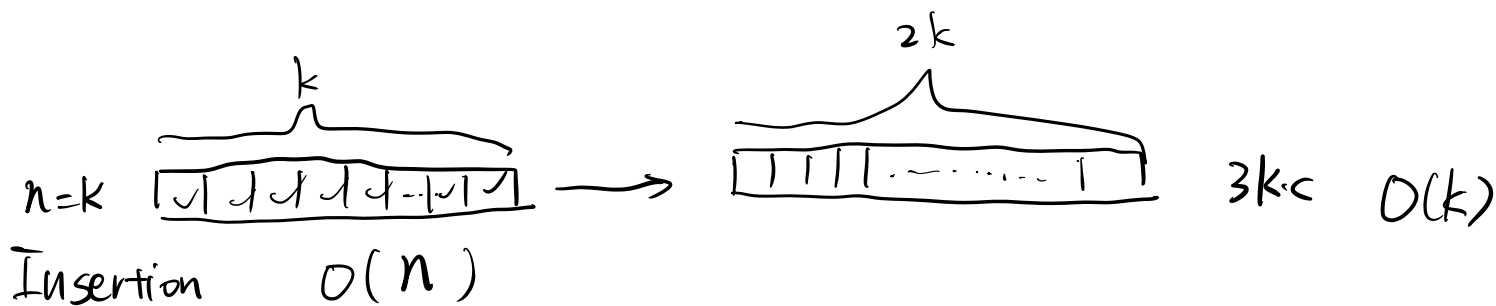
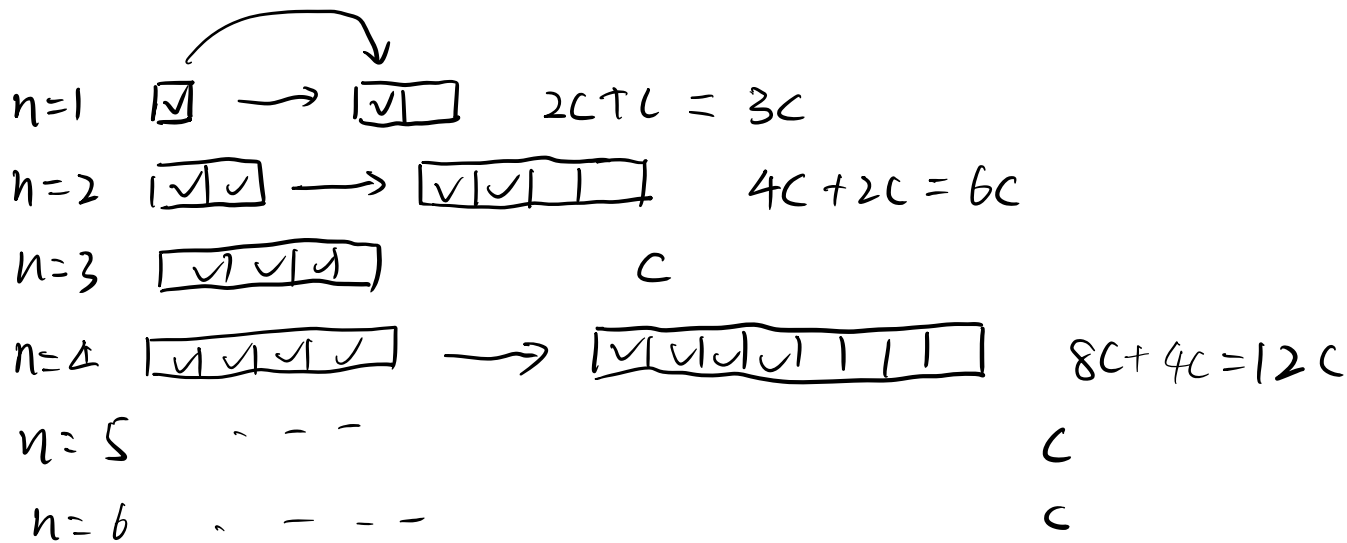


# dynamic array

- $A[i]$
- Insertion
- $O(n)$  space



$T(n)$ : cost of a sequence of  $n$  insertions in worst case. (from empty array)

$$\begin{aligned}
 T(n) &\leq Cn + 2^0 \cdot 3C + 2^1 \cdot 3C + \dots + 2^{(k)} \cdot 3C \\
 &\leq Cn + (2^{k+1} - 1) \cdot 3C \\
 &\leq Cn + (2 \cdot n - 1) \cdot 3C \\
 &\leq 7C \cdot n \quad O(n)
 \end{aligned}$$

largest  $k$   
 $2^k \leq n$

$$T(n)/n = O(1) \rightarrow \text{amortized cost}$$

Amortized cost  $\Lambda(n) = \frac{T(n)}{n}$  where  $T(n)$  is the worst-case

$$\text{Amortized cost } A(n) = \frac{T(n)}{n}$$

cost of all possible sequences of  $n$  operations.



similar

When there are multiple types of operations.

