



Competitive Programming

From Problem 2 Solution in $O(1)$

Combinatorial Game Theory

Game of Nim - Examples

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Nim with skip move

- Given the original nim game with extra rule
 - Skip turn: A player is allowed to say I will skip turn
 - Player 1 is allowed up to A skippings
 - Player 2 is allowed up to B skippings
 - Given the N piles, A, B who is the winner?
 - If $A = B$? Then whoever the winner can **cancel** the other player move to insure winning (**Move Cancellation**)
 - If $A > B$ and 1st wins in normal nim, then he just plays normally and cancels moves for 2nd if he skipped
 - If $A > B$ and 1st loses in normal nim, then he **skip the first time**, and then play as previous case to win too :)
 - So if $A > B$, 1st always win. If $A < B$, 2nd always win.

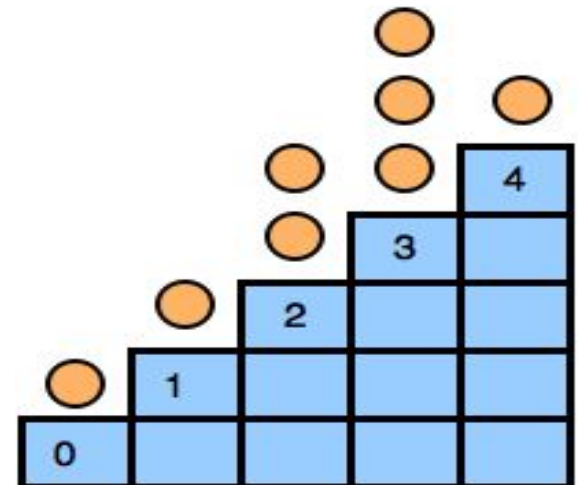
Dividing a number

- Given an integer N
 - Move: divide N by a prime power > 1
 - e.g. $3, 3^2, 3^3 \dots$
 - primes: $2, 3, 5, 7, 11 \dots$
 - Loser: $N = 1$
 - **Solution:** Represent N using its prime powers
 - $$n = 2^{a_1} 3^{a_2} 5^{a_3} 7^{a_4} \dots p_k^{a_k}$$
 - Then we have k piles, each has a_i stones
 - $N = 1440600 = 2*2*2*3*5*5*7*7*7*7$
 - So piles are $= \{3, 1, 2, 4\}$
 - E.g. we have 4 7s

Staircase Nim

■ Staircase Nim

- Staircase with n steps, each step has some coins
- Move: move some coins to the left step (except first step)
- Loser: Can't make a move (e.g. all coins at `arr[0]`)
- Intuition: Every step is a pile? No only odd positions

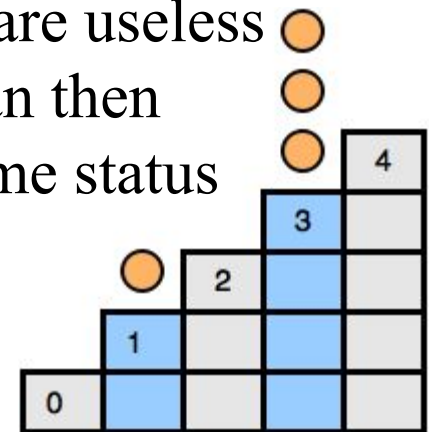


Src: <http://codeforces.com/blog/entry/44651>

Staircase Nim

■ Notes

- Movement from step 1 to step 0 \Rightarrow remove action
- If opponent moved 5 coins from step 2 to step 1, you can just move them to step 0, hence removed them too
- Recall: this is move **cancellation strategy**
- In general, movements on even position are useless
- Movements from odd position to even can then be considered as **removed**, but affect game status
- **Solution:** xor odd positions (the piles)



Even steps are useless!

Turning Turtles Game

- Given a horizontal line of N coins: Head/Tail
 - 1 2 3 4 5 6 7 8 9 10
 - T H T T H H T T T H
- Moves
 - Pick any head, and flip it to tail
 - Optionally, flip any coin on left of your chosen coin
 - Loser: No more heads to flip
- Example: THHH (let's make action on 4)
 - THHT (flip 4th only)
 - THTT (flip 4th and 3rd) TTHT (flip 4th and 2nd)
 - HHHT (flip 4th and 1st) So Position 4 => 4 moves

Turning Turtles Game

■ Observations

- Action only on Heads. So may be heads \sim Nim piles
- Head in k th position has k moves. May be Pile size = k
- E.g. THHTTH \Rightarrow {2, 3, 6} as Nim pile sizes!
- We need to verify 3 game moves as 2 nim possibilities
 - We can take whole nim pile
 - We can take ANY thing $<$ whole size

■ Verification for:

- Flip Head only
- Flip Head and Flip a left coin: Tail to Head
- Flip Head and Flip a left coin: Head to Tail

Turning Turtles Game

■ Verification

- Flip k-th Head only
 - If we did that, we cancelled the k moves from position
 - Seems as if we just removed while pile content!
- Flip k-th Head and t-th tail (recall $k > t$)
 - Head at k has k moves cancelled
 - But a tail at is now H and has t moves
 - As if pile of k stones reduced to t stones only
- Flip k-th Head and t-th head (recall $k > t$)
 - This is tricky. We already satisfied nim moves
 - Intuition: If this game is nim \Rightarrow this move is not a new case or a useless move

Turning Turtles Game

- Flip k-th Head and t-th head
 - This move cancels: my k moves and his t moves!!
 - As if it means, the other head will never change the final status situation (e.g. I have a winning for it)
 - Let's try example: THTTH \Rightarrow Nim {2, 5}
 - Recall the **move duplication strategy** to always win
 - Make them equal piles \Rightarrow (2, 2) \Rightarrow (0, 0) \Rightarrow 1st win
 - So flipping kth head means reduce k to t (e.g. 5 to 2)
 - Then his second pile equal to mine (2, 2)
 - So I will win anyway \Rightarrow e.g. the 2 H's are cancelled
- Overall: 2 equivalent games.

