



# Competitive Programming

From Problem 2 Solution in  $O(1)$

## Computational Geometry

### Simple and Convex Polygons

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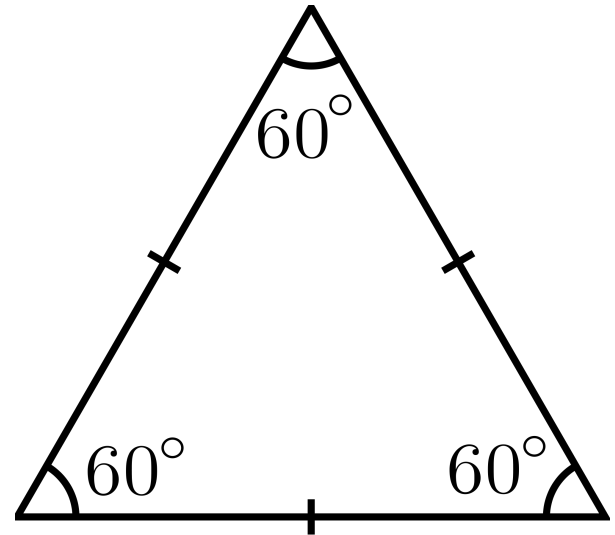
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# Triangle

## ■ Triangle is the **simplest polygon**

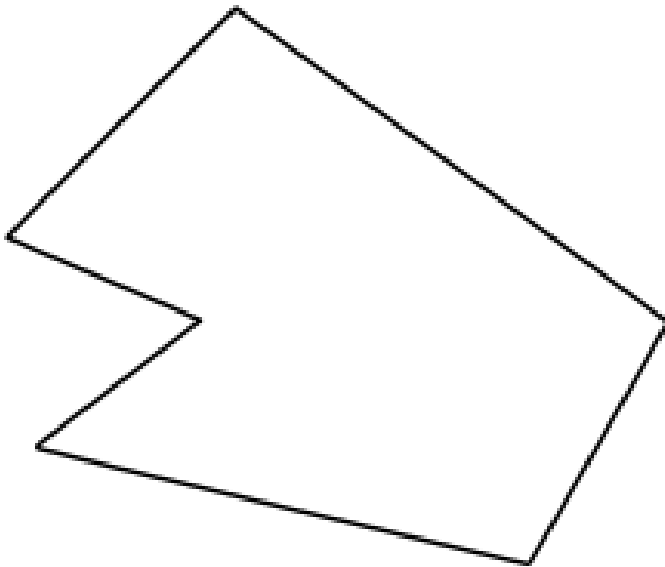
- has 3 points
- has 3 sides
- has 3 angles with sum = 180
- If all sides of equal length  
 $\Rightarrow$  equal angles =  $180 / 3 = 60$



Src: [https://upload.wikimedia.org/wikipedia/commons/thumb/9/96/Triangle\\_Equilateral.svg/2000px-Triangle\\_Equilateral.svg.png](https://upload.wikimedia.org/wikipedia/commons/thumb/9/96/Triangle_Equilateral.svg/2000px-Triangle_Equilateral.svg.png)

# Polygon

- Sequence of  $n$  points/vertices/corners ( $n \geq 3$ )
  - Terminologies to read

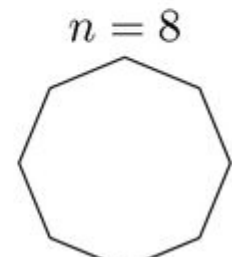
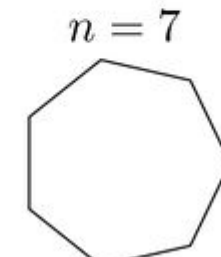
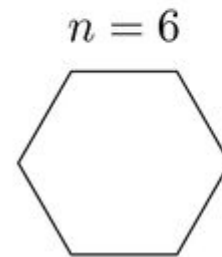
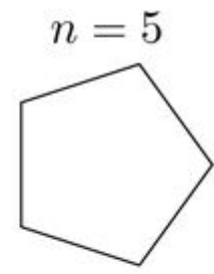
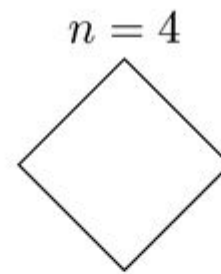
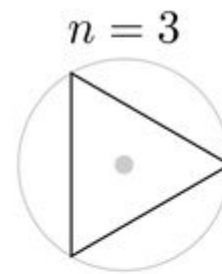


Src: <http://www.basic-mathematics.com/types-of-polygons.html>

Number of sides	Name of polygon
3	triangle
4	quadrilateral
5	pentagon
6	hexagon
7	heptagon
8	octagon
9	nonagon
10	decagon
12	dodecagon

