NAME

```
stat, fstat, lstat – get file status
```

SYNOPSIS

```
#include <sys/types.h>
#include <sys/stat.h>
#include <unistd.h>

int stat(const char *path, struct stat *buf);
int fstat(int fd, struct stat *buf);
int lstat(const char *path, struct stat *buf);
```

Feature Test Macro Requirements for glibc (see **feature_test_macros**(7)):

```
lstat(): _BSD_SOURCE || _XOPEN_SOURCE >= 500
```

DESCRIPTION

These functions return information about a file. No permissions are required on the file itself, but — in the case of **stat()** and **lstat()** — execute (search) permission is required on all of the directories in *path* that lead to the file.

stat() stats the file pointed to by *path* and fills in *buf*.

Istat() is identical to **stat**(), except that if *path* is a symbolic link, then the link itself is stat-ed, not the file that it refers to.

fstat() is identical to **stat()**, except that the file to be stat-ed is specified by the file descriptor fd.

All of these system calls return a stat structure, which contains the following fields:

```
struct stat {
    dev_t st_dev; /* ID of device containing file */
    ino_t st_ino; /* inode number */
    mode_t st_mode; /* protection */
    nlink_t st_nlink; /* number of hard links */
    uid_t st_uid; /* user ID of owner */
    gid_t st_gid; /* group ID of owner */
    dev_t st_rdev; /* device ID (if special file) */
    off_t st_size; /* total size, in bytes */
    blksize_t st_blksize; /* blocksize for file system I/O */
    blkcnt_t st_blocks; /* number of 512B blocks allocated */
    time_t st_atime; /* time of last access */
    time_t st_mtime; /* time of last modification */
    time_t st_ctime; /* time of last status change */
};
```

The *st_dev* field describes the device on which this file resides. (The **major**(3) and **minor**(3) macros may be useful to decompose the device ID in this field.)

The st_rdev field describes the device that this file (inode) represents.

The *st_size* field gives the size of the file (if it is a regular file or a symbolic link) in bytes. The size of a symlink is the length of the pathname it contains, without a trailing null byte.

The st_blocks field indicates the number of blocks allocated to the file, 512-byte units. (This may be smaller than $st_size/512$ when the file has holes.)

The st_blksize field gives the "preferred" blocksize for efficient file system I/O. (Writing to a file in smaller

chunks may cause an inefficient read-modify-rewrite.)

Not all of the Linux file systems implement all of the time fields. Some file system types allow mounting in such a way that file and/or directory accesses do not cause an update of the st_atime field. (See *noatime*, *nodiratime*, and *relatime* in **mount**(8), and related information in **mount**(2).) In addition, st_atime is not updated if a file is opened with the **O_NOATIME**; see **open**(2).

The field *st_atime* is changed by file accesses, for example, by **execve**(2), **mknod**(2), **pipe**(2), **utime**(2) and **read**(2) (of more than zero bytes). Other routines, like **mmap**(2), may or may not update *st_atime*.

The field *st_mtime* is changed by file modifications, for example, by **mknod**(2), **truncate**(2), **utime**(2) and **write**(2) (of more than zero bytes). Moreover, *st_mtime* of a directory is changed by the creation or deletion of files in that directory. The *st_mtime* field is *not* changed for changes in owner, group, hard link count, or mode.

The field *st_ctime* is changed by writing or by setting inode information (i.e., owner, group, link count, mode, etc.).

The following POSIX macros are defined to check the file type using the *st_mode* field:

```
S_ISREG(m) is it a regular file?
S_ISDIR(m) directory?
S_ISCHR(m) character device?
S_ISBLK(m) block device?
S_ISFIFO(m) FIFO (named pipe)?
S_ISLNK(m) symbolic link? (Not in POSIX.1-1996.)
S_ISSOCK(m) socket? (Not in POSIX.1-1996.)
```

The following flags are defined for the *st_mode* field:

1B 1 1. S_IFMT 0170000 bit mask for the file type bit fields S_IFSOCK 0140000 socket S_IFLNK 0120000 symbolic link S_IFREG 0100000 regular S_IFBLK 0040000 directory 0060000 block device S_IFDIR 0020000 character device S IFIFO0010000 FIFO S ISUID S IFCHR 0004000 set UID bit S_ISGID 0002000 set-group-ID bit (see below) S_ISVTX 0001000 sticky bit (see below) S IRWXU mask for file owner permissions S IRUSR 00400 owner has read permission S_IWUSR 00200 owner write permission has 00100 owner has execute permission S_IRWXG 00070 S_IXUSR mask for group permissions S IRGRP 00040 group has read permission S_IWGRP 00020 group write permission S IXGRP 00010 group has permission has execute S_IRWXO 00007 mask for permissions for others (not in group) S IROTH 00004 others have read permission S_IWOTH 00002 others have write permission S_IXOTH 00001 others have execute permission

