



Introducing Reflection into a Verification System

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Semester Project Report

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Abstract

Stainless is a tool for verifying Scala programs. During verification, tree reflection is sometimes needed. It can be used to send constraints to an underlying constraint solving during program execution. With tree reflection, the program can be described as an algebraic data type. You can then type check and interpret it within Stainless.

1 Introduction

The goal of this project was to implement tree reflection within Stainless. Tree reflection means to allow access to an expression as an algebraic data type. For example, if you had this expression:

`x + y * z`

the corresponding algebraic data type could be:

`Plus(Var("x"), Times(Var("y"), Var("z")))`

For this purpose, we defined which expressions we wanted to be able to reflect and how to describe them as algebraic data types. Furthermore, we defined some basic types to support type checking on our expressions. As the underlying expressions could be interpreted, we added an interpreter for our algebraic data types that returned the same result but wrapped in its reflection.

2 Implementation

The system is written in Pure Scala, thus allowing Stainless to verify it. As Stainless ensures lot of properties, the code has to be as simple as it could be, otherwise Stainless is not able to terminate the verification in a sustainable time.

2.1 Expressions

Expressions are algebraic data types. They represent the abstract syntax tree (AST) of a program that can be type checked and interpreted. It was implemented as an abstract class, *Expr*, and case classes that extend *Expr*.

The expressions that will be found in the leaves of the AST are the literals. They are the "basic unit" of expressions. For example, the literal for a character is defined as:

```
case class CharLiteral(value: Char) extends Expr
```

There are also literals that represent integers, booleans, strings and fractions. They are defined in a similar manner except for fractions which take a tuple of `BigInt`, the first element of the tuple representing the numerator and the second one the denominator. They all correspond to one of the basic types described in the following section.

Another type of value (leaf of the AST) is lambda definition:

```
case class Lambda(params: List[(Identifier, Type)],
                  body: Expr) extends Expr
```

It represents lambda-expressions such as

```
 $\lambda x: \text{BigInt}, \lambda y: \text{BigInt} \rightarrow x + y$ 
```

In the program, a list containing (x, BigInt) and (y, BigInt) represents the *params*. In fact, x and y would be `Identifiers`. `Identifiers` are defined as a case class that takes a string as parameter, the "name" of what needs to be identified. $x + y$ represents the *body*. In the body, x and y would be represented by *Variable*, which is an expression defined by:

```
case class Variable(id: Identifier) extends Expr
```

The nodes of the AST represent operations done on the leaves and on other nodes. For example, the addition operation is implemented as:

```
case class Plus(lhs: Expr, rhs: Expr) extends Expr
```

Here, *Plus* is a node in the AST and *lhs* and *rhs* are the children of this node. The other arithmetic operations (minus, unary minus, times, division, remainder and modulo) are defined in a similar way, except for the unary minus which takes only one argument.

In addition to arithmetic, there are also expressions to represent:

- String operations such as concatenation, computing the length and taking a subset of the given string expression
- Equality test on two expressions
- Comparisons on integers, fractions and characters expressions such as greater than ($>$), greater equals (\geq), less than ($<$) and less equals (\leq).
- Logical operations on booleans expressions such as logical and, logical or, logical implication and logical not
- If-then-else expression

- Let expression, an expression which gives a certain value to an identifier and then uses this definition in the subsequent expression

Another type of expression is *Application*:

```
case class Application(callee: Expr, args: List[Expr])
  extends Expr
```

It is used when we want to apply arguments to a *Lambda*. For example, if our lambda is

```
add = λ x: BigInt, λ y: BigInt → x + y
```

then *add*(1, 2) would correspond to an *Application* where *callee* is *add* and a list containing 1 and 2 represents *args*.

When there is a problem during the interpretation, the program outputs an expression which represents an error:

```
case class ErrorValue(error: String) extends Expr
```

2.2 Types

Types are represented as an abstract class *Type* and case classes that extend it and represent basic types of Stainless. There is a type for each kind of literal. For example, the type corresponding to a character is:

```
case class CharType() extends Type
```

In the same way, *IntegerType* is for integers, *BooleanType* for booleans, *StringType* for strings and *RealType* for fractions.

The type that corresponds to lambdas is:

```
case class FunctionType(from: List[Type], to: Type) extends Type
```

where *from* contains the types of the parameters *params* of *Lambda* and *to* represents the type of the body of *Lambda*.

2.3 DSL

To be able to define expressions in a more convenient way, the user can call some methods of the Domain Specific Language (DSL).

