Number Theory 2.7 CRT $n = p_1^{e_1} \cdots p_r^{e_r}$ The Well-Ordering Property $\emptyset \neq S \subseteq \mathbb{N} \Rightarrow \min S \in S$ Let $N_i = n/n_i$ for every $i \in [k]$, $\exists s_i, t_i, N_i s_i + n_i t_i = 1$ **Division Algorithm** Let $b = b_1 N_1 s_1 + \dots + b_k N_k s_k$, then 1.3 $x \equiv b \pmod{n}$ a = bq + r $0 \le r < b$ for unique q, r2.8 DLOG & CDH 1.4 Ideal of $\mathbb Z$ A nonempty set $I\subseteq \mathbb Z$ such that $f_{\text{DLOG}}(q, G, g; h) = \log_q h, f_{\text{CDH}}(q, G, g; A, B) = g^{ab}$ $a,b\in I\Rightarrow a+b\in I$ 2.9 Diffie-Hellman Key Exchange $a \in I, r \in \mathbb{Z} \Rightarrow ra \in I$ Alice: $a \leftarrow \mathbb{Z}_q, A = g^a$, send (q, G, g, A) to Bob I is an ideal of $\mathbb{Z} \Leftrightarrow I = d\mathbb{Z}$ Bob: $b \leftarrow \mathbb{Z}_q$, $B = g^b$, send B to Alice; output $k = A^b$ $I_1 + I_2 = \{x + y : x \in I_1, y \in I_2\}$ Alice: output $k = B^a$ $a\mathbb{Z} + b\mathbb{Z} = \gcd(a, b)\mathbb{Z}$ 3 Group Theory **Great Common Divisor** 1.5 3.1 Group gcd(a,b) = as + btClosure $\forall a, b \in G, a \star b \in G$ **Euler's Phi Function Associative** $\forall a, b, c \in G, a \star (b \star c) = (a \star b) \star c$ $\Phi(n) = |\mathbb{Z}_n^*|, \forall n \in \mathbb{Z}^+$ **Identity** $\exists e \in G, \forall a \in G, a \star e = e \star a = a$ If $n = p_1^{e_1} \cdots p_r^{e_r}$, then Inverse $\forall a \in G, \exists b \in G, a \star b = b \star a = e$ $\Phi(n) = \Phi(p_1^{e_1}) \cdots \Phi(p_r^{e_r}) = n(1 - p_1^{-1}) \cdots (1 - p_r^{-1})$ Commutative (Abelian Group) $\forall a, b \in G, a \star b = b \star a$ If n = pq, then 3.2 Field $(\mathbb{F},+,\cdot)$ 3.3 Polynomial $\Phi(n) = (p-1)(q-1)$ The Set \mathbb{Z}_n and \mathbb{Z}_n^* Let $f(X) = f_t X^t + \dots + f_1 X + f_0 \in \mathbb{Z}_p[X]$ and $\alpha \in \mathbb{Z}_p$, then $\exists q(X) = q_{t-1} X^{t-1} + \dots + q_0 \in \mathbb{Z}_p[X]$ s.t. $\mathbb{Z}_n = \{[0]_n, [1]_n, \cdots, [n-1]_n\}$ $\mathbb{Z}_n^* = \{[a]_n \in \mathbb{Z}_n : \gcd(a, n) = 1\}$ Euler's Theorem $f(X) = (X - \alpha)q(X) + f(\alpha)$ Let $n \geq 1$ and $\alpha \in \mathbb{Z}_n^*$, then $\alpha^{\Phi(n)} = 1$ $q_{t-1} = f_t$ 1.9 Fermat's Little Theorem If p is a prime and $\alpha \in \mathbb{Z}_p$, then $\alpha^{p-1} = 1$ $q_{t-2} = f_{t-1} + f_t \alpha$ 1.10 Wilson's Theorem If p is a prime, then $(p-1)! \equiv -1 \pmod{p}$ $q_0 = f_1 + f_2 \alpha + \dots + f_t \cdot \alpha^{t-2}$ $f(X) \in \mathbb{Z}_p[X] \text{ has } \leq \deg(f) \text{ roots in } \mathbb{Z}_p$ Cryptography 2.1 RSA 3.4 Order $(pk, sk) \leftarrow \mathbf{Gen}(1^n)$: The order of a group G is the cardinality of