B.4 Branch-and-bound algorithm

Algorithm B.2 Branch-and-bound for integer linear programming.

Integer linear programming algorithm.

```
struct node t{
                            /* Constraints. */
    int
                m
    int
                            /* Decision variables. */
                n
    int
                k
                            /* Parent branches on x_k. */
    int
                h
                            /* Branch on x_h. */
    double
                xh
                            /* x_h. */
    double
                ak
                            /* Parent a_k. */
    double
                bk
                            /* Parent b_k. */
                            /* Lower bounds. */ } + ๑ เ ๖ธกุลท
/* Upper bounds. */
    double
                min[n]
    double
                max[n]
                            /* A. */
    double
                a[m][n]
    double
                b[m];
                            /* b. */
    double
                            /* x. */
                x[n];
                            /* c. */
    double
                c[n];
    double
                            /* z. */
                Z;
}
                                                (Skapar en Mihal
(Nod med Strucku ovan)
function initial\_node(m, n, a, b, c)
begin
    auto p = allocate memory for a node
    p.a = new double [m+1][n+1]
    p.b = new double [m+1]
    p.c = new double [n+1]
    p.x = new double [n+1]
    p.min = new double [n]
    p.max = new double [n]
    p.m = m
    p.n = n
    copy a, b, and c parameters to p
    for (i = 0; i < n; i = i + 1)
        p.min[i] = -\infty
        p.max[i] = +\infty
    return p
end
```

```
function extend (p, m, n, a, b, c, k, ak, bk)
begin
           q = allocate memory for a node
   auto
   int
           i,j
   q.k = k
   q.ak = ak
   q.bk = bk
   if ak > 0 and p.max[k] < \infty then
       q.m = p.m
   else if ak < 0 and p.min[k] > 0 then
        q.m = p.m
   else
       q.m = p.m + 1
   q.n = p.n
   q.h = -1
   q.a = new double [q.m+1][q.n+1] // note normally q.m > m
   q.b = new double [q.m+1]
   q.c = new double [q.n+1]
   q.x = new double [q.n+1]
   q.min = new double [n]
   q.max = new double [n]
   copy p.min and p.max to q // each element and not only pointers
   copy m first rows of parameter a to q.a // each element
   copy m first elements of parameter b to q.b
   copy parameter c to q.c // each element
   if ak > 0 then
       if q.max[k] = \infty or bk < q.max[k] then
            q.max[k] = bk
   else if q.min[k] = -\infty or -bk > q.min[k] then
       q.min[k] = -bk
   for (i = m, j = 0; j < n; j = j + 1) {
       if q.min[j] > -\infty then
            q.a[i][j] = -1
           q.b[i] = -q.min[j]
           i += 1
       if q.max[j] < \infty then
           q.a[i][j] = 1
           q.b[i] = q.max[j]
           i += 1
   }
   return q
end
```

```
function is_integer(xp)
begin
    // xp is a pointer to a double
    double x = *xp
    double r = round(x) // ISO C lround
    if |r-x| < \epsilon then
        *xp = r
        return 1
    else
        return 0
end
function integer(p)
begin
    int i
    for (i = 0; i < p.n; i = i + 1)
        if !is_integer(&p.x[i]) then
            return 0
    return 1
end
procedure bound(p, h, zp, x)
    // zp is a pointer to max z found so far
    if p.z > *zp then
        zp = p.z
        copy each element of p.x to x // save best x
        remove and delete all nodes q in h with q.z < p.z
end
function isfinite(x)
begin
    // ISO C function
    if x is a NaN or |x| = \infty then
        \textbf{return} \ 0
    else
        return 1
end
```

```
begin
    double min, max
   if q.z < z then
        return 0
   for (h = 0; h < q.n; h = h + 1)
        \textbf{if} ~! \textit{is\_integer}(\&q.x[h]) ~\textbf{then} \\
            if q.min[h] = -\infty then
                min = 0
            else
                min = q.min[h]
            max = q.max[h]
            if \lfloor q.x[h] \rfloor < min \text{ or } \lceil q.x[h] \rceil > max \text{ then}
            q.h = h
            q.xh = q.x[h]
            delete each of a,b,c,x of q // or recycle in other way
   return 0
end
procedure succ(p, h, m, n, a, b, c, k, ak, bk, zp, x)
   auto q = extend(p,m,n,a,b,c,k,ak,bk)
   if q = null then 3 gith clet => return
                                                      foraldom med
        return
                                                       Myn villuor
   q.z = simplex(q.m, q.n, q.a, q.b, q.c, q.x, 0)
                                          hollar om qz inte
air so ether Nav => fanns inte en losning.
   if is finite (q.z) then
        if integer(q) then
            bound(q,h,zp,x)
                                          Om q. 2 ar en int=>
        else if branch(q, *zp) then
                                             di keur vi kara bort en
            add q to h
                                              massa saker, vill inte
            return
                                              under som dessa.
   delete q
                                          Annurs Kollar vi om vi
end
                                             Skn forgrena
```

function branch(q, z)

```
function intopt(m, n, a, b, c, x)
begin
     auto p = initial node(m,n,a,b,c)
     \mathbf{set} \ \mathbf{h} = \{p\} \quad -
     double z = -\infty // best integer solution found so far
     p.z = simplex(p.m, p.n, p.a, p.b, p.c, p.x, 0)
     if integer(p) or !isfinite(p.z) then
           z = p.z
                                          > so or Naw
           if integer(p) then
                 copy p.x to x
          delete p

return z

nch(p,z) \leftarrow clelw \cup PP

duvi brancher

ile h \neq \emptyset

take p from h

succ(p, h, m, n, a, b, c, p.h, 1, [p.xh], &z, x)

succ(p, h, m, n, a, b, c, p.h, -1, -[p.xh], &z, x)

Shuper vi ett niger och

delete p

vänstr barn
     branch(p,z) \leftarrow delay upp
while h \neq \emptyset
     while h \neq \emptyset
     if z = -\infty then
           return NaN // not-a-number
     else
           return z
end
```