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CS 536: Science of Programming

Homework 7: Separation Logic, Nondeterminism, and Parallelism

1 Linked list and separation logic

Task 1.1 (Written, 10 points)

```
S := head, := head;  
   $x_0 := !(head, +1);$   
   $x_1 := !(x_0 + 1);$   
   $head_2 := !(x_1, +1);$   
   $!(x_1, +1) := nil$ 
```

Task 1.2 (Written, 12 points)

$$\begin{aligned} & \{ \text{List}(L, \text{head}, m+3) \} \\ & \{ \text{List}(L_1; L_2, \text{head}, 3+m) \} \\ & \{ \text{List}(a_1; a_2; a_3; L_2, \text{head}, 3+m) \} \end{aligned}$$

$$\Rightarrow \{ \text{head} \mapsto a_1, i_1 * i_1 \mapsto a_2, i_2 * i_2 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$\text{head}_1 := \text{head}; \quad \{ \text{head}_1 \mapsto a_1, i_1 * i_1 \mapsto a_2, i_2 * i_2 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$x_0 := !(\text{head}_1 + 1); \quad \{ x_0 = i_1 \wedge \text{head}_1 \mapsto a_1, x_0 * i_1 \mapsto a_2, i_2 * i_2 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$\Rightarrow \{ \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, i_2 * i_2 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$x_1 := !(x_0 + 1); \quad \{ x_1 = i_2 \wedge \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, x_1 * i_2 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$\Rightarrow \{ \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, x_1 * x_1 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$\text{head}_2 := !(x_1 + 1); \quad \{ \text{head}_2 = i_3 \wedge \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, x_1 * x_1 \mapsto a_3, i_3 * \text{List}(L_2, i_3, m) \}$$

$$\Rightarrow \{ \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, x_1 * x_1 \mapsto a_3, \text{head}_2 * \text{List}(L_2, \text{head}_2, m) \}$$

$$!(x_1 + 1) := \text{nil} \quad \{ \text{head}_1 \mapsto a_1, x_0 * x_0 \mapsto a_2, x_1 * x_1 \mapsto a_3, \text{nil} * \text{List}(L_2, \text{head}_2, m) \}$$

$$\Rightarrow \{ \text{List}(a_1; a_2; a_3, \text{head}_1, 3) * \text{List}(L_2, \text{head}_2, m) \}$$



## 2 Resource Logic

### Task 2.1 (Written, 10 points)

a)  $x \mapsto \boxed{x}$

b)  $x \mapsto \boxed{2} \quad \boxed{2} \leftarrow y$

c)  $x \mapsto \begin{array}{|c|} \hline x+1 \\ \hline x \\ \hline \end{array}$

d)  $x \mapsto \boxed{y} \mapsto \boxed{2} \leftarrow w$

e)  $y \mapsto \boxed{2} \quad \text{where } x=2$

### Task 2.2 (Written, 6 points)

a)  $x \mapsto \boxed{x} \leftarrow y$

$x \mapsto \boxed{x} \leftarrow z$

b)  $x \mapsto \boxed{2} \quad \boxed{2} \leftarrow y$

$x \mapsto \boxed{2}$   
 $y \mapsto \boxed{2}$

c)  $x \mapsto \boxed{5}$

$x \mapsto \boxed{7}$

### Task 2.3 (Written, 8 points)

a)  $(x \mapsto 2 * y \mapsto 2) * z \mapsto 2 \Rightarrow (z \mapsto 2 * y \mapsto 2) * x \mapsto 2$

$$(x \mapsto 2 * y \mapsto 2) * z \mapsto 2$$

$$\Leftrightarrow x \mapsto 2 * (y \mapsto 2 * z \mapsto 2) \quad \text{Rule 2}$$

$$\Leftrightarrow x \mapsto 2 * (z \mapsto 2 * y \mapsto 2) \quad \text{Rule 1}$$

$$\Leftrightarrow (z \mapsto 2 * y \mapsto 2) * x \mapsto 2 \quad \text{Rule 1}$$

Hence  $(x \mapsto 2 * y \mapsto 2) * z \mapsto 2 \Rightarrow (z \mapsto 2 * y \mapsto 2) * x \mapsto 2$

b)  $(x \mapsto 3 \wedge y \mapsto 3) \Rightarrow x = y$

By definition,  $e \mapsto e'$  is considered imprecise assertion because  $e \mapsto e'$  just needs to appear somewhere in the heap. Therefore we cannot say for sure that  $x = y$  since  $x \mapsto 3$  may appear in one place and  $y \mapsto 3$  may appear somewhere else.

Counter example:

$$x \mapsto \boxed{3}$$

$$\boxed{3} \mapsto y$$

Here we have  $x \mapsto 3 \wedge y \mapsto 3$  but  $x \neq y$ .

