString("Hello, NVMM!"));
    JNVM.faEnd();
  Simple(long addr) {
    this.addr = addr;
    this.resurrect();
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```
// transformed code (continued)
  long addr; // the persistent data structure
  long size() { return 12; }
  PString getMsg() { return (PString)
          JNVM.readPObject(addr, 0); }
  void setMsg(PString v) {
          JNVM.writePObject(addr, 0, v); }
  int getX() {return JNVM.readInt(addr, 8);}
  void setX(int v) {JNVM.writeInt(addr, 8, v);}
```

- class-wide off-heap layout
- generate constructor, re-constructor
- replace (non-transient) field accesses
- wrap non-private methods
- generate or transform field accessors

J-PFA

Automatic crash-consistent update usage = JNVM.faStart() some code JNVM.faEnd()

Per-thread persistent redo-log (inspired by Romulus)

Log new, free and updates granularity = a block of PMEM

Do *not* log updates to "new" persistent objects (e.g. allocated within the FA-block)

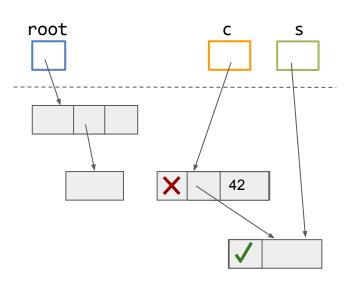
J-PDT + Low-level interface

J-PDT

drop-in replacement for (part of) JDK e.g., string, native array, map.

Low-level interface

- unsafe.{pwb,pfence, psync}
- NVMM block allocator
- recovery time GC (à la Makalu)
- validation = 1 bit in object header
 - makes atomic reclamation easier
 - allows deferring object liveness
 - interpreted on recovery to reclaim reachable invalid objects



```
Complex c = new Complex(s);
root.put("Complex", c);
c.validate();
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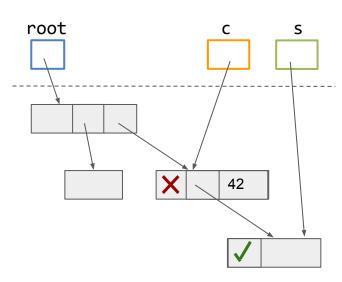
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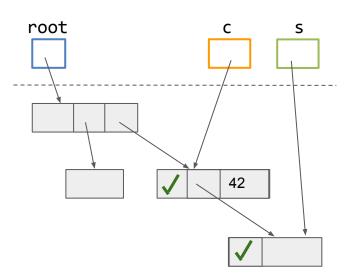
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- prior works
- overview

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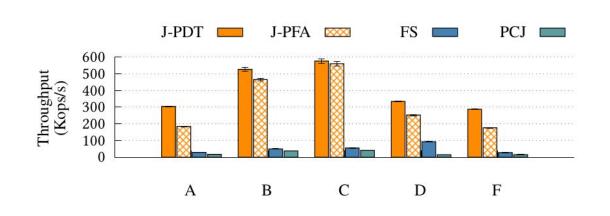
- programming model
 - persistent objects
 - code generator
- JPFA
- JPDT

Evaluation

- YCSB benchmark
- recovery

Conclusion

YCSB Benchmark



Durable backends for Infinispan:

- PCJ = HashMap from Persistent Collections Java (JNI + PMDK)
- FS: ext4-dax

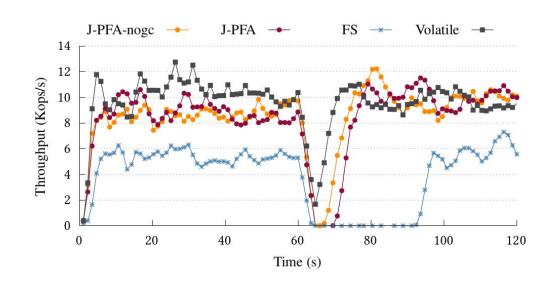
Hardware used:

4 Intel CLX 6230 HT 80-core 128GB DDR4, 4x128GB Optane (gen1)

Takeaways:

- J-NVM up to 10.5x (resp. 22.7x) than FS (resp. PCJ)
- no need for volatile cache

Recovery



TPC-B like benchmark 10M accounts (140 B each) client-server setting SIGKILL after 1 min

Takeaways:

- J-NVM is more than 5x faster to recover than FS
- no-need for graph traversal in some cases (e.g., only FA blocks)

Conclusion

J-NVM = off-heap persistent objects

Each persistent object is composed of

- a persistent data structure: unmanaged, allocated off-heap (NVMM)
- a proxy: managed, allocated on-heap (DRAM)

Pros:

- unique data representation (no data marshalling)
- recovery-time GC (not at runtime, does not scale)
- consistently faster than external designs (JNI, FS)

Cons:

- explicit free <u>but</u> common for durable data
- limited code re-use <u>but</u> safer programming model