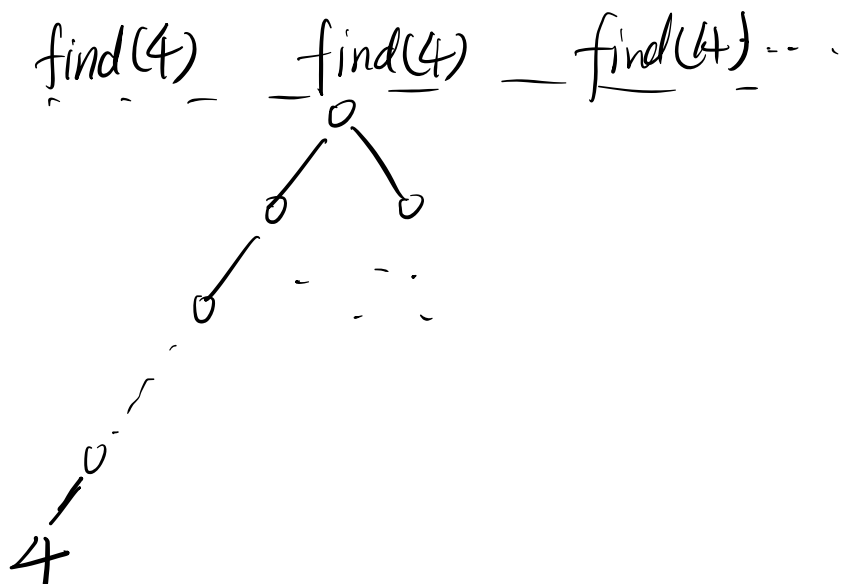


AVL

- ∴ $O(\lg n)$ in worst case
- ∴ complicated implementation
 - extra space h_L, h_R
 - non-adaptive.

Splay tree.

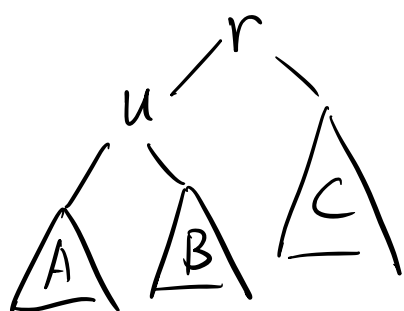
- ∴ $O(n)$ in worst case
- ∴ $O(\lg n)$ amortized cost
- ∴ easy implementation
 - no extra space
 - adaptive



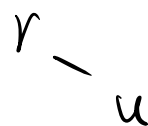
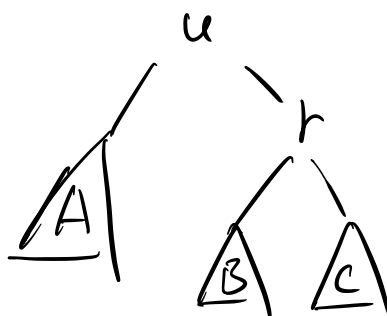
splay(u) = move u to the root by rotations.

If u is the root
done

else if u is a child of the root

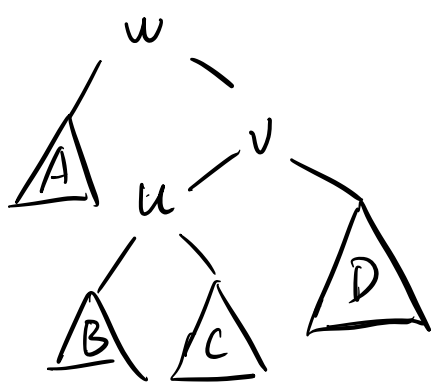


single
rotation

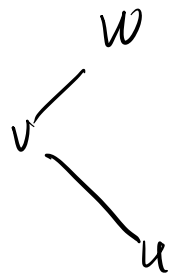
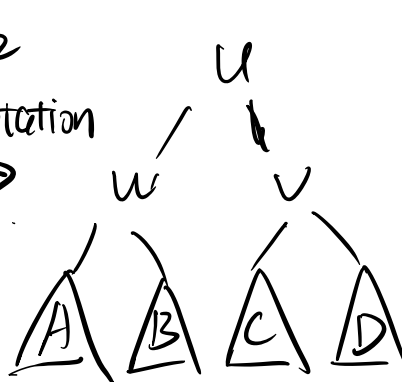


else // u has a grandparent.

