



Competitive Programming

From Problem 2 Solution in $O(1)$

Computational Geometry

Line Sweep - Segments Intersection

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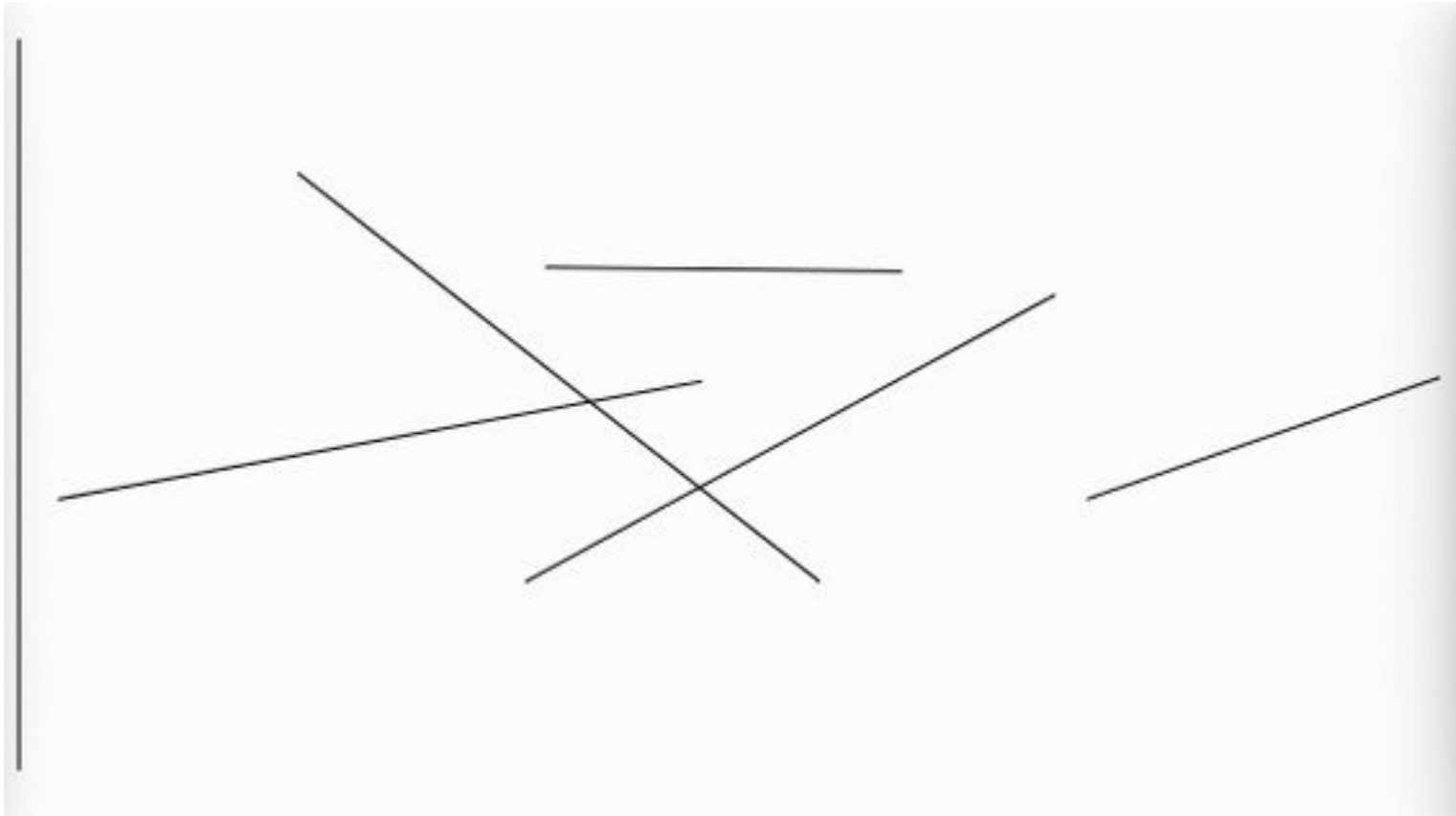
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Line Sweep

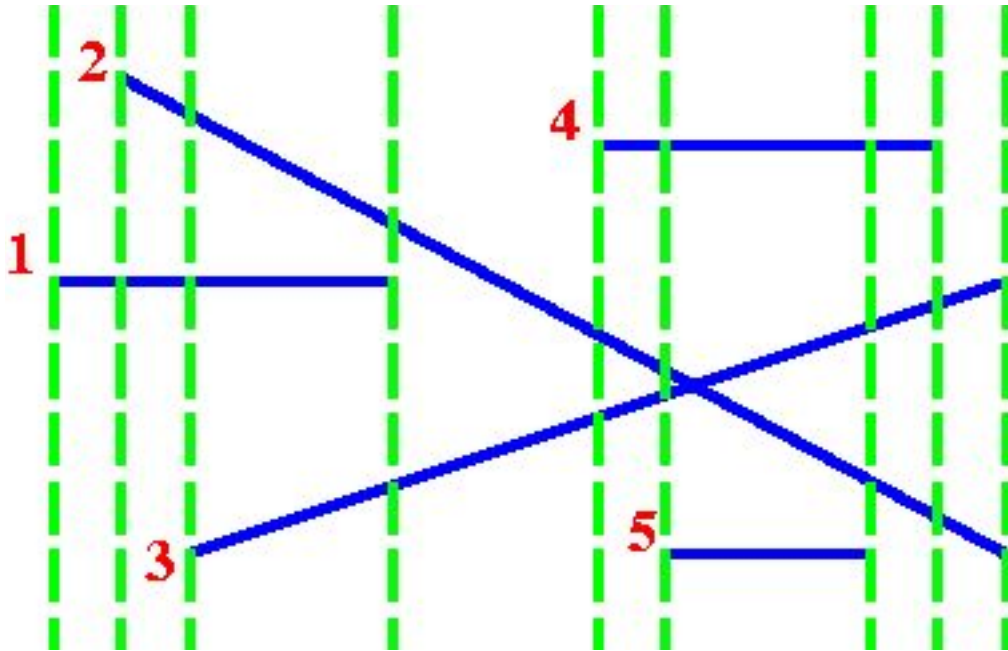
- Giving a plane of objects and some task
 - E.g. Set of segments to intersect
 - E.g. Set of rectangles to compute intersection
- One can consider every pair of objects!
 - However, every object (e.g. segment) interacts (e.g. intersect) with some surrounding ones **NOT ALL**
- Imagine a **vertical line** that is **swept** across the plane, specially **at discrete points** (events)
 - E.g. start/end of segment or a rectangle
 - We will use them to identify the **surrounding objects**

