3.2. AUTÓMATAS APILADORES (AA)

DEFINICIÓN

$$A = (Q, \Sigma, \Gamma, \delta, q_0, z_0, F)$$

Q : conjunto finito de estados.

 Σ : alfabeto de entrada.

 Γ : alfabeto de la pila.

 δ : función de transición.

$$\delta$$
: Q × (Σ ∪ {ε}) × Γ → Q × Γ*

q₀: estado inicial.

 $q_0 \in Q$

z₀ : símbolo inicial de la pila.

 $z_0 \in \Gamma$

F: conjunto de estados finales o de aceptación.

 $F \subset Q$

Ejemplo 1:

$$A = (\{q_1, q_2, q_3, q_4\}, \{a, b\}, \{A, B\}, \delta, q_1, A, \{q_4\})$$

$$\delta(q_1, a, A) = \{(q_2, BA), (q_4, A)\}$$

$$\delta(q_1, \varepsilon, A) = \{(q_4, \varepsilon)\}$$

$$\delta(q_2, a, B) = \{(q_2, BB)\}$$

$$\delta(q_2, b, B) = \{(q_3, \varepsilon)\}\$$

$$\delta(q_3, \epsilon, A) = \{(q_4, A)\}$$

$$\delta(q_3, b, B) = \{(q_3, \varepsilon)\}$$

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Ejemplo 2:

$$\begin{split} &A = (\{q_1, q_2, q_3, q_4, q_5, q_6\}, \{a, b, c\}, \{A, B, Z_0\}, \delta, q_1, Z_0, \{q_6\}) \\ &\delta(q_1, a, A) = \{(q_1, AA)\} \\ &\delta(q_1, a, Z_0) = \{(q_1, AZ_0), (q_2, Z_0)\} \\ &\delta(q_1, b, A) = \{(q_3, A)\} \\ &\delta(q_2, b, B) = \{(q_2, BB)\} \\ &\delta(q_2, a, Z_0) = \{(q_2, Z_0)\} \\ &\delta(q_2, b, Z_0) = \{(q_2, BZ_0)\} \\ &\delta(q_2, c, B) = \{(q_4, \epsilon)\} \\ &\delta(q_3, b, A) = \{(q_3, A)\} \\ &\delta(q_3, c, A) = \{(q_5, \epsilon)\} \\ &\delta(q_4, c, B) = \{(q_6, \epsilon)\} \\ &\delta(q_5, c, A) = \{(q_5, \epsilon)\} \\ &\delta(q_5, c, A) = \{(q_6, \epsilon)\} \end{split}$$

MOVIMIENTOS

$$\begin{array}{ll} (q,\alpha)\in\delta(p,a,z)\Rightarrow(p,a\omega,z\beta)\vdash(q,\omega,\alpha\beta) & p,q\in Q\\ &a\in(\Sigma\cup\{\epsilon\})\\ &\omega\in\Sigma^*\\ &Z\in\Gamma\\ &\alpha,\beta\in\Gamma^* \end{array}$$

$$\vdash \mbox{movimiento en un paso.} \\ i\\ \vdash \mbox{movimiento en i pasos.} \\ *\\ \vdash \mbox{movimiento en cero o más pasos.} \\ +\\ \vdash \mbox{movimiento en uno o más pasos.} \end{array}$$

3.2.1. AUTÓMATAS APILADORES DETERMINISTAS (AAD)

$$A = (Q, \Sigma, \Gamma, \delta, q_0, z_0, F)$$

1.
$$\delta(q, \varepsilon, z) \neq \emptyset \Rightarrow \delta(q, \sigma, z) = \emptyset$$
 $\forall q \in Q, \sigma \in \Sigma, z \in \Gamma$
2. $\# \delta(q, a, z) \leq 1$ $\forall q \in Q, a \in (\Sigma \cup \{\epsilon\}), z \in \Gamma$

Ejemplo:

$$A = (\{q_0, q_1\}, \{a, b\}, \{A, B, Z_0\}, \delta, q_0, Z_0, \{q_0\})$$

$$\delta(q_0, a, Z_0) = \{(q_1, AZ_0)\}$$

$$\delta(q_0, b, Z_0) = \{(q_1, BZ_0)\}$$

$$\delta(q_1, a, A) = \{(q_1, AA)\}$$

$$\delta(q_1, b, B) = \{(q_1, BB)\}$$

$$\delta(q_1,\,a,\,B)\ = \{(q_1,\,\epsilon)\}$$

$$\delta(q_1, b, A) = \{(q_1, \varepsilon)\}\$$

$$\delta(q_1, \, \epsilon, \, Z_0) = \{(q_0, \, Z_0)\}$$

$$L(A) = \left\{\omega \in \left\{a, b\right\}^* / |\omega|_a = |\omega|_b\right\}$$