The set-group-ID bit (**S_ISGID**) has several special uses. For a directory it indicates that BSD semantics is to be used for that directory: files created there inherit their group ID from the directory, not from the effective group ID of the creating process, and directories created there will also get the **S_ISGID** bit set. For a file that does not have the group execution bit (**S_IXGRP**) set, the set-group-ID bit indicates mandatory file/record locking.

The sticky bit (**S_ISVTX**) on a directory means that a file in that directory can be renamed or deleted only by the owner of the file, by the owner of the directory, and by a privileged process.

RETURN VALUE

On success, zero is returned. On error, -1 is returned, and *errno* is set appropriately.

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ERRORS

EACCES

Search permission is denied for one of the directories in the path prefix of *path*. (See also **path_resolution**(7).)

EBADF

fd is bad.

EFAULT

Bad address.

ELOOP

Too many symbolic links encountered while traversing the path.

ENAMETOOLONG

File name too long.

ENOENT

A component of *path* does not exist, or *path* is an empty string.

ENOMEM

Out of memory (i.e., kernel memory).

ENOTDIR

A component of the path prefix of *path* is not a directory.

EOVERFLOW

(stat()) path refers to a file whose size cannot be represented in the type off_t. This can occur when an application compiled on a 32-bit platform without -D_FILE_OFFSET_BITS=64 calls stat() on a file whose size exceeds (2<<31)-1 bits.

CONFORMING TO

These system calls conform to SVr4, 4.3BSD, POSIX.1-2001.

Use of the *st_blocks* and *st_blksize* fields may be less portable. (They were introduced in BSD. The interpretation differs between systems, and possibly on a single system when NFS mounts are involved.)

POSIX does not describe the **S_IFMT**, **S_IFSOCK**, **S_IFLNK**, **S_IFREG**, **S_IFBLK**, **S_IFDIR**, **S_IFCHR**, **S_IFIFO**, **S_ISVTX** bits, but instead demands the use of the macros **S_ISDIR**(), etc. The **S_ISLNK**() and **S_ISSOCK**() macros are not in POSIX.1-1996, but both are present in POSIX.1-2001; the former is from SVID 4, the latter from SUSv2.

Unix V7 (and later systems) had **S_IREAD**, **S_IWRITE**, **S_IEXEC**, where POSIX prescribes the synonyms **S_IRUSR**, **S_IWUSR**, **S_IXUSR**.

Other Systems

Values that have been (or are) in use on various systems:

11111. hex	name ls	octal	description	n f000	S_IFMT	1700	000 mask for file	
type 0000		000000	SCO	out-of-ser	vice i	node; B	SD unknown	
		type;	SVID-v	2 a	ınd	XPG2	have both	
		0	and	0100000) f	or or	dinary file	
1000 S_IFIFC	pp 01000	0 FIFO (n	amed pipe)	2000	S_IFCHR	c	020000 char-	
acter special (V7) 3000 S	_IFMPC		030000) multiple:	xed characte	er special (V7)	
4000 S_IFDIF	d/	040000	directory				(V7)	
5000 S_IFNA	M	050000	XENIX	n	amed	speci	al file	
		with	two	sub	types,	distingu	iished by	
		st_rdev v	alues 1, 2 00	001 S_	INSEM	S	000001 XENIX	
semaphore subtype of IFNAM 0002 S_INSHD m 000002 XENIX shared data subtype of								
schiaphore subty	of it is the or)0 <u>=</u>	IDIID	***	000002	ZLI 11ZL SHALC	a data subtype of	
IFNAM		FBLK		060000 b		specia		
1 71	6000 S_I	FBLK		060000 b			l (V7)	

