

# 浙江大学 2017 - 2018 学年 春夏 学期

## 《数据库系统》课程期末考试试卷 (A 卷)

### 参考答案及评分细则

#### Answers of Problem 1:

(12 points, 3 points per part)

(1)            **select c1.cno, c1.name**  
                 **from (card as c1) natural join (detail as d1),**  
                 **detail as d2**  
                 **where c1.depart= 'CS' and d2.cno ='c0002' and**  
                 **d1.cdate = d2.cdate and d1.pno=d2.pno;**

*another answer:*

**select c1.cno, c1.name**  
**from card as c1, detail as d1, detail as d2**  
**where c1.cno=d1.cno and**  
**c1.depart= 'CS' and d2.cno ='c0002'**  
**d1.cdate = d2.cdate and d1.pno=d2.pno;**

评分细则:

每个 where 条件错扣 1 分, select 错扣 1 分, 直至扣完

(2)             $\Pi_{c1.cno, c1.name} (\sigma_{d1.cdate=d2.cdate \wedge d1.pno=d2.pno}$   
                  $( (\sigma_{c1.depart='CS'} (\rho_{c1}(card))) \bowtie \sigma_{(\rho_{d1}(detail)) \times}$   
                  $(\sigma_{d2.cno='c0002'} (\rho_{d2}(detail))) ) )$

(3)            **select cno**  
                 **from detail natural join pos**  
                 **where year(detail.cdate)=2018**  
                 **group by cno**  
                 **having count(distinct campus)=1;**

*another answer:*

**select \***  
**from card c1**  
**where exists(**  
                 **select \***

```

        from (detail natural join pos) as r1
        where r1.cno=c1.cno )
and not exists(
    select *
    from    (detail natural join pos) as r1,
            (detail natural join pos) as r2
    where   r1.cno=c1.cno and r2.cno=c1.cno and
            year(r1.cdate)=2018 and year(r2.cdate)=2018
            and r1.campus<> r2.campus
    )

```

(4)

```

select cno
from detail natural join pos
where pos.campus='紫金港' and year(detail.cdate)=2018
group by pno
having sum(amount)>=all (
    select sum(amount)
    from detail natural join pos
    where pos.campus=' 紫金港' and year(detail.cdate)=2018
    group by pno
)

```

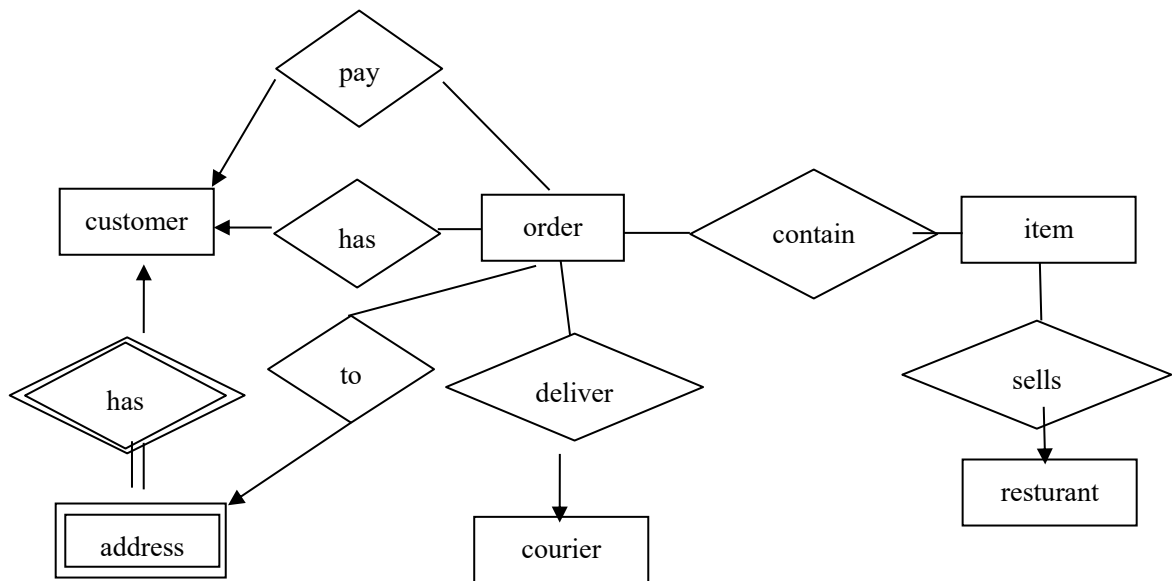
(5)

```

update card set balance = balance -20 where cno='c0002';
insert into detail(cno, pno, cdate, ctime, amount)
    values('c0002', 'p001','2018-07-02', '08:08:08', 20);
commit;

```

**Answers of Problem 2:**  
**(11 points)**



**Answers of Problem 3:**  
**(12 points, 3 points per part)**

(1)  $F_c = \{A \rightarrow C, C \rightarrow B, B \rightarrow DE\}$

(2)  $(B)^+ = (B, D, E)$

(3)  $R_1(B, D, E), F_1 = \{B \rightarrow DE\}$   
 $R_2(C, B), F_2 = \{C \rightarrow B\}$   
 $R_3(A, C), F_3 = \{A \rightarrow C\}$

(4) The decomposition is dependency preserving,  
 because  $(F_1 \cup F_2 \cup F_3)^+ = F^+$

