

NAME

stat, fstat, lstat – get file status

SYNOPSIS

```
#include <sys/types.h>
```

```
#include <sys/stat.h>
```

```
#include <unistd.h>
```

```
int stat(const char *path, struct stat *buf);
```

```
int fstat(int fd, struct stat *buf);
```

```
int lstat(const char *path, struct stat *buf);
```

Feature Test Macro Requirements for glibc (see **feature_test_macros(7)**):

```
lstat(): _BSD_SOURCE || _XOPEN_SOURCE >= 500
```

DESCRIPTION

These functions return information about a file. No permissions are required on the file itself, but — in the case of **stat()** and **lstat()** — execute (search) permission is required on all of the directories in *path* that lead to the file.

stat() stats the file pointed to by *path* and fills in *buf*.

lstat() is identical to **stat()**, except that if *path* is a symbolic link, then the link itself is stat-ed, not the file that it refers to.

fstat() is identical to **stat()**, except that the file to be stat-ed is specified by the file descriptor *fd*.

All of these system calls return a *stat* structure, which contains the following fields:

```
struct stat {
    dev_t    st_dev;    /* ID of device containing file */
    ino_t    st_ino;    /* inode number */
    mode_t    st_mode;  /* protection */
    nlink_t   st_nlink; /* number of hard links */
    uid_t    st_uid;    /* user ID of owner */
    gid_t    st_gid;    /* group ID of owner */
    dev_t    st_rdev;   /* device ID (if special file) */
    off_t     st_size;   /* total size, in bytes */
    blksize_t st_blksize; /* blocksize for file system I/O */
    blkcnt_t  st_blocks; /* number of 512B blocks allocated */
    time_t    st_atime;  /* time of last access */
    time_t    st_mtime;  /* time of last modification */
    time_t    st_ctime;  /* time of last status change */
};
```

The *st_dev* field describes the device on which this file resides. (The **major(3)** and **minor(3)** macros may be useful to decompose the device ID in this field.)

The *st_rdev* field describes the device that this file (inode) represents.

The *st_size* field gives the size of the file (if it is a regular file or a symbolic link) in bytes. The size of a symlink is the length of the pathname it contains, without a trailing null byte.

The *st_blocks* field indicates the number of blocks allocated to the file, 512-byte units. (This may be smaller than *st_size*/512 when the file has holes.)

The *st_blksize* field gives the "preferred" blocksize for efficient file system I/O. (Writing to a file in smaller

chunks may cause an inefficient read-modify-rewrite.)

Not all of the Linux file systems implement all of the time fields. Some file system types allow mounting in such a way that file and/or directory accesses do not cause an update of the *st_atime* field. (See *noatime*, *nodiratime*, and *relatime* in **mount(8)**, and related information in **mount(2)**.) In addition, *st_atime* is not updated if a file is opened with the **O_NOATIME**; see **open(2)**.

The field *st_atime* is changed by file accesses, for example, by **execve(2)**, **mknod(2)**, **pipe(2)**, **utime(2)** and **read(2)** (of more than zero bytes). Other routines, like **mmap(2)**, may or may not update *st_atime*.

The field *st_mtime* is changed by file modifications, for example, by **mknod(2)**, **truncate(2)**, **utime(2)** and **write(2)** (of more than zero bytes). Moreover, *st_mtime* of a directory is changed by the creation or deletion of files in that directory. The *st_mtime* field is *not* changed for changes in owner, group, hard link count, or mode.

The field *st_ctime* is changed by writing or by setting inode information (i.e., owner, group, link count, mode, etc.).

The following POSIX macros are defined to check the file type using the *st_mode* field:

S_ISREG(m)	is it a regular file?
S_ISDIR(m)	directory?
S_ISCHR(m)	character device?
S_ISBLK(m)	block device?
S_ISFIFO(m)	FIFO (named pipe)?
S_ISLNK(m)	symbolic link? (Not in POSIX.1-1996.)
S_ISSOCK(m)	socket? (Not in POSIX.1-1996.)

