

Lecture 2: Queues

EECG142- Data Structures

FIFO

Textbook:

Data Structures via C++: Objects by Evolution
by A. Michael Berman

First year - EECE Department
Spring 2025

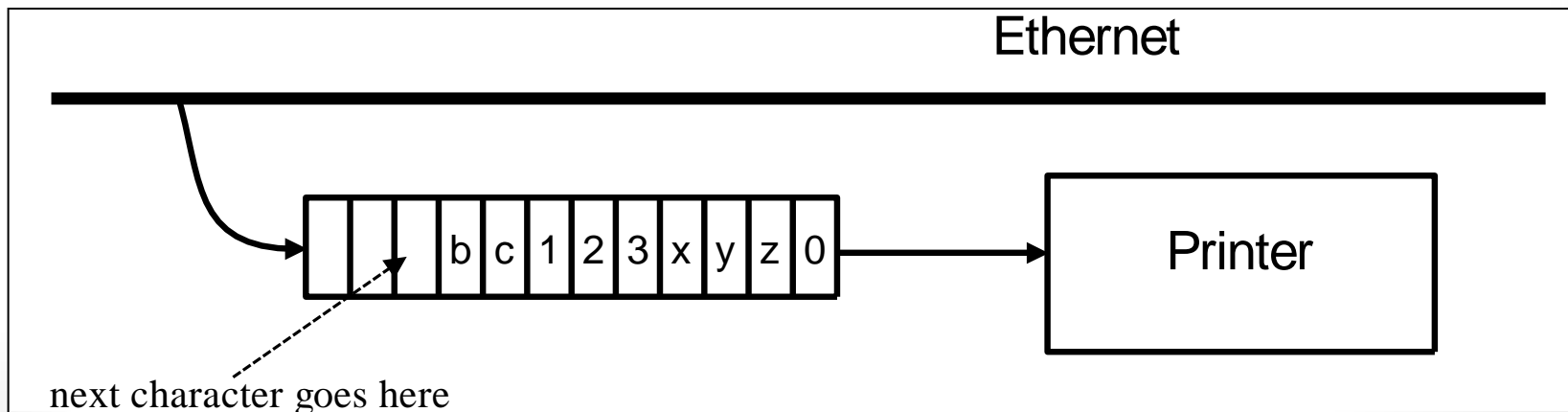
This lecture covers

1. Queue definition and examples
2. Queue implementation using linear arrays
3. Queue implementation using circular arrays
4. Queue implementation using linked-lists

Queue Example

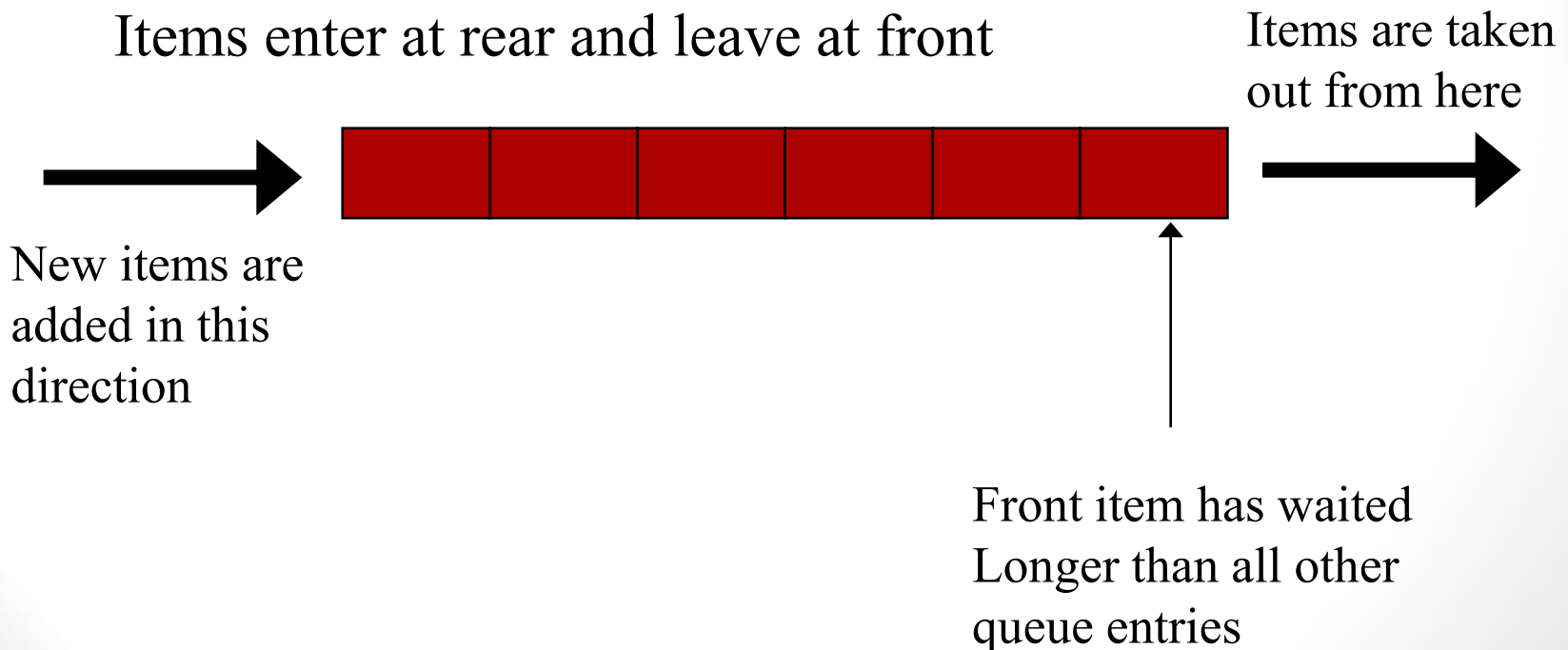
Examples of queues include:

- people waiting in line for a movie or a shop.
- jobs waiting to be executed by a processors
- documents needed to be printed by a printer
- data in a buffer waiting to be transmitted



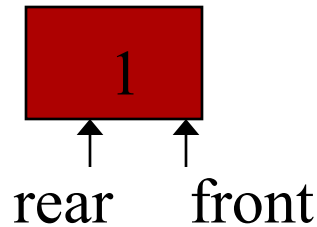
Queues

Unlike stacks, queues have a **First in, first out (FIFO)** property. **Items are added to the rear and removed from the front**

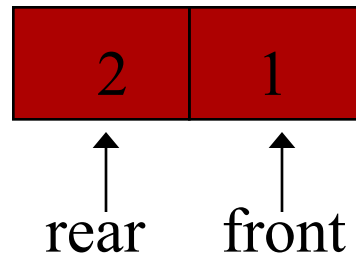


Simple Queue Operations

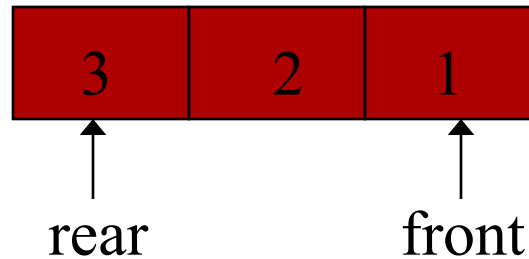
q.enqueue(1)



q.enqueue(2)



