Parallel Computing for Science & Engineering Spring 2013: MPI point-to-point 1

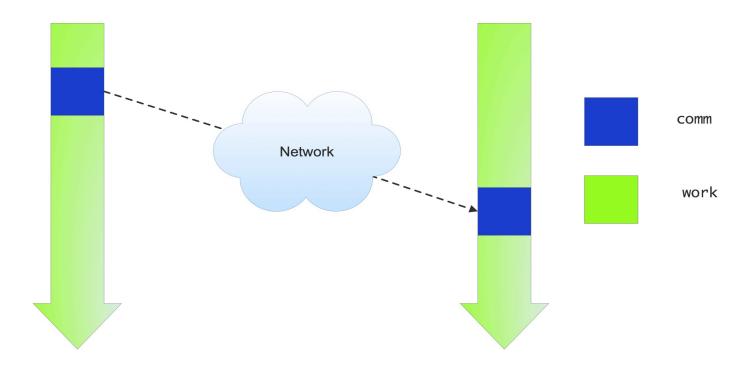
Instructors:

Victor Eijkhout, Research Scientist, TACC Kent Milfeld, Research Associate, TACC



Life would be simple if....

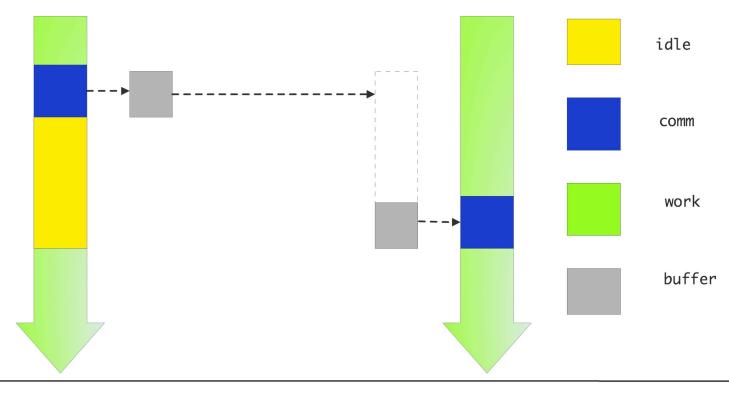
 Processors would just send and receive, and the network would DWIM





Unfortunately

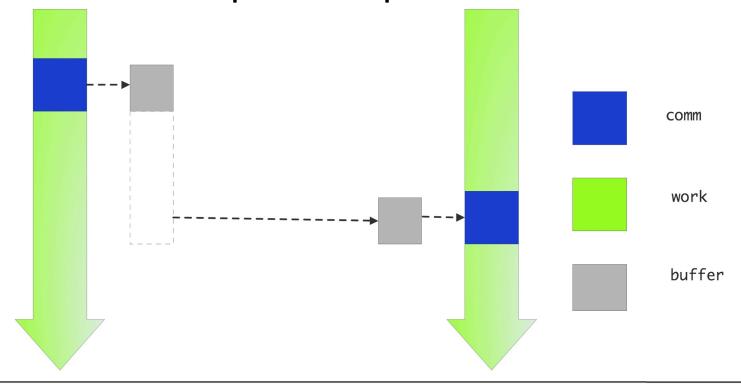
 Data has to be somewhere: on one process or the other





Non-blocking solution

 Create a buffer and let the send data sit there until someone picks it up



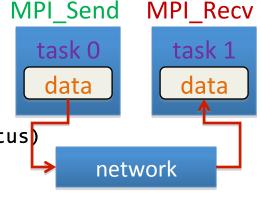


Blocking Send/Receive

Generic Syntax

- MPI_Send(buf, count, datatype, dest, tag, comm)
- MPI_Recv(buf, count, datatype, source, tag, comm, status)
- When MPI sends a message, it doesn't just send the contents; it also sends an envelope describing the contents:

Argument	Description	
buf	initial address of send/receive buffer (reference)	
count	number of items to send (integer)	
datatype	MPI data type of items to send/receive	
dest	MPI rank of task receiving the data (integer)	
source	MPI rank of task sending the data (integer)	
tag	message ID (integer)	
comm	MPI communicator (set of exchange processors)	
status	returns information on the message received	



Parts of a P-2-P Communication:

