

Multipath variants for congestion control in context of WebRTC

Student:
Ardelean Daniel-George

Coordinator:
Sterca Ioan-Adrian

Table of contents

01

Motivation

02

Research goals

03

Congestion
control

04

Why Multipath

05

Methodology

06

System
architecture

07

Simulation topology

08

Path selection
strategies

09

Evaluated metrics

10

Results

11

Trade-off
s

12

Future work

13

Conclusions

Motivation

- Increased demand for real-time communication [1]
- WebRTC dominant technology for video streaming [2]
- Single path limitations



Google Meet



Microsoft Teams



WhatsApp



zoom



Jitsi Meet

Research goals

- Survey existing congestion control algorithms [3]
- Explore benefits of multipath transmission in WebRTC
- Evaluate the performance across various metrics

WebRTC & Congestion control

- Peer-to-peer communication over UDP
- No built-in congestion control in UDP
- Existing algorithms: GCC, SReAM, NADA [5][6][7]
- Limitations: single-path assumptions

04

Why Multipath?

- Send across multiple network routes at the same time [4]
- Advantages:
 - Increased resilience
 - Redundancy and fault tolerance
 - Efficient bandwidth use
 - Seamless mobility for mobile devices

Multipath in practice



MPTCP: Multipath TCP [8]

- Built into the protocol
- Multipath subflows



UDP lacks native congestion control support

Congestion control is not built into the protocol itself



Challenges

- Packet reordering
- Integration with existing architectures
- Path selection strategies

Methodology

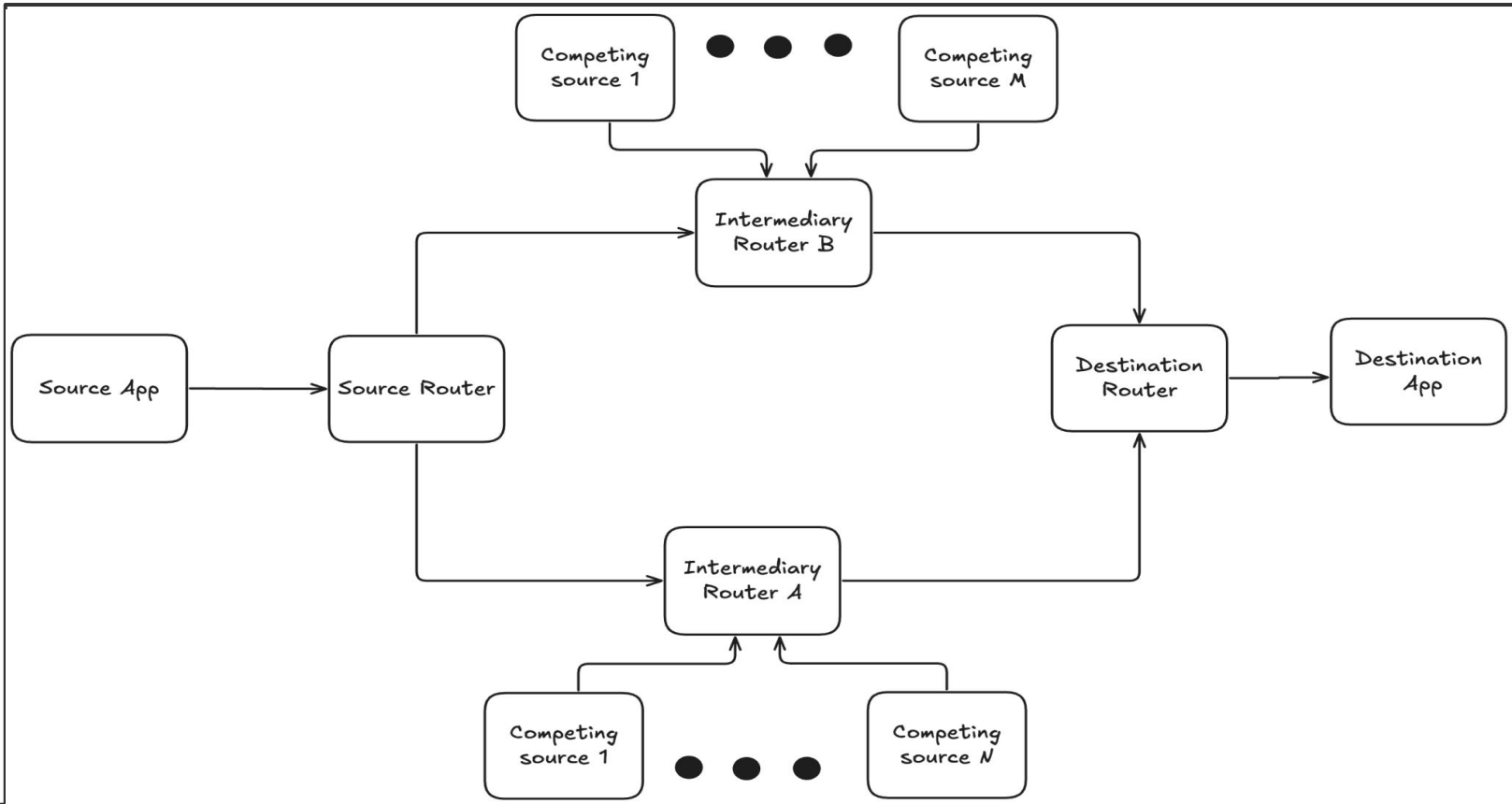
- Create a realistic multipath simulation scenario
- Implement a custom multipath capable NADA that can switch between multiple path selection strategies
- Compare the behaviour of a classic NADA [7] and its multipath counterpart