G. Choose two *n*-bit primes $p \neq q$, N = pqWhen $|G| < \infty, \forall a \in G$, the order of a is the least integer Choose e, d s.t. $0 \le e, d < \Phi(N), \gcd(e, \Phi(N)) = 1$ l > 0 s.t. $a^l = 1$ (la = 0 for additive group) $d = e^{-1} \bmod \Phi(N)$ $\forall a \in G, a^{|G|} = 1$ output pk = (N, e), sk = (N, d)3.5 Cyclic Group $c \leftarrow \mathbf{Enc}(pk, m)$: Abelian group (G, \cdot) is a cyclic group if $\exists g \in G$ s.t. $G = \langle g \rangle$ output $c = m^e \mod N, 0 \le c < N$ Combinatorics $m \leftarrow \mathbf{Dec}(sk, c)$: **Functions** output $m = c^d \mod N, 0 \le m < N$ Let $A, B \neq \emptyset$ be two sets. A function(map) $f: A \rightarrow B$ assigns Arithmetic Operations a unique element $b \in B$ for all $a \in A$ $a = (a_{k-1} \cdots a_1 a_0)_2, \ b = (b_{l-1} \cdots b_1 b_0)_2$ $\ell(a) = \begin{cases} \lfloor \log_2(|a|) \rfloor + 1, & a \neq 0, \\ 1, & a = 0 \end{cases}$ $k = \ell(a), l = \ell(b)$ **injective** $f(a) = f(b) \Rightarrow a = b$ surjective f(A) = Bbijective injective and surjective Cantor's Diagonal Argument Addition & Subtraction a + b or a - b: O(k) $|A| \neq |\mathbb{Z}^+|$ Multiplication a * b: $O(k^2)$ Cantor's Theorem 4.3**Division** a/b: $O((k-l+1) \cdot l)$ Let A be any set, then $|A| < |\mathcal{P}(A)|$ Arithmetic Module N $(a \pm b) \mod N$: $O(\ell(N))$, 4.4 The Halting Problem (ab) mod N: $O(\ell(N)^2)$ There is no Turing machine computing Square-and-Multiply $\mathbf{HALT}(P,I) = \begin{cases} \text{``halts''} & \text{if } P(I) \text{ halts;} \\ \text{``loops forever.''} & \text{if } P(I) \text{ loops forever.} \end{cases}$ Square $x_0 = a$ $x_{k-1} = x_{k-2}^2 \bmod N = a^{2^{k-1}} \bmod N$ Multiply $a^e \bmod N = (x_0^{e_0} \cdot x_1^{e_0} \cdots x_{k-1}^{e_0}) \bmod N$ Countable an Uncountable A set A is countable if $|A| < \infty$ or $|A| = |\mathbb{Z}^+|$; 2.3 EA otherwise, it is said to be uncountable Compute $d = \gcd(a, b)$ 4.6 Schröder-Bernstein Theorem 2.4 EEA If $|A| \leq |B|$ and $|B| \leq |A|$, then |A| = |B|Compute $d = \gcd(a, b) = as + bt$: $O(\ell(a)\ell(b))$ $\aleph_0 = |\mathbb{Z}^+| < |\mathcal{P}(\mathbb{Z}^+)| = 2^{\aleph_0} = |[0,1)| = |(0,1)| = |\mathbb{R}| = c$ 2.5 Prime Number Theorem 4.7 Permutations of Set For $x \in \mathbb{R}^+$, $\pi(x) = \sum_{p \le x} \text{Numbers of primes}$ An *n*-element set has $P(n,r) = \frac{n!}{(n-r)!}$ and has n^r different $|\mathbb{P}_n| \ge \frac{2^n}{n \ln 2} (\frac{1}{2} + O(\frac{1}{n})) \text{ when } n \to \infty$ 2.6 Linear Congruence Equation r-permutations with repetition. 4.8 Combinations of Set $\binom{n}{r} = \frac{n!}{r!(n-r)!}$ $ax \equiv b \pmod{n}$