**Answers of Problem 4:**

**(12 points, 4 points per part)**

```
(1) <campus_cards>
      <pos pno="p003">
        <campus> 玉泉 </campus>
        <location> 四食堂 </location>
      </pos>
      <card cno="c0003" >
        <name> 王浩</name>
        <depart> CS </depart>
        <balance> 300</balance>
        <detail pno=" p003">
          <cdate> 2018-07-03</cdate>
          <ctime> 08:10:10 </ctime>
          <amount>25 </amount>
          <remark>餐饮</remark>
        </detail>
      </card>
    </campus_cards>
```

**(2)**

**/campus\_cards/card[name="张帅"]/detail[amount=50] /id(pno) /location/text()**

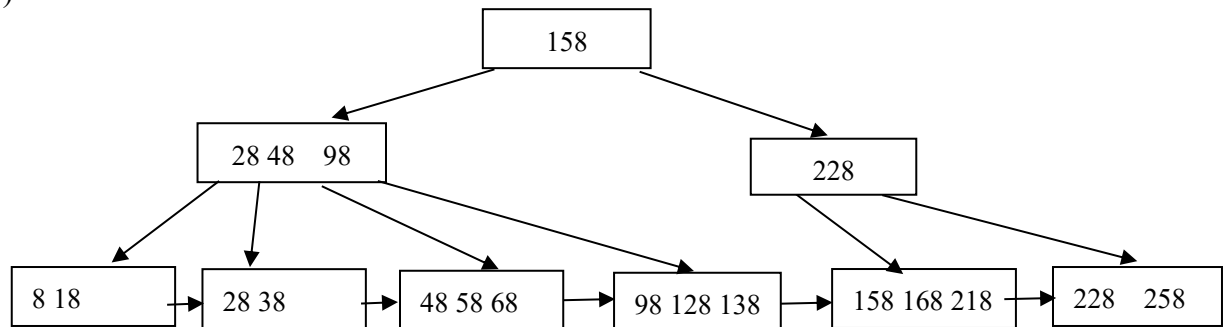
**(3)**

```
for $x in /campus_cards/card/detail
  $y in /campus_cards/card[name="张帅"]/detail
  where $x/@pno=$y/@pno and $x/cdate=$y/cdate and
    $x/@cno<>$y/@cno
  return  <cno> {$x/cno } </cno>
```

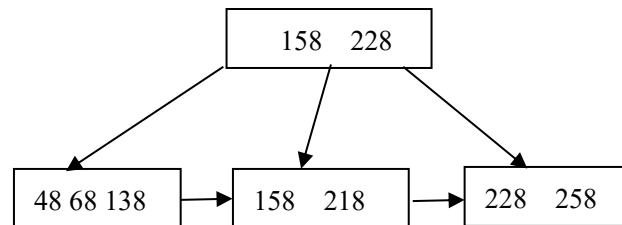
**Answers of Problem 5:**

**(12 points, 4 points per part)**

1)



2)



3)  $3+2+2+1=8$

**Answers of Problem 6:**

**(16 points, 4 points per part)**

1)  $(100000000 * 100000000) / (100 * 365)) * 3/12 = 684.93M$

- 2) Record number per block of card =  $4096/25 = 163$   
 Blocks of card =  $10000/163 = 61.3 \rightarrow 62$   
 Record number per block of detail =  $4096/29 = 141.24 \rightarrow 141$   
 Blocks of detail =  $10000000/141 = 70922$
- 3) Fan-out rate n of the B+-tree =  $(4096-4)/(5+4) + 1 = 455$
- 4) Min height of B+tree =  $\log_{455} (10000) \rightarrow 2$  (向上取整)  
 Max height of B+tree =  $\log_{228} (10000/2) + 1 = \rightarrow 2$  (向下取整)  
 So height of B+tree = 2

5) Cost for evaluating  $\sigma$  operation (2分, 各1分)

block transfer =  $62 t_T$

seek time =  $1 t_S$

cost for the natural join operation (2分,  $t_S$  和  $t_T$  各1分)

return number of  $\sigma$  name = '张帅' (card) =  $(10000/5000) = 2$

block number for each card cno in detail =  $(10000000/10000)/141 = 7.09 = 8$

cost for the natural join operation =  $2*(2 t_S + 2 t_T + 1 t_S + 8 t_T)$

=  $2*(3 t_S + 10 t_T) = 6 t_S + 20 t_T$

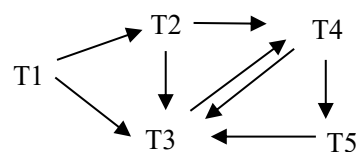
pipeline evaluation:

Total cost =  $(1 t_S + 62 t_T) + (6 t_S + 20 t_T) = 7 t_S + 82 t_T$

**Answers of Problem 7:**

**(12 points, 3 points per part)**

1)



2)

S is not serializable, because there are cycles in the graph :  $T3 \sim T4 \sim T3$ ;

T3~T4~T5~T3

3) no, because every schedule generated by 2PL is serializable.

4) w3(B) , or w4(B), , or w3(C)

**Answers of Problem 8:**

**(10 points, 2 points per part)**

1) 1006 (2分)

2) 1005 (2分)

3)

PageID	PageLSN	RecLSN
8001	1010	1010
8002	1014	1006
8003	1015	1015

4) “8002.1” = 55 (1分)

“8002.2” = 99 (1分)

5) 1017: <T3, 8002.1, 66>

1018: <T3, 8002.1, 55>

1019: <T3 abort>