System architecture

- Simulations realised using Network Simulator 3 (ns-3) [9]
- Extended existing UDP client classes
- Simulate realistic video streams by implementing custom video packets
- Simulate a receiver application to consume the video stream

Simulation topology

- Two paths: different bandwidth/latency
- Competing TCP flows to simulate real traffic and achieve statistical multiplexing over the network
- Wrapper classes over Source/Destination connections
- Custom video buffer and receiver app to simulate a video player



Path selection strategies

- Dynamic Weights
- Best Path
- Round Robin
- Frame-Aware
- Buffer-Aware

Congestion signals



Agg signal

Packet loss

Packet loss is the primary congestion indicator

One-way delay

The one-way delay between sender indicates network queuing and potential congestion buildup

Receiving / Sending rate

NADA compares the receiving rate with the send rate to detect throughput reductions

Explicit congestion notifications

Packets are marked as congested by routers

Evaluated metrics

<u>Throughput</u>	Measures the effective data rate at the receiver, calculated as total received bytes divided by simulation time (in Mbps)
<u>Delay</u>	The average time taken by packets to travel from sender to receiver, including all transmission, propagation, and queuing delays
<u>Packet loss</u>	The percentage of packets that were sent but never received at the destination
<u>Jitter</u>	Measures the variability in packet arrival times, indicating how uneven the delivery intervals are
<u>Buffer length</u>	Indicates how much video content (in seconds) is currently preloaded in the receiver's playback buffer
<u>Buffer underruns</u>	Counts the number of times the playback buffer is empty and causes a video freeze or rebuffering
MOS (Mean Opinion Score)	An estimated user experience score (0 to 5) derived from received video bitrate using a logistic mapping
Stability	Reflects the consistency of throughput over time; large deviations indicate unstable performance
Efficiency	Represents how many of the transmitted bytes actually reached the destination, indicating delivery effectiveness

10

Results





60.8% - 219.4%

Throughput increase

88.9%–99.9%
reduction

Packet loss

Up to 86.7% lower

Jitter on stable networks

84.3 - 218.5%

Variability in throughput

0.01% – 0.88%

Improvement in MOS score

84.8 - 414.6%

In buffer length (ms)

Trade-offs

- Increased throughput variability
- Adaptive strategies have increased jitter
- Path selection adds complex overhead

Future work

- Adaptive filtering and congestion prediction
- Real deploys on real networks and evaluating real user experience
- Testing on 5G/6G, IoT, edge computing scenarios
- Machine learning for strategy selection

Conclusions

- The Multipath NADA approach improves real-time performance
- Modularity of the proposed architecture can easily be extended for future research
- Best for hybrid topologies
- Buffer-aware / Frame-aware are the most stable

references

- [9].A. B. Massimo Antonio Tozzato, Fabio Cazzin. (2025) Network simulator 3. [Online]. Available: <https://www.nsnam.org/>

**THANK
YOU**



Q&A