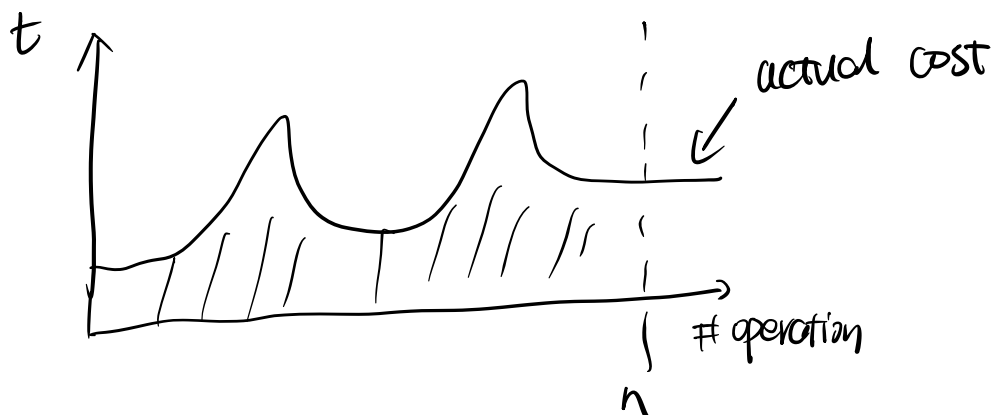
## Stack

	Worst-case	Amortized
- Push( $S, x$ )	$O(1)$	$O(1)$
- pop( $S$ ) if $S$ is empty error	$O(1)$	$O(1)$
- MULTIPOP( $S, k$ )	$O(k)$	$O(1)$
while $S$ is not empty and $k > 0$ pop( $S$ ) $k = k - 1$		