The following flags are defined for the *st_mode* field:

S_IFMT	0170000	bit mask for the file type bit fields
S_IFSOCK	0140000	socket
S_IFLNK	0120000	symbolic link
S_IFREG	0100000	regular file
S_IFBLK	0060000	block device
S_IFDIR	0040000	directory
S_IFCHR	0020000	character device
S_IFIFO	0010000	FIFO
S_ISUID	0004000	set UID bit
S_ISGID	0002000	set-group-ID bit (see below)
S_ISVTX	0001000	sticky bit (see below)
S_IRWXU	00700	mask for file owner permissions
S_IRUSR	00400	owner has read permission
S_IWUSR	00200	owner has write permission
S_IXUSR	00100	owner has execute permission
S_IRWXG	00070	mask for group permissions
S_IRGRP	00040	group has read permission
S_IWGRP	00020	group has write permission
S_IXGRP	00010	group has execute permission
S_IRWXO	00007	mask for permissions for others (not in group)
S_IROTH	00004	others have read permission
S_IWOTH	00002	others have write permission
S_IXOTH	00001	others have execute permission

The set-group-ID bit (**S_ISGID**) has several special uses. For a directory it indicates that BSD semantics is to be used for that directory: files created there inherit their group ID from the directory, not from the

effective group ID of the creating process, and directories created there will also get the **S_ISGID** bit set. For a file that does not have the group execution bit (**S_IXGRP**) set, the set-group-ID bit indicates mandatory file/record locking.

The sticky bit (**S_ISVTX**) on a directory means that a file in that directory can be renamed or deleted only by the owner of the file, by the owner of the directory, and by a privileged process.

RETURN VALUE

On success, zero is returned. On error, -1 is returned, and *errno* is set appropriately.

ERRORS

EACCES

Search permission is denied for one of the directories in the path prefix of *path*. (See also **path_resolution(7)**.)

EBADF

fd is bad.

EFAULT

Bad address.

ELOOP

Too many symbolic links encountered while traversing the path.

ENAMETOOLONG

File name too long.

ENOENT

A component of *path* does not exist, or *path* is an empty string.

ENOMEM

Out of memory (i.e., kernel memory).

ENOTDIR

A component of the path prefix of *path* is not a directory.

EOVERFLOW

(**stat()**) *path* refers to a file whose size cannot be represented in the type *off_t*. This can occur when an application compiled on a 32-bit platform without *-D_FILE_OFFSET_BITS=64* calls **stat()** on a file whose size exceeds $(2^{31}-1)$ bits.

CONFORMING TO

These system calls conform to SVr4, 4.3BSD, POSIX.1-2001.

Use of the *st_blocks* and *st_blksize* fields may be less portable. (They were introduced in BSD. The interpretation differs between systems, and possibly on a single system when NFS mounts are involved.)

POSIX does not describe the **S_IFMT**, **S_IFSOCK**, **S_IFLNK**, **S_IFREG**, **S_IFBLK**, **S_IFDIR**, **S_IFCHR**, **S_IFIFO**, **S_ISVTX** bits, but instead demands the use of the macros **S_ISDIR()**, etc. The **S_ISLNK()** and **S_ISSOCK()** macros are not in POSIX.1-1996, but both are present in POSIX.1-2001; the former is from SVID 4, the latter from SUSv2.

Unix V7 (and later systems) had **S_IREAD**, **S_IWRITE**, **S_IEXEC**, where POSIX prescribes the synonyms **S_IRUSR**, **S_IWUSR**, **S_IXUSR**.

Other Systems

Values that have been (or are) in use on various systems:

hex	name	ls	octal	description
f000	S_IFMT		170000	mask for file type
0000			000000	SCO out-of-service inode; BSD unknown type; SVID-v2 and XPG2 have both 0 and 0100000 for ordinary file

