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## 22BIO211: Intelligence of Biological Systems - 2

# INTRODUCTION TO DYNAMIC PROGRAMMING

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# The Change Problem

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- How can a cashier make change using the fewest number of coins?
- Different currencies have different possible coin values, or denominations.
- The heuristic used by cashiers all over the world to make change, which we call GreedyChange, iteratively selects the largest coin denomination possible.

*GreedyChange(money)*

```
change ← empty collection of coins
while money > 0
    coin ← largest denomination that is less than or equal to money
    add a coin with denomination coin to the collection of coins change
    money ← money - coin
return change
```

# The Change Problem

- Does GreedyChange always return the minimum possible number of coins?
- GreedyChange is suboptimal for some denominations!
- Example
  - Denomination - (120, 40, 30, 24, 20, 10, 5, 4, 1)
  - we want to change 48 units of currency
  - Greedychange selection-(40,5,1,1,1)
  - Minimum Selection - (24,24)
- Since GreedyChange is incorrect, we need to come up with a different approach.

# The Change Problem

- We can represent coins from  $d$  arbitrary denominations by an array of integers
  - $\text{Coins} = (\text{coin}_1, \dots, \text{coin}_d)$
  - where the values  $\text{coin}_i$  are given in decreasing order.
- Example : Consider the following denominations of coins
  - $\text{Coins} = (120, 40, 30, 24, 20, 10, 5, 4, 1)$
- We can represent another array of  $d$  positive integers ( $\text{change}_1, \dots, \text{change}_d$ ) such that
  - $\text{coin}_1 \cdot \text{change}_1 + \dots + \text{coin}_d \cdot \text{change}_d = \text{money}$
- Example:
  - Money = 48
  - Coins = (120, 40, 30, 24, 20, 10, 5, 4, 1)
  - Change = (0, 1, 0, 0, 0, 0, 1, 0, 3) corresponding to (40,5,1,1,1)
  - or Change = (0, 0, 0, 2, 0, 0, 0, 0, 0) corresponding to (24,24)

# Recursive Solution to the Change Problem

- A recursive solution to this problem involves breaking down the problem into smaller subproblems and solving them recursively.
- For each coin, you have two choices: include the coin or exclude the coin.
- If you include the coin, reduce the amount by the coin's value and keep the number of coins the same.
  - *Because the money is decremented and the coin can be repeated as a choice*
- If you exclude the coin, keep the amount the same but reduce the number of coins considered by one.
  - *No effect on money and the coin can not be considered as a choice*

# Recursive Solution to the Change Problem

- Let  $\text{Coins} = (5, 4, 1)$ , then for a money  $m$ ,

$$\text{MINNUMCOINS}(m) = \min \begin{cases} \text{MINNUMCOINS}(m - 5) + 1 \\ \text{MINNUMCOINS}(m - 4) + 1 \\ \text{MINNUMCOINS}(m - 1) + 1 \end{cases}$$