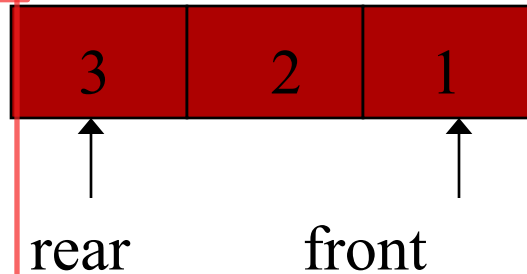
q.enqueue(3)



Simple Queue Operations

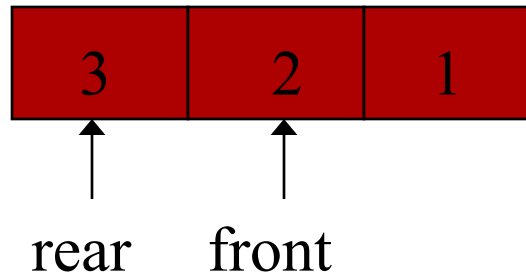
```
cout << q.front();
```

1



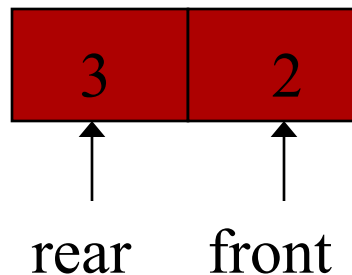
```
cout << q.dequeue();
```

1



```
cout << q.front();
```

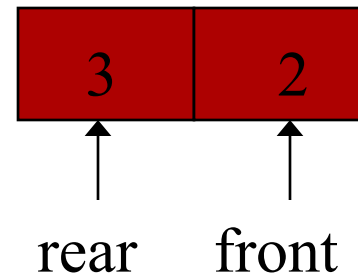
2



Simple Queue Operations

```
if (q.isEmpty())  
    cout << "empty" << endl;  
else  
    cout << "not empty" << endl;
```

not empty



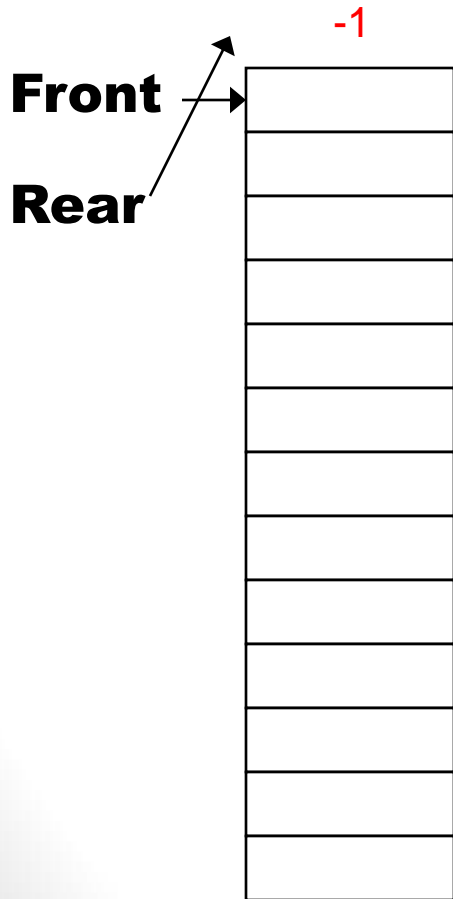
Array-based Implementation

- You need to keep track of the queue's front and rear.
- The rear index of the queue will normally be higher than its front index with the exceptions
 - Empty queue ($\text{rear} < \text{front}$)
 - One-item queue ($\text{rear} = \text{front}$)
- A full queue has its rear index = array size-1

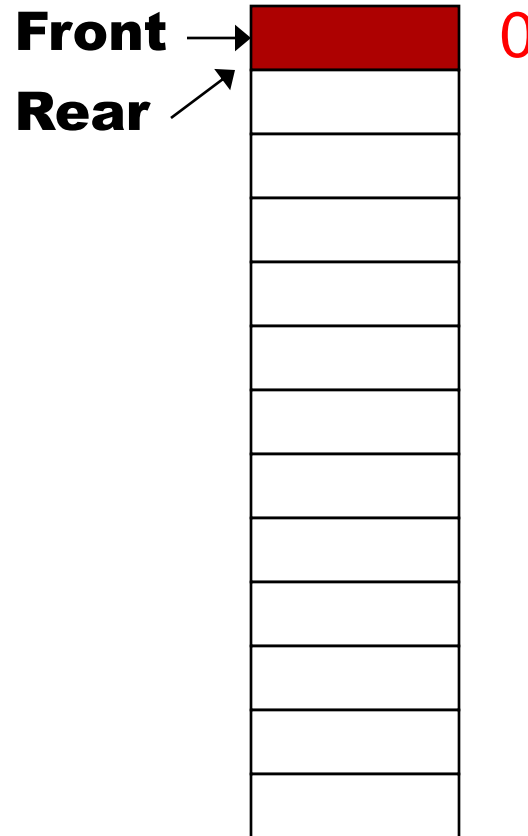
front -> entrance	3 cases	
rear -> exit	1) $\text{rear} < \text{front}$	empty
front -> first element	2) $\text{rear} > \text{front}$	not empty
rear -> last element	2.1) $\text{rear} = \text{front}$	one element (always true)
	2.2) $\text{rear} = \text{max}-1$	full

