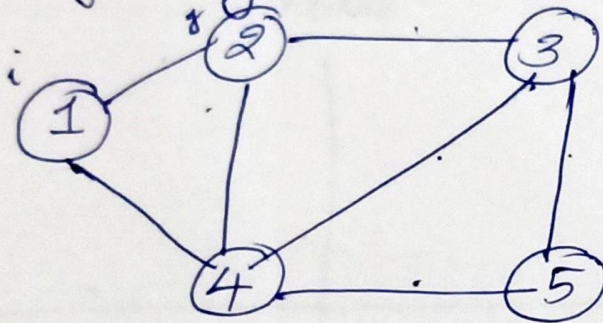


S₃ - AIE

Data Structures & Algorithms - 2

Graph Representation

- (1) Adjacency Matrix (3) Incidence matrix
(2) Adjacency list



Matrix $m \times n$
↓
rows → cols.

Adj. matrix

$n \times n$ where $n \rightarrow$ no. of vertices.

$i \rightarrow$	1	2	3	4	5
1	0	1	0	1	0
2	1	0	1	1	0
3	0	1	0	1	1
4	1	1	1	0	1
5	0	0	1	1	0

space complexity $\Rightarrow O(n^2)$
better \rightarrow if dense graph.

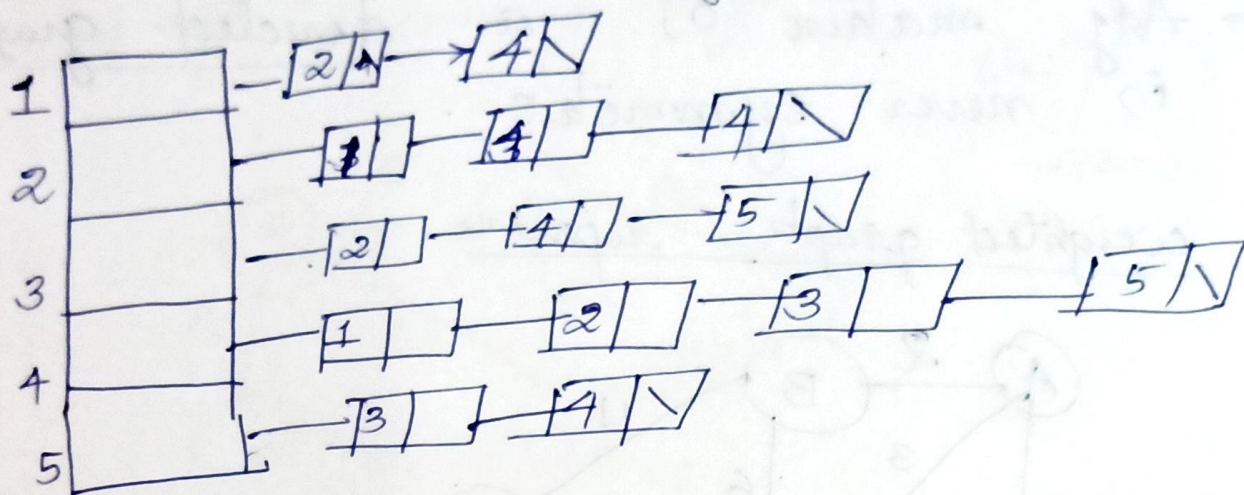
5x5

Defn

It is a matrix $A[n][n]$ where n is the no. of vertices.
 $a[i][j] = \begin{cases} 1, & \text{if } i \& j \text{ are adjacent} \\ 0, & \text{otherwise} \end{cases}$

Adjacency list

→ for each vertex one ^{LH} would be maintained. That ^{LH} will have contain the adj node to that node



space complexity $O(n + 2e)$

$n \rightarrow$ vertices

$e \rightarrow$ edges

eg: $(1 \rightarrow 2 \text{ \& } 2 \rightarrow 1)$

$2e \rightarrow$ edges we are taking 2 times.

dense \rightarrow adjacency matrix
 sparse \rightarrow adjacency list
 → each node is connected to every other nodes.
 → very few no. of edges.

Adjacency matrix

A graph $G = (V, E)$ where

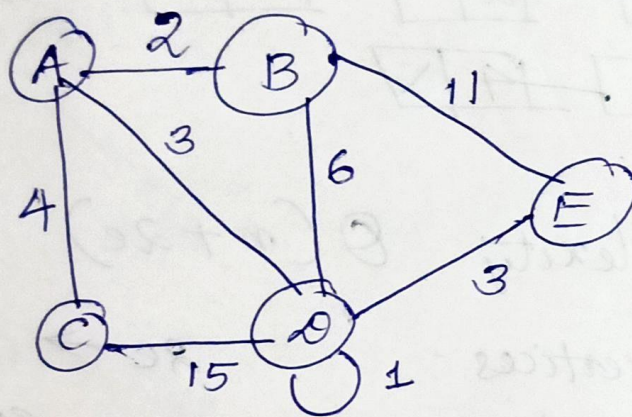
$V = \{0, 1, \dots, n-1\}$ can be represented using dimensional array of size $n \times n$.

→ Adj matrix of an undirected graph is always a symmetric matrix.

i.e., an edge (i, j) implies edge (j, i)

→ Adj matrix of a directed graph is never symmetric.

weighted graph represen.



	A	B	C	D	E
A	0	2	4	3	0
B	2	0	0	6	11
C	4	0	0	15	0
D	3	6	15	1	3
E	0	11	0	3	0

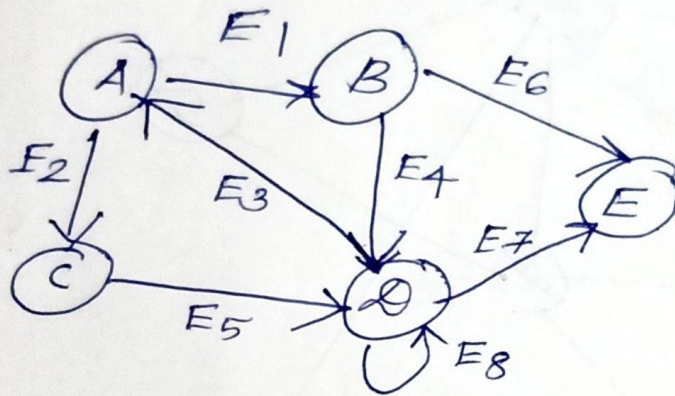
Incidence matrix

→ In this matrix, rows represent vertices, columns represent edges.

→ This matrix is filled with either 0 or 1 or -1.

- $0 \Rightarrow$ row edge is not connected to column.
 $1 \Rightarrow$ row edge is connected to out ~~going~~ edge of column.
 $-1 \Rightarrow$ incoming edge.

Outgoing +1
 incoming -1



	E_1	E_2	E_3	E_4	E_5	E_6	E_7	E_8
A	1	1	-1	0	0	0	0	0
B	-1	0	0	1	0	1	0	0
C	0	-1	0	0	1	0	0	0
D	0	0	1	-1	-1	0	1	1
E	0	0	0	0	0	-1	-1	0