of r-combinations of an n element set with repetition, # of natural number solutions of the equation

$$x_1 + x_2 + \dots + x_n = r$$
 are $\binom{n+r-1}{r}$

4.9 Multiset

 $A = \{n_1 \cdot a_1, n_2 \cdot a_2, \dots, n_k \cdot a_k\}$ is an $(n_1 + n_2 + \dots + n_k)$ -multiset

4.10 Permutations of Multiset

A has exactly $\frac{(n_1 + n_2 + \cdots + n_k)!}{n_1! n_2! \cdots n_k!}$ permutations

4.11 Combination of Multiset

An r-subset(multiset) of A is r-combination of A

4.12 Shortest Path

of shortest paths from (0,0) to (p,q) is $\frac{(p+q)!}{p!a!}$

There is a T-route from $A = (a, \alpha)$ to $B = (b, \beta)$ only if (1) b > a; (2) $b - a \ge |\beta - \alpha|$ (3) $2 | (b + \beta - a - \alpha)$

4.14 Numbers of T-Routes

of T-routes from $A = (a, \alpha)$ to $B = (b, \beta)$ is

$$\frac{(b-a)!}{(\frac{b-a}{2} + \frac{\beta-\alpha}{2})!(\frac{b-a}{2} - \frac{\beta-\alpha}{2})!}$$

of T-routes that intersect the x-axis is

$$\frac{(b-a)!}{(\frac{b-a}{2}+\frac{\beta+\alpha}{2})!(\frac{b-a}{2}-\frac{\beta+\alpha}{2})!}$$

4.15 Solution of Bertrand's Ballot Problem The sequence $x_1x_2 \dots x_{2n}$ is a ballot

The probability that A never trials B is $p_n = C_n / {2n \choose n}$

4.16 Catalan Number

of solutions of the equation system

$$\begin{cases} x_1 + x_2 + \dots + x_{2n} = n \\ x_1 + x_2 + \dots + x_i \le i/2, i = 1, 2, \dots, 2n - 1 \\ x_i \in \{0, 1\}, i = 1, 2, \dots, 2n \end{cases}$$

$$C_n = \frac{(2n)!}{n!(n+1)!}$$

4.17 Inverse Binomial Transform The binomial transform of $\{a_n\}_{n\geq s}$ is $\{b_n\}_{n\geq s}$ s.t.

$$b_n = \sum_{k=s}^{s} \binom{n}{k} a_k$$

The inverse binomial transform of $\{b_n\}_{n\geq s}$ is $\{a_n\}_{n\geq s}$ s.t.

$$a_n = \sum_{k=1}^{s} (-1)^{n-k} \binom{n}{k} b_k$$

4.18 Distribution Problems Type 1 n labeled $\rightarrow k$ labeled: $|S| = k^n$

$$n \to i: \ N_1 = \frac{n!}{n_1! n_2! \cdots n_k!}$$

Type 2 n unlabeled $\rightarrow k$ labeled: $|S| = \binom{n+k-1}{n}$

Type 3 n labeled $\rightarrow k$ unlabeled: $|S| = \sum_{i=1}^{n} S_2(n, j)$

Type 4 n unlabeled $\rightarrow k$ unlabeled: $|S| = \sum_{i=1}^{n} p_j(n)$

4.19 Stirling number of the second kind $S_2(n, j)$

$$S_2(n,j) = \frac{1}{j!} \sum_{i=0}^{j-1} (-1)^i \binom{j}{i} (j-i)^n \text{ when } n \ge j \ge 1$$
$$S_2(n,j) = S_2(n-1,j-1) + jS_2(n-1,j)$$

4.20 Partitions of Integers

For $n \in \mathbb{Z}^+$, $p_j(n+j) = \sum_{k=1}^{J} p_k(n)$,

 $p_k(n) = p_{k-1}(n-1) + p_k(n-k)$ **4.21 Characteristic Roots**Characteristic equation: $r^k - c_1 r^{n-1} - c_2 r^{n-2} - \dots - c_k = 0$

4.22 LHRR

 $a_n = \sum_{i=1}^{n} c_i a_{n-i}$, where $n \geq k, \{c_i\}_{i=1}^k$ are constants, $c_k \neq 0$

No multiple roots $\{r_1,\ldots,r_k\}$: $x_n=\sum_{i=1}^n\alpha_ir_j^n$

Multiple roots $\{m_1 \cdot r_1, \dots m_t \cdot r_t\}$: $x_n = \sum_{j=1}^t (\sum_{\ell=0}^{m_j-1} \alpha_{j,\ell} n^\ell) r_j^n$

4.23 LNRR

 $a_n = \sum_{i=0}^{n} c_i a_{n-i} + F(n), \{c_i\}_{i=1}^k$ are constants, $c_k, F(n) \neq 0$

Particular Solutions $F(n) = (f_l n^l + \dots + f_1 n + f_0) s^n = f(n) s^n$

s: a root of characteristic equation, m: multiplication of s

$$x_n = (p_l n^l + \dots + p_1 n + p_0) s^n n^m$$

 $x_n=(p_ln^l+\cdots+p_1n+p_0)s^nn^m$ General Solutions Particular solution of LNRR + General solution of the associated LHRR

4.24 Generating Function

$$A(x) = \sum_{r=0}^{\infty} a_r x^r, B(x) = \sum_{r=0}^{\infty} b_r x^r$$

$$A(x) \pm B(x) = \sum_{r=0}^{\infty} (a_r \pm b_r) x^r, A(x) \cdot B(x) = \sum_{r=0}^{\infty} (\sum_{i=0}^{r} a_i b_{r-i}) x^r$$

A(x) has an inverse iff $a_0 \neq 0$.

