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Camila Loiola Brito Maia, Thiago do Nascimento Ferreira, Fabrício Gomes de Freitas and Jerffeson Teixeira de Souza

Optimization in Software Engineering Group (GOES.UECE)

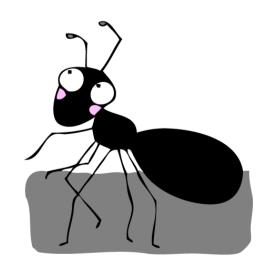
State University of Ceará, Brazil

The following topics will be discussed

- Introduction
- Problem definition
- Problem encoding
- The proposed ACO-based algorithm
- Conclusions and future works



INTRODUCTION



Test Case Prioritization

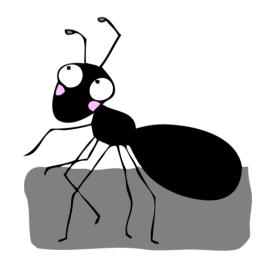
- The more important ones has to be tested before
 - What is more important for the client?
 - Are there any constraints for the prioritization?

ACO applied to TCP problem

- Few works on TCP
- Precedence of test cases not considered yet
- Application in other problems in our group
- A new approach considering the precedence of test cases



PROBLEM DEFINITION



Requirements

- Set R of requirements
 - N elements
- Attributes
 - importance:
 - volatility

Test Cases

- Set C of test cases
 - M elements
- Attributes
 - precedence;
 - coverage;
 - executionTime_j
 - $score_{j} = \sum ((importance_{i} * weight1) + (volatility_{i} * weight2)),$ $\forall requirement_{i} \in coverage_{j}$

Mathematical Formulation

$$Maximize \sum_{j=1}^{M} \left(\frac{score_j}{executionTime_j} * P_j \right)$$

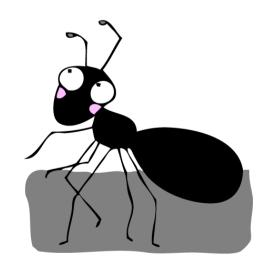
Subject to:

$$\forall \ t_{j1}, t_{j2} \in \mathcal{C}, \qquad \left(precedence_{t_{j2}} = \ t_{j1}\right) \rightarrow (q_{t_{j1}} < q_{t_{j2}})$$

- Where $P_j = M q_j + 1$
- q_j is the position of test case j in the ordered suite



PROBLEM ENCODING



Overview

- Directed graph G = (V,E)
- Each test case represents a vertice on the graph
 - V has M elements
- All vertices are connected with all others
- Pheromone equally distributed

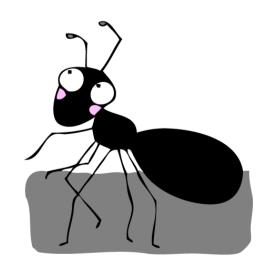
The Proposed Approach

- Each ant has the following information
 - nextNode
 - visitedNodes
 - allowedNodes
- Heuristic Function