# Regular Polygon

- It is equiangular and equilateral
  - Equal Angles, Equal Sides
- Internal angle =  $(n-2) * 180 / n$



Src: <http://www.mathguide.com/lessons/RegularPoly/regularpoly.gif>

# Polygon: Implementation

- Typically polygon of  $n$  points is represented with vector of  $n$  points
- Then  $(p[i], p[i+1])$  is an edge:  $i \Rightarrow [0 - n-1]$
- Last edge is problematic
  - One way use:  $(p[i], p[(i+1)\%n])$  to wrap to begin
  - Note: Mod is expensive...and code wise sometimes annoying
  - Other trick: add  $p[0]$  to the vector of points ( $n+1$  points)
  - Iterate till  $n$  exactly (not  $p.size()$ )
  - From problem to another, determine what fits more

# Simple Polygon

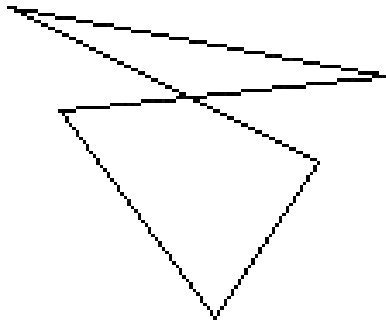
- A flat shape consisting of straight, **non-intersecting** line segments



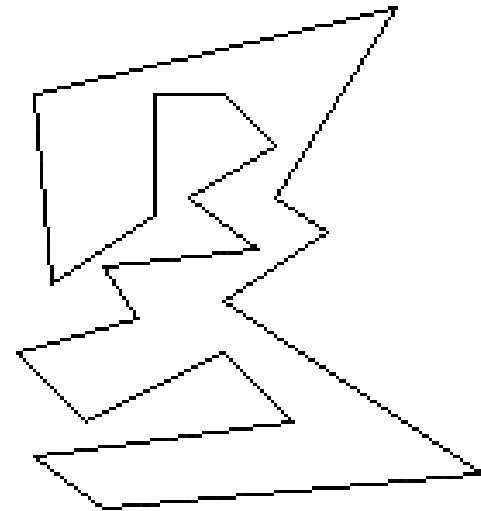
Src: [https://upload.wikimedia.org/wikipedia/commons/thumb/3/3c/Polygons\\_Examples\\_of\\_polygons.png/220px-Polygons\\_Examples\\_of\\_polygons.png](https://upload.wikimedia.org/wikipedia/commons/thumb/3/3c/Polygons_Examples_of_polygons.png/220px-Polygons_Examples_of_polygons.png)

# Non Simple Polygon

- Intersecting edges



Not a Simple Polygon



Simple Polygon

Src: <http://www.personal.kent.edu/~rmuhamma/Compgeometry/MyCG/Definitions/definition.htm>

# Is Simple Polygon Implementation

- BruteForcing it is easy in  $O(n^2)$ 
  - just make sure no 2 edges intersects
  - Note 1: Don't check consecutive edges
  - Note 2: Remember last edge, edge 1 are consecutive
  - Use ccw/cp to test intersection
- Other advanced ways
  - $O(n \log n)$  using line sweep
  - $O(n)$  using very complex algorithm

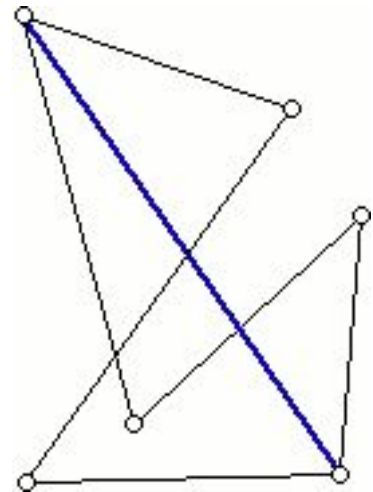
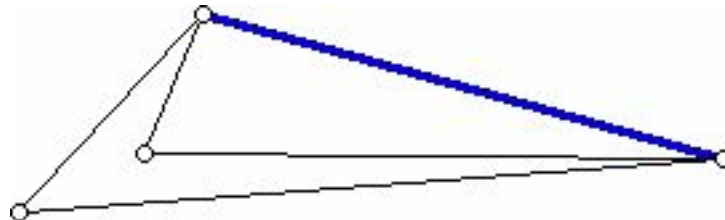
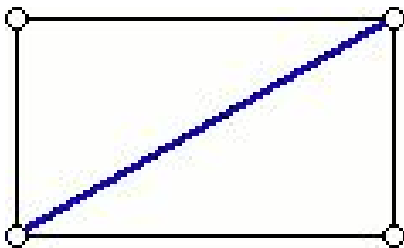


