

Competitive Programming From Problem 2 Solution in O(1)

Number Theory Totient and Möbius Functions

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Euler's totient function

- \bullet $\phi(n)$, the Phi Function
 - Count integers i < n such that gcd(i, n) = 1
 - gcd(a, b) = 1 => then coprimes: gcd(5, 7), gcd(4, 9)
 - gcd(prime, i) = 1 for i < prime
- $\phi(10) = 4 \Rightarrow 1, 3, 7, 9$
- $\phi(5) = 4 \Rightarrow 1, 2, 3, 4 \dots \phi(prime) = prime-1$
- If a, b, c are pairwise coprimes, then
 - $\varphi(a*b*c) = \varphi(a) * \varphi(b) * \varphi(c)$
- if $k \ge 1$ $\varphi(p^k) = p^k p^{k-1} = p^{k-1}(p-1) = p^k \left(1 \frac{1}{p}\right)$.

Euler's totient numbers

- Online Sequence
- φ(n) = 1, 1, 2, 2, 4, 2, 6, 4, 6, 4, 10, 4, 12, 6, 8, 8, 16, 6, 18, 8, 12, 10, 22, 8, 20, 12, 18, 12, 28, 8, 30, 16, 20, 16, 24, 12, 36, 18, 24, 16, 40, 12
 - $\phi(1) = \phi(2) = 1. \ \phi(5) = 4$
 - $\varphi(n)$ is even for n > 2
 - $\operatorname{sqrt}(n) \le \varphi(n) \le n \operatorname{sqrt}(n)$: Except 2, 6

 - $n = \sum_{i} \varphi(di)$ where d are the divisors of n

Menon's identity

$$\sum_{\substack{1 \le k \le n \\ \gcd(k,n)=1}} \gcd(k-1,n) = \varphi(n)d(n),$$

d(n): # of n divisors

The identity is not important in competitive:)

Phi Code: Brute Force

```
int phi(int n) {
   int ret = 0;
   for (int i = 1; i <= n; i++)
       if (gcd(i, n) == 1)
        ret++;
   return ret;
}</pre>
```

Phi Code: Prime Factorization

```
// Factorize and use fact p^(n-1)*(p-1)
int phi(int n)
    int p to k, relative primes = 1;
    for (int i = 2, d = 1; i*i <= n; i += d, d = 2) {
        if(!(n % i)) {
            p to k = 1;
            while (!(n % i))
                p to k *= i, n /= i;
            relative primes *= (p to k/i) * (i-1);
    if (n != 1)
        relative primes *= (n-1);
    return relative primes;
```

Phi Code: Range Generator

```
void phi generator() // using seive
    const int MAX = 1000000;
    char primes[MAX];
    int phi[MAX];
   memset(primes, 1, sizeof(primes));
    for (int k = 0; k < MAX; ++k)
        phi[k] = 1;
    for (int i = 2; i \le MAX; ++i) {
        if (primes[i]) {
            phi[i] = i-1; // ph(prime) = p-1
            for (int j = i * 2; j \le MAX; j += i) {
                primes[j] = 0;
                int n = j, pow = 1;
                while (!(n % i))
                    pow *= i, n /= i;
               phi[i] *= (pow/i) * (i-1);
```

Phi Factorial Code

Can you prove?

```
// phi(N!) = (N is prime ? N-1 : N ) * phi((N-1)!)
ll phi_factn(int n)
{
    ll ret = 1;
    for (int i = 2; i <= n; ++i)
        ret = ret * (isprime(i) ? i-1 : i);
    return ret;
}</pre>
```