Special Cases of a Queue

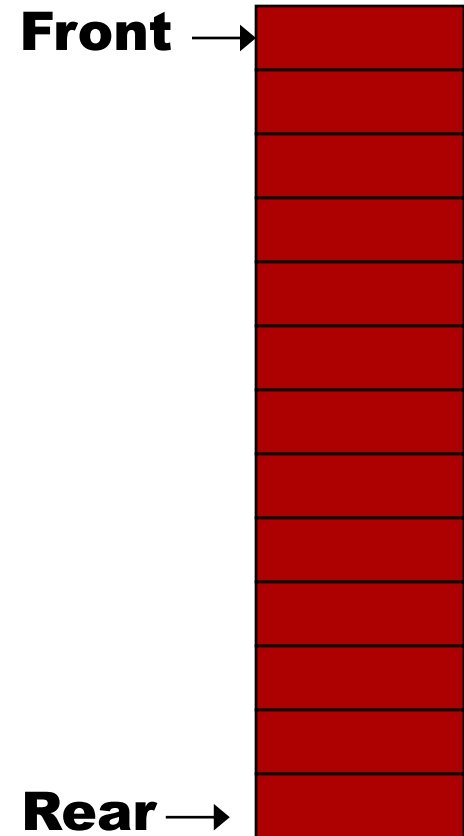
Empty queue



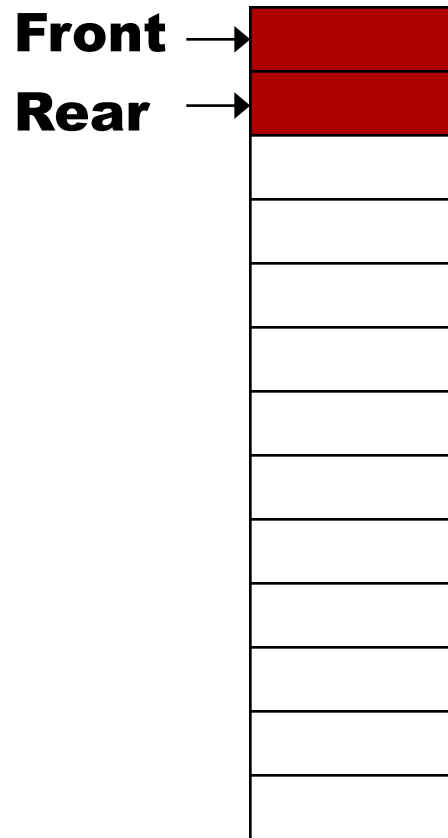
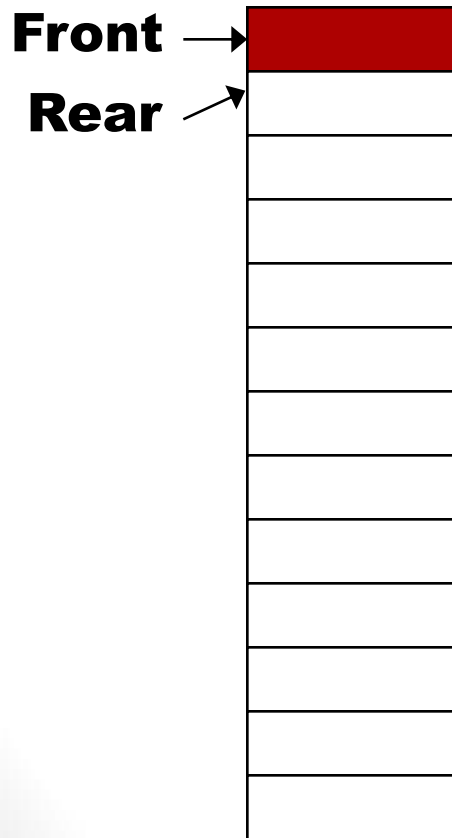
Single element queue