Hence  $(x \mapsto 3 \wedge y \mapsto 3) \not\Rightarrow x = y$



$$c) \text{ emp} * ((\exists x. y \mapsto x) * w \mapsto 2) \Rightarrow \exists x. (y \mapsto x * w \mapsto 2)$$

$$\text{emp} * ((\exists x. y \mapsto x) * w \mapsto 2)$$

$$\Leftrightarrow ((\exists x. y \mapsto x) * w \mapsto 2) * \text{emp} \quad \text{Rule 1}$$

$$\Leftrightarrow ((\exists x. y \mapsto x) * w \mapsto 2) \quad \text{Rule 3}$$

$$\Leftrightarrow \exists x. (y \mapsto x * w \mapsto 2) \quad \text{Rule 6}$$

$$\text{Hence } \text{emp} * ((\exists x. y \mapsto x) * w \mapsto 2) \Rightarrow \exists x. (y \mapsto x * w \mapsto 2)$$

$$d) (\exists x. y \mapsto x) * x \mapsto 2 \not\Rightarrow \exists x. (y \mapsto x * x \mapsto 2)$$

As per rule 6, if  $x$  is not free in  $P_2$ , then only can we apply rule 6.

However, in this example,  $x$  is free since it is not bound to the  $x$  in  $\exists x$ .

Counter example:

$$x \mapsto \boxed{2} \quad \boxed{6} \vdash \dots y \quad \text{where } x=6$$

In the above counter example, we have  $(\exists x. y \mapsto x) * x \mapsto 2$  but not  $\exists x. (y \mapsto x * x \mapsto 2)$ .

$$\text{Hence } (\exists x. y \mapsto x) * x \mapsto 2 \not\Rightarrow \exists x. (y \mapsto x * x \mapsto 2)$$

### 3 Non-determinism

Task 3.1 (Written, 10 points)

a)  $S_1 \triangleq \text{while } \{x > 0 \rightarrow y := x/y \mid x \leq 0 \rightarrow y := y * x\}$

$$\sigma \triangleq \{x=0, y=1\}$$

$$\langle S_1, \{x=0, y=1\} \rangle$$

$$\rightarrow \langle y := x/y; S_1, \{x=0, y=1\} \rangle$$

$$\rightarrow^2 \langle S_1, \{x=0, y=0\} \rangle$$

$$\rightarrow \langle y := y * x; S_1, \{x=0, y=0\} \rangle$$

$$\rightarrow^2 \langle S_1, \{x=0, y=0\} \rangle$$

$$\rightarrow \langle y := y * x; S_1, \{x=0, y=0\} \rangle$$

$$\rightarrow^2 \langle S_1, \{x=0, y=0\} \rangle$$

...

b)  $\Sigma_0 = \{\{x=0, y=1\}\}$

$$\Sigma_1 = M(y := x/y, \{x=0, y=1\}) \vee M(y := y * x, \{x=0, y=1\})$$

$$= \{\{x=0, y=0\}\} \vee \{\{x=0, y=0\}\}$$

$$= \{\{x=0, y=0\}\}$$

$$\Sigma_2 = M(y := x/y, \{x=0, y=0\}) \vee M(y := y * x, \{x=0, y=0\})$$

$$= \{\perp\} \vee \{\{x=0, y=0\}\}$$

$$= \{\perp, \{x=0, y=0\}\}$$

#### 4 Parallel programs

Task 4.1 (Written, 10 points)

$$\langle [x := y; y := y + 1 \parallel y := y + 3], (\{x = 3, y = 2\}, \{ \}) \rangle$$

$$\langle [y := y + 1 \parallel y := y + 3], (\{x = 2, y = 2\}, \{ \}) \rangle$$

$$\langle [x := y; y := y + 1 \parallel \text{skip}], (\{x = 3, y = 5\}, \{ \}) \rangle$$

$$\langle [\text{skip} \parallel y := y + 3], (\{x = 2, y = 3\}, \{ \}) \rangle$$

$$\langle [y := y + 1 \parallel \text{skip}], (\{x = 2, y = 5\}, \{ \}) \rangle$$

$$\langle [y := y + 1 \parallel \text{skip}], (\{x = 5, y = 5\}, \{ \}) \rangle$$

$$\langle [\text{skip} \parallel \text{skip}], (\{x = 2, y = 6\}, \{ \}) \rangle$$

$$\langle [\text{skip} \parallel \text{skip}], (\{x = 5, y = 6\}, \{ \}) \rangle$$

$$m(S_2, (\{x = 3, y = 2\}, \{ \})) = \{ (\{x = 2, y = 6\}, \{ \}), (\{x = 5, y = 6\}, \{ \}) \}$$

5 One more wrap-up question

Task 5.1 (Written, 0 points).

15 hours