Data
Send to/Recv from
Message ID



Details

buffer	data (address in C, name of array/value in Fortran)		
count	Length of source array (in elements, 1 for scalars)		
datatype	Data Type: e.g. mpi_int (C), mpi_integer (F90), mpi_double_precision (F90), mpi_double (C), etc.		
source	Rank (proc #) of source in communicator group		
tag	Message identifier (arbitrary integer)		
communicator	Group of processors		
status	Information about message		
ierr	Error (argument in Fortran, returned in C)		

	С		Fortran	
status	MPI_Status	mystat;	<pre>integer mystat(MPI_STATUS_SIZE)</pre>	
datatype	MPI_Datatype	mytype;	integer mytype	
comm.	MPI_Comm	mycomm;	integer mycomm	



Language Example

```
    C
        ierr=MPI_Send(&a[0],cnt,type,dest,tag,com);
    F
        call MPI_Send(a, cnt,type,dest,tag,com,ierr)
    C
        ierr=MPI_Recv(&b[0],cnt,type, src,tag,com,&status);
    F
        call MPI_Recv(b, cnt,type, src,tag,com, status,ierr)
```

 Call blocks until send data of a has been sent or copied to a buffer. Recv's block until data is in b.



P-2-P Example

```
#include <mpi.h>
int main(int argc, char* argv[]){
MPI Comm Comm=MPI COMM WORLD;
int npes,iam=-1,ierr;
ierr=MPI Init(&argc, &argv);
ierr=MPI Comm size(Comm, &npes);
ierr=MPI Comm rank(Comm, &iam);
ierr=MPI Finalize();
printf("iam=%d\n",iam);
```



P-2-P Example

```
#include <mpi.h>
int main(int argc, char* argv[]){
MPI Comm Comm=MPI COMM WORLD;
MPI Status status;
int npes, iam=-1, ierr, irec=-1;
ierr=MPI Init(&argc, &argv);
ierr=MPI Comm size(Comm, &npes);
ierr=MPI Comm rank(Comm, &iam);
if(iam==0)
        ierr=MPI Send(&iam, 1,MPI INT, 1,9, Comm);
if(iam==1)
        ierr=MPI Recv(&irec,1,MPI INT, 0,9, Comm,&status);
ierr=MPI Finalize();
printf("iam=%d, received=%d\n",iam,irec);
```



The 6 Basic MPI Call Summary

- MPI is used to create parallel programs based on message passing
- Usually the same program is run on multiple processors
- The 6 basic calls in MPI are

```
MPI_Init(&argc, &argv);
MPI_Comm_Rank(MPI_COMM_WORLD, &myid);
MPI_Comm_Size(MPI_COMM_WORLD, &numprocs);
MPI_Send(buffer,count, MPI_TYPE, dest,tag, MPI_COMM_WORLD);
MPI_Recv(buffer,count, MPI_TYPE, src, tag, MPI_COMM_WORLD, &stat);
MPI_Finalize();
```

MPI_TYPE is an MPI Parameter or User Data Type buffer is passed by reference



MPI_SendRecv

- Initiates send and receive at the same time.
- Completes when both send and receive buffers are safe to use
- Useful for communications patterns where each node sends and receives messages (two-way communication). Good for avoiding deadlock, implementing shifts/rings.
- Executes a **standard mode** send & receive operation for dest and src, respectively.
- The send and receive operations use the same communicator, but have distinct tags.



Bidirectional Communication with MPI Sendrecv

• C

Fortran



Blocking vs Non-blocking

Blocking

- A blocking send routine will only return after it is safe to modify the data area.
- Safe means that modifications in the data area will not affect the data to be sent.
- A Safe send does not imply that the data was actually received.
- A blocking send can be either synchronous or asynchronous.

Non-blocking

- Send/receive routines return immediately.
- Non-blocking operations request the MPI library to perform the operation when possible.
- It is unsafe to modify the data area until the requested operation has been performed. There are wait routines used to do this (MPI_Wait)
- Primarily used to overlap computation with communication



Blocking vs non-Blocking Routines

Description	Syntax for C bindings
Blocking send	MPI_Send(buf, count, datatype, dest, tag, comm)
Non-blocking send	<pre>MPI_Isend(buf,count, datatype, dest, tag, comm, request)</pre>
Blocking receive	<pre>MPI_Recv(buf, count, datatype, source, tag, comm, status)</pre>
Non-blocking receive	<pre>MPI_Irecv(buf,count, datatype, source, tag, comm, request)</pre>
Wait for completion	MPI_Wait(request, status)

request: used by non-blocking send and receive operation.