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9000	S_IFCMP		110000	VxFS				c	ompressed
9000	S_IFNWK	n	110000	network		spe	cial		(HP-UX)
a000	S_IFLNK	1@	120000	symbolic			link		(BSD)
b000	S_IFSHAD		130000	Solaris	sha	dow	inode	for	ACL
			(not		seen		by		userspace)
c000	S_IFSOCK	s=	140000	socket	(BSD;	also	"S_IFSOC"	on	VxFS)
d000	S_IFDOOR	D>	150000	Solaris do	or e000	S_IFWH	IT w%	160	0000 BSD
whiteout (not used for inode) 0200 S_ISVTX 001000 sticky bit: save swapped text									
even			г	ıfter use	(V7)				reserved
(SVID-	v2)			Or	noi	n-directories	s: don't	cach	e this
			file (S	SunOS)				On o	directories:
restricte	ed de	eletion					flag	(S	VID-v4.2)
0400	S_ISGID		002000	set-group-	·ID	on	execut	ion	(V7)
			for	directori	es:	use	BSD se	emantics	s for
			propaga	tion of G	D 0400	S_ENF	МT	00	2000 Sys-
tem V	V file lockin	g enforc	ement	(shared				with	S_ISGID)
0800	S_ISUID		004000	set-user-II)	on	executi	on	(V7)
0800	S_CDF	004000	director	y is a conte	xt depen	dent			file
(HP-UX	X)								

A sticky command appeared in Version 32V AT&T UNIX.

NOTES

Since kernel 2.5.48, the *stat* structure supports nanosecond resolution for the three file timestamp fields. Glibc exposes the nanosecond component of each field using names either of the form *st_atim.tv_nsec*, if the **_BSD_SOURCE** or **_SVID_SOURCE** feature test macro is defined, or of the form *st_atimensec*, if neither of these macros is defined. On file systems that do not support sub-second timestamps, these nanosecond fields are returned with the value 0.

On Linux, lstat() will generally not trigger automounter action, whereas stat() will.

For most files under the /proc directory, **stat**() does not return the file size in the st_size field; instead the field is returned with the value 0.

Underlying kernel interface

Over time, increases in the size of the *stat* structure have led to three successive versions of **stat**(): $sys_stat()$ (slot $_NR_oldstat)$, $sys_newstat()$ (slot $_NR_stat)$, and $sys_stat64()$ (new in kernel 2.4; slot $_NR_stat64$). The glibc **stat**() wrapper function hides these details from applications, invoking the most recent version of the system call provided by the kernel, and repacking the returned information if required for old binaries. Similar remarks apply for **fstat**() and **lstat**().

EXAMPLE

The following program calls **stat**() and displays selected fields in the returned *stat* structure.

```
#include <sys/types.h>
#include <sys/stat.h>
#include <time.h>
#include <stdio.h>
#include <stdlib.h>

int
main(int argc, char *argv[])
{
   struct stat sb;
   if (argc != 2) {
```

```
fprintf(stderr, "Usage: %s <pathname>\n", argv[0]);
             exit(EXIT_FAILURE);
           }
           if (stat(argv[1], \&sb) == -1) {
             perror("stat");
             exit(EXIT_SUCCESS);
           printf("File type:
                                      ");
           switch (sb.st_mode & S_IFMT) {
           case S_IFBLK: printf("block device\n");
                                                           break:
           case S_IFCHR: printf("character device\n");
                                                            break;
           case S_IFDIR: printf("directory\n");
                                                        break;
           case S IFIFO: printf("FIFO/pipe\n");
                                                          break;
           case S_IFLNK: printf("symlink\n");
                                                          break;
           case S_IFREG: printf("regular file\n");
                                                         break;
           case S_IFSOCK: printf("socket\n");
                                                         break;
           default:
                      printf("unknown?\n");
                                                      break;
           }
           printf("I-node number:
                                         %ld\n", (long) sb.st_ino);
           printf("Mode:
                                      %lo (octal)\n",
                (unsigned long) sb.st_mode);
           printf("Link count:
                                       %ld\n", (long) sb.st_nlink);
           printf("Ownership:
                                       UID=%ld GID=%ld\n",
                (long) sb.st_uid, (long) sb.st_gid);
           printf("Preferred I/O block size: %ld bytes\n",
                (long) sb.st blksize);
                                     % lld bytes \n'',
           printf("File size:
                (long long) sb.st_size);
           printf("Blocks allocated:
                                         %lld\n",
                (long long) sb.st_blocks);
           printf("Last status change:
                                         %s", ctime(&sb.st_ctime));
           printf("Last file access:
                                       %s", ctime(&sb.st atime));
           printf("Last file modification: %s", ctime(&sb.st_mtime));
           exit(EXIT_SUCCESS);
SEE ALSO
        access(2), chmod(2), chown(2), fstatat(2), readlink(2), utime(2), capabilities(7), symlink(7)
```

COLOPHON

This page is part of release 3.22 of the Linux *man-pages* project. A description of the project, and information about reporting bugs, can be found at http://www.kernel.org/doc/man-pages/.