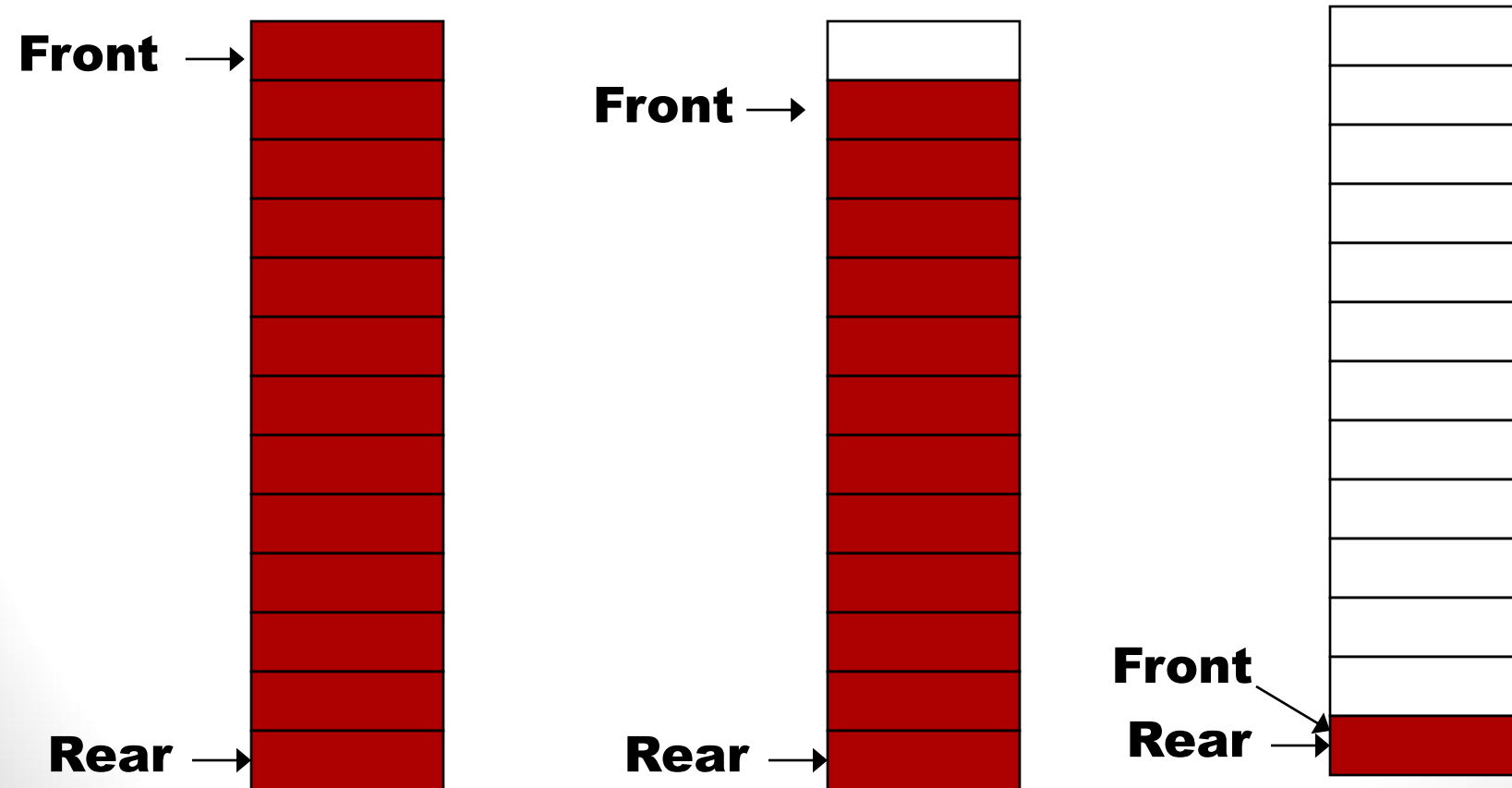
Full queue



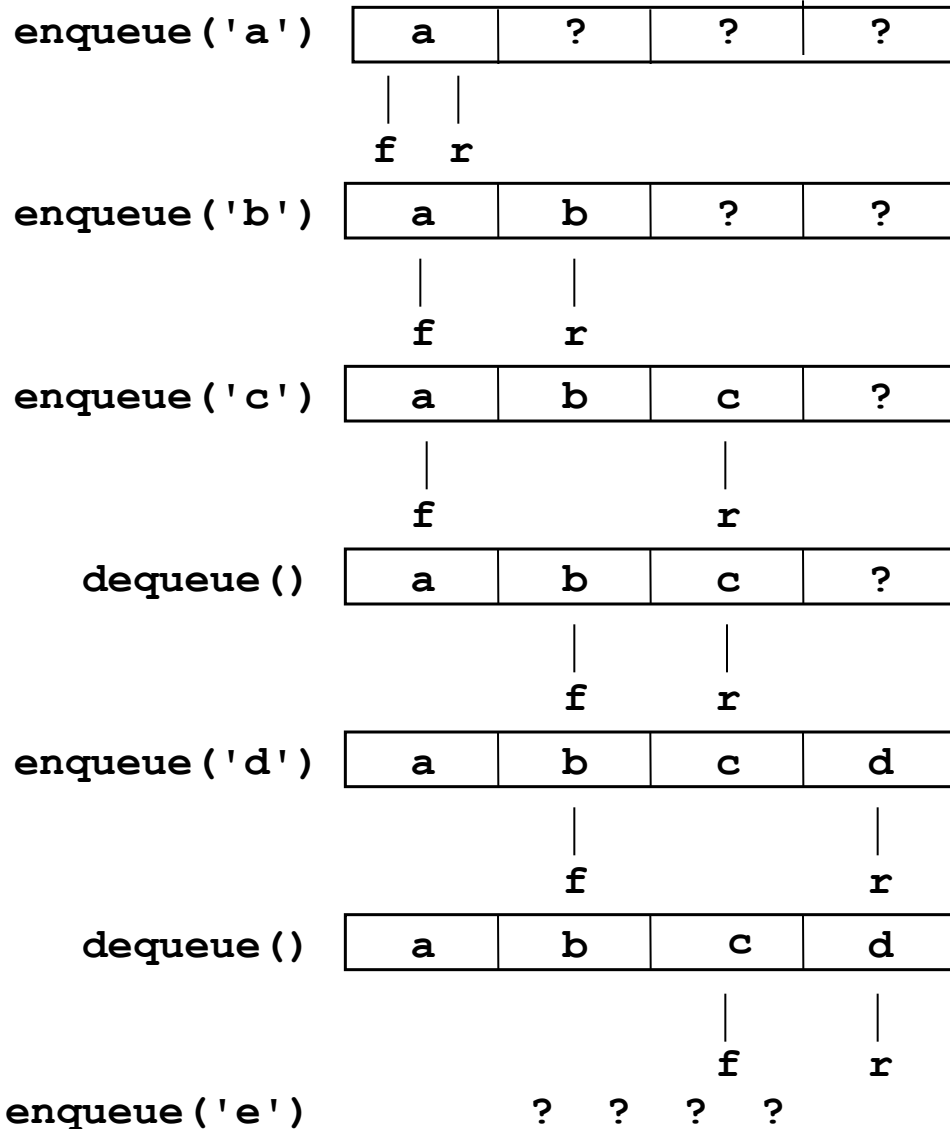
Enqueueing Items



Dequeuing Items

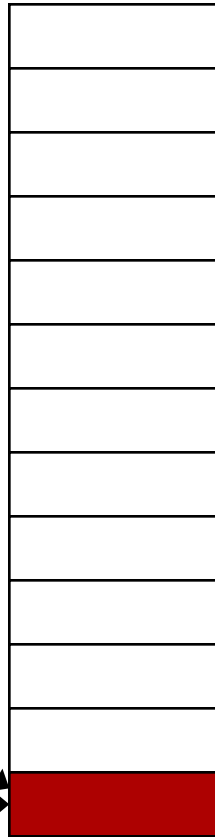


Implementation Problems



Implementation Problems

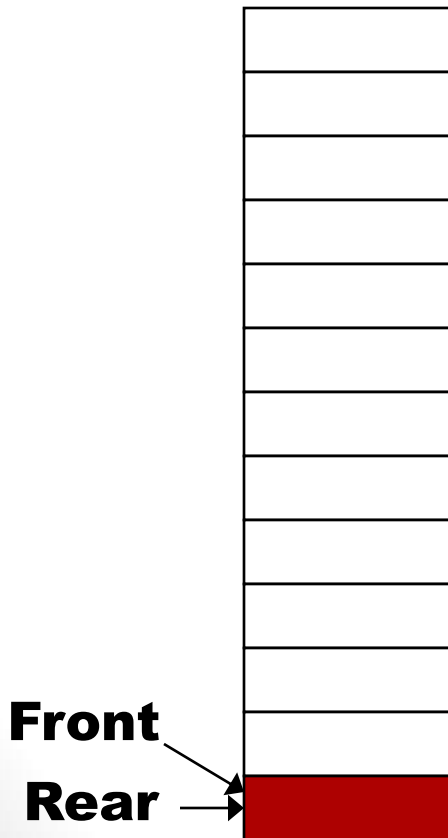
This single element
queue is still considered
full



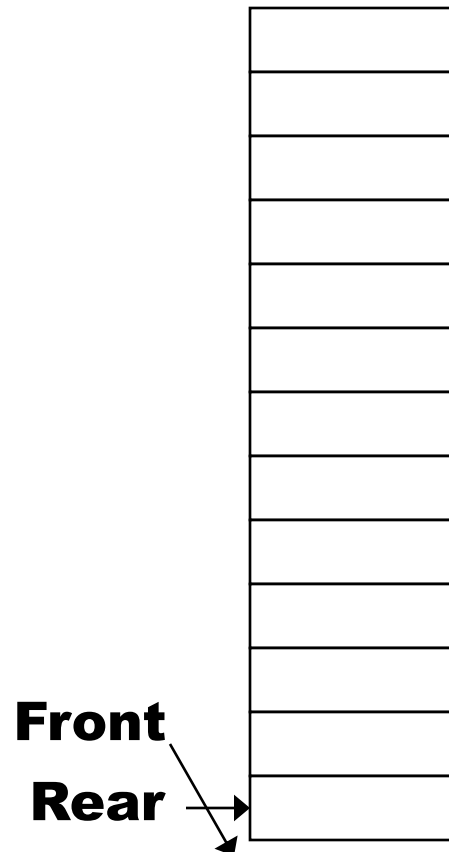
- Once the rear reaches the last index of the array, the queue is considered full and cannot accept more items even if there are empty cells after dequeuing
- We need to wait till the queue becomes empty and then reset the rear index.
- Such way is inefficient and unacceptable in many applications

Implementation Problems

Single element
queue



Is the queue empty or full?