Non-blocking Communication

- Non-blocking send
 - send call returns immediately
 - send actually occurs later
- Non-blocking receive
 - receive call returns immediately
 - when received data is needed, call a wait subroutine
- Non-blocking communication used to overlap communication with computation (and communication with communication!).
- Can be used to prevent deadlock.



Non-blocking Send with MPI_Isend

C

Fortran

- request is the id for the message call
- Don't use data area until communication is complete



Non-blocking Receive with MPI_Irecv

C

Fortran

- request is an id for communication
- Note: There is no status parameter.
- Don't use **data** area until communication is complete



MPI_Wait Used to Complete Communication

- request from MPI_Isend or MPI_Irecv
 - the completion of a send operation indicates that the sender is now free to update the data in the send buffer
 - the completion of a receive operation indicates that the receive buffer contains the received message
- MPI_Wait blocks until message specified by request completes



MPI_Wait Usage

```
MPI Request request;
  MPI Status status;
  ierr = MPI Wait(&request, &status)

    Fortran

  integer request
  integer status(MPI STATUS SIZE)
  call MPI Wait ( request, status, ierr)
```



Nonblocking Examples



One-way Send/Recv Communication

Blocking Send & Non-blocking Recv

```
if (rank==0) then
    call MPI_send(    sendbuf, count, MPI_REAL, 1, tag, MPI_COMM_WORLD,
    ierr)
elseif (rank==1) then
    call MPI_Irecv(    recvbuf, count, MPI_REAL, 0, tag, MPI_COMM_WORLD,
    req, ierr)
    call MPI_Wait(    req,status,ierr)
endif
```

Non-blocking Send & Non-blocking Recv



Two-way Communication: Deadlock

Deadlock 1 (always deadlocks)

```
if (rank==0) then
    call MPI_Recv(    recvbuf,count,MPI_REAL,1,tag,MPI_COMM_WORLD,status,ierr)
    call MPI_Send(    sendbuf,count,MPI_REAL,1,tag,MPI_COMM_WORLD,ierr)
elseif (rank==1) then
    call MPI_Recv(    recvbuf,count,MPI_REAL,0,tag,MPI_COMM_WORLD,status,ierr)
    call MPI_Send(    sendbuf,count,MPI_REAL,0,tag,MPI_COMM_WORLD,ierr)
endif
```

Deadlock 2 (deadlocks when system buffer is too small)



Two-way Communication: Solutions

Solution 1 (but this doesn't allow bidirectional communication)

```
if (rank==0) then
    call MPI_Send(    sendbuf,count,MPI_REAL, 1,tag,MPI_COMM_WORLD,ierr)
    call MPI_Recv(    recvbuf,count,MPI_REAL, 1,tag,MPI_COMM_WORLD,status,ierr)
elseif (rank==1) then
    call MPI_Recv(         recvbuf,count,MPI_REAL, 0,tag,MPI_COMM_WORLD,status,ierr)
    call MPI_Send(         sendbuf,count,MPI_REAL, 0,tag,MPI_COMM_WORLD,ierr)
endif
```

Solution 2



Two-way Communication: Solutions

```
Solution 3
if (rank==0) then
   call MPI_ISend( sendbuf,count,MPI_REAL,
1,tag,MPI_COMM_WORLD,req1,ierr)
   call MPI_IRecv( recvbuf,count,MPI_REAL,
0,tag,MPI_COMM_WORLD,req2,ierr)
elseif (rank==1) then
   call MPI_ISend( sendbuf,count,MPI_REAL,
0,tag,MPI_COMM_WORLD,req1,ierr)
   call MPI IRecv( recvbuf, count, MPI_REAL,
   1, tag, MPT COMM WORLD, req2, ierr)
endif
         call MPI Wait( reg1, status, ierr)
         call MPI Wait( reg2, status, ierr)
Solution 4
if (rank==0) then
   call MPI BSend(
                          sendbuf,count,MPI REAL,1,tag,MPI COMM WORLD,ierr)
                            recvbuf, count, MPI REAL,
   call MPI Recv(
   0,tag,MPI COMM WORLD,status,ierr)
elseif (rank==1) then
   call MPI BSend(
                          sendbuf,count,MPI REAL,1,tag,MPI COMM WORLD,ierr)
   call MPI Recv( recvbuf, count
0, tag, MPI_COMM_WORLD, status, ierr)
                            recvbuf, count, MPI REAL.
endif
```