Square-free integer

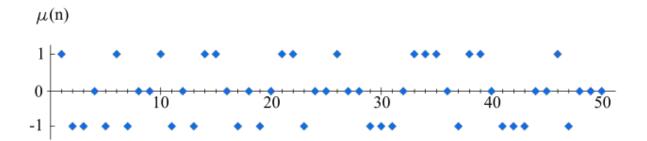
- Is not divisible by perfect square (except 1)
 - perfect square: sqrt(n) = is integer. sqrt(16) = 4
 - SQ: e.g. not divisible by 16=4x4...or 49=7x7...etc
- In other words, no prime number occurs more than once: e.g. n = 2*5*11 is square free, but n = 2*3*3*3*7 is not (divisible by 9 = 3x3)
- I-th square free: 1, 2, 3, 5, 6, 7, 10, 11, 13, 14, 15, 17, 19, 21, 22, 23, 26, 29, 30, 31, 33, 34
 - F(13) = 19

Möbius function

- $\mu(1) = 1$
- $\mu(n) = 1$ if n is a **square-free** positive integer with an **even number** of prime factors.
 - E.g. $\mu(2*3*5*7) = 1$
- $\mu(n) = -1$ if n is a square-free positive integer with an odd number of prime factors.
 - E.g. $\mu(2*3*5) = -1$
- $\mu(n) = 0$ if n has a squared prime factor.
 - E.g. $\mu(2*3*3*7) = 0$

Möbius sequence

- $\mu(n) = 1, -1, -1, 0, -1, 1, -1, 0, 0, 1, -1, 0, -1, 1, 1, 0, -1, 0, -1, 0, -1, 0, -1, 0, 1, 1, -1, 0, 0, 1, 0, 0, -1, -1, -1, 0, 1, 1, 1, 1, 0, -1, 1, 1, 0, -1, -1, 0, 0, 1, -1, 0$
- $\underline{\mu(n)} + 1 = 2, 0, 0, 1, 0, 2, 0, 1, 1, 2, 0, 1, 0, 2$



Möbius function code

```
int moebius(int n)
    int mebVal = -1;
    for(int i = 2; i * i <= n; i++)
        if(n % i == 0)
            if(n %(i*i)) == 0)
                return 0;
            n \neq i, mebVal = -mebVal;
    if(n)
        mebVal = - mebVal;
    return mebVal;
}
```

Möbius generator code

Möbius and Inclusion Exclusion

- Recall, in IE we compute all subsets, and add odd subsets and remove negative subsets
- Assume generating all subsets of primes but in implicit way (e.g. iterate on numbers), Möbius can tell you if number is odd subset or even
- Typically, ignoring numbers with repeated prime factors is target
- Then Möbius(n) plays perfect role in that

Möbius and Inclusion Exclusion

- Given square free number, find its index?
 - E.g. F reverse(n = 19) = 13
- Reverse thinking: Can we remove non SFree?
- In range n, remove who divides by 2x2, 3x3, 4x4, 5x5, 6x6...etc
 - 4x4 and 6x6 already computed by previous ones
 - Ignore duplicate primes (4x4)...use IE for others F(2)+F(3)+F(5)-F(6)

```
ll val = 19, idx = val;
for (ll i = 2; i*i<=val; i++)
   idx -= moebius[i]*(val/(i*i));</pre>
```

Möbius and Inclusion Exclusion

- Count the triples (a,b,c) such $a,b,c \le n$, and gcd(a,b,c) = 1
 - Reverse thinking, total (# triples gcd > 1)
 - How many triples with gcd multiple of 2: $(n/2)^3$
 - How many triples with gcd multiple of 3: $(n/3)^3$
 - and 4? Ignore any numbers of internal duplicate primes
 - and 6? already computed in 2, 3. Remove it: $-(n/6)^3$
 - This is Inclusion Exclusion

```
int n = 4;
ll sum = n*n*n;
for (ll i = 2; i <= n; ++i)
    sum -= moebius[i]*(n/i)*(n/i)*(n/i);</pre>
```

Totient and Möbius connection

Sum over divisors d of n

$$\sum_{d} d \mu \left(\frac{n}{d} \right) = \phi (n),$$

Read Totient functions proves

تم بحمد الله

علمكم الله ما ينفعكم

ونفعكم بما تعلمتم

وزادكم علمأ

Problems

- UVA 417, 11417, 10179, 10820, 10990, 11426, 10299, 11327
- TIMUS: 1673
- http://www.geeksforgeeks.org/eulers-totientfunction/