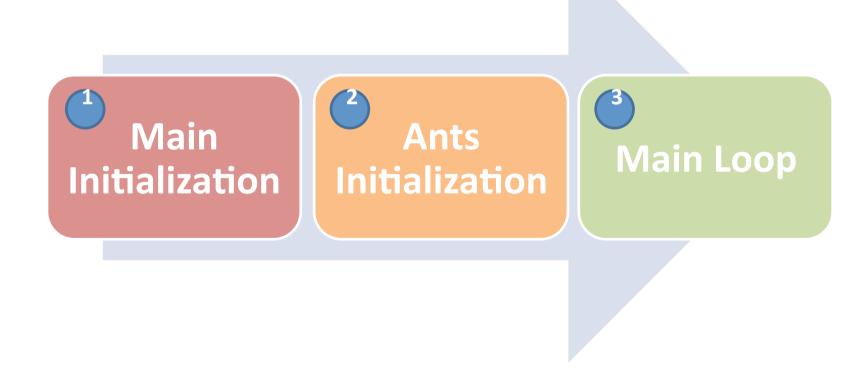
$$\frac{score_j}{executionTime_j}$$



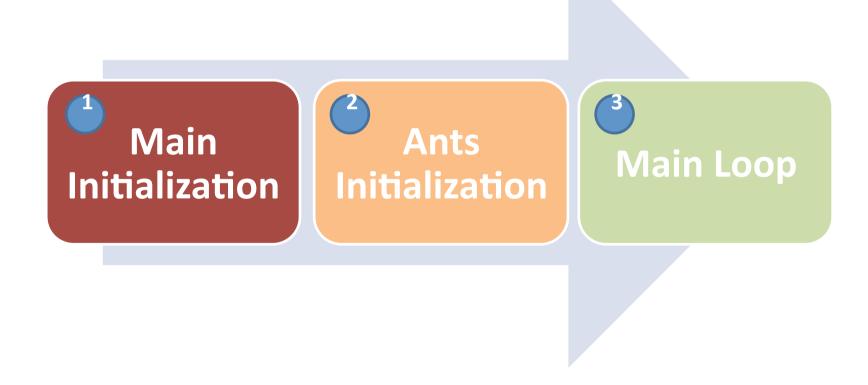
THE PROPOSED ACO-BASED ALGORITHM



Steps of the Algorithm



Steps of the Algorithm



ITERATION = 0

Read test case precedence information Generate directed graph Initialize pheromone

ITERATION = 0

Read test case precedence information

Generate directed graph

Initialize pheromone

ITERATION = 0

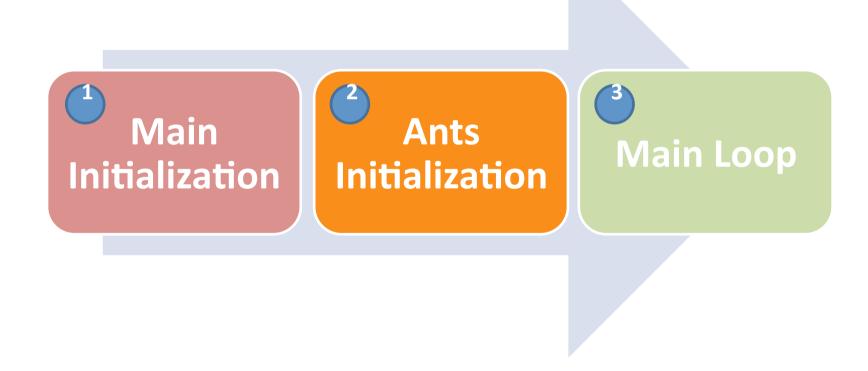
Read test case precedence information

Generate directed graph

Initialize pheromone

ITERATION = 0
Read test case precedence information
Generate directed graph
Initialize pheromone

Steps of the Algorithm



FOR ALL ants

FOR ALL vertices t_j of V, visited_j \leftarrow False ant_k.allowedNodes = UPDATE_ALLOWED_NODES() ant_k.nextNode = FIND_INITIAL_NODE() ant_k.visitedNode.add(nextNode)

FOR ALL ants

FOR ALL vertices t_i of V, visited_i \leftarrow False

 ant_k .allowedNodes = UPDATE_ALLOWED_NODES()

ant_k.nextNode = FIND_INITIAL_NODE()

ant_k.visitedNode.add(nextNode)

FOR ALL ants

FOR ALL vertices t_i of V, visited_i \leftarrow False

ant_k.allowedNodes = UPDATE_ALLOWED_NODES()

 $ant_k.nextNode = FIND_INITIAL_NODE()$

ant_k.visitedNode.add(nextNode)