# Simple Polygon Implementation

```
// assume p0 is added again to p (n+1 points)
bool isSimplePolygon(vector<point> &p) {
    for (int i = 0; i < sz(p) - 1; i++) {
        for (int j = i + 2; j < sz(p) - 1; j++)
            if (intersect(p[i], p[i + 1], p[j], p[j + 1])
                && (i != 0 || j != sz(p) - 2) /* last/first edges are consecutive */)
                return false; //Segments must not share vertices
    }
    return true;
}
```

# Polygon Diagonals

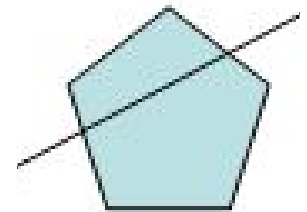
- A polygon's diagonals are line segments from one vertex (point) to another, but not the sides.
- The number of diagonals of an n-sided polygon is:  $n(n - 3) / 2$



Src: <http://www.ams.org/samplings/feature-column/fcarc-diagonals2>

# Convex Polygon

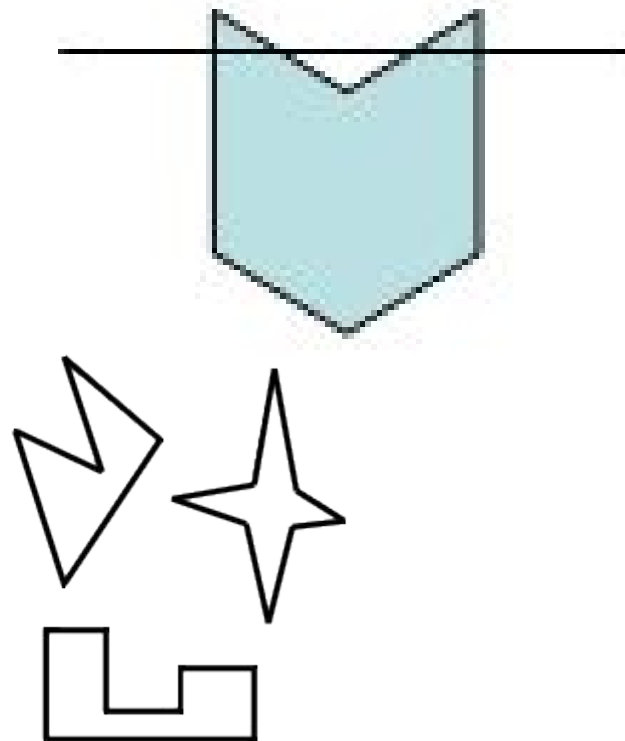
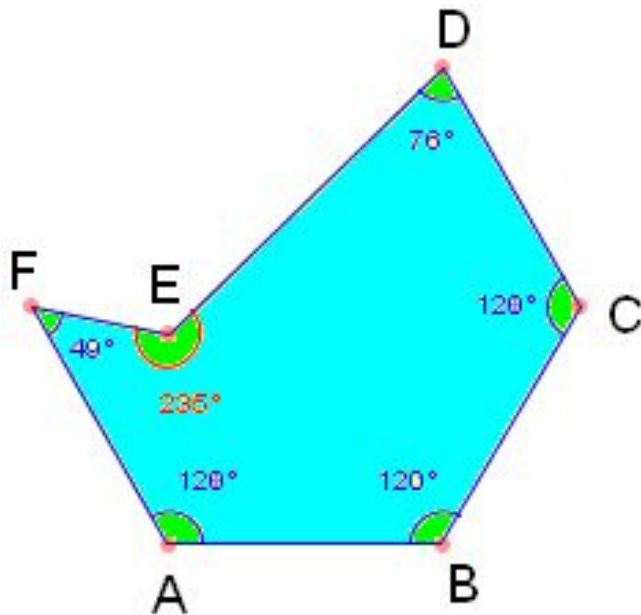
- All interior angles  $\leq 180$ 
  - E.g. all the vertices of the polygon will point outwards, away from the interior of the shape
- A line drawn through a convex polygon will intersect the polygon exactly twice
- All the diagonals of a convex polygon lie entirely inside the polygon



