



# Competitive Programming

From Problem 2 Solution in  $O(1)$

## Linear Algebra

## Gaussian Elimination - 2

**Mostafa Saad Ibrahim**

PhD Student @ Simon Fraser University



**Code**

# Gaussian Apps

- Solve system of linear Equations
- Solve system of linear Equations under Mod
- Compute Matrix Inverse
- Compute Determinant
- Some specific apps:
  - Patterns: Compute nth polynomial degree parameters
  - Solving Recurrence (useful if cyclic)

# Gaussian and Prime Mod

- Solve following system under prime = 5
  - $x + y = 5 \pmod{5}$
  - $3x + 6y = 1 \pmod{5}$
  - Solution:  $x = 3$  and  $y = 2$
- Gaussian operations are finally  $+$ ,  $-$ ,  $*$ ,  $/$ . So we can apply  $\%$  (use Mod inverse)
- Better flip all input matrix to +ve values  $\% p$
- If Prime = 2. No division operation.  
Subtraction of 2 equations under mod 2 is just **xor** of values. So we can have faster code.

**Code**

# Gaussian and Prime Mod

- As a note, solving
  - $x + y = 5 \pmod{5}$
  - $3x + 6y = 1 \pmod{5}$
- Same as solving
  - $x + y = 5 + 5k_1$
  - $3x + 6y = 1 + 5k_2$
- So don't be cheated with the extra 2 variables  $k_1, k_2$ ...think like take all  $\pmod{5}$

# Gaussian and Matrix Inverse

$$\begin{pmatrix} 1 & 2 & 1 \\ 2 & 9 & 2 \\ 1 & 7 & 2 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 2 & 1 \\ 0 & 5 & 0 \\ 0 & 5 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ -1 & 0 & 1 \end{pmatrix} \text{R}_2 = \text{R}_2 - 2 * \text{R}_1 \text{ AND } \text{R}_3 = \text{R}_3 - \text{R}_1$$

$$\begin{pmatrix} 1 & 2 & 1 \\ 0 & 1 & 0 \\ 0 & 5 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ -0.4 & 0.2 & 0 \\ -1 & 0 & 1 \end{pmatrix} \text{R}_2 = \text{R}_2 / 5$$

$$\begin{pmatrix} 1 & 2 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ -0.4 & 0.2 & 0 \\ 1 & -1 & 1 \end{pmatrix} \text{R}_3 = \text{R}_3 - 5 * \text{R}_2$$

$$\begin{pmatrix} 1 & 2 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 & -1 \\ -0.4 & 0.2 & 0 \\ 1 & -1 & 1 \end{pmatrix} \text{R}_1 = \text{R}_1 - \text{R}_3$$

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 0.8 & 0.6 & -1 \\ -0.4 & 0.2 & 0 \\ 1 & -1 & 1 \end{pmatrix} \text{R}_1 = \text{R}_1 - 2 * \text{R}_2$$

# Determinant

- Recall...direct computations is expensive

$$\begin{vmatrix} a & b \\ c & d \end{vmatrix} = ad - bc.$$

$$\begin{aligned} \begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix} &= a \begin{vmatrix} e & f \\ h & i \end{vmatrix} - b \begin{vmatrix} d & f \\ g & i \end{vmatrix} + c \begin{vmatrix} d & e \\ g & h \end{vmatrix} \\ &= a(ei - fh) - b(di - fg) + c(dh - eg) \\ &= aei + bfg + cdh - ceg - bdi - afh. \end{aligned}$$



# Determinant and Elementary Operations

2	4
10	6

$$A = 2 * 6 - 4 * 10 = -28...$$

What if we swapped 2 rows?

What if we divided by factor?

What if we added 2 rows?

10	6
2	4

Swapping:  $A = 4 * 10 - 2 * 6 = 28...$  Just multiply by -ve

1	2
10	6

Divide first row by 2

$$1 * 6 - 2 * 10 = -14 \Rightarrow \text{Original Divided by 2}$$

2	4
8	2

$$E2 \leftarrow E1$$

$$2 * 2 - 4 * 8 = -28 \Rightarrow \text{NO effect}$$

# Gaussian and Determinant

- Given that the elementary row operations are valid and have controllable effect, gaussian algorithm can be **trivially changed**
- Swapping two rows multiplies the determinant by  $-1$
- Multiplying a row by a nonzero scalar multiplies the determinant by the same scalar
- Adding to one row a scalar multiple of another does not change the determinant.

# Gaussian and Determinant

- Given the division operation, computations will be saved in doubles...and may be we face precision problem
- Sol1: be very careful when casting  $|A|$  value
- Sol2: Use fractions instead of doubles, if problem says fraction values won't overflow
- Sol3: Use Big Integers/Big Decimals
- Sol4: Compute % prime if result won't overflow

# Gaussian and Determinant

- **Compute % prime** if result won't overflow
- Assume we know final results fit in 32 bits.
- Find some problems  $p_1 p_2 \dots p_n$  such that
  - $p_1 * p_2 \dots p_n > \text{MAX\_ANSWER}$
  - $p_1 * p_2 \dots p_n < \text{Data type limit (e.g. long long)}$
  - E.g. primes: 257, 263, 269, 271
- Compute determinant %  $P_i$  for every prime
- Use Chinese remainder theorem to compute final answer

**Code**

# Gaussian and Patterns

- Assume we have some function:
- $f(n) = 5, 13, 24, 38, 55, 75, 98, \dots$ 
  - You are informed that it has form:
  - $f(n) = an^2 + bn + c$
  - Can you find  $f(n)$ ?
- We have 3 unknowns..can we get 3 equations?
- Evaluate  $f(1), f(2), f(3) \Rightarrow 3$  equations
- Solve them:
  - $a = 3/2, b = 7/2, c = 0$
  - $f(n) = (3n^2 + 7n)/2$  starting with  $n = 1$

# Gaussian and Recurrence

- Recall Fibonacci?
  - $F(n) = F(n-1) + F(n-2)$  and  $F(0) = F(1) = 1$
- If  $n = 4$ , we can think  $F(i)$  is variable / 5 vars
- Each  $F(i)$  is equation
  - E.g.  $F(4) = F(3) + F(2) \Rightarrow 0 = F(3) + F(2) - F(4)$
  - E.g. matrix row for  $F(4)$  can be:  $[-1 \ 1 \ 1 \ 0 \ 0 \mid 0]$
- Now, just solve the system and compute  $F(4)$
- As matrix is sparse, this is  $O(n^2)$
- This is used in Dynamic Programming when recurrence are cyclic and can't be coded

# Your Todo

- Implement  $O(n)$  solution for solving **Tridiagonal matrix** instead of Gaussian

$$\begin{bmatrix} b_1 & c_1 & & & 0 \\ a_2 & b_2 & c_2 & & \\ & a_3 & b_3 & \ddots & \\ & & \ddots & \ddots & c_{n-1} \\ 0 & & & a_n & b_n \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} d_1 \\ d_2 \\ d_3 \\ \vdots \\ d_n \end{bmatrix}.$$



# تم بحمد الله

علمكم الله ما ينفعكم

ونفعكم بما تعلمتم

وزادكم علماً

# Problems

- UVA 684, 10109, 10766, 10828, 11319, 11542, 11755
- [HDU 4418 - Time travel](#), its [solution](#)
- <http://codeforces.com/blog/entry/2536>
- SPOJ([NWERC04H](#))