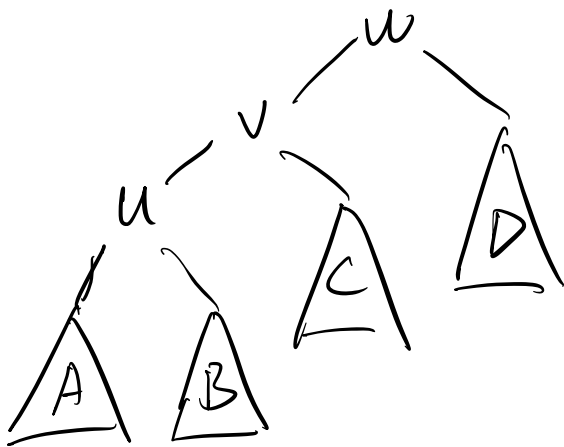
case 1. zig-zag



rotate u x 2
double rotation

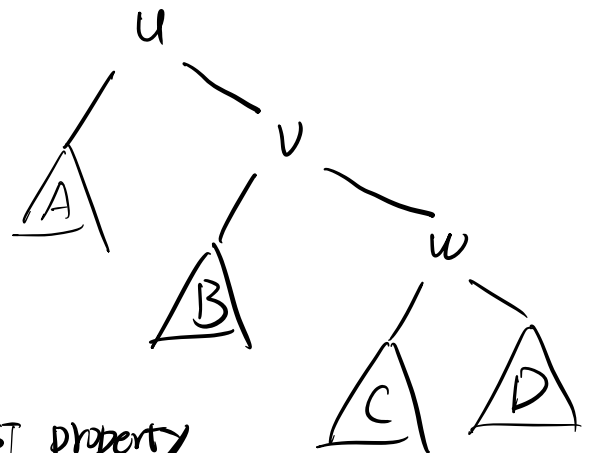


case 2 zig-zig



rotate v
rotate u

double rotation



preserve the BST property

Find key:

1. find as in BST
2. splay the node to the root

$$\begin{aligned}
 & \left\{ \begin{array}{l} r \\ \vdots \\ u \end{array} \right\} + \left(\begin{array}{l} c' \cdot \# \text{rotations} \\ + \\ c \cdot \# \text{rotations} \end{array} \right) \\
 & \left(\begin{array}{l} \# \text{rotations} \geq \ell/2 \\ c \cdot \# \text{rotations} \end{array} \right)
 \end{aligned}$$

Insert:

1. insert as in BST
2. splay the newly inserted node to the root $c \cdot \# \text{rotations}$

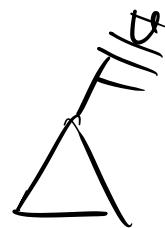
Delete:

$c \cdot \# \text{rotations}$

1. find the node u to be deleted and splay it to the root.

2. if u has most one child

delete u directly

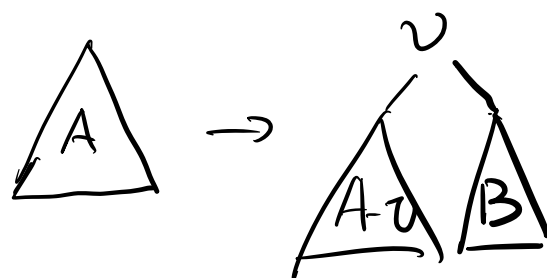
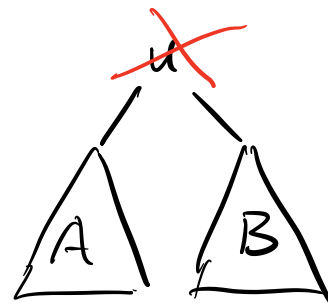