$$T(n) \leq O(n) + O(n) = O(n)$$

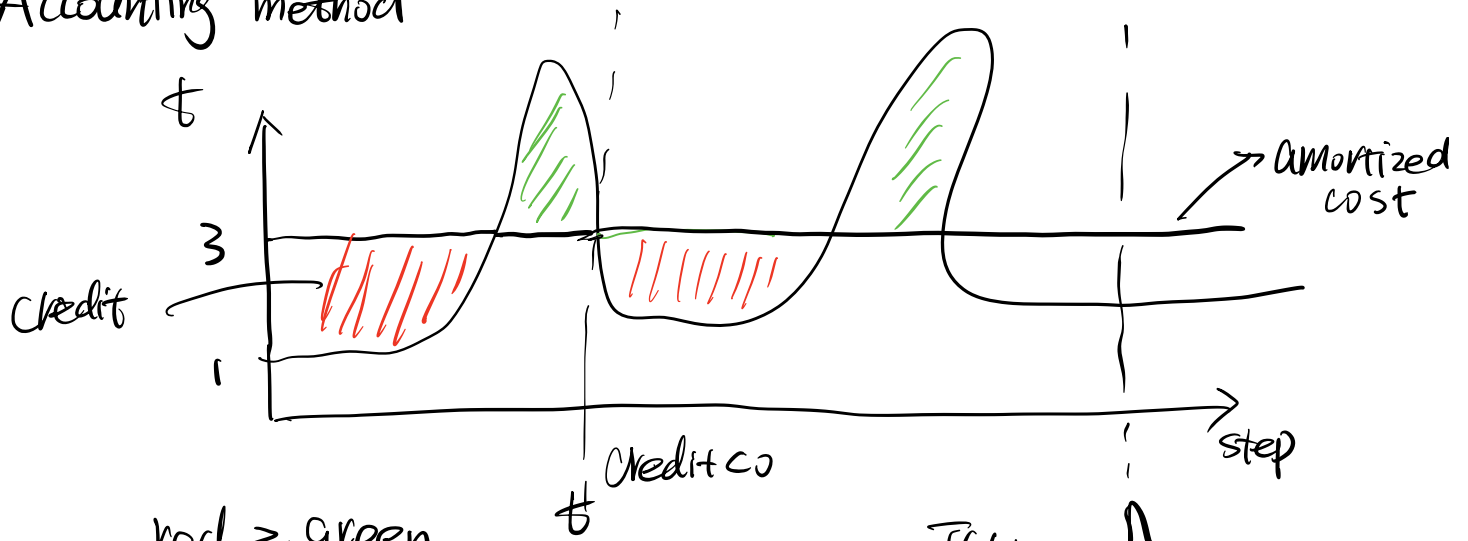
$$T(n)/n = O(1)$$

## Aggregate analysis



$$T(n)$$

## Accounting method



red  $\geq$  green

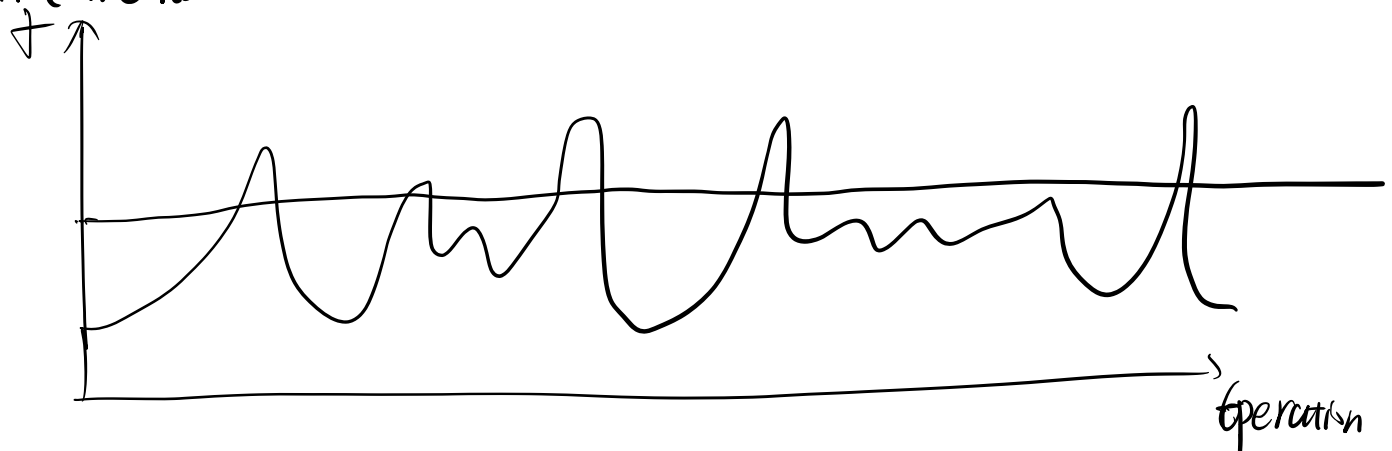
$\Rightarrow$  rect  $\geq$  curve

$\Rightarrow \sum \text{amortized cost} \geq \sum \text{actual cost}$

$$A(t) = \frac{\bar{n}(t)}{t}$$

	Actual cost	Amortized cost
PUSH(S, x)	$c$	$\geq c$ (associate the extra $c$ with $x$ )
POP(S)	$c$	0
Multipop(S, k)	$c \cdot \min( S , k)$	0

## Potential method.



hard to analyze directly that red  $\geq$  green

$\Phi(i) = \# \text{ credit left after the } i\text{th operation}$        $\Phi(0) = \# \text{ credits at the beginning}$   
 potential function energy

$$(2) \quad \underset{\substack{\downarrow \\ \text{amortized}}}{\hat{c}_i} = \underset{\substack{\downarrow \\ \text{actual cost}}}{c_i} + \Phi(i) - \Phi(i-1)$$

$$(1) \quad \underline{\Phi(n) \geq \Phi(0)} \Rightarrow \text{red} \geq \text{green} \Rightarrow \sum_i \hat{c}_i \geq \sum_i c_i$$

$$\Rightarrow \sum_i \hat{c}_i = \sum_i c_i + \Phi(n) - \Phi(0)$$


---

$\Phi(i) = c \cdot \# \text{ objects in the stack after } i\text{th operation}$

$$\Phi(0) = 0 \quad \Phi(n) \geq 0$$

PUSH:

$$\hat{c}_i = c_i + \Phi(i) - \Phi(i-1) = c + c = 2c$$

POP

$$\hat{c}_i = c_i + \Phi(i) - \Phi(i-1) = c - c = 0$$

MULTIPOP

$$\dots = c \cdot \# \text{ pop} - c \cdot \# \text{ pop} = 0$$