4.25 $(1 + \alpha x)^u$

The extended binomial coefficient

$$(1+x)^u = \sum_{r=0}^{\infty} \binom{u}{r} x^r$$

Counting Combinations with GFs

$$a_r = |\{(r_1, \dots, r_n) : r_i \in R_i, r_1 + \dots + r_n = r\}|$$

$$a_r = |\{(r_1, \dots, r_n) : r_i \in R_i, r_1 + \dots + r_n = r\}|$$

$$\sum_{r=0}^{\infty} a_r x^r = \prod_{i=1}^n \sum_{r_i \in R_i} x^{r_i}$$

$$a_r = \sum_{r_1 \in R_1, \dots, r_n \in R_n, r_1 + \dots r_n = r} \frac{r!}{r_1! \dots r_n!}$$

4.27 Counting Permutations with GFs
$$a_r = \sum_{r_1 \in R_1, \dots, r_n \in R_n, r_1 + \dots r_n = r} \frac{r!}{r_1! \dots r_n!} \sum_{r=0}^{\infty} \frac{a_r}{r!} x^r = \prod_{i=1}^n \sum_{r_i \in R_i} \frac{x^{r_n}}{r_n!}$$

Partial Fraction Decomposition

Let Q(x), P(x) be two polynomial s.t. deg(Q) > deg(P). If $Q(x) = (1 - r_1 x)^{m_1} \cdots (1 - r_t x)^{m_t}$ for distinct non-zero numbers r_1, \ldots, r_t and integers $m_1, \ldots, m_t \geq 1$, then there exist unique coefficients $\{\alpha_{j,u}: j \in [t], u \in [m_j]\}$ such that

$$\frac{P(x)}{Q(x)} = \sum_{j=1}^{t} \sum_{u=1}^{m_j} \frac{\alpha_{j,u}}{(1 - r_j x)^u}$$

4.29 Principle of IE

Let S be a finite set, $A_1, A_2, A_n \subseteq S$, then

$$\left| \bigcup_{i=1}^{n} A_{i} \right| = \sum_{t=1}^{n} (-1)^{t-1} \sum_{1 \le i_{1} < \dots < i_{t} \le n} \left| A_{i_{1}} \cap \dots \cap A_{i_{t}} \right|$$

$$\left| \bigcap_{i=1}^{n} A_{i} \right| = \sum_{t=1}^{n} (-1)^{t-1} \sum_{1 \le i_{1} < \dots < i_{t} \le n} \left| A_{i_{1}} \cup \dots \cup A_{i_{t}} \right|$$

A cover of a finite set A is a family $\{A_1, A_2, \dots A_n\}$ of subsets of A such that $\bigcup_{i=1}^n A_i = A$.

4.31 Pigeonhole Principle

Let A be a set with $\geq N$ elements. Let $\{A_1, A_2, \dots, A_n\}$ be a cover of A, then $\exists k \in [n], |A_k| \geq \lceil N/n \rceil$.

5 Propositional Logic

5.1 Logical Operators

p	$\neg p$
Т	F
F	Т

p	q	$p \wedge q$	$p \lor q$	$p \rightarrow q$	$p \leftrightarrow q$
Т	Т	Т	Т	Т	Т
\mathbf{T}	F	F	${ m T}$	\mathbf{F}	F
\mathbf{F}	Τ	F	${ m T}$	${ m T}$	F
\mathbf{F}	F	F	F	${ m T}$	${ m T}$

5.2 Type of WFFs

Tautology: truth value is **T** for all truth assignment **Contradiction**: truth value is **F** for all truth assignment **Contingency**: neither tautology or contradiction

5.3 Logical Equivalence

Proving $A \equiv B$ (1) $A^{-1}(\mathbf{T}) = B^{-1}(\mathbf{T})$ (2) $A^{-1}(\mathbf{F}) = B^{-1}(\mathbf{F})$ (3) $A \leftrightarrow B$ is a tautology

5.4 Tautological Implications

Proving $A \Rightarrow B$ (1) $A^{-1}(\mathbf{T}) \subseteq B^{-1}(\mathbf{T})$ (2) $B^{-1}(\mathbf{F}) \subseteq A^{-1}(\mathbf{F})$ (3) $A \to B$ is a tautology (4) $A \land \neg B$ is a contradiction