Line Sweep



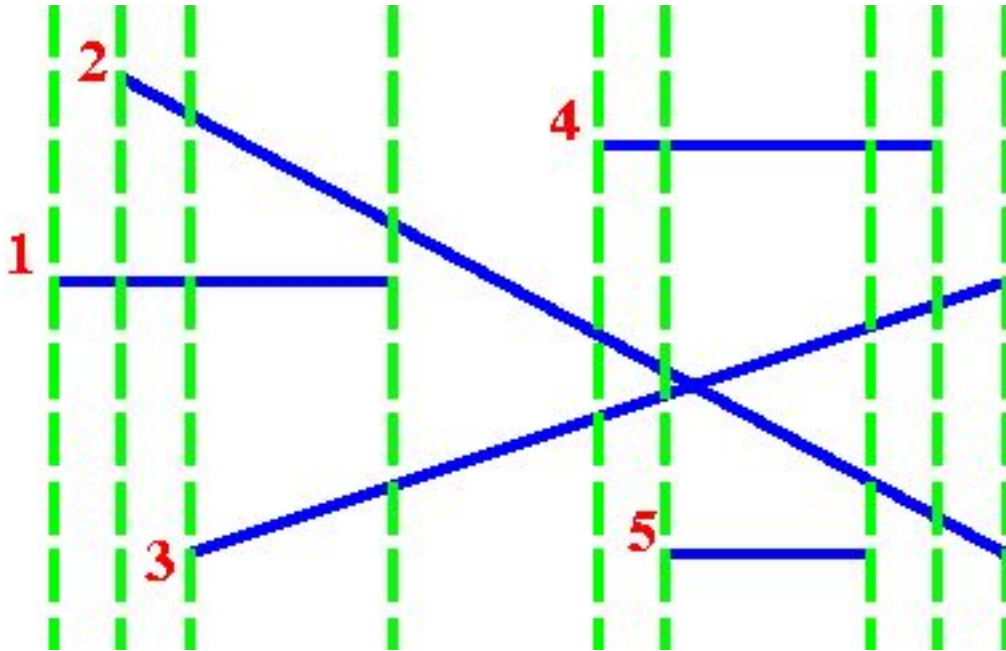
Src: <http://cglab.ca/~morin/teaching/2402/notes/planesweep.pdf>

Line Sweep (using vertical line)



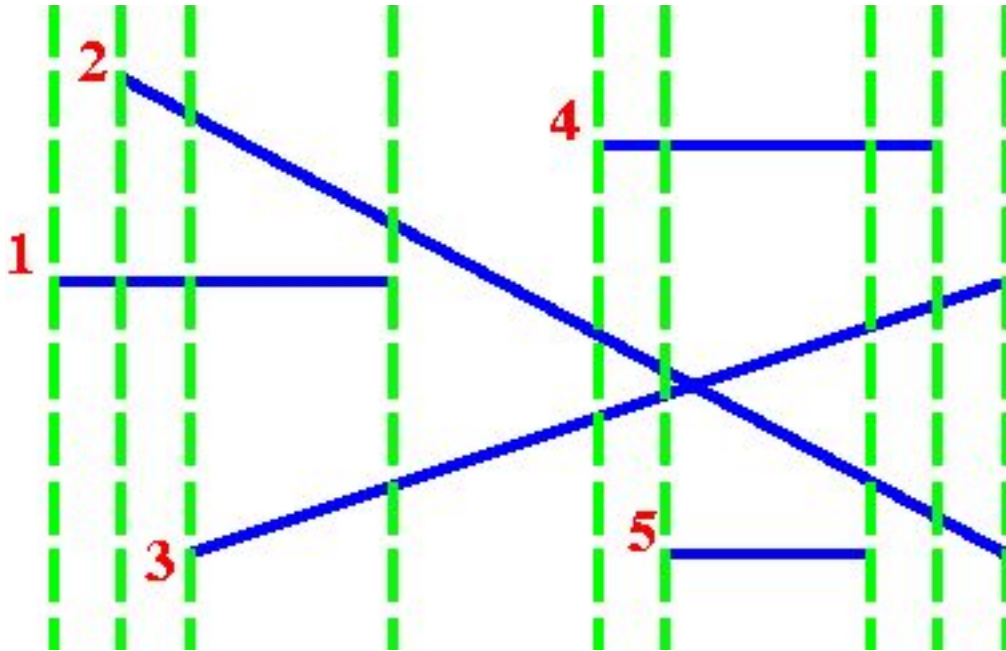
- Given N segments, find their **intersections**
- Naive solution is $O(n^2)$: intersect every pair
- 5 segments (blue)
- 10 end points
- 10 sweep **events** (vertical lines)
- Each endpoint is **event**
- Event: start / end of segment

Line Sweep (using vertical line)



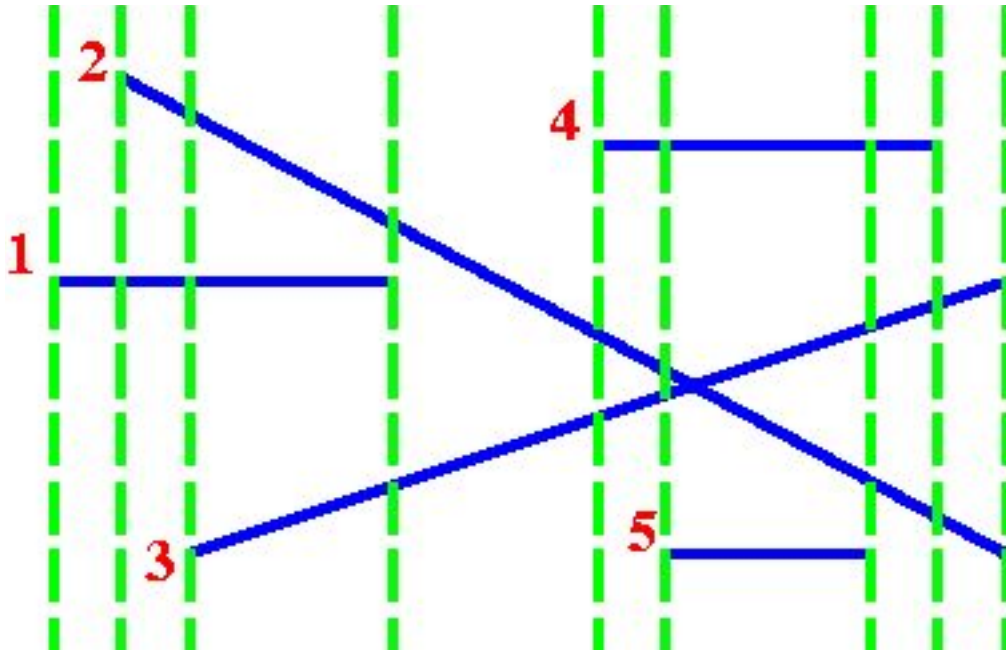
- Instead of considering ALL other segments
- We need to identify related subset only
- Let's call them the **active set**
- Whenever event occurs
 - It only consider the active set
 - it may use subset of it or all of it
 - If start event = add to active set
 - If end event = remove from active set
 - Processing of start event might be different from end event
- Let 4S (start of segment 4 event)
- Let 4E (end of segment 4 event)
- And so on for other segments

Line Sweep (using vertical line)