3. else // u has two children

4. delete u

5. splay the largest element v of A to the root of A

6. attach B to v



Observation

actual cost of each operation is $c \cdot \# \text{rotations}$
(find, Ins, Del) \downarrow some constant

$$\Phi(n) \geq \Phi(0)$$

$$\Phi(i) - \Phi(i-1) \leq c \lg n - c \cdot \# \text{rotations}$$

$$\hat{C}_i = C_i + \Phi(i) - \Phi(i-1) = c \cdot \cancel{\# \text{rotations}} + c \lg n - c \cdot \cancel{\# \text{rotation}} \\ = c \lg n = O(\lg n)$$

Given a tree T ,

• for each node $u \in T$

$$\text{rank} \quad \text{size}(u) = \# \text{ nodes in } T_u$$

$$r(u) = \lg(\text{size}(u))$$

$$\Phi(T) = \sum_{u \in T} r(u)$$

Claim: let T be a splay tree. Let $u \in T$ be a node.
let T' be the tree after $\text{splay}(u)$.

$$\Phi(T') - \Phi(T) \leq \underbrace{3c \cdot (r'(u) - r(u))}_{\substack{\text{rank of} \\ u \text{ in } T'}} - 2c \cdot (\# \text{rotations} - 1)$$

Find key:

1. find as in BST

$$\Delta_1 = 0$$

2. $\text{splay}(u)$

$$\Delta_2 = 3c \cdot (r'(u) - r(u)) - 2c \cdot (\# \text{rotations} - 1)$$

$$\leq r'(u) \leq \lg n$$

$$\text{amortized cost} = \underbrace{c \cdot \# \text{rotations}}_{\text{actual cost}} + \Delta_2 \leq 3c \cdot \lg n - c \cdot \# \text{rotations} + 2c$$

$$\leq 3c \cdot \lg n + 2c$$

$$= O(\lg n)$$

Insertion

1. insert as in BST

$$\Delta_1 \leq h \leq 2 \# \text{rotations}$$

2. $\text{splay}(u)$

$$\Delta_2 \leq 3c \cdot (r'(u) - r(u)) - 2c \cdot (\# \text{rotations} - 1)$$

h
 u

$$\text{amortized cost} = c \cdot \# \text{rotations} + A_1 + A_2$$

$$\leq 3c \cdot (r'(u) - r(u)) - (c-2) \cdot \# \text{rotations} + 2c$$

assume $c \geq 2$

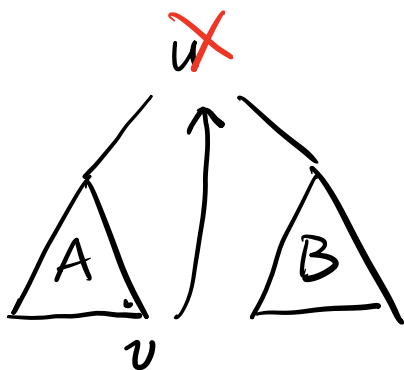
$$\leq 3c \cdot \frac{(r'(u) - r(u))}{\lg n} + 2c = O(\lg n)$$

Del:

1. find u and $\text{splay}(u)$
2. delete u
3. find v and $\text{splay}(v)$ (to the root of A)
4. attach B to v

$$\begin{aligned} \underline{A_1} &\leq 3c \cdot (r'(u) - r(u)) \\ \underline{A_2} &= -2c \cdot (\# \text{rotations} - 1) \\ &\quad \text{in step 1} \end{aligned}$$

$$\begin{aligned} \underline{A_3} &\leq 3c \cdot (r'(v) - r(v)) \\ \underline{A_4} &= -2c \cdot (\# \text{rotation} - 1) \\ &\quad \text{in step 3} \end{aligned}$$