5.5 Rules of Replacement

$$\neg (P \land Q) \equiv (\neg P) \lor (\neg Q) \quad \neg (P \lor Q) \equiv (\neg P) \land (\neg Q)$$

$$P \lor (P \land Q) \equiv P \land (P \lor Q) \equiv P$$

$$P \to Q \equiv \neg P \lor Q \quad P \leftrightarrow Q \equiv (P \to Q) \land (Q \to P)$$

6 Predicate Logic

6.1 From Natural Language to WFFs

 $\forall x (P(x) \to Q(x)) \ \exists x (P(x) \land Q(x))$

6.2 Type of WFFs

A WFF is **logically valid** if it is **T** in every interpretation A WFF is **unsatisfiable** if it is **F** in every interpretation A WFF is **satisfiable** if it is **T** in some interpretation

6.3 Logical Equivalence

 $A \equiv B$ iff $A \leftrightarrow B$ is logically valid

6.4 De Morgan's Laws for Quantifiers

 $\neg \forall x P(x) \equiv \exists x \neg P(x), \neg \exists x P(x) \equiv \forall x \neg P(x)$

6.5 Distributive Laws for Quantifiers

$$\forall x (P(x) \land Q(x)) \equiv \forall x P(x) \land \forall x Q(x)$$
$$\exists x (P(x) \lor Q(x)) \equiv \exists x P(x) \lor \exists x Q(x)$$

6.6 Rules of Substitution

 $(P) \land (Q) \Rightarrow P \land Q \qquad \qquad \text{Conjunction} \\ P \land Q \Rightarrow P \qquad \qquad \text{Simplification} \\ P \Rightarrow P \lor Q \qquad \qquad \text{Addition} \\ P \land (P \rightarrow Q) \Rightarrow Q \qquad \qquad \text{Modus Ponens} \\ \neg Q \land (P \rightarrow Q) \Rightarrow \neg P \qquad \qquad \text{Modus Tollens} \\ \neg P \land (P \lor Q) \Rightarrow Q \qquad \qquad \text{Disjunctive Syllogism} \\ (P \rightarrow Q) \land (Q \rightarrow R) \Rightarrow (P \rightarrow R) \qquad \text{Hypothetical Syllogism} \\ (P \lor Q) \land (\neg P \lor R) \Rightarrow Q \lor R \qquad \text{Resolution} \\ P \Rightarrow Q \rightarrow R \equiv P \land Q \Rightarrow R \qquad \qquad \text{Conclusion Premise}$

6.7 Tautological Implications

 $\forall x P(x) \lor \forall x Q(x) \Rightarrow \forall x (P(x) \lor Q(x))$ $\exists x (P(x) \land Q(x)) \Rightarrow \exists x P(x) \land \exists x Q(x)$ $\forall x (P(x) \rightarrow Q(x)) \Rightarrow \forall x P(x) \rightarrow \forall x Q(x)$ $\forall x (P(x) \rightarrow Q(x)) \Rightarrow \exists x P(x) \rightarrow \exists x Q(x)$ $\forall x (P(x) \leftrightarrow Q(x)) \Rightarrow \forall x P(x) \leftrightarrow \forall x Q(x)$ $\exists x (P(x) \leftrightarrow Q(x)) \Rightarrow \exists x P(x) \leftrightarrow \exists x Q(x)$ $\forall x (P(x) \rightarrow Q(x)) \land \forall x (Q(x) \rightarrow R(x)) \Rightarrow \forall x (P(x) \rightarrow R(x))$ $\forall x (P(x) \rightarrow Q(x)) \land P(a) \Rightarrow Q(a)$

6.8 Rules of Inference for \forall , \exists

 $\forall x P(x) \Rightarrow P(a)$ Universal Instantiation $P(a) \Rightarrow \forall x P(x)$ Universal Generalization $\exists x P(x) \Rightarrow P(a)$ Existential Instantiation $P(a) \Rightarrow \exists x P(x)$ Existential Generalization

7 Graph

7.1 Types of Graph

7.2 Handshaking Theorem

Let G = (V, E) be an undirected graph, then $2|E| = \sum_{v \in V} \deg(v)$ and $|\{v \in V : \deg(v) \text{ is odd}\}|$ is even. Let G = (V, E) be an directed graph, then

$$\sum_{v \in V} \deg^{-}(v) = \sum_{v \in V} \deg^{+}(v) = |E|$$

7.3 Edge Contraction

Let G = (V, E) be a simple graph and $e = \{u, v\} \in E$ Define G/e = (V', E'), where $V' = (V - \{u, v\}) \cup \{w\}$ and $E' = \{e' \in E : e' \cap e = \emptyset\} \cup \{\{w, x\} : \{u, x\} \in E \text{ or } \{v, x\} \in E\}$