3.2.2. AUTÓMATAS APILADORES NO DETERMINISTAS (AAN)

Ejemplo 1:

$$A = (\{q_0, q_1, q_2\}, \{a, b\}, \{A, B, Z_0\}, \delta, q_0, Z_0, \emptyset)$$

$$\delta(q_0, a, Z_0) = \{(q_1, AZ_0)\}$$

$$\delta(q_0, b, Z_0) = \{(q_1, BZ_0)\}$$

$$\delta(q_0, \, \epsilon, \, Z_0) = \{(q_2, \, \epsilon)\}$$

$$\delta(q_1, a, A) = \{(q_1, AA)\}$$

$$\delta(q_1, b, B) = \{(q_1, BB)\}$$

$$\delta(q_1, a, B) = \{(q_1, \epsilon)\}$$

$$\delta(q_1, b, A) = \{(q_1, \epsilon)\}$$

$$\delta(q_1, \, \epsilon, \, Z_0) = \{(q_0, \, Z_0)\}$$

$$N(A) = \{\omega \in \{a, b\}^* / |\omega|_a = |\omega|_b\}$$

Ejemplo 2:

$$A = (\{q_1, q_2\}, \{a, b\}, \{A, B, Z\}, \delta, q_1, Z, \{q_2\})$$

$$\delta(q_1, \varepsilon, Z) = \{(q_2, Z)\}$$

$$\delta(q_1, a, Z) = \{(q_1, AZ)\}$$

$$\delta(q_1, b, Z) = \{(q_1, BZ)\}$$

$$\delta(q_1, a, A) = \{(q_1, AA)\}$$

$$\delta(q_1, b, A) = \{(q_1, \epsilon)\}$$

$$\delta(q_1, a, B) = \{(q_1, \varepsilon)\}$$

$$\delta(q_1, b, B) = \{(q_1, BB)\}\$$

$$L(A) = \{ \omega \in \{a, b\}^* / |\omega|_a = |\omega|_b \}$$

 $\omega = abba$

3.2.3. LENGUAJE ACEPTADO MEDIANTE ESTADO FINAL (L(A))

$$A = (Q, \Sigma, \Gamma, \delta, q_0, z_0, F)$$

 $F \neq \emptyset$

$$L(A) = \{ \omega \in \Sigma^* / (q_0, \omega, z_0) \vdash (q, \varepsilon, \gamma), q \in F \land \gamma \in \Gamma^* \}$$

Ejemplo 1:

$$A = (\{q_0, q_1, q_2\}, \{0, 1, c\}, \{0, 1, Z_0\}, \delta, q_0, Z_0, \{q_2\})$$

$$\delta(q_0, 0, Z_0) = \{(q_0, 0Z_0)\}$$

$$\delta(q_0, 1, Z_0) = \{(q_0, 1Z_0)\}$$

$$\delta(q_0, 0, 0) = \{(q_0, 00)\}$$

$$\delta(q_0, 0, 1) = \{(q_0, 01)\}$$

$$\delta(q_0, 1, 0) = \{(q_0, 10)\}$$

$$\delta(q_0, 1, 1) = \{(q_0, 11)\}$$

$$\delta(q_0, c, Z_0) = \{(q_1, Z_0)\}$$

$$\delta(q_0, c, 0) = \{(q_1, 0)\}$$

$$\delta(q_0, c, 1) = \{(q_1, 1)\}$$

$$\delta(q_1, 0, 0) = \{(q_1, \epsilon)\}$$

$$\delta(q_1, 1, 1) = \{(q_1, \varepsilon)\}$$

$$\delta(q_1, \, \epsilon, \, Z_0) = \{(q_2, \, Z_0)\}$$

$$L(A) = \{\omega c \omega^{R} / \omega \in \{0, 1\}^{*}\}$$

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Ejemplo 2:

$$A = (\{q_0,\,q_1,\,q_2\},\,\{0,\,1\},\,\{0,\,1,\,Z_0\},\,\delta,\,q_0,\,Z_0,\,\{q_2\})$$