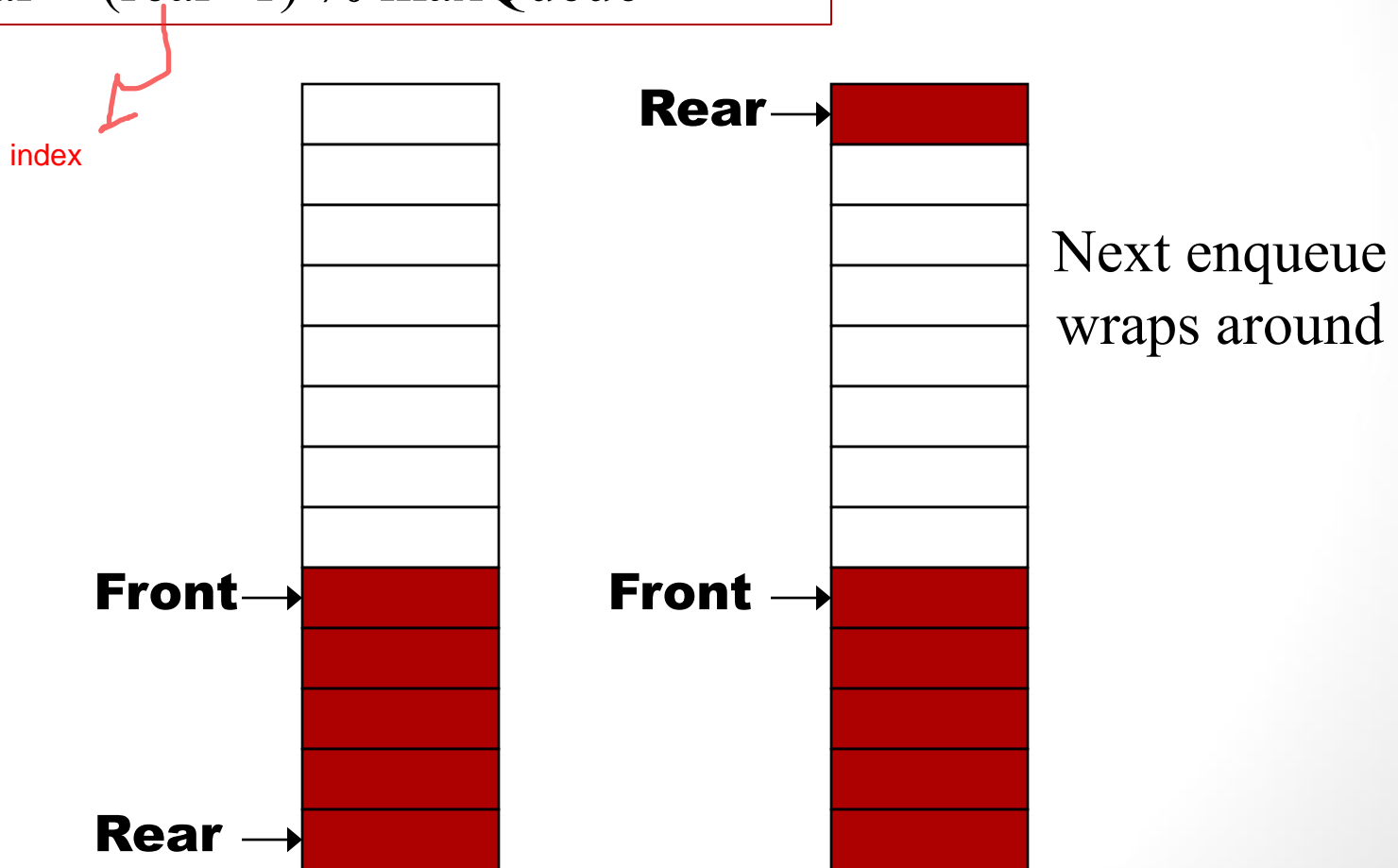
Our definition of empty is $\text{rear} < \text{front}$, so it is empty.

But, rear has reached its limit. It cannot go beyond the array, Therefore, the queue is full!

Solution: Wrapping Around

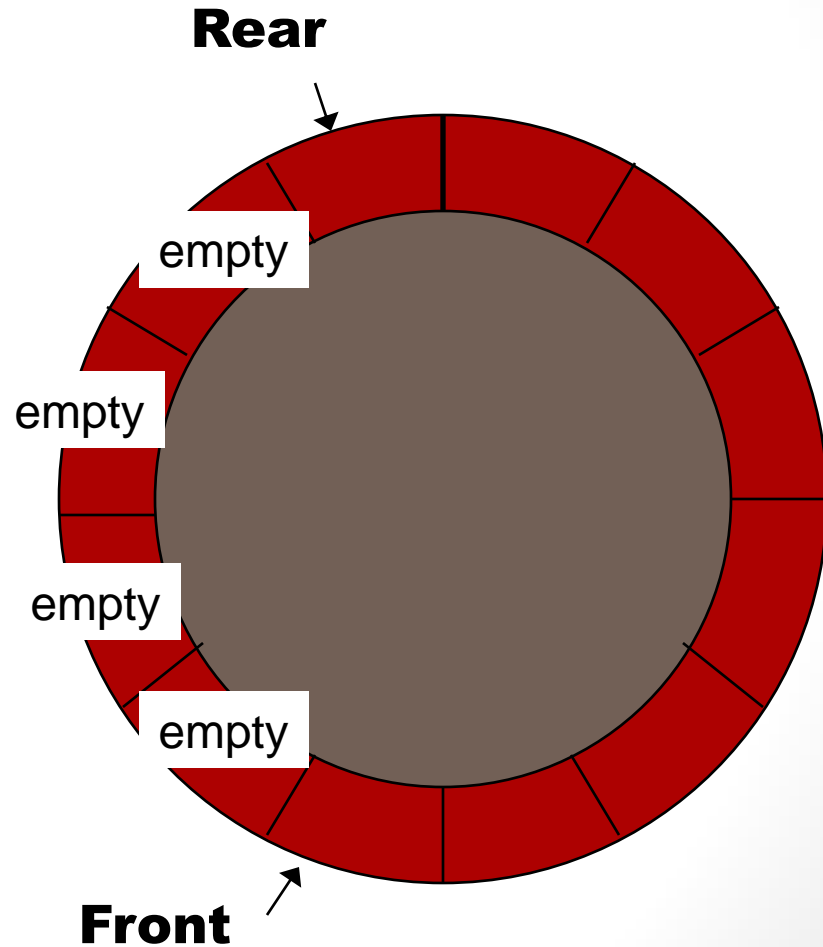
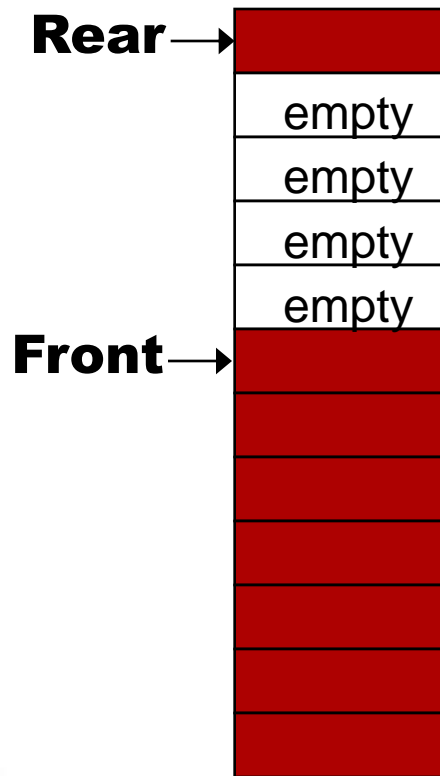
¹³
If $\text{rear} + 1 > \text{maxQueue} - 1$, then $\text{rear} = 0$