1000	S_IFIFO	p	010000	FIFO (named pipe)
2000	S_IFCHR	c	020000	character special (V7)
3000	S_IFMPC		030000	multiplexed character special (V7)
4000	S_IFDIR	d/	040000	directory (V7)
5000	S_IFNAM		050000	XENIX named special file with two subtypes, distinguished by <i>st_rdev</i> values 1, 2
0001	S_INSEM	s	000001	XENIX semaphore subtype of IFNAM
0002	S_INSHD	m	000002	XENIX shared data subtype of IFNAM
6000	S_IFBLK	b	060000	block special (V7)
7000	S_IFMPB		070000	multiplexed block special (V7)
8000	S_IFREG	-	100000	regular (V7)
9000	S_IFCMP		110000	VxFS compressed
9000	S_IFNWK	n	110000	network special (HP-UX)
a000	S_IFLNK	l@	120000	symbolic link (BSD)
b000	S_IFSHAD		130000	Solaris shadow inode for ACL (not seen by userspace)
c000	S_IFSOCK	s=	140000	socket (BSD; also "S_IFSOC" on VxFS)
d000	S_IFDOOR	D>	150000	Solaris door
e000	S_IFWHT	w%	160000	BSD whiteout (not used for inode)
0200	S_ISVTX		001000	sticky bit: save swapped text even after use (V7) reserved (SVID-v2) On non-directories: don't cache this file (SunOS) On directories: restricted deletion flag (SVID-v4.2)
0400	S_ISGID		002000	set-group-ID on execution (V7) for directories: use BSD semantics for propagation of GID
0400	S_ENFMT		002000	System V file locking enforcement (shared with S_ISGID)
0800	S_ISUID		004000	set-user-ID on execution (V7)
0800	S_CDF		004000	directory is a context dependent file (HP-UX)

A sticky command appeared in Version 32V AT&T UNIX.

NOTES

Since kernel 2.5.48, the *stat* structure supports nanosecond resolution for the three file timestamp fields. Glibc exposes the nanosecond component of each field using names either of the form *st_atim.tv_nsec*, if the **_BSD_SOURCE** or **_SVID_SOURCE** feature test macro is defined, or of the form *st_atimensec*, if neither of these macros is defined. On file systems that do not support sub-second timestamps, these nanosecond fields are returned with the value 0.

On Linux, **lstat()** will generally not trigger automounter action, whereas **stat()** will.

For most files under the */proc* directory, **stat()** does not return the file size in the *st_size* field; instead the field is returned with the value 0.

Underlying kernel interface

Over time, increases in the size of the *stat* structure have led to three successive versions of **stat()**: *sys_stat()* (slot *__NR_oldstat*), *sys_newstat()* (slot *__NR_stat*), and *sys_stat64()* (new in kernel 2.4; slot *__NR_stat64*). The glibc **stat()** wrapper function hides these details from applications, invoking the most recent version of the system call provided by the kernel, and repacking the returned information if required for old binaries. Similar remarks apply for **fstat()** and **lstat()**.

EXAMPLE

The following program calls **stat()** and displays selected fields in the returned *stat* structure.

```
#include <sys/types.h>
#include <sys/stat.h>
#include <time.h>
#include <stdio.h>
#include <stdlib.h>

int
main(int argc, char *argv[])
{
    struct stat sb;

    if (argc != 2) {
        fprintf(stderr, "Usage: %s <pathname>\n", argv[0]);
        exit(EXIT_FAILURE);
    }

    if (stat(argv[1], &sb) == -1) {
        perror("stat");
        exit(EXIT_SUCCESS);
    }

    printf("File type:          ");

    switch (sb.st_mode & S_IFMT) {
    case S_IFBLK: printf("block device\n");      break;
    case S_IFCHR: printf("character device\n");    break;
    case S_IFDIR: printf("directory\n");           break;
    case S_IFIFO: printf("FIFO/pipe\n");           break;
    case S_IFLNK: printf("symlink\n");             break;
    case S_IFREG: printf("regular file\n");         break;
    case S_IFSOCK: printf("socket\n");             break;
    default:      printf("unknown?\n");            break;
    }

    printf("I-node number:      %ld\n", (long) sb.st_ino);

    printf("Mode:                %lo (octal)\n",
           (unsigned long) sb.st_mode);

    printf("Link count:          %ld\n", (long) sb.st_nlink);
    printf("Ownership:          UID=%ld  GID=%ld\n",
           (long) sb.st_uid, (long) sb.st_gid);

    printf("Preferred I/O block size: %ld bytes\n",
           (long) sb.st_blksize);
    printf("File size:              %lld bytes\n",
           (long long) sb.st_size);
    printf("Blocks allocated:       %lld\n",
           (long long) sb.st_blocks);

    printf("Last status change:     %s", ctime(&sb.st_ctime));
}
```

```
printf("Last file access:      %s", ctime(&sb.st_atime));  
printf("Last file modification: %s", ctime(&sb.st_mtime));  
  
exit(EXIT_SUCCESS);  
}
```

SEE ALSO

access(2), chmod(2), chown(2), fstatat(2), readlink(2), utime(2), capabilities(7), symlink(7)

COLOPHON

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