Src: <http://www.mathopenref.com/polygonconvex.html>

# Concave Polygon

- Non convex one



Src: <http://image.tutorvista.com/cms/images/38/concave-polygon.jpg>

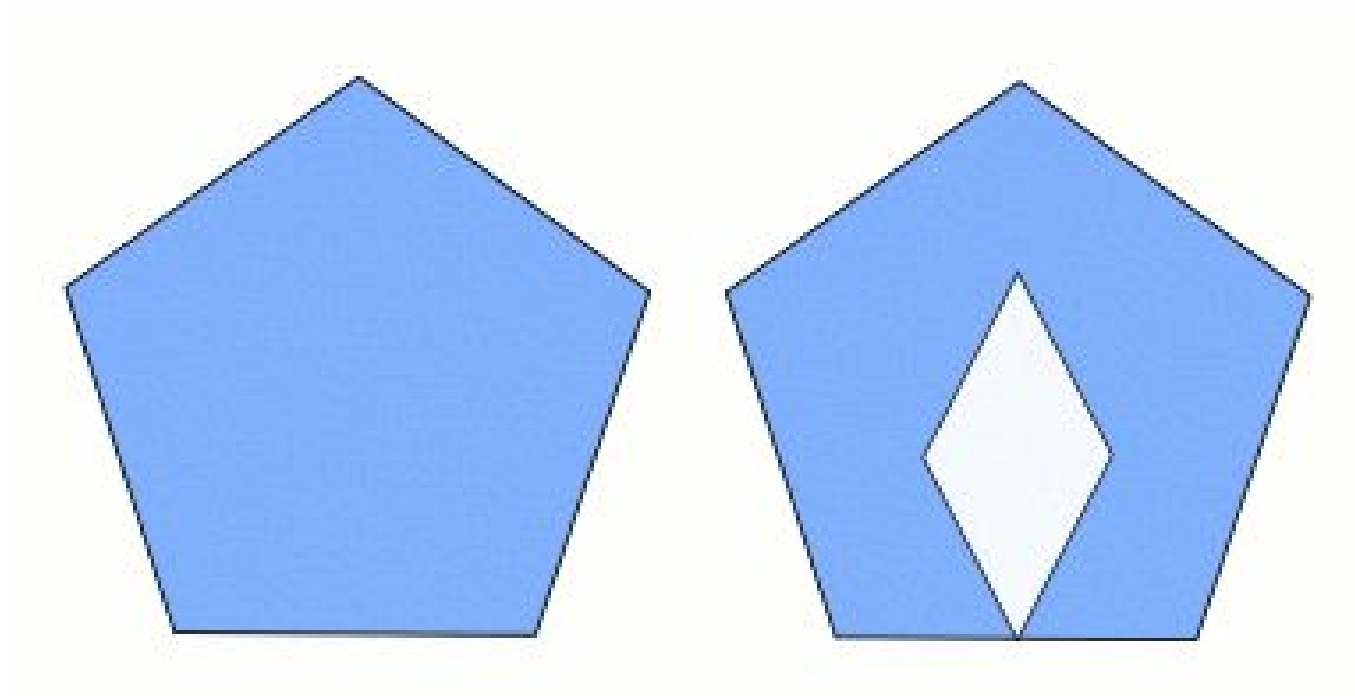
# Is Convex Polygon Implementation

- One can just iterate / compute each angle
  - if all  $\leq 180 \Rightarrow$  convex
  - This involves double computations
- What about using CCW test?
- Let's assume the polygon points are ordered
  - E.g. all are in ccw order or cw
- If every 2 edges have same direction (e.g. all are ccw), then this is convex
  - If there mix of signs, this is concave

# Convex Polygon Implementation

```
bool isConvexPolygon(vector<point> &p) {  
    // all polygon 3 consecutive points must have same sign (ccw or cw)  
    p.push_back(p[0]), p.push_back(p[1]); // wrap points for simplicity  
  
    int sign = ccw(p[0], p[1], p[2]);  
    bool ok = true;  
    for (int i = 1; ok && i < sz(p) - 2; i++) {  
        if (ccw(p[i], p[i + 1], p[i + 2]) != sign)  
            ok = false;  
    }  
    p.pop_back(), p.pop_back();  
    return ok;  
}  
  
// isConcave = !isConvex  
// isConcave = 2 different signs
```

# Polygon with holes (ignore)



Src: [https://plus.maths.org/issue43/features/kirk/polygon\\_hole.gif](https://plus.maths.org/issue43/features/kirk/polygon_hole.gif)

# Critical Notes

## ■ Polygon input

- **Coordinates** are integer or real?
- Points are **ordered**? yes: ccw? cw? on angle with P0?
- Caution: Sometimes input has **duplicate** points
- Caution: Sometimes input has **collinear** points
- Caution: for invalid input ( $n < 3$ ), special cases:  $n = 3$

## ■ Polygon type

- Remember keywords: **simple/!simple - convex/concave**
- Some algorithms for simple and/or convex polygons are very fast than generic case



# تم بحمد الله

علمكم الله ما ينفعكم

ونفعكم بما تعلمتم

وزادكم علماً