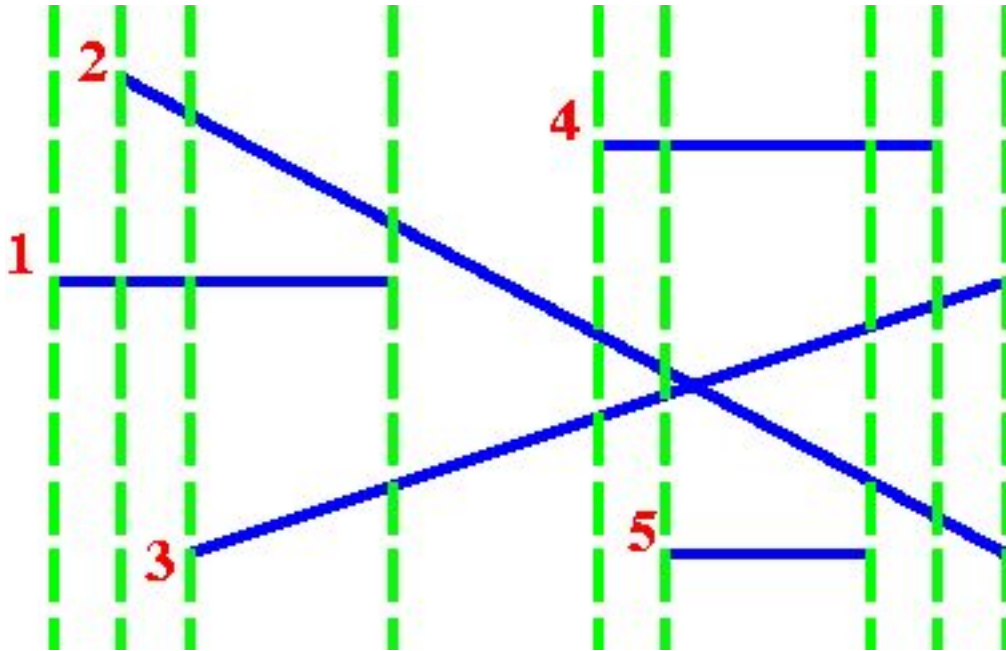
- Let 4S (start of segment 4 event)
- Let 4E (end of segment 4 event)
- And so on for other segments
- Given that we sweep vertically
- Then we can think of sorting segments based on X-axis
- **Events sorting:**
- 1S, 2S, 3S, 1E, 4S, 5S, 5E, 4E, 3E, 2E

Line Sweep (using vertical line)



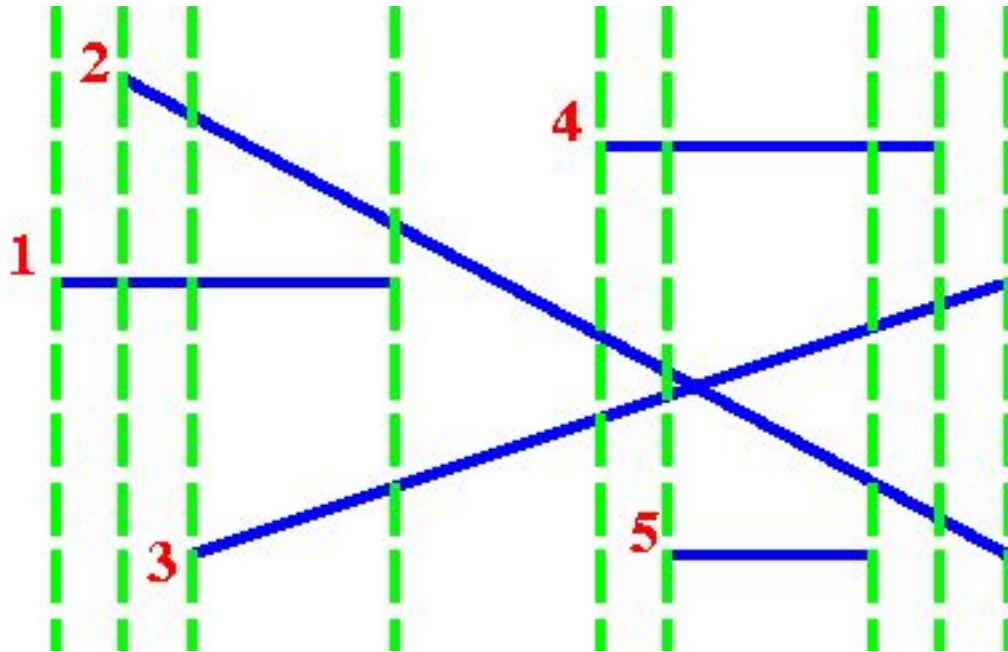
- 1S, 2S, 3S, 1E, 4S, 5S, 5E, **4E**, 3E, 2E
- What might be **active set** for event **4E**?
 - Ignore all event after it
 - Remove events that started/ended
- 1S, 2S, 3S, 1E, 4S, 5S, 5E, **4E**, 3E, 2E
- 1S, 2S, 3S, 1E, 4S, 5S, 5E, **4E**
- , 2S, 3S, , 4S, , , **4E**
- So only segments {2, 3} are active for 4E
- Look at 4E in the figure **vertically**

Line Sweep (using vertical line)



- 1S, 2S, 3S, 1E, 4S, 5S, 5E, 4E, 3E, 2E
- What might be **active set** for event 1S?
 - Segments $\{\}$ \Rightarrow Empty
- What might be **active set** for event 2S?
 - Segments $\{1\}$
- What might be **active set** for event 1E?
 - Segments $\{2, 3\} \Rightarrow$ Not same as 1S
- What might be **active set** for event 5S?
 - Segments $\{2, 3, 4\}$
- What might be **active set** for event 3E?
 - Segments $\{2\}$
- What might be **active set** for event 2E?
 - Segments $\{\}$ \Rightarrow Empty

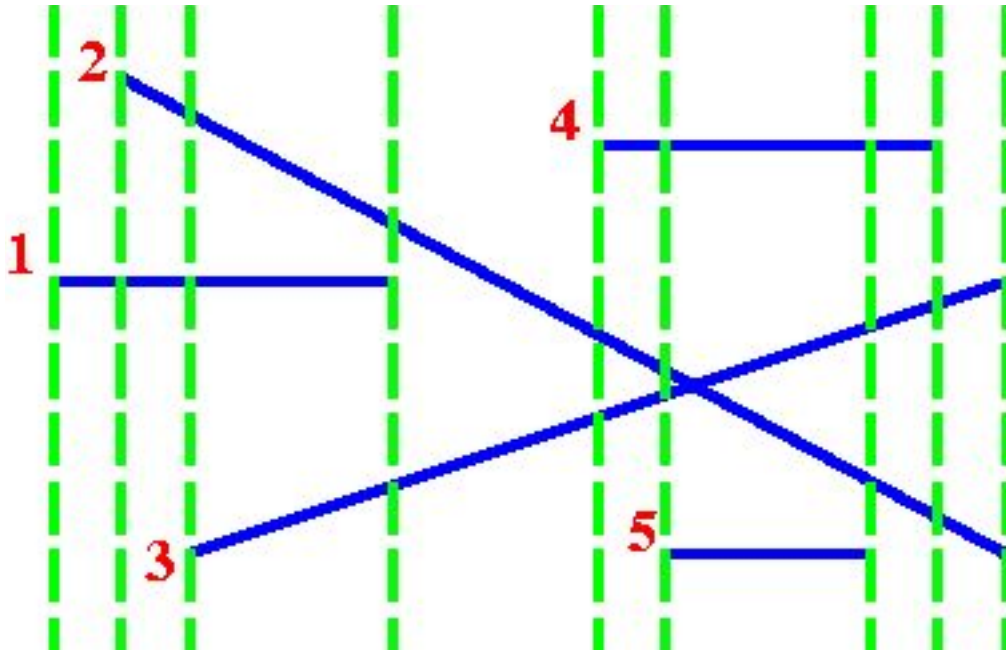
Line Sweep (using vertical line)