For literals, the methods are named with the first letter of the literal in capital. For example, for a *CharLiteral*:

```
def C(c : Char) = CharLiteral(c)
```

For expressions that need an Identifier as argument, the user only has to provide a string which represents the name of this Identifier. For example, to define a *Let* expression:

```

def let(name: String, tpe: Type, value: Expr)
  (body: Identifier => Expr): Expr = {
    val id = Identifier(name)
    Let(id, tpe, value, body(id))
  }

```

For arithmetic, string concatenation, comparisons and logical operators, the methods are defined in the abstract class *Expr* using operator overloading. For example, for *Plus*:

```
def +(rhs: Expr) = Plus(this, rhs)
```

Then, if you need to express the addition of 1 and 2, you can write it as:

```
I(1) + I(2)
```

and it will represent

```
Plus(I(1), I(2))
```

Unary operators such as *UMinus* and *Not* were defined by:

```

def unary_- = UMinus(this)
def unary_! = Not(this)

```

Then, to express -1, the user only needs to write:

```
-I(1)
```

and to express the negation of the boolean true:

```
! B(true)
```

2.4 Type checker

The type checker was implemented in a recursive manner. It is a method that takes as arguments an expression and an environment which is represented by a Map of Identifier to Type. It does a pattern match on the given expression and applies an inference rule to obtain either *None()* if the expression does not type check or *Some(resulting type)* if the expression type checks.

The inference rules are (E representing the environment):

For literals:

$$\begin{array}{l}
 \text{CharLiteral} \frac{c \text{ is a Char}}{E \vdash \text{CharLiteral}(c) : \text{CharType}} \\
 \text{IntegerLiteral} \frac{i \text{ is a BigInt}}{E \vdash \text{IntegerLiteral}(i) : \text{IntegerType}} \\
 \text{BooleanLiteral} \frac{b \text{ is a Boolean}}{E \vdash \text{BooleanLiteral}(b) : \text{BooleanType}} \\
 \text{StringLiteral} \frac{s \text{ is a String}}{E \vdash \text{StringLiteral}(s) : \text{StringType}}
 \end{array}$$

$$\text{FractionLiteral} \frac{n \text{ is a } \mathit{BigInt}, d \text{ is a } \mathit{BigInt}}{E \vdash \mathit{FractionLiteral}((n, d)) : \mathit{RealType}}$$

$$\text{Variable} \frac{v : T \in E}{E \vdash v : T}$$

$$\text{Let} \frac{E \vdash \text{value} : T1 \quad E, id : T1 \vdash \text{body} : T2}{E \vdash \mathit{Let}(id, T1, \text{value}, \text{body}) : T2}$$

$$\text{If-then-else} \frac{E \vdash \text{cond} : \mathit{BooleanType} \quad E \vdash \text{thenn} : T \quad E \vdash \text{elze} : T}{E \vdash \mathit{IfExpr}(\text{cond}, \text{thenn}, \text{elze}) : T}$$

For lambdas:

$$\text{Lambda} \frac{E_{[id1:=t1, \dots, idN:=tN]} \vdash \text{body} : T}{E \vdash \mathit{Lambda}([(id1, t1), \dots, (idN, tN)], \text{body}) : \mathit{FunctionType}([t1, \dots, tN], T)}$$

$$\text{Application} \frac{E \vdash \text{callee} : \mathit{FunctionType}([t1, \dots, tN], T) \quad E \vdash a_1 : t_1 \dots E \vdash a_N : t_N}{E \vdash \mathit{Application}(\text{callee}, [a_1, \dots, a_N]) : T}$$

For arithmetic expressions:

$$\text{General arithmetic} \frac{E \vdash \text{lhs} : T \quad E \vdash \text{rhs} : T}{E \vdash \text{op}(\text{lhs}, \text{rhs}) : T} \quad \text{Unary minus} \frac{E \vdash e : T}{E \vdash \mathit{UMinus}(e) : T}$$

The general arithmetic rule is applied by *Plus*, *Minus*, *Times*, *Division*, *Modulo* and *Remainder* where *op* represents the corresponding operation. For the last two, *T* can only be *IntegerType*. For the others, *T* is either *IntegerType* or either *RealType*.

The Unary minus rule is applied by *UMinus* and *T* can be *IntegerType* or *RealType*.

For string operations:

$$\begin{aligned} \text{Strings concatenation} & \frac{E \vdash \text{lhs} : \mathit{StringType} \quad E \vdash \text{rhs} : \mathit{StringType}}{E \vdash \mathit{StringConcat}(\text{lhs}, \text{rhs}) : \mathit{StringType}} \\ \text{Substring} & \frac{E \vdash e : \mathit{StringType} \quad E \vdash \text{start} : \mathit{IntegerType} \quad E \vdash \text{end} : \mathit{IntegerType}}{E \vdash \mathit{SubString}(e, \text{start}, \text{end}) : \mathit{IntegerType}} \\ \text{String length} & \frac{E \vdash e : \mathit{StringType}}{E \vdash \mathit{StringLength}(e) : \mathit{IntegerType}} \end{aligned}$$