¹³
 $\text{rear} = (\text{rear} + 1) \% \text{maxQueue}$



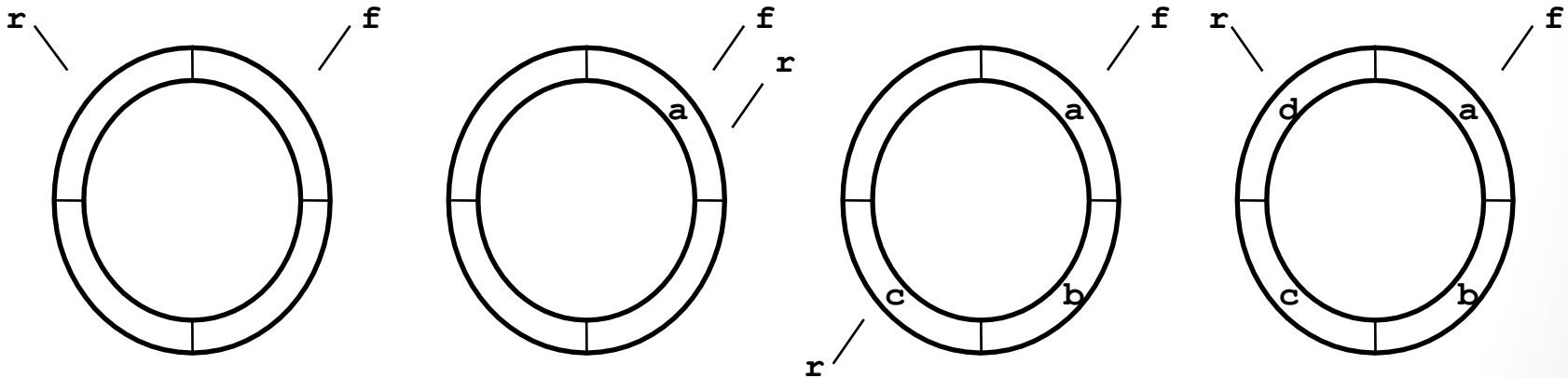
The Circular Queue

Circular queue

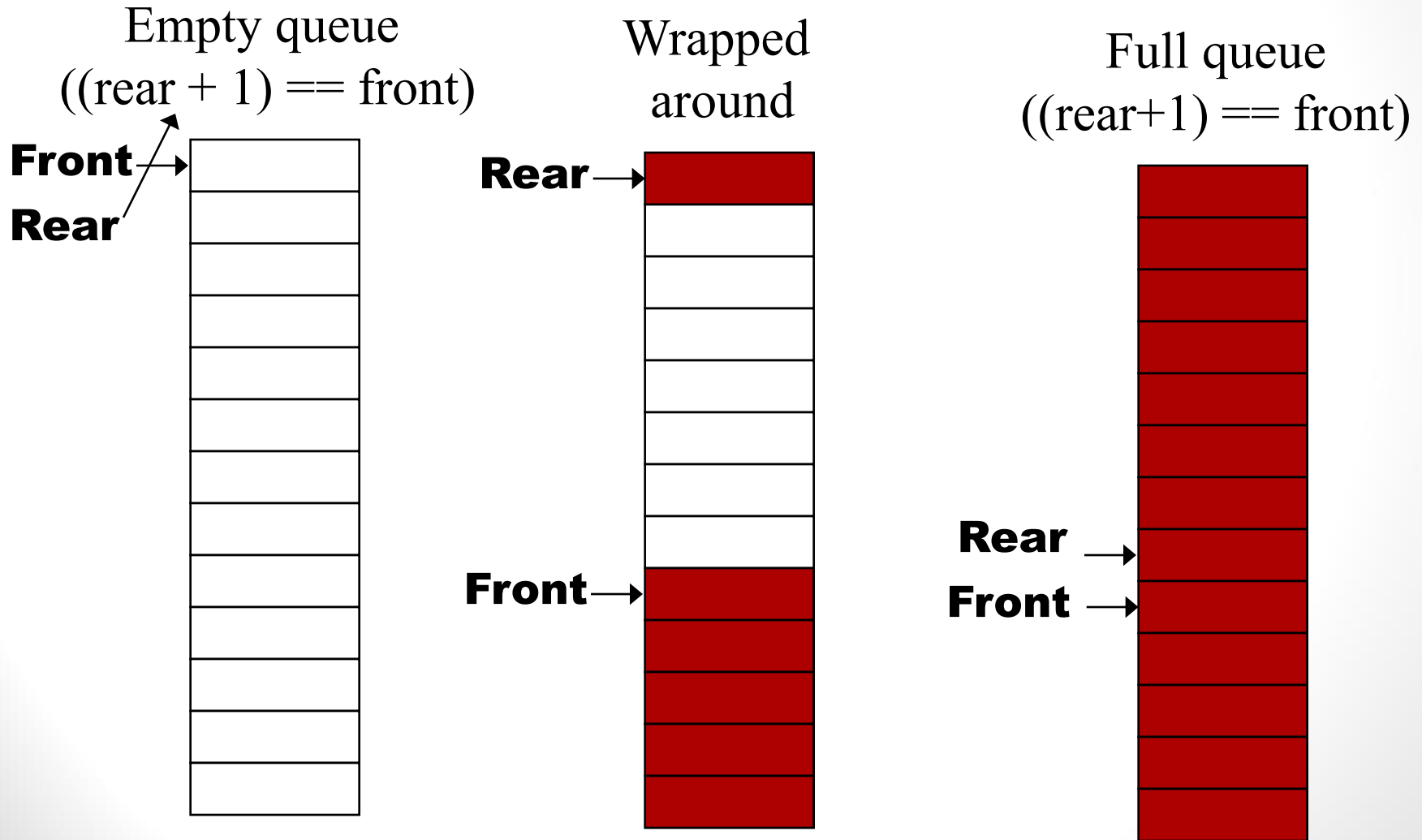


Is this queue empty or full?

We cannot use ' $\text{rear} < \text{front}$ ' as a test of empty circular queue because of the wrapping around.