- **active set** for event **5S**? Segments {2, 3, 4}
- **active set** for event **5E**? Segments {2, 3, 4}
- **Questions** to think about?!
- Q1) Should we check intersections
 - at start event only?
 - end event? Or both?
- Q2) When we try to intersect: Should we try with **all** active set elements or a **subset**?
- Q3) What to check at each event? Is it **same** checkings for both endpoints?

Finding the right questions & answers is the major concern in Line Sweep Algorithms

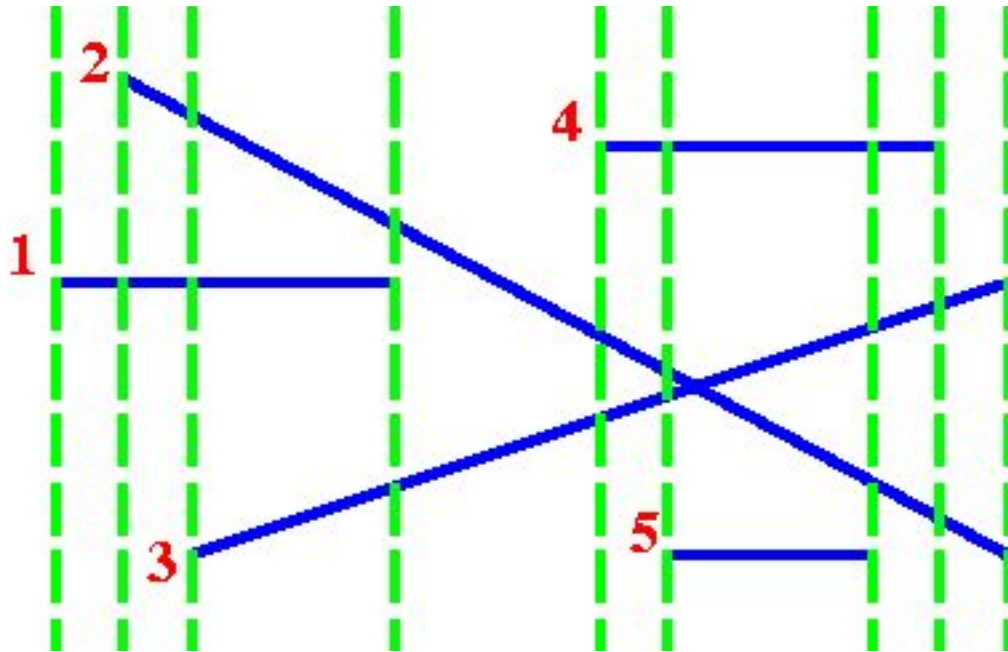
Line Sweep (using vertical line)



Bentley-Ottman Algorithm

- Let current segment is C
- From active set, let segment above C is A and segment below C is B
- **For Start Event**
 - Check intersect(Current, Above)
 - Check intersect(Current, Below)
 - Add C to active Set
- **For End Event**
 - Remove C from active Set
 - Check intersect(Above, Below)

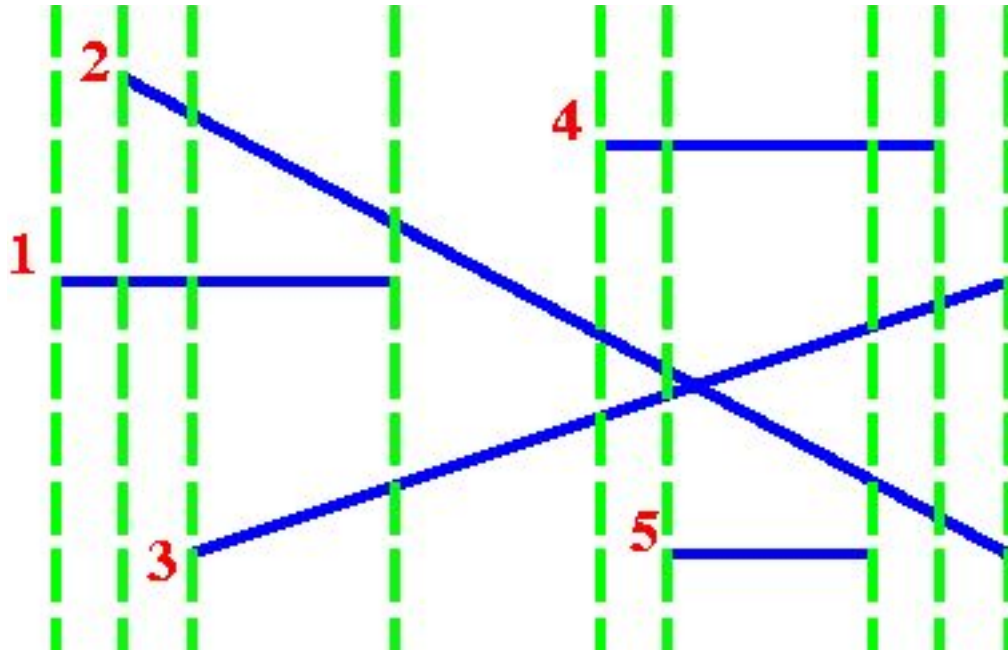
Line Sweep (using vertical line)



Events Processing

- Active List = {}
- Current Event **1S**
 - above = {}
 - below = {}
 - No intersection tests
- Add 1 to active list