For comparisons expressions:

$$\text{Equals} \frac{E \vdash \text{lhs} : T1 \quad E \vdash \text{rhs} : T2}{E \vdash \mathit{Equals}(\text{lhs}, \text{rhs}) : \mathit{BooleanType}}$$

$$\text{Other comparisons} \frac{E \vdash lhs : T \quad E \vdash rhs : T}{E \vdash op(lhs, rhs) : T}$$

For equals, $T1$ and $T2$ can be of any type but lhs and rhs must type check.
 For other comparisons, T must be either *IntegerType*, *RealType* or *CharType*.
 This rule is applied by *LessThan*, *GreaterThan*, *LessEquals* and *GreaterEquals*.

For logical operations:

$$\begin{array}{c} \text{Logical operators} \frac{E \vdash lhs : BooleanType \quad E \vdash rhs : BooleanType}{E \vdash op(lhs, rhs) : BooleanType} \\ \text{Not} \frac{E \vdash e : BooleanType}{E \vdash Not(e) : BooleanType} \end{array}$$

For logical operators, op can be *And*, *Or* or *Implies*.

2.5 Interpreter

The interpreter is a small-step interpreter. It has two main methods: *interpret*, which takes an expression and returns the fully interpreted expression (either a literal, a lambda or an error value) and *next*, which takes an expression and returns an option, either *Some(the next small-step of interpretation of the expression)* or either *None()* if it is stuck. Being stuck can happen for two reasons. First, the expression is already fully evaluated, thus we cannot progress. Second, the expression does not make sense, thus there are no inference rule to progress in the interpretation.

The interpreter substitutes the variables in the rest of the expression as soon as their value is defined. For example, if you have this pseudo-code below:

```
let x = 1; let x = x + 2; x + 4
```

substitution of the first x will return:

```
let x = 1; let x = 1 + 2; x + 4
```

because the second x masks the first one in the $x + 4$ statement. Then, after the second x has been evaluated, the substitution would return:

```
let x = 1; let x = 3; 3 + 4
```

The *next* method follows these inference rules, where \rightarrow means a step of evaluation and v represents a fully interpreted expression (a literal, a lambda or an error):

$$\begin{array}{c} \text{For Let:} \\ \frac{e \rightarrow e'}{Let(id, t, e, body) \rightarrow Let(id, t, e', body)} \\ \frac{v \text{ is a value}}{Let(id, t, v, body) \rightarrow body_{[id:=v]}} \end{array}$$

For Application:

$$\begin{array}{c}
\frac{\text{callee} \rightarrow \text{callee}'}{\text{Application}(\text{callee}, \text{args}) \rightarrow \text{Application}(\text{callee}', \text{args})} \\
\frac{\text{c is a Lambda, } e_i \rightarrow e'_i}{\text{Application}(c, [v_1, \dots, v_{i-1}, e_i, e_{i+1}, \dots, e_N]) \rightarrow \text{Application}(c, [v_1, \dots, v_{i-1}, e'_i, e_{i+1}, \dots, e_N])} \\
\frac{\text{c is a Lambda}([(id_1, t_1), \dots, (id_N, t_N)], \text{body}), v_1, \dots, v_N \text{ are values}}{\text{Application}(c, [v_1, \dots, v_N]) \rightarrow \text{body}_{[id_1:=v_1, \dots, id_N:=v_N]}}
\end{array}$$

For if-then-else expressions:

$$\begin{array}{c}
\frac{\text{cond} \rightarrow \text{cond}'}{\text{IfExpr}(\text{cond}, \text{thenn}, \text{elze}) \rightarrow \text{Let}(\text{cond}', \text{thenn}, \text{elze})} \\
\frac{v \text{ is BooleanLiteral}(\text{true})}{\text{IfExpr}(v, \text{thenn}, \text{elze}) \rightarrow \text{thenn}} \quad \frac{v \text{ is BooleanLiteral}(\text{false})}{\text{IfExpr}(v, \text{thenn}, \text{elze}) \rightarrow \text{elze}}
\end{array}$$