$$r(v) \leq r(u)$$

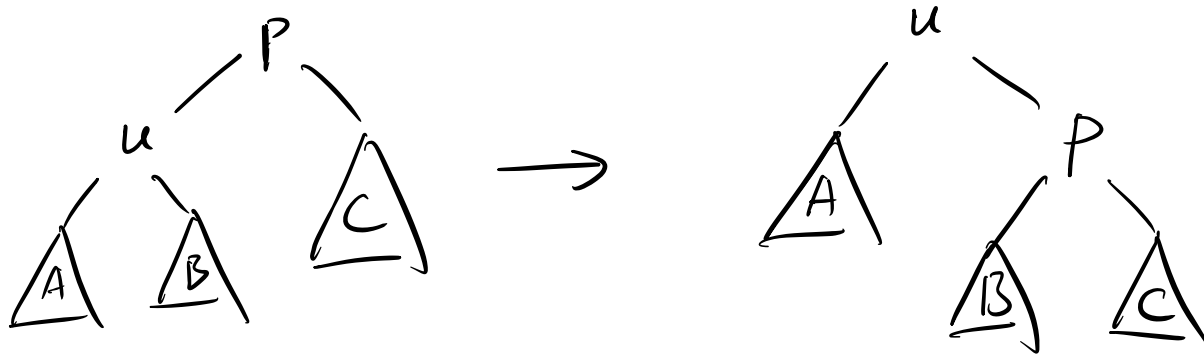
$$\Rightarrow A_2 + A_4 \leq 0$$

$$\begin{aligned} \text{amortized cost} &= \cancel{c} \cdot \# \text{rotations} + 3c \cdot (r'(u) - r(u)) + 3c \cdot (r'(v) - r(v)) \\ &\quad + \cancel{-2c} \cdot \# \text{rotations} + 4c \end{aligned}$$

$$\leq 6c \cdot \lg n + 4c \quad O(\lg n)$$

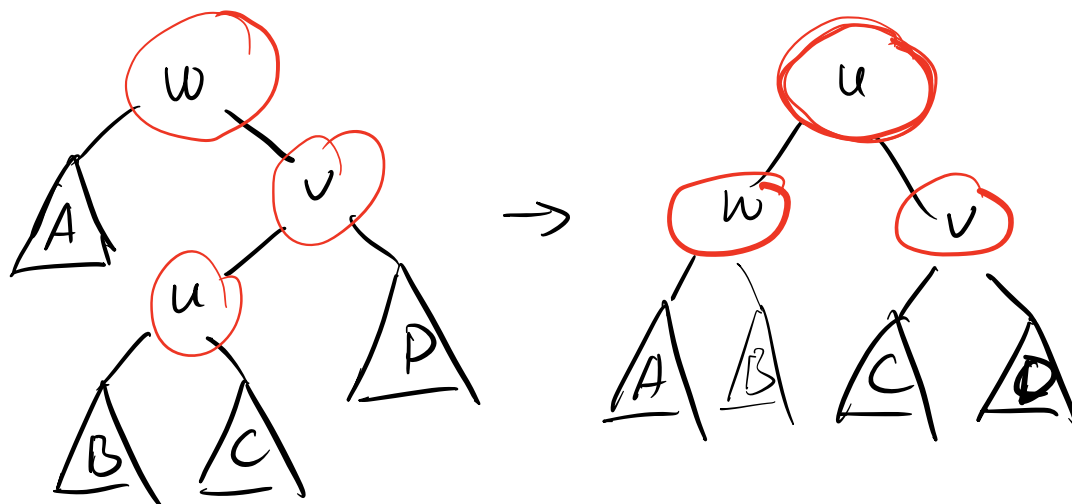
$$\Phi(T') - \Phi(T) \leq 3(r'(u) - r(u)) - 2(\# \text{rotations} - 1)$$

Case 1. single rotation



$$\begin{aligned} \frac{\Phi(T') - \Phi(T)}{C} &= r'(u) - r(u) + r'(p) - r(p) \\ &\leq r'(u) - r(u) \\ &\leq 3(r'(u) - r(u)) \end{aligned}$$