Your turn: Twins game

- Same as Turning Turtles Game, but
 - You must flip 2 coins, not optionally
 - Then k th position has $k-1$ move NOT k
 - Also, it means we can't take whole pile move!
 - Actually both previous notes lead to same thing
 - As in terms of nim, pile size = 1 is losing condition
 - So $F(n)$ in this game = **$F(n-1)$ in normal nim**
 - So k th position is pile of size $n-1$
 - So THHTTH $\Rightarrow \{1, 2, 5\}$ NOT $\{2, 3, 6\}$ sizes
 - If you **indexed** game as 0-based, then $F(n) = n$
 - Sometimes one of the 2 indexings makes computations/patterns easier (later)

Silver Dollar Game

■ From judge

- Array of N cells, that has coins (at most 1 coin in a cell)
- Move: Pick a coin and move to **any left square: BUT**
 - **Can't jump** over other coins
 - **Can't put** two coins in one cell
- Loser: Can't do a move
- Example. for 4 coins in position {2, 5, 7, 10}
- | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|----|----|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 1 |

Silver Dollar Game

■ Thoughts

- Issue: each coin is **not independent** game any more
- **Intuition:** Each coin is constrained with its left one.
 - Distance between every 2 consecutive is pile size
 - E.g. $\{2, 5, 7, 11\} \Rightarrow$ piles $\{1, 2, 1, 3\}$
 - E.g. between 7 and 11 = 3 steps
 - **Wrong equivalence:** Moving 7 to 6: decrease in pile and add in another (e.g. generates piles: $\{1, 2, 0, 4\}$)

Silver Dollar Game

- Let's understand the game more
 - Input={5}: First player will win directly
 - Input={5, 6}: Second player will always win directly
 - 1st will have to move 5 to 4
 - 2nd can just return them consecutive: e.g. from 6 to 5
 - So overall even # of steps for any input $\{x, x+1\}$
 - Clearly this input is always like $\{1, 2\}$: pile size = 0
 - Input={5, 8}
 - In similar manner, a move from 5 to 4 can be **canceled** by move from 8 to 7 to have same difference (poker)
 - So, the 2 players should **focus** on moving 2nd not 1st
 - Clearly this input is always like $\{1, 4\}$: pile size = 2

Silver Dollar Game

- Let's understand the game more
 - Input = {2, 4, 6}
 - One can force (4, 6) as a pile
 - And (0, 2) as a pile (e.g. 0 from left boundary)
 - And handle each pair with same winning strategy for 2 values. so total piles {1, 1} \Rightarrow xor = 0 (2nd can win)
 - Given that we identified a winning strategy that we can definitely apply (either first if will win or 2nd), we can safely consider it
 - You don't need to think of what other ways of winning: having one way to force winning is enough
 - Game using our grouping/cancellation tricks = Nim

Silver Dollar Game

- Let's understand the game more
 - Input = {2, 4, 6}
 - What about grouping (2, 4) \Rightarrow pile size = 1
 - Now 6 is free: it doesn't correspond to something in Nim
 - Let's restrict it with its max possible moves to left
 - 6 has 2 numbers before it: so its pile size = $6 - 2 - 1 = 3$
 - E.g. piles (1, 3) \Rightarrow $\text{xor} \neq 0 \Rightarrow$ 1st win \Rightarrow where mistake?
 - Although 6 has up to 3 moves to left, if one directly moved 6 to 5, now it corresponds to empty pile
 - When 4 moves to 3: 6 pile grows up again
 - So our mapping is not nim equivalent = wrong approach

Silver Dollar Game

■ Overall solution

- From most right number to left:
- Group each 2 consecutive numbers
- If they are odd, most left number will be grouped with 0
- $\{2, 5, \underline{7}, \underline{11}\} \Rightarrow \{2, 5\}, \{7, 11\} \Rightarrow (2, 3)$
- $\{2, 5, 7, \underline{11}, \underline{20}\} \Rightarrow \{0, 2\}, \{5, 7\}, \{11, 20\} \Rightarrow (1, 1, 8)$

■ For other ways of explaining it: [See1](#), [See2](#)

- Mainly, they map faster to nim and consider it poker nim
- I cancelled a move in the game \Rightarrow Direct nim
- Think **much**: Duplication, Cancellation, Actual game, Win strategy, Correct equivalence

Pots of gold game

- Given set of coins (e.g. 10, 7, 8, 12).
 - Move: Take coins in a cell from one of ends (e.g. 10 | 12)
 - The **winner** is the player who collected a **higher number of coins** at the end of game.
- Solution:
 - Don't be cheated. This is NOT nim-related game
 - Winning is not about **LAST** move, but **overall** choices
 - This is a classical search/minimax/dp problem
- **Lesson:** Remember properties of Nim

Final Notes

- Make sure game is impartial
 - If it is not, think in search/backtrack directions
- Remember Nim game nature
 - With exception, the independance feature of sub-games
- Dependent games are tricky/hard to analyze
 - Understand & analyze the game
 - Try to decouple to independent sub-games
 - Suggest mappings and verify it carefully
 - Let duplication/cancellation strategies be your friends

تم بحمد الله

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