For arithmetic expressions:

$$\begin{array}{c}
\text{Arithmetic 1} \quad \frac{\text{lhs} \rightarrow \text{lhs}'}{\text{Expr}(\text{lhs}, \text{rhs}) \rightarrow \text{Expr}(\text{lhs}', \text{rhs})} \\
\text{Arithmetic 2} \quad \frac{v \text{ is an IntegerLiteral, } \text{rhs} \rightarrow \text{rhs}'}{\text{Expr}(v, \text{rhs}) \rightarrow \text{Expr}(v, \text{rhs}')} \\
\text{Arithmetic 3} \quad \frac{v \text{ is a FractionLiteral, } \text{rhs} \rightarrow \text{rhs}'}{\text{Expr}(v, \text{rhs}) \rightarrow \text{Expr}(v, \text{rhs}')} \\
\text{Arithmetic 4} \quad \frac{v1, v2 \text{ are IntegerLiterals}}{\text{Expr}(v1, v2) \rightarrow v1 \text{ op } v2} \\
\text{Arithmetic 5} \quad \frac{v1, v2 \text{ are FractionLiterals}}{\text{Expr}(v1, v2) \rightarrow v1 \text{ op } v2} \\
\frac{e \rightarrow e'}{UMinus(e) \rightarrow UMinus(e')} \quad \frac{v \text{ is an IntegerLiteral or a FractionLiteral}}{UMinus(v) \rightarrow -v}
\end{array}$$

The rules Arithmetic 1, 2, 3, 4 and 5 are used by the *Expr*: Plus, Minus, Times, Division and their corresponding operation *op*: +, -, *, / The rules Arithmetic 1, 2 and 4 are implemented by the *Expr*: Remainder, Modulo and their corresponding operation *op*: %, mod.

For operations on strings:

$$\begin{array}{c}
\frac{\text{lhs} \rightarrow \text{lhs}'}{\text{StringConcat}(\text{lhs}, \text{rhs}) \rightarrow \text{StringConcat}(\text{lhs}', \text{rhs})} \\
\frac{v \text{ is a StringLiteral, } \text{rhs} \rightarrow \text{rhs}'}{\text{StringConcat}(v, \text{rhs}) \rightarrow \text{StringConcat}(v, \text{rhs}')}
\end{array}$$

$$\begin{array}{c}
\frac{v1, v2 \text{ are StringsLiterals}}{StringConcat(v1, v2) \rightarrow v1 + v2} \\
\frac{e \rightarrow e'}{SubString(e, start, end) \rightarrow SubString(e', start, end)} \\
\frac{v \text{ is a StringLiteral}, start \rightarrow start'}{SubString(v, start, end) \rightarrow SubString(v, start', end)} \\
\frac{v1 \text{ is a StringLiteral}, v2 \text{ is an IntegerLiteral}, end \rightarrow end'}{SubString(v1, v2, end) \rightarrow SubString(v1, v2, end')} \\
\frac{v1 \text{ is a StringLiteral}, v2, v3 \text{ are IntegerLiterals}}{SubString(v1, v2, v3) \rightarrow v1.bigSubstring(v2, v3)} \\
\frac{e \rightarrow e'}{StringLength(e) \rightarrow StringLength(e')} \quad \frac{v \text{ is a StringLiteral}}{StringLength(v) \rightarrow v.bigLength}
\end{array}$$

For comparisons:

$$\begin{array}{c}
\frac{lhs \rightarrow lhs'}{Expr(lhs, rhs) \rightarrow Expr(lhs', rhs)} \\
\frac{v \text{ is a lit}, rhs \rightarrow rhs'}{Expr(v, rhs) \rightarrow Expr(v, rhs')} \quad \frac{v1, v2 \text{ are lits}}{Expr(v1, v2) \rightarrow v1 \text{ op } v2}
\end{array}$$

For the *Expr* Equals, *v1* and *v2* can be any type of value, *v1* and *v2* can even be different type of literals and then *op* is ==. For the *Expr* LessThan, LessEquals, GreaterThan, GreaterEquals, *lit* are IntegerLiteral, FractionLiteral and CharLiteral. *v1* and *v2* must be the same kind of literals. The corresponding *op* are: <, ≤, >, ≥.

For logical operators:

$$\begin{array}{c}
\text{General logic } \frac{lhs \rightarrow lhs'}{Expr(lhs, rhs) \rightarrow Expr(lhs', rhs)} \\
\text{And and Implies 1 } \frac{v \text{ is BooleanLiteral}(true), rhs \rightarrow rhs'}{Expr(v, rhs) \rightarrow rhs'} \\
\text{And and Implies 2 } \frac{v1 \text{ is BooleanLiteral}(true), v2 \text{ is a BooleanLiteral}}{Expr(v1, v2) \rightarrow v2} \\
\frac{v \text{ is a BooleanLiteral}(false), rhs \rightarrow rhs'}{Or(v, rhs) \rightarrow rhs'} \\
\frac{v1 \text{ is a BooleanLiteral}(false), v2 \text{ is a BooleanLiteral}}{Or(v, v2) \rightarrow v2} \\
\frac{v \text{ is a BooleanLiteral}(false)}{And(v, rhs) \rightarrow BooleanLiteral(false)} \\
\frac{v \text{ is a BooleanLiteral}(true)}{Or(v, rhs) \rightarrow BooleanLiteral(true)}
\end{array}$$

$$\begin{array}{c}
\frac{v \text{ is a } \textit{BooleanLiteral}(\textit{false})}{\textit{Implies}(v, \textit{rhs}) \rightarrow \textit{BooleanLiteral}(\textit{true})} \\
\frac{e \rightarrow e'}{\textit{Not}(e) \rightarrow \textit{Not}(e')} \quad \frac{v \text{ is a } \textit{BooleanLiteral}}{\textit{Not}(v) \rightarrow !v}
\end{array}$$

General logic is used by And, Or and Implies.