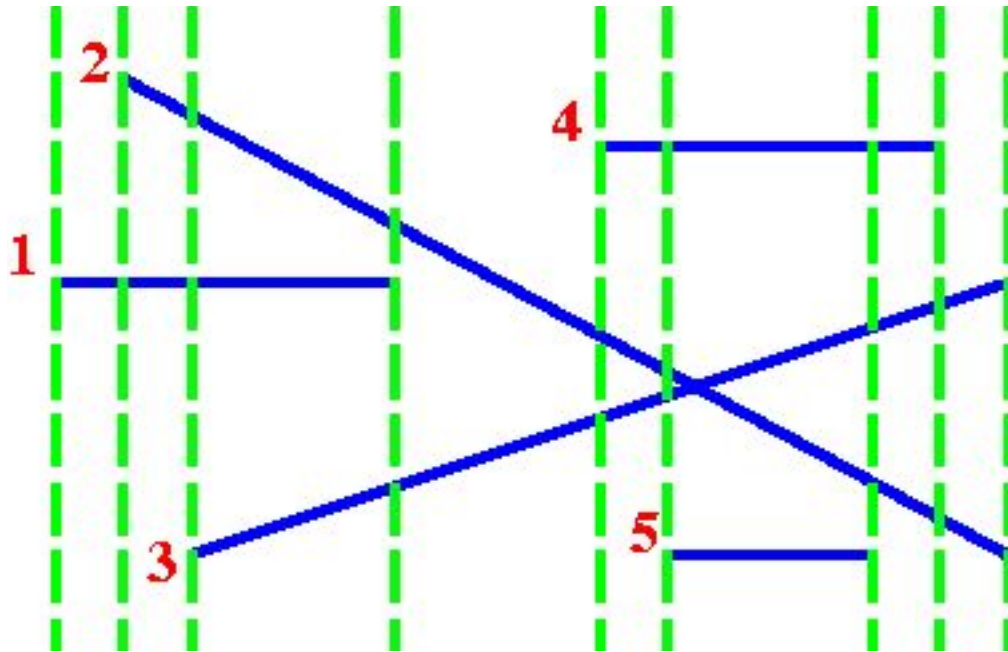
Line Sweep (using vertical line)



Events Processing

- Active List = {1}
- Current Event **2S**
 - above = {}
 - below = {1}
 - intersection(2, 1) = False
- Add 2 to active list

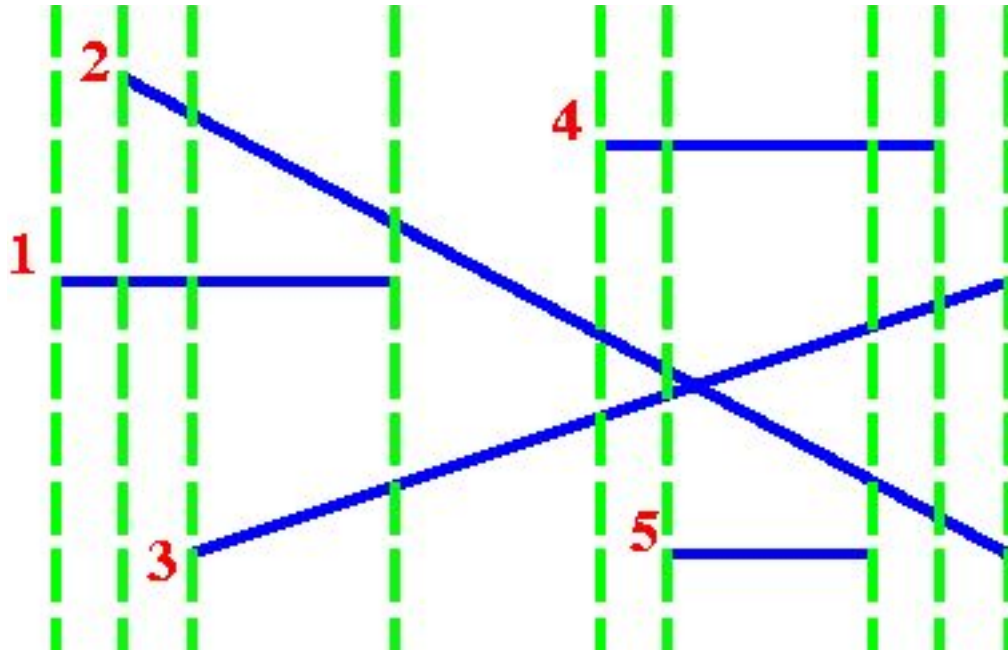
Line Sweep (using vertical line)



Events Processing

- Active List = {1, 2}
- Current Event **3S**
 - above = {1}
 - below = {}
 - intersection(3, 1) = False
- Add 3 to active list

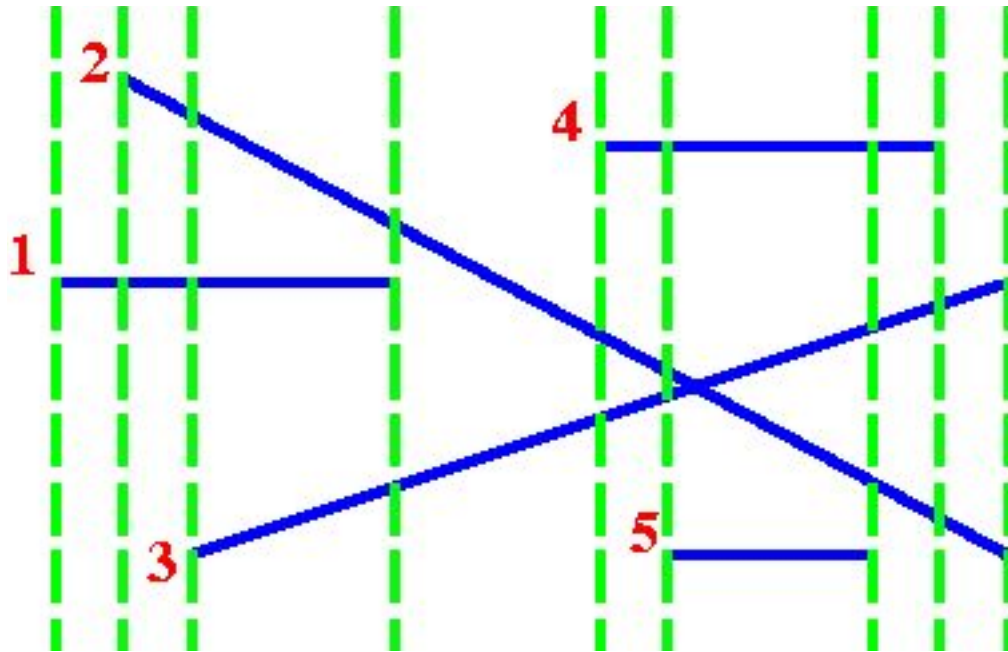
Line Sweep (using vertical line)



Events Processing

- Active List = {1, 2, 3}
- Current Event **1E**
 - above = {2}
 - below = {3}
 - intersections(2, 3) = True
- **Remove** 1 from active list

