

$$\begin{aligned}
& \int \pi ::= \\
& \int \tau ::= \\
& \pi | (\tau_1, \tau_2) | \tau \\
& 1, v_2) | \{v_1, \dots, v_k\} \\
& ?? \\
& 1, e_2) pair \\
& \{e_1, \dots, e_k\}_{k \geq 1} primitive sequence \\
& \{\} \tau empty sequence of type \tau \\
& 1 e_2 let - binding \\
& e built - in function call \\
& e_1 x e_0 general comprehension \\
& e_1 e_0 restricted comprehension \\
& f(e_1, \dots, e_k) user - defined function call \\
& d ::= \\
& \times \\
& function f(x_1 : \\
& \tau_1, \dots, x_k : \\
& \tau_k) : \\
& \tau = \\
& e user - defined function \\
& e_1 \\
& \times \\
& using \\
& \times \\
& mkseq \\
& \times \\
& seq \\
& ?? \\
& ?? \\
& \Gamma = [x_1 || - > \tau_1, \dots, x_i || - > \tau_i] \\
& \Sigma \\
& \Sigma = [f_1 || - > \tau_1 \times \dots \times \tau_k \tau, \dots, f_i || - > \tau'_1 \times \dots \times \tau'_m \tau'] \\
& \vdash_{\Sigma} \\
& \mathcal{E} : \\
& \pi) \vdash_{\Sigma} a : \pi((x) = \tau) \vdash_{\Sigma} x : \tau \\
& \vdash_{\Sigma} e_1 : \tau_1 \vdash_{\Sigma} e_2 : \tau_2 \vdash_{\Sigma} (e_1, e_2) : (\tau_1, \tau_2) \\
& \vdash_{\Sigma} e_1 : \tau \dots \vdash_{\Sigma} e_k : \tau \vdash_{\Sigma} \{e_1, \dots, e_k\} : \{\tau\} \vdash_{\Sigma} \{\} \tau : \{\tau\} \\
& \vdash_{\Sigma} e_1 : \tau_1 [x || - > \tau_1] \vdash_{\Sigma} e_2 : \tau \vdash_{\Sigma} x e_1 e_2 : \tau \\
& \vdash_{\Sigma} e_1 : \tau_1 \dots \vdash_{\Sigma} e_k : \tau_k (\Sigma()) = \tau_1 \times \dots \times \tau_k \tau) \vdash_{\Sigma} e : \tau \\
& \vdash_{\Sigma} e_0 : \{\tau_0\} [x || - > \tau_0, j x_i || - > \tau_i] \vdash_{\Sigma} e_1 : \tau (j(x_i) = \tau_i \in \{\pi, (\pi_1, \pi_2)\}) \vdash_{\Sigma} e_1 x e_0 : \tau \\
& \vdash_{\Sigma} e_0 : [j x_i || - > \tau_i] \vdash_{\Sigma} e_1 : \tau (j(x_i) = \tau_i) \vdash_{\Sigma} e_1 e_0 : \tau \\
& \vdash_{\Sigma} e_1 : \tau_1 \dots \vdash_{\Sigma} e_k : \tau_k (\Sigma(f) = \tau_1 \times \dots \times \tau_k \tau) \vdash_{\Sigma} f(e_1, \dots, e_k) : \tau \\
& \int \\
& v_1 \tau_1 v_2 \tau_2 (v_1, v_2) (\tau_1, \tau_2) v_1 \tau \dots v_k \tau v \{\tau\} Typing rules of values \\
& e_1 \\
& x_1, \dots, x_j \\
& k \\
& k \geq \\
& 0 \\
& k \\
& k? \\
& \rho \\
& \rho = [x_1 || - > v_1, \dots, x_i || - > v_i] \\
& \Phi = [f || - > (v_1, \dots, v_k) v] \\
& \rho \vdash_{\Phi} \\
& ev \rho \vdash_{\Phi} aa(\rho(x) = v) \rho \vdash_{\Phi} x v \rho \vdash_{\Phi} e_1 v_1 \rho \vdash_{\Phi} e_2 v_2 \rho \vdash_{\Phi} (e_1, e_2)(v_1, v_2) \\
& \rho \vdash_{\Phi} e_1 v_1 \dots \rho \vdash_{\Phi} e_k v_k \rho \vdash_{\Phi} \{e_1, \dots, e_k\} \{v_1, \dots, v_k\} \rho \vdash_{\Phi} \{\} \tau \{\} \\
& \rho \vdash_{\Phi} e_1 v_1 \rho [x || - > v_1] \vdash_{\Phi} e_2 v \rho \vdash_{\Phi} e_1 x e_2 v \\
& \rho \vdash_{\Phi} e_1 v_1 \dots \rho \vdash_{\Phi} e_k v_k ((\Phi()) = (v_1, \dots, v_k) v) \rho \vdash_{\Phi} ev \\
& \rho \vdash_{\Phi} \\
& e_0 \{v_1, \dots, v_k\} (\rho [x || - > v_i] \vdash_{\Phi} e_1 v'_i)_{i=1}^k \rho \vdash_{\Phi} e_1 x e_0 v' \\
& \rho \vdash_{\Phi} e_0 \rho \vdash_{\Phi} e_1 e_0 \{\} \rho \vdash_{\Phi} e_0 \rho \vdash_{\Phi} e_1 v_1 \rho \vdash_{\Phi} e_1 x \{v_1\} \\
& \rho \vdash_{\Phi} e_1 v_1 \dots \rho \vdash_{\Phi} e_k v_k ((\Phi(f) = (v_1, \dots, v_k) v) \rho \vdash_{\Phi} f ev Semantic of \\
& bound sequence \\
& e_0 \\
& k \\
& e_1 \\
& e_1 \\
& k \\
& k \\
& guard \\
& e_0 \\
& e_1 \\
& ?? \\
& ?? \\
& \times \\
& mqseq \\
& \times \\
& zip \\
& \times \\
& scan \\
& \times \\
& reduce \\
& \vdash \\
& \otimes \\
& ?? \\
& \oplus \times
\end{aligned}$$