2.6 Soundness theorem

The soundness theorem can be stated as: "If a program type checks, its evaluation does not get stuck"(1). To show that the system (or at least a subset of the Expressions) was a sound system, we proved with Stainless two lemmas, progress and preservation, on some of the Expressions.

2.6.1 Progress

Progress can be stated as: "If a program type checks, it is not stuck"(1). In the program, it is translated by a method which takes as arguments an expression *expr* and a type *t*. It has a precondition:

```
require (!isValue(expr) &&
        typecheck(expr, Map[Identifier, Type]()) == Some(t))
```

and a post condition:

```
next(expr).nonEmpty
```

where the *isValue* method returns true if *expr* is a literal, a *Lambda* or an *ErrorValue*. To ensure the post condition, we used the method *holds* of Stainless.

Stainless was not able to prove this as it was stated but fortunately, by using additional lemmas (*check* method of Stainless), Stainless did verify progress on a subset of expressions. The lemmas were added depending on the given expression. To do so, we used pattern matching. For expressions such as general arithmetic, string operations, comparisons, logical operators and if-then-else, we added an equality test on the type of the expression and checked that the fields of each expression had a consistent type. Furthermore, we pattern matched on the fields and checked the progress of the first non-evaluated field. For example, here are the added lemmas for *Plus*:

```
case Plus(lhs, rhs) => {
  check((t == IntegerType() || t == RealType()) &&
        typecheck(lhs, Map[Identifier, Type]()) == Some(t) &&
        typecheck(rhs, Map[Identifier, Type]()) == Some(t))
  (lhs, rhs) match{
    case (IntegerLiteral(_), IntegerLiteral(_)) => true
    case (FractionLiteral(_), FractionLiteral(_)) => true
    case (IntegerLiteral(_), _) => check(progress(rhs, t))
    case (FractionLiteral(_), _) => check(progress(rhs, t))
    case (_, _) => check(progress(lhs, t))
  }
```

```

    }
  }
}

```

For *Let*, we added:

```

case Let(id, tValue, value, body) =>
  value match{
    case _ if(isValue(value)) => true
    case _ => check(progress(value, tValue))
  }

```

For *Application*, Stainless could not verify it if it had an unbounded number of parameters, but we were able to prove progress on an *Application* which had only one or two parameters. To be able to do so, we defined:

```

case class FunctionType1(from: Type, to: Type) extends Type
case class FunctionType2(from1: Type, from2: Type, to: Type)
  extends Type

case class Lambda1(id: Identifier, t: Type, body: Expr)
  extends Expr
case class Lambda2(id1: Identifier, t1: Type,
  id2: Identifier, t2: Type, body: Expr) extends Expr

case class Application1(callee: Expr, arg: Expr) extends Expr
case class Application2(callee: Expr, arg1: Expr, arg2: Expr)
  extends Expr

```

Lambda1, *Lambda2* and *Application1*, *Application2* follow the same rules for type checking and interpretation as *Lambda* and *Application* but with respectively one and two parameters. The lemmas that were added to prove progress on *Application1* are:

```

case Application1(callee, arg) =>
  callee match {
    case Lambda1(id, t, body) => arg match{
      case _ if(isValue(arg)) => true
      case _ => check(progress(arg, t))
    }
    case _ =>{
      val tArg = typecheck(arg,
        Map[Identifier, Type]())
      check(tArg.nonEmpty &&
        progress(callee, FunctionType1(tArg.get, t)))
    }
  }
}

```

It is similar for *Application2*, except it checks progress on the first non-evaluated *arg* and type checks both *args*.

Due to the duration of verifying progress with Stainless, we could only verify together:

- Arithmetic, string operations and if-then-else expressions (approximately 4 hours for the post condition of progress only and 5 hours to prove all the others lemmas in progress)
- Let expressions, comparisons and logical operators (approximately 4 hours for the post condition of progress only and 5 hours to prove all the other lemmas in progress)
- Application1 (less than a minute to prove progress)
- Application2 (less than a minute to prove progress)

The duration may seem extremely long but Stainless had to control the exhaustiveness of every pattern matching and that there were no divisions, remainders or modulus by zero. It also verified every clause in each *check* method and then that the overall lemma was correct. Furthermore, it had to verify the post-condition of progress, which took the longest time. The output of Stainless for each verification is shown in the annexes.