FOR ALL ants

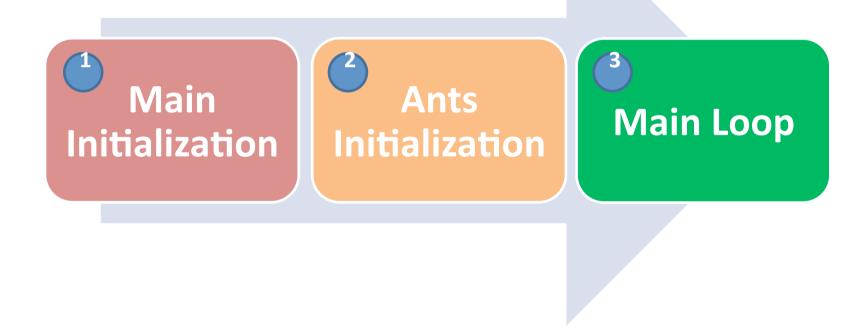
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FOR ALL ants

FOR ALL vertices t_j of V, visited_j \leftarrow False ant_k.allowedNodes = UPDATE_ALLOWED_NODES() ant_k.nextNode = FIND_INITIAL_NODE() ant_k.visitedNode.add(nextNode)

Steps of the Algorithm



FOR EACH antk

```
WHILE ant_k.allowedNodes.size() > 0 k = ant_k.actualNode ant_k.nextNode = ant_k.FIND_NEXT_NODE() j = ant_k.actualNode ant_k.visitedNodes.add(ant_k.nextNode) ant_k.allowedNodes = UPDATE_ALLOWED_NODES() Update\ pheromone\ in\ edge\ (k,j),with \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
```

FOR EACH ant_k

WHILE ant_k.allowedNodes.size() > 0

```
\begin{aligned} & \text{k} = \text{ant}_k.\text{actualNode} \\ & \text{ant}_k.\text{nextNode} = \text{ant}_k.\text{FIND\_NEXT\_NODE()} \\ & \text{j} = \text{ant}_k\text{actualNode} \\ & \text{ant}_k.\text{visitedNodes.add(ant}_k.\text{nextNode)} \\ & \text{ant}_k.\text{allowedNodes} = \text{UPDATE\_ALLOWED\_NODES()} \\ & \text{Update pheromone in edge (k,j),with} \\ & \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0 \end{aligned}
```

```
FOR EACH ant<sub>k</sub>

WHILE ant<sub>k</sub>.allowedNodes.size() > 0
```

 $k = ant_k.actualNode$

```
ant<sub>k</sub>.nextNode = ant<sub>k</sub>.FIND_NEXT_NODE()
j = ant<sub>k</sub>actualNode
ant<sub>k</sub>.visitedNodes.add(ant<sub>k</sub>.nextNode)
ant<sub>k</sub>.allowedNodes = UPDATE_ALLOWED_NODES()
Update pheromone in edge (k,j),with
```