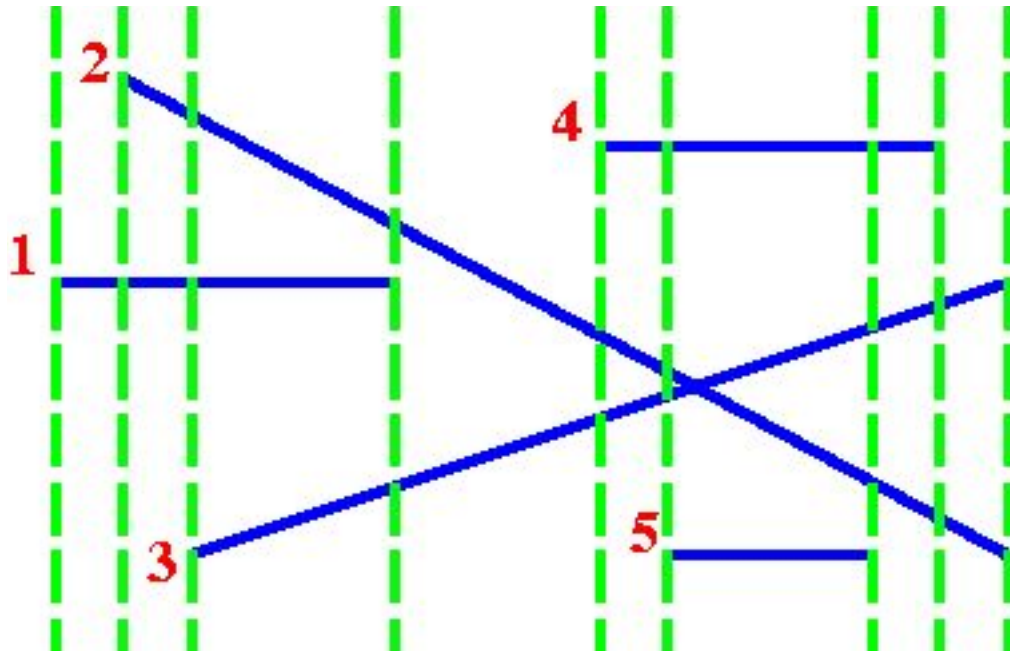
Line Sweep (using vertical line)



Events Processing

- Active List = {2, 3}
- Current Event **S4**
 - above = {2}
 - below = {3}
 - intersection(4, 2) = False
 - intersection(4, 3) = False
- Add 4 to active list

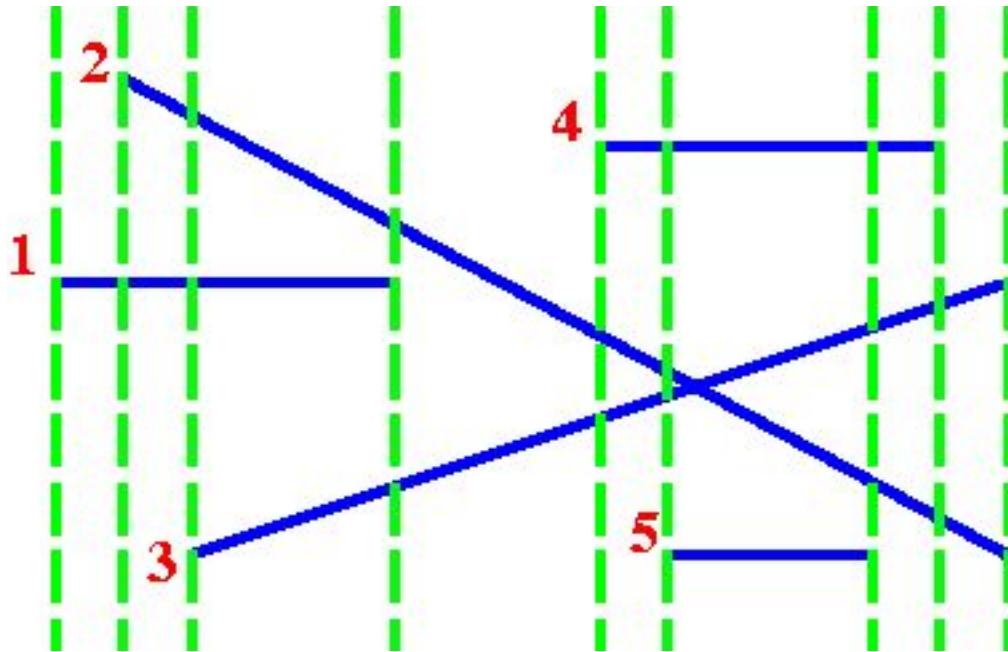
Line Sweep (using vertical line)



Events Processing

- Active List = {2, 3, 4}
- Current Event **S5**
 - above = {3}
 - below = {}
 - intersection(5, 3) = False
- Add 5 to active list

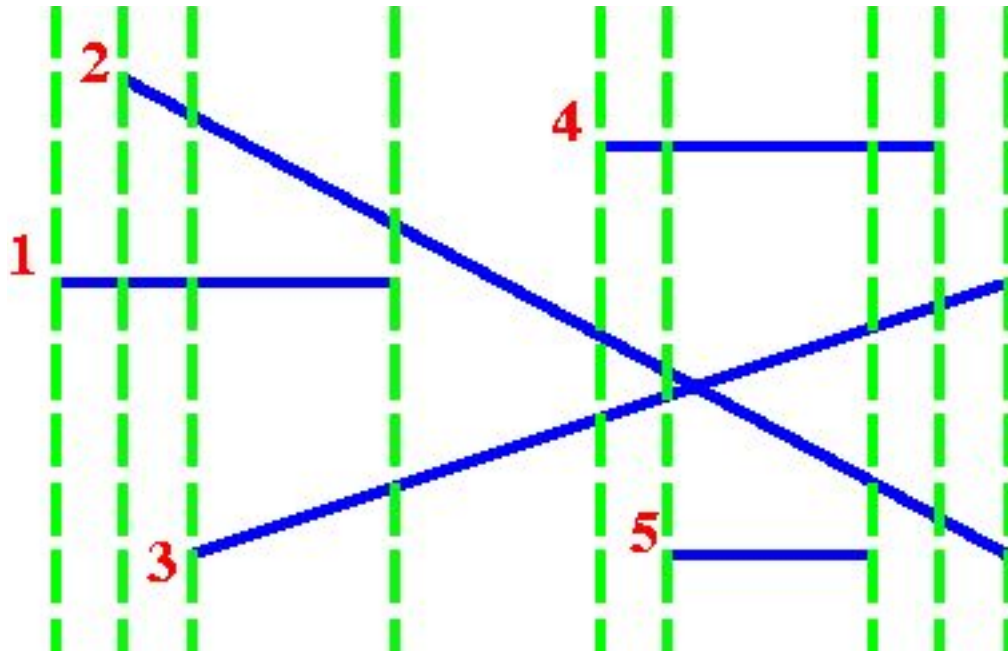
Line Sweep (using vertical line)



Events Processing

- Active List = {2, 3, 4, 5}
- Current Event **E5**
 - above = {3}
 - below = {}
 - No intersection tests
- Remove 5 from active list