2.6.2 Preservation

Preservation can be stated as: "If a program type checks and makes one [next] step [with the Interpreter], then the result again type checks"(1). In the program, it is translated by a method which takes as arguments an expression *e1* and a type *t*. It has a precondition:

```
require(typecheck(e1, Map[Identifier, Type]()) == Some(t) &&
      next(e1).nonEmpty)
```

and a post condition:

```
typecheck(e2, Map[Identifier, Type]()) == Some(t)
```

where *e2* represents *next(e1)*. To ensure the post condition, we used the method *holds* of Stainless.

Unfortunately, as it is, Stainless is not able to prove preservation on any of the expressions (or at least not in less than ten hours). We tried to add lemmas with the *check* method of Stainless but due to the opacity of the verification system, it is too complicated to find which lemmas can improve the verification, or at least not worsen it.

2.7 Tests

We made tests to check if the type checker and the interpreter had a normal behaviour. By side effect, it also tested the DSL. Some of the tests were made using the method *holds* of Stainless. For example:

```
def testInterpretPlusInteger(): Boolean = {  
    interpret(I(1) + I(2)) == I(3)  
}.holds
```

Due to the duration of the verification, beside these tests, we used the ScalaTest library with the FunSuite class.

3 Conclusion

This project was challenging due to the opacity and duration of the verification in Stainless. Sometimes, I had to run Stainless for more than 8 hours to receive a result. The system did not give enough feedback to know if and where it was stuck and what could help it progress. Without the help of my supervisor, Romain Edelmann, I would not have found some of the tricks to help Stainless do the verification, like using pattern matching on list instead of if-then-else expressions to verify the ADT invariant of recursive functions.

I have also lost a lot of time creating a system too complex for Stainless to verify, groping toward the goal and following wrong paths.

Despite these complications, this project was a thrilling experience. It taught me to tackle a problem step-by-step and to use a verification system, which existence I was not aware of.

References

- [1] LARA, EPFL, *Computer Language Processing, Lecture 9, CS-320, Edition 2018* in <http://lara.epfl.ch/cc18:top>, Date of access: 06.06.19.

Annexes

stainless summary					
areBothBoolean	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:145:17	0.020
areBothInt	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:161:17	0.020
areBothIntOrReal	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:178:17	0.017
areBothIntOrRealOrChar	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:196:17	0.022
interpret	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:39:25	0.322
isValue	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:16:17	0.019
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:56:17	0.223
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:62:59	0.226
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:63:62	0.224
progress	postcondition	valid	nativez3	src/main/scala/Progress.scala:26:9	0.897
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:28:17	0.571
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:38:33	0.107
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:39:70	0.051
progress	precond. (call check(progress(arg, t)))	valid	nativez3	src/main/scala/Progress.scala:41:59	0.563
progress	precond. (call check(arg, t))	valid	nativez3	src/main/scala/Progress.scala:41:65	0.999
progress	precond. (call check(res))	valid	nativez3	src/main/scala/Progress.scala:45:49	0.857
progress	precond. (call progress(callee, FunctionType1(get[Type] ...))	valid	nativez3	src/main/scala/Progress.scala:45:72	0.725
progress	precond. (call get[Type](tArg))	valid	nativez3	src/main/scala/Progress.scala:45:103	0.572
subst	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:352:17	0.019
typecheck	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:20:17	0.024
typecheck	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:58:33	0.044
total: 21 valid: 21 (0 from cache) invalid: 0 unknown: 0 time: 6.522					

Figure 1: Verification of progress on Lambda1 and Application1

stainless summary					
areBothBoolean	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:145:17	0.158
areBothInt	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:161:17	0.020
areBothIntOrReal	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:178:17	0.038
areBothIntOrRealOrChar	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:196:17	0.024
interpret	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:39:25	0.238
isValue	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:16:17	0.029
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:56:17	0.046
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:70:66	0.049
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:71:73	0.049
next	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:72:69	0.039
progress	postcondition	valid	nativez3	src/main/scala/Progress.scala:26:9	0.941
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:28:17	0.253
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:49:33	0.057
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:50:81	0.044
progress	match exhaustiveness	valid	nativez3	src/main/scala/Progress.scala:51:77	0.066
progress	precond. (call check(progress(arg2, t2)))	valid	nativez3	src/main/scala/Progress.scala:53:67	0.771
progress	precond. (call progress(arg2, t2))	valid	nativez3	src/main/scala/Progress.scala:53:73	1.728
progress	precond. (call check(progress(arg1, t1)))	valid	nativez3	src/main/scala/Progress.scala:55:59	0.691
progress	precond. (call progress(arg1, t1))	valid	nativez3	src/main/scala/Progress.scala:55:65	1.836
progress	precond. (call check(res))	valid	nativez3	src/main/scala/Progress.scala:60:49	1.094
progress	precond. (call progress(callee, FunctionType2(get[Type] ...))	valid	nativez3	src/main/scala/Progress.scala:61:57	1.169
progress	precond. (call get[Type](tArg1))	valid	nativez3	src/main/scala/Progress.scala:61:88	0.621
progress	precond. (call get[Type](tArg2))	valid	nativez3	src/main/scala/Progress.scala:61:99	0.686
subst	match exhaustiveness	valid	nativez3	src/main/scala/Interpreter.scala:352:17	0.053
typecheck	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:20:17	0.035
typecheck	match exhaustiveness	valid	nativez3	src/main/scala/Typechecker.scala:65:33	0.066
total: 26 valid: 26 (0 from cache) invalid: 0 unknown: 0 time: 10.801					