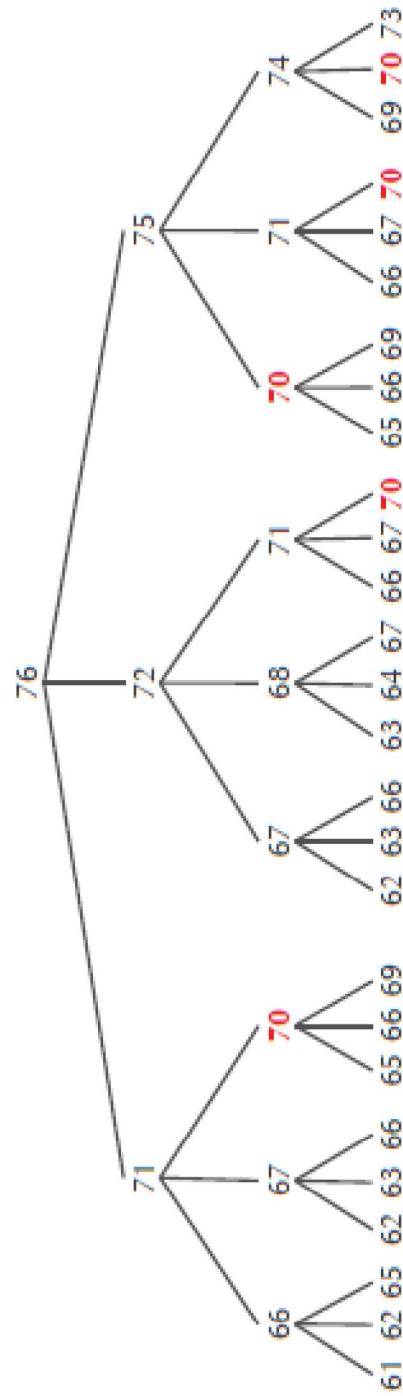
- We can form a recurrence relation from this
  - For the general denominations  $\text{Coins} = (coin_1, \dots, coin_d)$ ,  $\text{MinNumCoins}(money)$  is equal to the minimum of  $d$  numbers:

$$\text{MINNUMCOINS}(money) = \min \begin{cases} \text{MINNUMCOINS}(money - coin_1) + 1 \\ \vdots \\ \text{MINNUMCOINS}(money - coin_d) + 1 \end{cases}$$

# Recursive Solution to the Change Problem

- Recursive algorithms are impractical because it recalculates the optimal coin combination for a given value of money over and over again.

- A tree illustrating Recursive change for money = 76 with Coins = (5, 4, 1).



## Recursive Solution to the Change Problem

- `MinNumCoins(70)` gets computed six times, five of which are shown in the figure .
- When `MinNumCoins(70)` is computed for the sixth time (corresponding to the path  $76 \rightarrow 75 \rightarrow 74 \rightarrow 73 \rightarrow 72 \rightarrow 71 \rightarrow 70$ ), `RecursiveChange` has already been called hundreds of times.
- `MinNumCoins(30)` will be computed billions of times!

# Recursive Solution to the Change Problem

```
RecursiveChange(money, Coins)

if money = 0
    return 0

MinNumCoins ← ∞

for i ← 0 to |Coins| - 1

    if money ≥ coini
        NumCoins ← RecursiveChange(money - coini, Coins)

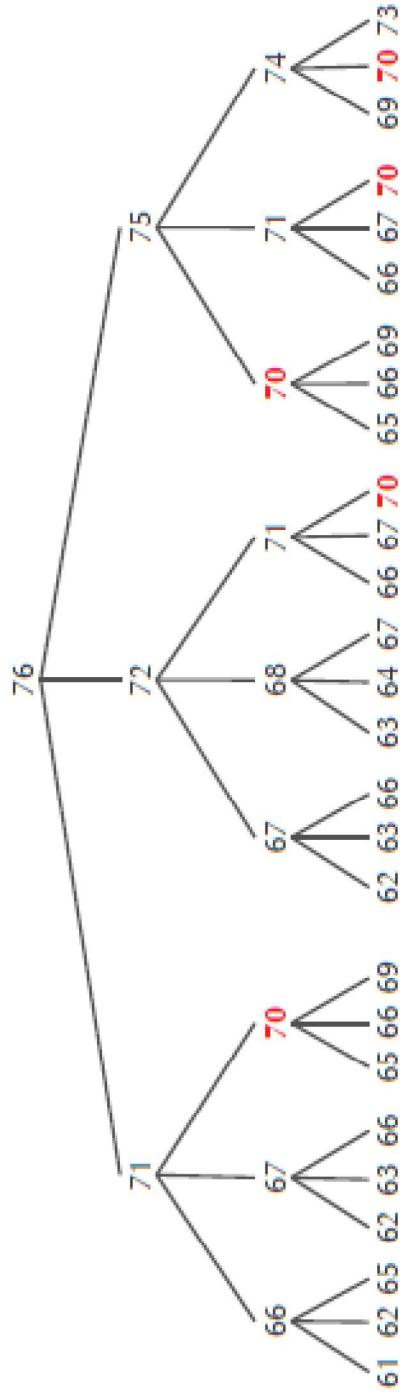
        if NumCoins + 1 < MinNumCoins
            MinNumCoins ← NumCoins + 1

return MinNumCoins
```