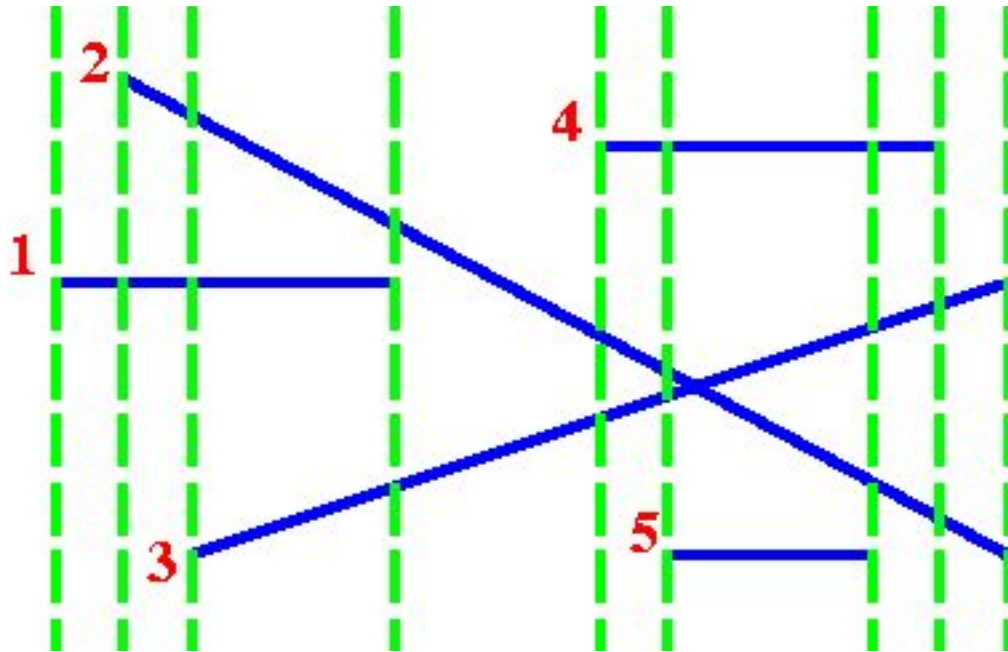
Line Sweep (using vertical line)



Events Processing

- Active List = {2, 3, 4}
- Current Event **E4**
 - above = {2}
 - below = {3}
 - intersections(2, 3) = YES
- Remove 4 from active list

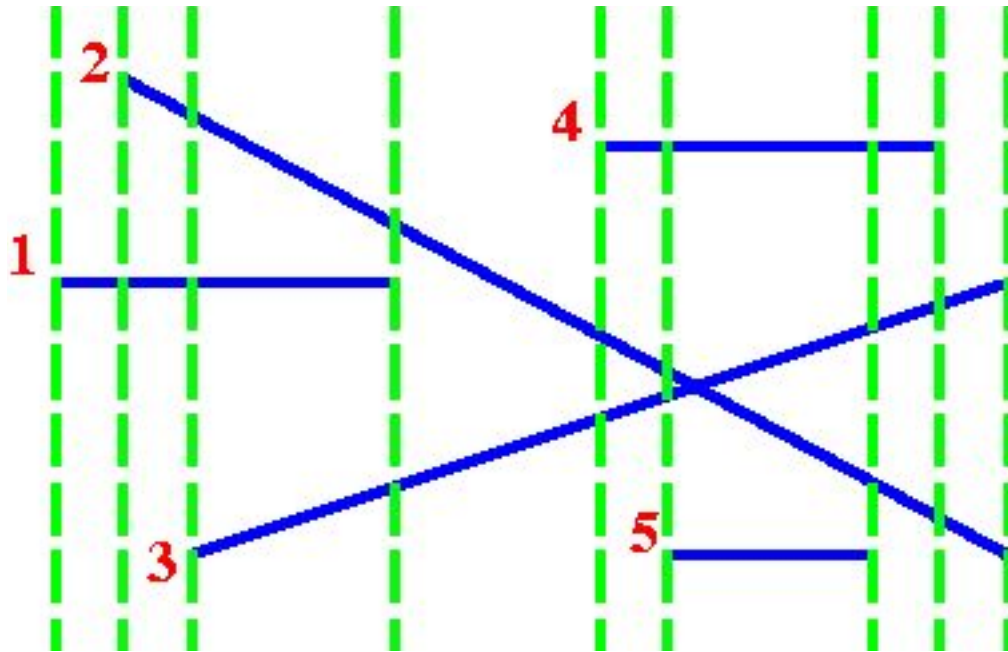
Line Sweep (using vertical line)



Events Processing

- Active List = {2, 3}
- Current Event **E3**
 - above = {2}
 - below = {}
 - No intersection tests
- Remove 3 from active list
- Similarly, remove 2 from active list

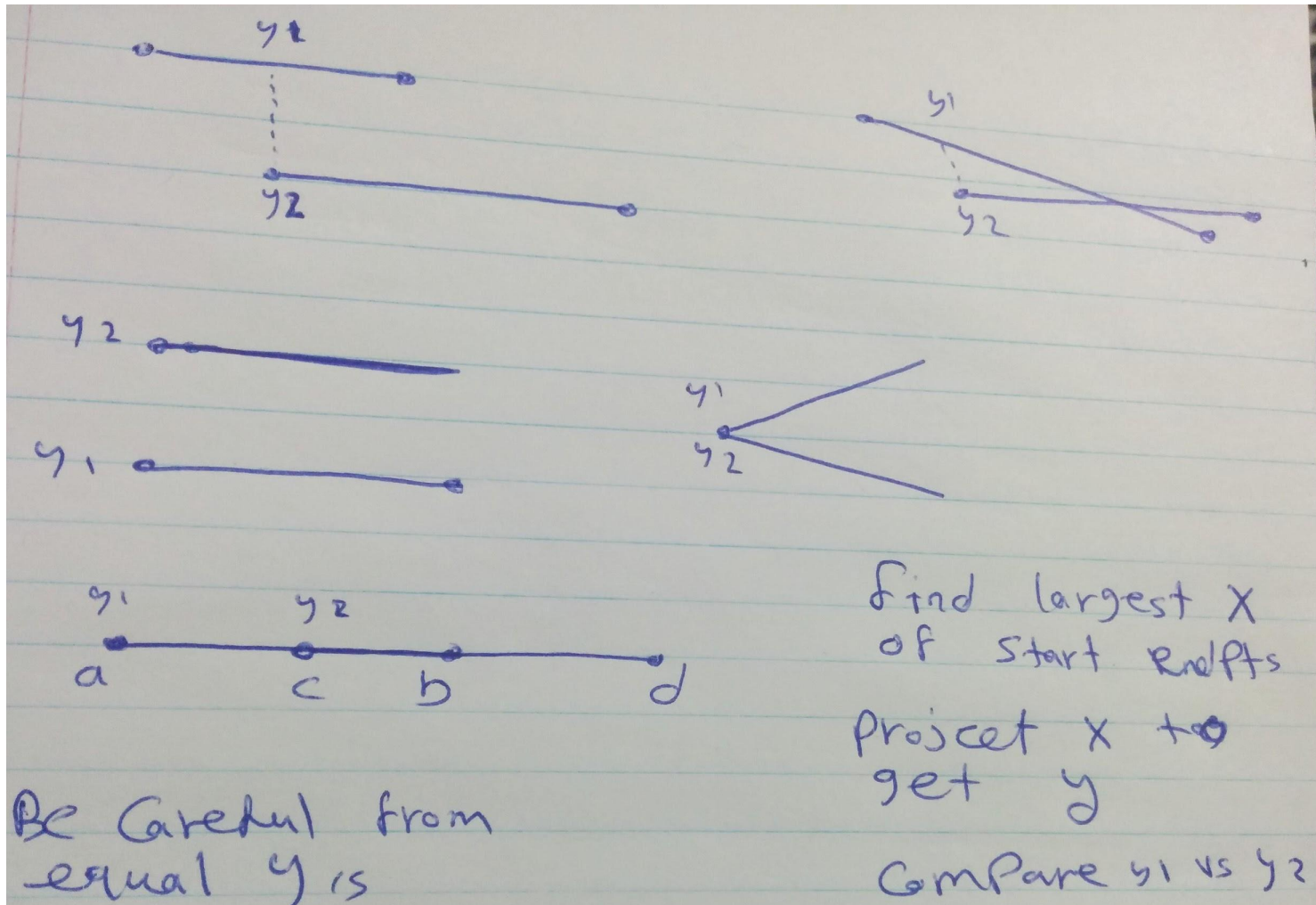
Line Sweep (using vertical line)