$$\tau_{kj} = (1 - \varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0$$

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FOR EACH ant<sub>k</sub>.allowedNodes.size() > 0 k = \operatorname{ant_k}.\operatorname{actualNode} \operatorname{ant_k}.\operatorname{nextNode} = \operatorname{ant_k}.\operatorname{FIND_NEXT_NODE}() j = \operatorname{ant_k}\operatorname{actualNode} \operatorname{ant_k}.\operatorname{visitedNodes.add}(\operatorname{ant_k}.\operatorname{nextNode}) \operatorname{ant_k}.\operatorname{allowedNodes} = \operatorname{UPDATE\_ALLOWED\_NODES}() \operatorname{Update\ pheromone\ in\ edge\ }(k,j),\operatorname{with} \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
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FOR EACH ant<sub>k</sub>.allowedNodes.size() > 0 k = \operatorname{ant}_k.\operatorname{actualNode}  \operatorname{ant}_k.\operatorname{nextNode} = \operatorname{ant}_k.\operatorname{FIND\_NEXT\_NODE}() j = \operatorname{ant}_k\operatorname{actualNode}  \operatorname{ant}_k.\operatorname{visitedNodes.add}(\operatorname{ant}_k.\operatorname{nextNode})  \operatorname{ant}_k.\operatorname{allowedNodes} = \operatorname{UPDATE\_ALLOWED\_NODES}()  \operatorname{Update\ pheromone\ in\ edge\ }(k,j),\operatorname{with}  \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
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 \operatorname{ant}_k.\operatorname{allowedNodes} = \operatorname{UPDATE\_ALLOWED\_NODES}() 
 \operatorname{Update\ pheromone\ in\ edge\ }(k,j),\operatorname{with} 
 \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
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FOR EACH ant<sub>k</sub>.allowedNodes.size() > 0 k = \operatorname{ant}_k.\operatorname{actualNode} \operatorname{ant}_k.\operatorname{nextNode} = \operatorname{ant}_k.\operatorname{FIND\_NEXT\_NODE}() j = \operatorname{ant}_k\operatorname{actualNode} \operatorname{ant}_k.\operatorname{visitedNodes.add}(\operatorname{ant}_k.\operatorname{nextNode}) \operatorname{ant}_k.\operatorname{allowedNodes} = \operatorname{UPDATE\_ALLOWED\_NODES}() \operatorname{Update\ pheromone\ in\ edge\ }(k,j),\operatorname{with} \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
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```
FOR EACH ant_k ant_k allowedNodes.size() > 0  k = \mathrm{ant}_k.\mathrm{actualNode}  ant_k.\mathrm{nextNode} = \mathrm{ant}_k.\mathrm{FIND\_NEXT\_NODE}()  j = \mathrm{ant}_k\mathrm{actualNode}  ant_k.\mathrm{visitedNodes.add}(\mathrm{ant}_k.\mathrm{nextNode})  ant_k.\mathrm{allowedNodes} = \mathrm{UPDATE\_ALLOWED\_NODES}() Update pheromone in edge (k,j), with  \tau_{kj} = (1-\varphi) \cdot \tau_{kj} + \varphi \cdot \tau_0
```

currentSolution = EVALUATE_BEST_SOLUTION()

IF ((bestSolution.value() is null)

or (bestSolution.value() < currentSolution.value())) **THEN** bestSolution = currentSolution

ITERATION ++

currentSolution = EVALUATE_BEST_SOLUTION()

IF ((bestSolution.value() is null)

or (bestSolution.value() < currentSolution.value())) THEN

bestSolution = currentSolution

ITERATION ++

currentSolution = EVALUATE_BEST_SOLUTION()
IF ((bestSolution.value() is null)

or (bestSolution.value() < currentSolution.value())) **THEN**

bestSolution = currentSolution

ITERATION ++

The bestSolution variable is the result of the algorithm

ant_k.allowedNodes.clear()

FOR ALL vertices $t_i \in V$

IF (t_i respects the precedence constraint and is not yet in solution) ant_k.allowedNodes.add(t_i)

ant_k.allowedNodes.clear()

FOR ALL vertices $t_i \in V$

IF (t_i respects the precedence constraint and is not yet in solution) ant_k.allowedNodes.add(t_i)

ant_k.allowedNodes.clear()

FOR ALL vertices $t_i \in V$

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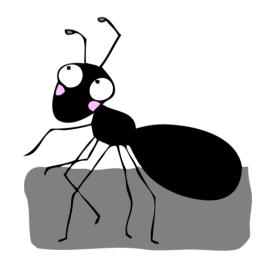
antk.FIND_NEXT_NODE()

Move ant k to a vertex t_i with probability p_{ij}^k or considering $max(\tau_{ij} \cdot w_j^\beta)$, considering only nodes in allowedNodes set

Return t_i



CONCLUSIONS AND FUTURE WORKS



Conclusions

- Currently
 - Implementation phase
 - Preliminary results not already to present
- Future
 - Evaluate and compare to other search-based algorithms



Camila Maia

camila.maia@gmail.com