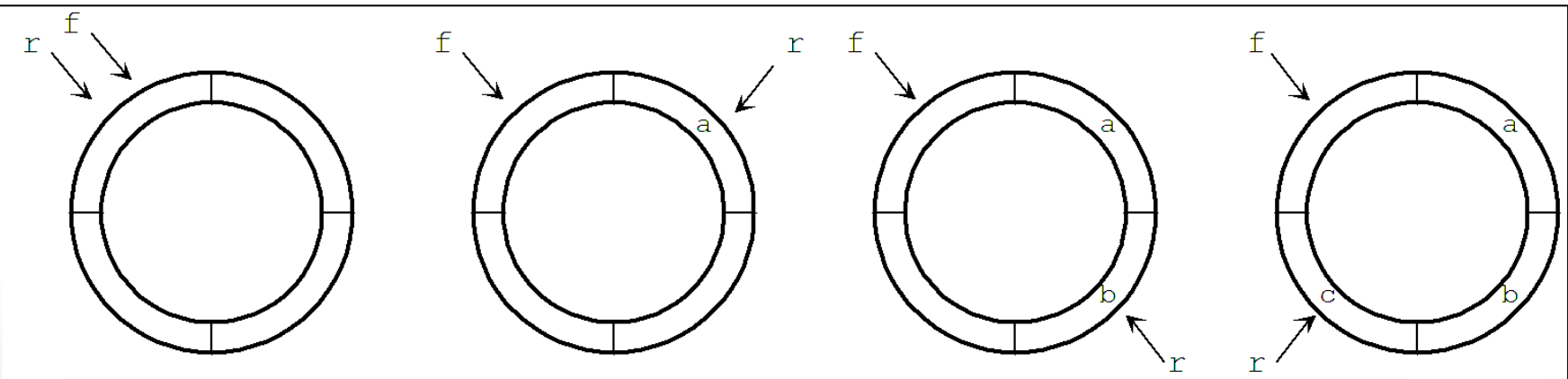


Explanation using Conventional Arrays



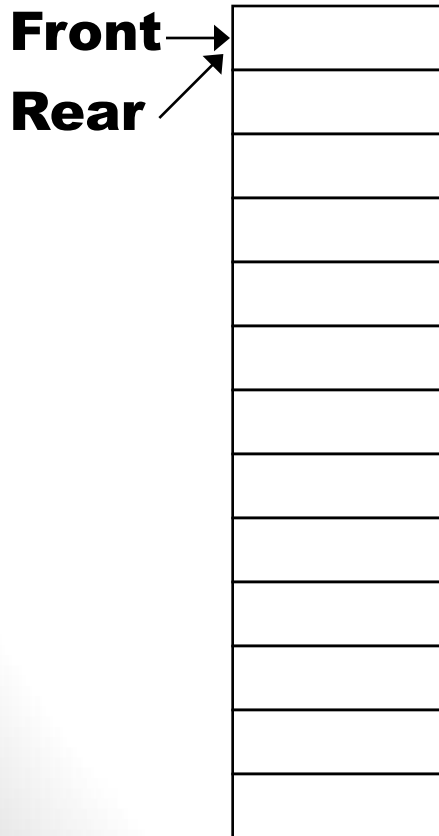
Empty or Full fix

- Let the front always point to an empty cell.
- Elements are added at $\text{nextPos}(\text{rear})$
- An empty queue is defined by $\text{rear} == \text{front}$
- A full queue is defined by $\text{nextPos}(\text{rear}) == \text{front}$

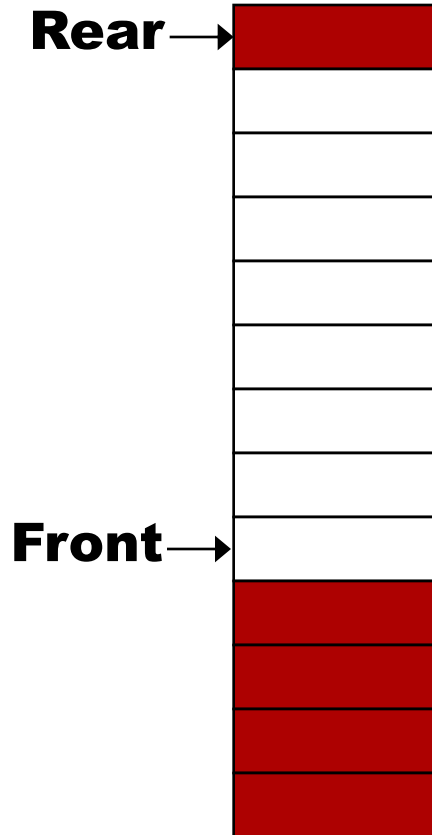


Explanation using Conventional Arrays

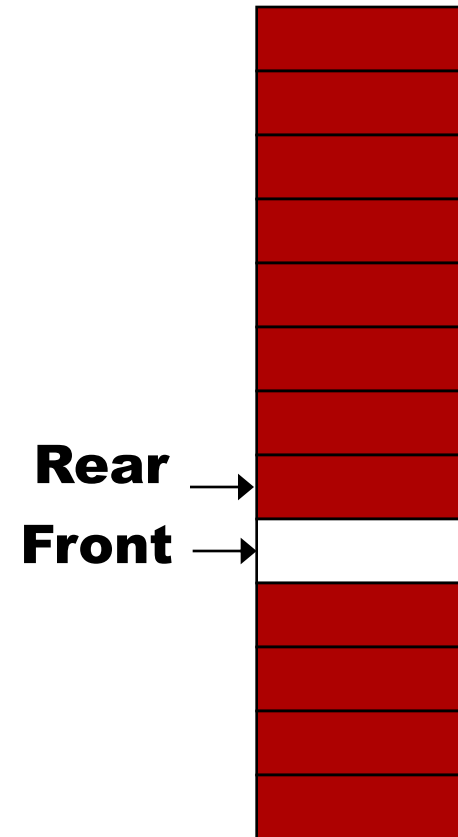
Empty queue
($\text{rear} == \text{front}$)



Wrapped
around



Full queue
($((\text{rear} + 1) == \text{front})$)



Queue Header file

```
const int maxQueue = 200;
template < class queueElementType >
class Queue {
public:
    Queue();
    void enqueue(queueElementType e);
    queueElementType dequeue();
    queueElementType front();
    bool isEmpty();
private:
    int f; // marks the front of the queue
    int r; // marks the rear of the queue
    queueElementType elements[maxQueue];
```

empty :
Rear = front

full:
nextPos(Rear) = front

front :
before the start

Rear :
point to the end

enqueue :
1) check not full
2) rear++
3) put the element

dequeue:
1) check not empty
2) front ++

any place occupied by front called before start
and the place it left considered empty

Queue Header file (cntd.)