Notes

- Any intersecting pairs will eventually **become adjacent**...hence detected
- The **variability** of checks at start and end event surrounds the different cases
- E.g. During S3, 3 couldn't check intersect with above (2), because 1 is barrier in-between.
- However, when 1E occurs, it check above and below, which found intersection
- In other words, if at both start & end we check current segment against above/below we will miss this case.

Is segment A below B?



Data Structures

■ Segment Structure

- 2 points of a segment
- sorting segments based on how is lower (in y)

■ Insert/Find/Delete Segments Structure

- To do that efficiently, one needs Balanced Tree (e.g. AVL)
- One can use C++ set/Priority Queue most of time
- However, in some problem, you must use your Tree

■ Events structure

- endpoint point, event type, parent segment
- sorting events based on X, if tie **ENTRY** events first

Algorithm Flow

- Read Segments...Create & Sort Events
- Iterate on events
- If entry event
 - Find BTree find below & above segments of current
 - intersect(current, below)? intersect(current, above)?
 - Insert it.
- If exist event
 - Find BTree find below & above segments of current
 - intersect(above, below)?
 - Delete it.

Segment Struct

```
struct segment {
    point p, q;
    int seg_idx;

    segment() {seg_idx = -1;}
    segment(point p_, point q_, int seg_idx_) {
        if (q_ < p_)
            swap(p_, q_);
        p = p_, q = q_, seg_idx = seg_idx_;
    }

    double CY(int x) const {
        if (dcmp(p.X, q.X) == 0)
            return p.Y; // horizontal

        double t = 1.0 * (x - p.X)/(q.X - p.X);
        return p.Y + (q.Y - p.Y)*t;
    }
    // operator< is very tricky and can cause 100 WAs.
    bool operator<(const segment& rhs) const {
        if(same(p, rhs.p) && same(q, rhs.q))
            return false;

        int maxX = max(p.X, rhs.p.X);
        int yc = dcmp(CY(maxX), rhs.CY(maxX));

        if (yc == 0) // critical condition
            return seg_idx < rhs.seg_idx;
        return yc < 0;
    }
};
```


Event Structure

```
int ENTRY = +1, EXIT = -1;           // entry types
const int MAX_SEGMENTS = 50000 + 9;
const int MAX_EVENTS = MAX_SEGMENTS * 2;

struct event {
    point p;
    int type, seg_idx;
    // smaller X first. If tie: ENTRY event first. Last on smaller Y
    bool operator <(const event & rhs) const {
        if (dcmp(p.X, rhs.p.X) != 0)
            return dcmp(p.X, rhs.p.X) < 0;
        if (type != rhs.type)
            return type > rhs.type;
        return dcmp(p.Y, rhs.p.Y) < 0;
    }
};

int n;
segment segments[MAX_SEGMENTS];
event events[MAX_EVENTS];
set<segment> sweepSet;
typedef set<segment>::iterator ITER;
```

Read Segments..Create/Order Events

```
cin >> n;
lp(i, n)
{
    cin >> x >> y;    point p1 = point(x, y);
    cin >> x >> y;    point p2 = point(x, y);
    segments[i] = segment(p1, p2, i);
}
bentleyOttmann_lineSweep();
```

```
void bentleyOttmann_lineSweep() {    // O( (k+n) logn )
    // Prepare events
    lp(i, n)
    {
        events[2*i] = {segments[i].p, ENTRY, i};
        events[2*i+1] = {segments[i].q, EXIT, i};
    }
    sort(events, events+2*n);
```

Sweeping over the events

```
lp(i, 2*n) {  
    if (events[i].type == ENTRY) {  
        auto status = sweepSet.insert(segments[events[i].seg_idx]);  
        ITER cur = status.first, below = before(cur), above = after(cur);  
  
        if(!status.second) {  
            FoundIntersection(cur->seg_idx, events[i].seg_idx); // Duplicate  
        } else {  
            if(intersectSeg(cur, above))  
                FoundIntersection(cur->seg_idx, above->seg_idx);  
            if(intersectSeg(cur, below))  
                FoundIntersection(cur->seg_idx, below->seg_idx);  
        }  
    } else {  
        ITER cur = sweepSet.find(segments[events[i].seg_idx]);  
  
        if(cur == sweepSet.end())  
            continue; // e.g. Duplicate  
  
        ITER below = before(cur), above = after(cur);  
  
        if(intersectSeg(above, below))  
            FoundIntersection(above->seg_idx, below->seg_idx);  
        sweepSet.erase(cur);  
    }  
}
```

Little Utilities

```
bool intersectSeg(ITER seg1Iter, ITER seg2Iter) {
    if (seg1Iter == sweepSet.end() || seg2Iter == sweepSet.end())
        return false;
    return intersect(seg1Iter->p, seg1Iter->q, seg2Iter->p, seg2Iter->q);
}

ITER after(ITER cur) {
    return cur == sweepSet.end() ? sweepSet.end() : ++cur;
}

ITER before(ITER cur) {
    return cur == sweepSet.begin() ? sweepSet.begin() : --cur;
}

void FoundIntersection(int i, int j) {
    printf("%d %d\n", i + 1, j + 1);
}
```

Simple Apps

■ Test if polygon is simple

- Recall that, a polygon is simple if it doesn't have intersecting segments
- Note, consecutive polygon edges intersect at endpoint
- We implemented before $O(n^2)$ solution
- One can do trivial adjustment to this code to check if any 2 edges intersect or not

■ Test if two polygons intersect

- Again, direct test if any of edges intersect
- Again, special care for endpoints

Harder Apps

- Intersection (or union, or difference) of two simple polygons
 - Determine all intersection points, and use them as new vertices to construct the answer
- Polygon Triangulation
 - [See1](#), [See2](#)

تم بحمد الله

علمكم الله ما ينفعكم

ونفعكم بما تعلمتم

وزادكم علماً