# Dynamic Programming solution to the Change Problem

- The change problem is the problem of finding the optimal distribution of change among a given amount of money and a set of coin denominations.
- Dynamic programming is an optimization technique used to solve complex problems by breaking them down into smaller, simpler subproblems.
- Dynamic programming can be applied to the change problem by using an array or table to store the solutions to subproblems.
- In the context of the change problem, dynamic programming is often used to determine the minimum number of coins required to make a given amount of change.

# Dynamic Programming solution to the Change Problem



- Instead of computing  $\text{MinNumCoins}(m)$  for every value of  $m$  from 76 downward toward  $m = 0$  via recursive calls,
  - we will invert our thinking and compute  $\text{MinNumCoins}(m)$  from  $m = 0$  upward toward 76,
  - storing all these values in an array so that we only need to compute  $\text{MinNumCoins}(m)$  once for each value of  $m$ .

# Dynamic Programming solution to the Change Problem

- $\text{MinNumCoins}(m)$  is still computed via the same recurrence relation:

$$\text{Coins} = (5, 4, 1)$$

$$\text{MINNUMCOINS}(m) = \min \begin{cases} \text{MINNUMCOINS}(m - 5) + 1 \\ \text{MINNUMCOINS}(m - 4) + 1 \\ \text{MINNUMCOINS}(m - 1) + 1 \end{cases}$$

$$\text{MINNUMCOINS}(6) = \min \begin{cases} \text{MINNUMCOINS}(1) + 1 = 2 \\ \text{MINNUMCOINS}(2) + 1 = 3 \\ \text{MINNUMCOINS}(5) + 1 = 2 \end{cases} = 2.$$
$$\text{MINNUMCOINS}(7) = \min \begin{cases} \text{MINNUMCOINS}(2) + 1 = 3 \\ \text{MINNUMCOINS}(3) + 1 = 4 \\ \text{MINNUMCOINS}(6) + 1 = 3 \end{cases} = 3.$$

$$\begin{array}{cccccccccc} m & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 \\ \text{MINNUMCOINS}(m) & 0 & 1 & 2 & 3 & 1 & 1 & 2 & 3 & 2 & 2 & 2 & 3 & 3 \end{array}$$

# Dynamic Programming solution to the Change Problem

- Notice that  $\text{MinNumCoins}(2)$  is used in the computation of both  $\text{MinNumCoins}(6)$  and  $\text{MinNumCoins}(7)$ 
  - *but instead of draining computational resources by having to compute this value from scratch both times, we can simply consult the pre-computed value in an array*

$m$	0	1	2	3	4	5	6	7	8	9	10	11	12
$\text{MINNUMCOINS}(m)$	0	1	2	3	1	1	2	3	2	2	2	3	3

- If  $\text{money} = 10^9$ , this algorithm requires a huge array of size  $10^9$ .

# Dynamic Programming solution to the Change Problem

```
DPChange(money, coins)

MinNumCoins(0) ← 0
for m ← 1 to money
    MinNumCoins(m) ← ∞
    for i ← 0 to |coins| - 1
        if m ≥ coini
            if MinNumCoins(m - coini) + 1 < MinNumCoins(m)
                MinNumCoins(m) ← MinNumCoins(m - coini) + 1
output MinNumCoins(money)
```

# Summary

- The Change Problem
- Recursive Solution to the Change Problem
- Dynamic Programming solution to the Change Problem