Figure 2: Verification of progress on Lambda2 and Application2

stainless summary				
areBothBoolean	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:145:17 0.214
areBothInt	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:161:17 0.014
areBothIntOrReal	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:178:17 0.015
areBothIntOrRealOrChar	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:196:17 0.018
interpret	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:39:25 0.059
isValue	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:16:17 0.016
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:56:17 0.058
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:85:59 0.069
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:93:48 0.088
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:94:60 0.081
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:99:67 0.042
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:107:49 0.097
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:108:60 0.075
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:113:67 0.087
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:121:43 0.021
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:127:49 0.018
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:128:60 0.018
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:133:67 0.020
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:141:52 0.015
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:142:60 0.016
next	division by zero	valid	native23	src/main/scala/Interpreter.scala:144:88 0.216
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:148:67 0.019
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:157:53 0.021
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:158:60 0.018
next	remainder by zero	valid	native23	src/main/scala/Interpreter.scala:160:88 0.217
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:167:50 0.033
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:168:60 0.029
next	modulo by zero	valid	native23	src/main/scala/Interpreter.scala:170:107 0.189
next	modulo by zero	valid	native23	src/main/scala/Interpreter.scala:171:88 0.196
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:180:56 0.030
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:182:41 0.023
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:190:58 0.019
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:192:41 0.015
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:194:57 0.025
next	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:206:49 0.015
progress	postcondition	valid	native23	src/main/scala/Progress.scala:26:9 9885.492
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:28:17 120.678
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:67:33 135.991
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:70:33 0.037
progress	precond. (call check(progress(thenm, t)))	valid	native23	src/main/scala/Progress.scala:72:78 1508.200
progress	precond. (call check(progress(thenm, t)))	valid	native23	src/main/scala/Progress.scala:72:84 136.448
progress	precond. (call check(progress(else, t)))	valid	native23	src/main/scala/Progress.scala:74:79 1499.534
progress	precond. (call check(progress(else, t)))	valid	native23	src/main/scala/Progress.scala:74:85 136.516
progress	precond. (call check(progress(cond, BooleanType())))	valid	native23	src/main/scala/Progress.scala:75:59 1512.209
progress	precond. (call check(progress(cond, BooleanType())))	valid	native23	src/main/scala/Progress.scala:75:65 135.164
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:80:33 139.125
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:83:33 0.033
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:86:72 1499.809
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:86:78 129.967
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:87:73 1501.996
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:87:79 129.705
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:88:56 1505.890
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:88:62 128.896
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:92:33 132.570
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:95:33 0.037
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:98:72 1640.960
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:98:78 129.119
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:99:73 1652.393
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:99:79 128.216
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:100:56 1641.018
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:100:62 127.853
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:104:33 130.821
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:106:33 0.031
progress	precond. (call check(progress(e, t)))	valid	native23	src/main/scala/Progress.scala:109:53 1641.897
progress	precond. (call check(progress(e, t)))	valid	native23	src/main/scala/Progress.scala:109:59 129.649
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:113:33 131.454
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:116:33 0.038
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:119:72 1652.756
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:119:78 128.983
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:120:73 1654.182
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:120:79 128.187
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:121:56 1640.785
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:121:62 128.137
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:125:33 130.193
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:128:33 0.035
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:131:72 1638.195
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:131:78 128.205
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:132:73 1649.007
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:132:79 129.011
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:133:56 1639.504
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:133:62 129.182
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:137:33 131.356
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:140:33 0.034
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:142:72 1653.870
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:142:78 129.254
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:143:56 1640.766
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:143:62 129.671
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:147:33 132.224
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:150:33 0.036
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:152:72 1640.855
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:152:78 128.834
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:153:56 1648.842
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:153:62 129.052
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:159:33 126.900
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:162:33 0.038
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:164:71 1655.203
progress	precond. (call check(progress(rhs, t)))	valid	native23	src/main/scala/Progress.scala:164:77 127.417
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:165:56 1610.654
progress	precond. (call check(progress(lhs, t)))	valid	native23	src/main/scala/Progress.scala:165:62 128.015
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:169:33 127.397
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:173:33 0.038
progress	precond. (call check(progress(end, IntegerType())))	valid	native23	src/main/scala/Progress.scala:176:49 1446.098
progress	precond. (call check(progress(end, IntegerType())))	valid	native23	src/main/scala/Progress.scala:176:55 127.868
progress	precond. (call check(progress(start, IntegerType())))	valid	native23	src/main/scala/Progress.scala:178:49 1403.438
progress	precond. (call check(progress(start, IntegerType())))	valid	native23	src/main/scala/Progress.scala:178:55 128.198
progress	precond. (call check(progress(e, t)))	valid	native23	src/main/scala/Progress.scala:179:59 1360.192
progress	precond. (call check(progress(e, t)))	valid	native23	src/main/scala/Progress.scala:179:65 127.965
progress	precond. (call check(res))	valid	native23	src/main/scala/Progress.scala:183:33 127.507
progress	match exhaustiveness	valid	native23	src/main/scala/Progress.scala:185:33 0.034
progress	precond. (call check(progress(e, StringType())))	valid	native23	src/main/scala/Progress.scala:187:53 1298.050
progress	precond. (call check(progress(e, StringType())))	valid	native23	src/main/scala/Progress.scala:187:59 128.050
subst	match exhaustiveness	valid	native23	src/main/scala/Interpreter.scala:352:17 0.021
typecheck	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:20:17 0.021
typecheck	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:40:45 0.012
typecheck	match exhaustiveness	valid	native23	src/main/scala/Typechecker.scala:41:41 0.015
total: 115 valid: 115 (0 from cache) invalid: 0 unknown: 0 time: 55658.039				