Two-way Communications Summary

	CPU 1	CPU 2
Deadlock 1 Recv/Send		Recv/Send
Deadlock 2	Send/Recv	Send/Recv
Solution 1	Send/Recv	Recv/Send
Solution 2	SendRecv	SendRecv
Solution 3	Isend/Irecv/Wait	Isend/Irecv/Wait
Solution 4	Bsend/Recv	Bsend/Recv



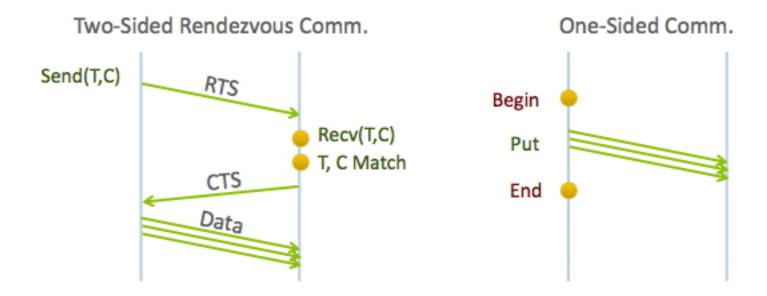
Wait types

- MPI_Wait : wait for one request
- MPI_Waitall: wait for an array of requests, good for load balanced tasks, or when all needed
- MPI_Waitany: wait for one in an array of requests, good for unbalanced tasks, or if they can be processed individually
- MPI_Waitsome: wait for any number in an array, much like MPI Waitany



One-Sided or RMA

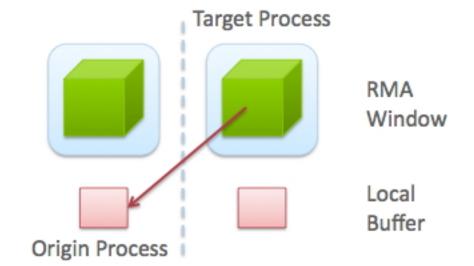
- It would be nice to avoid that two-way orchestration: just write into another process' memory or read from it
- Less overhead, easier to code





One-Sided concepts

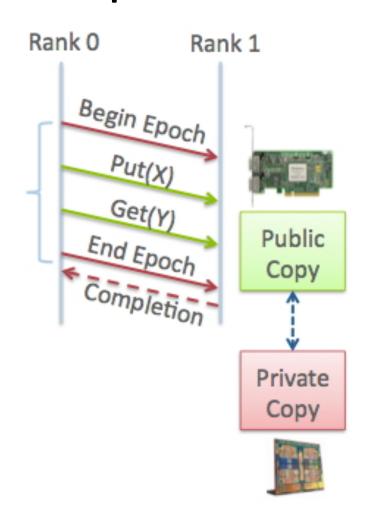
- Target & origin
 processes: origin
 issues the call, target
 does nothing explicit
- Window & local memory: window is accessible to others





More RMA concepts

- Origin vs Target, window vs local mem
- Actions: Put, Get, Accumulate
- Epoch: just like MPI_Wait: you have to make sure data has arrived





RMA routines

```
int MPI_Win_create(void *base, MPI_Aint size, int disp_unit, MPI_Info info, MPI_Comm comm, MPI_Win *win)
```

```
MPI_Get( origin_addr, origin_count, origin_datatype, target_rank, target_disp, target_count, target_datatype, win)
```

MPI_Win_fence(assert,win)

Use of fences is one way to synchronize. There are more.



Fences

- MPI_Win_fence(assert,win)
- Assertions:
 - MPI_MODE_NOSTORE
 - MPI_MODE_NOPUT
 - MPI MODE NOPRECEDE
 - MPI MODE NOSUCCEED
- Example:

```
MPI_Win_fence(
    (MPI_MODE_NOSTORE|MPI_MODE_NOPRECEDE), win);
```



RMA limitations