هالام

```
int nextPos(int p)
```

```
{
```

```
if (p == maxQueue - 1) // at end of circle
```

```
    return 0;
```

```
else
```

```
    return p+1;
```

```
}
```

```
};    int nextpos( int p){ return (p+1)%MaxSize;}
```

wrapping around effect

neve never never
do this
f++ r++

always use nextPos funtion

Queue Implementation

```
template < class queueElementType >
Queue < queueElementType >::Queue()
{ // start both front and rear at 0
    f = 0;
    r = 0;
}
template < class queueElementType > bool
Queue < queueElementType >::isEmpty()
{ // return true if the queue is empty
    return bool(f == r);
}
```

Queue Implementation

```
template < class queueElementType >
Void Queue < queueElementType > ::
    enqueue(queueElementType e)
{ // add e to the rear and advance r
    assert(nextPos(r) != f);
    r = nextPos(r);
    elements[r] = e;
}
```


Queue Implementation

```
template < class queueElementType >
queueElementType
Queue < queueElementType >::dequeue()
{ // advance the front and return the value at the front
  assert(f != r);
  f = nextPos(f);
  return elements[f];
}
```

Queue Implementation

```
template < class queueElementType >
queueElementType
Queue < queueElementType >::front()
{
    // return value of element at the front
    assert(f != r);
    return elements[nextPos(f)];
}
```

Header for Queue as Linked-list

```
template < class queueElementType >
class Queue {
public:
    Queue();
    void enqueue(queueElementType e);
    queueElementType dequeue();
    queueElementType front();
    bool isEmpty();
```

Private section

private:

```
Struct Node;  
typedef Node * nodePtr;  
struct Node {  
    queueElementType data;  
    nodePtr next;  
};  
nodePtr f; head  
nodePtr r; tail  
};
```

linked list

Linked-lists based implementation

```
template < class queueElementType >
Queue < queueElementType >::Queue()
{ // set both front and rear to null pointers
f = NULL;
    r = NULL;
}

template < class queueElementType > bool
Queue < queueElementType >::isEmpty()
{ // true if the queue is empty -- when f is a null pointer
    return bool(f == NULL);
}
```

Linked-lists based implementation

```
template < class queueElementType >
queueElementType
Queue < queueElementType >::front()
{
    assert(f); not pointing to null
    return f->data;
}
```

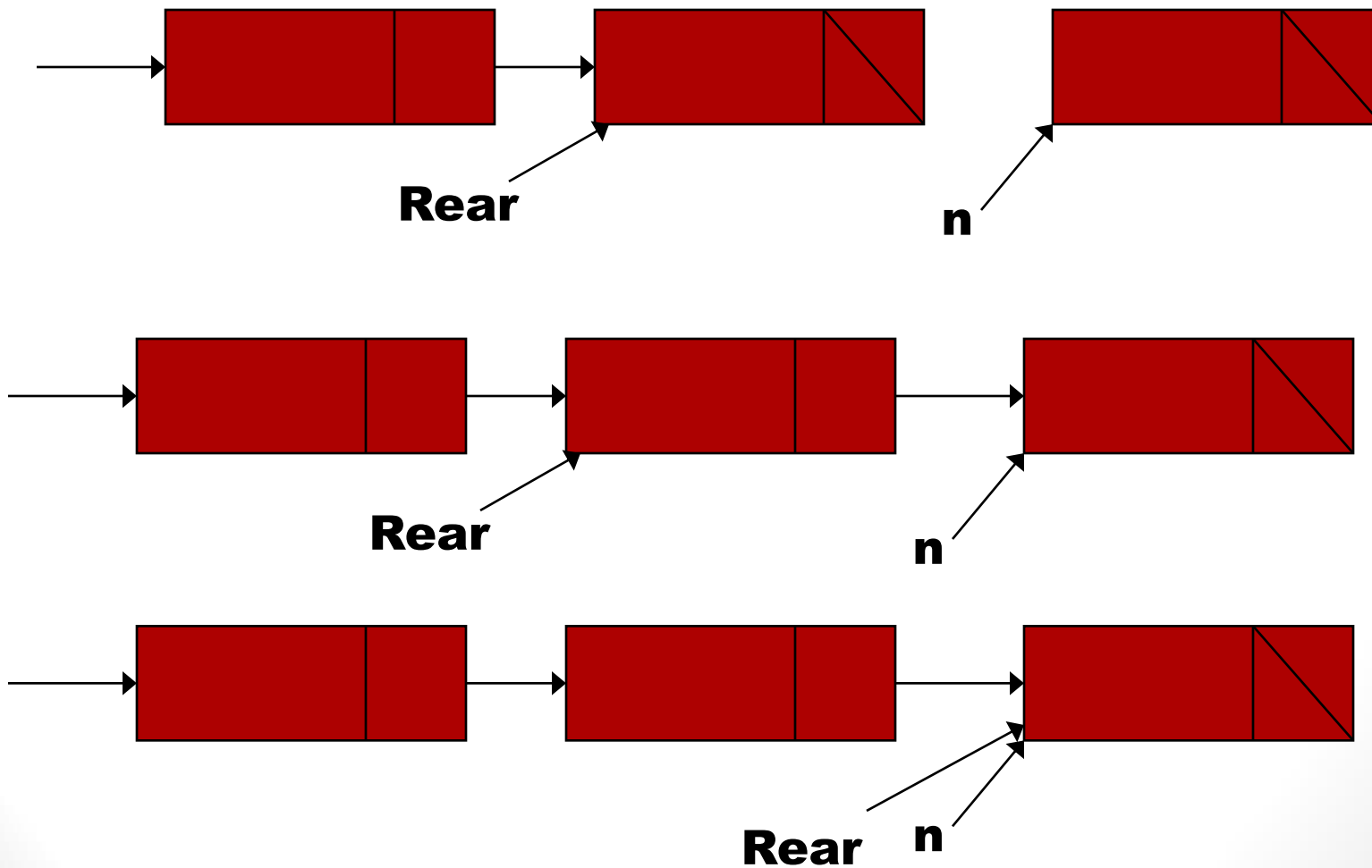
Linked-lists based implementation

```
template < class queueElementType >
void Queue < queueElementType > ::enqueue(queueElementType e)
{ // create a new node, insert it at the rear of the queue
  nodePtr* n = new Node;
  assert(n);
  n->next = NULL;
  n->data = e;
  if (f != NULL) { // existing queue is not empty
    r->next = n; // add new element to end of list
    r = n;
  } else { // adding first item in the queue
    f = n; // so front, rear must be same node
    r = n;
  }
}
```

enqueue is the same as insert function

هالام

enqueueing



dequeue()

```
template < class queueElementType > queueElementType
```

```
Queue < queueElementType >::dequeue()
```

```
{ assert(f); // make sure queue is not empty
```

```
queueElementType frontElement = f->data;
```

```
    nodePtr n=f;
```

```
    f = f->next;
```

front move forward when dequeue
rear move forward when enqueue

```
    delete n;
```

```
    if (f == NULL) // we're deleting last node
```

```
    rear = NULL;    in linked list implimentation its impossible that rear lags the front
```

```
    return frontElement;
```

```
}
```

dequeuing