Figure 3: Verification of progress on if-then-else expressions, general arithmetic and string operations

Figure 4: Verification of progress on Let expressions, comparisons and logical operators

stainless summary					
areBothBoolean	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:144:17	0.201
areBothInt	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:160:17	0.154
areBothIntOrReal	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:177:17	0.253
areBothIntOrRealOrChar	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:195:17	0.028
helperArgs	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:72:49	0.018
helperArgs	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:78:80	0.189
helperArgs	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:82:83	0.152
helperBody	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:90:49	0.016
helperNewEnv	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:53:49	0.023
interpret	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:40:25	0.185
isValue	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:18:17	0.034
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:57:17	0.088
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:63:59	0.428
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:97:41	0.409
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:113:59	0.570
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:121:48	0.259
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:122:60	0.333
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:127:67	0.102
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:135:49	0.258
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:136:60	0.166
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:141:67	0.047
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:149:43	0.046
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:155:49	0.052
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:156:60	0.050
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:161:67	0.064
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:169:52	0.074
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:170:60	0.041
next	division by zero	valid	nativev3	src/main/scala/Interpreter.scala:172:88	0.430
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:176:67	0.053
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:185:53	0.072
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:186:60	0.039
next	remainder by zero	valid	nativev3	src/main/scala/Interpreter.scala:188:88	0.335
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:195:50	0.046
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:196:60	0.041
next	modulo by zero	valid	nativev3	src/main/scala/Interpreter.scala:198:107	0.166
next	modulo by zero	valid	nativev3	src/main/scala/Interpreter.scala:199:88	0.064
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:208:56	0.038
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:210:41	0.046
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:218:58	0.041
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:220:41	0.039
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:222:57	0.049
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:234:49	0.051
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:246:52	0.046
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:247:60	0.057
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:252:67	0.035
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:257:57	0.039
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:265:55	0.040
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:266:60	0.070
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:271:67	0.083
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:276:57	0.038
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:284:54	0.050
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:285:60	0.031
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:290:67	0.042
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:295:57	0.031
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:303:57	0.030
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:304:60	0.040
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:309:67	0.044
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:314:57	0.042
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:325:47	0.044
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:328:41	0.044
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:337:46	0.036
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:340:41	0.036
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:349:51	0.035
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:352:41	0.040
next	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:361:40	0.042
subst	match exhaustiveness	valid	nativev3	src/main/scala/Interpreter.scala:380:17	0.030
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:21:17	0.015
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:41:45	0.078
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:42:41	0.255
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:63:33	0.109
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:67:64	0.089
typecheck	precond. (call get[Type](aType))	valid	nativev3	src/main/scala/Typechecker.scala:67:102	2616.159
typecheck	match exhaustiveness	valid	nativev3	src/main/scala/Typechecker.scala:105:33	0.088
total: 73 valid: 73 (0 from cache) invalid: 0 unknown: 0 time: 2623.568					

Figure 5: Verification of the files Interpreter, Typechecker, Types, Expressions and Identifiers with all the expressions