Case 2 Zig-zag



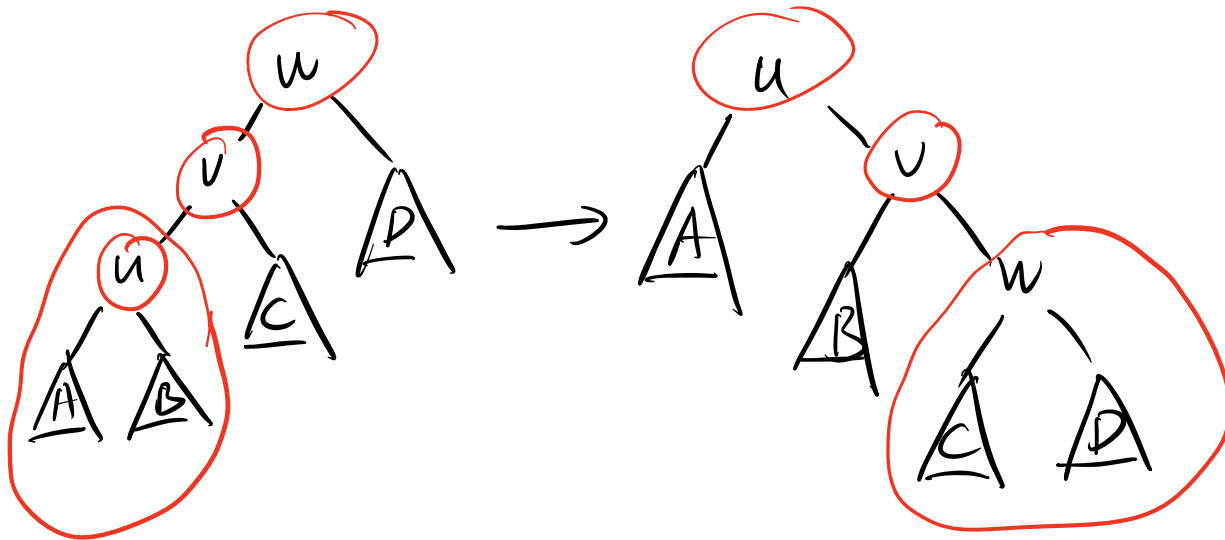
$$\begin{aligned} \frac{\Phi(T') - \Phi(T)}{C} &= \cancel{r'(w) - r(w)} + r'(v) - \cancel{r(v)} + r'(w) - \cancel{r(w)} \\ &\leq -r(u) + r'(v) - r(u) + r'(w) \quad r(w) > r(u) \\ &= -2r(u) + r'(v) + r'(w) \end{aligned}$$

$$\leq -2r(u) + 2r'(u) - 2 \leq 3(r'(u) - r(u)) - 2$$

$$\text{size}'(w) + \text{size}'(v) + 1 = \text{size}'(u)$$

$$r'(w) + r'(v) \leq 2r'(u) - 2$$

case 3. zig-zig



$$\begin{aligned} \frac{\Phi(T') - \Phi(T)}{c} &= \cancel{r(u)} - r(u) + \cancel{r'(v) - r(v)} + \cancel{r'(w) - r(w)} \\ &\leq -r(u) + r'(u) - r(u) + r'(w) \quad \because r'(v) \leq r'(u) \\ &\quad r(v) \geq r(u) \\ &= -2r(u) + r'(u) + r'(w) \\ &= -3r(u) + \cancel{r'(u)} + \cancel{r'(w)} + r(u) \leq 3(r'(u) - r(u)) - 2 \end{aligned}$$

$$\text{size}(u) + \text{size}'(w) = A + B + C + D + 2 \leq \text{size}'(u) - 2$$

$$r(u) + r'(w) \leq 2r'(u) - 2$$

splay(u):

$$\frac{\Phi(T') - \Phi(T)}{c} = 3 \cdot (r'(u) - r(u)) - 2 \cdot (\# \text{ double rotations})$$

$$\geq \# \text{ rotations} - 1$$

$$\leq 3 \cdot (r'(u) - r(u)) - 2 \cdot (\# \text{rotations} - 1)$$