$$\delta(q_0, 0, Z_0) = \{(q_0, 0Z_0)\}\$$

$$\delta(q_0, 1, Z_0) = \{(q_0, 1Z_0)\}$$

$$\delta(q_0, 0, 0) = \{(q_0, 00)\}$$

$$\delta(q_0, 0, 1) = \{(q_0, 01)\}$$

$$\delta(q_0, 1, 0) = \{(q_0, 10)\}$$

$$\delta(q_0, 1, 1) = \{(q_0, 11)\}$$

$$\delta(q_0,\,\epsilon,\,Z_0)=\{(q_1,\,Z_0)\}$$

$$\delta(q_0, \epsilon, 0) = \{(q_1, 0)\}$$

$$\delta(q_0, \varepsilon, 1) = \{(q_1, 1)\}$$

$$\delta(q_1, 0, 0) = \{(q_1, \varepsilon)\}$$

$$\delta(q_1, 1, 1) = \{(q_1, \epsilon)\}$$

$$\delta(q_1,\,\epsilon,\,Z_0)=\{(q_2,\,Z_0)\}$$

$$L(A) = \{\omega\omega^{R} / \omega \in \{0, 1\}^*\}$$

$$\omega = 1111$$

3.2.4. LENGUAJE ACEPTADO MEDIANTE AGOTAMIENTO DE PILA (N(A))

$$A = (Q, \Sigma, \Gamma, \delta, q_0, z_0, \varnothing)$$

$$*$$

$$N(A) = \{\omega \in \Sigma^* / (q_0, \omega, z_0) \vdash (q, \epsilon, \epsilon), q \in Q\}$$

$$Ejemplo 1:$$

$$A = (\{q_1, q_2\}, \{0, 1, c\}, \{R, B, G\}, \delta, q_1, R, \varnothing)$$

$$\delta(q_1, 0, R) = \{(q_1, BR)\}$$

$$\delta(q_1, 1, R) = \{(q_1, GR)\}$$

$$\delta(q_1, 0, B) = \{(q_1, BB)\}$$

$$\delta(q_1, 0, B) = \{(q_1, BB)\}$$

$$\delta(q_1, 1, B) = \{(q_1, GB)\}$$

$$\delta(q_1, 0, G) = \{(q_1, BG)\}$$

$$\delta(q_1, 0, G) = \{(q_1, GG)\}$$

$$\delta(q_1, c, R) = \{(q_2, R)\}$$

$$\delta(q_1, c, R) = \{(q_2, R)\}$$

$$\delta(q_1, c, G) = \{(q_2, E)\}$$

$$\delta(q_2, 0, B) = \{(q_2, \epsilon)\}$$

$$\delta(q_2, 0, B) = \{(q_2, \epsilon)\}$$

$$\delta(q_2, \epsilon, R) = \{(q_2, \epsilon)\}$$

$$N(A) = \{\omega c \omega^R / \omega \in \{0, 1\}^*\}$$

$$Ejemplo 2:$$

$$A = (\{q_1, q_2\}, \{0, 1\}, \{R, B, G\}, \delta, q_1, R, \varnothing)$$

$$\delta(q_1, 0, R) = \{(q_1, BR)\}$$

$$\delta(q_1, 1, R) = \{(q_1, GR)\}$$

$$\delta(q_1, 0, B) = \{(q_1, BB), (q_2, \epsilon)\}$$

$$\delta(q_1, 0, G) = \{(q_1, BG)\}$$

$$\delta(q_1, 1, B) = \{(q_1, GB)\}$$

$$\delta(q_1, 1, G) = \{(q_1, GG), (q_2, \epsilon)\}$$

$$\delta(q_2, 0, B) = \{(q_2, \epsilon)\}$$

$$\delta(q_2, 0, B) = \{(q_2, \epsilon)\}$$

$$\delta(q_1, \epsilon, R) = \{(q_2, \epsilon)\}$$

$$\delta(q_1, \epsilon, R) = \{(q_2, \epsilon)\}$$

 $\omega = 001100$

 $N(A) = \{\omega\omega^{R} / \omega \in \{0, 1\}^*\}$

3.2.5. EQUIVALENCIAS

$$A = (\{q_1, q_2, q_3\}, \{a, b\}, \{a, z\}, \delta, q_1, z, \{q_3\})$$

$$\delta(q_1, a, z) = \{(q_1, az)\}$$

$$\delta(q_1, b, z) = \{(q_2, \epsilon)\}$$

$$L(A) = \{aa\}$$

$$N(A) = \{b\}$$

$3.2.5.1. L(A) \Rightarrow N(A')$

 $\delta(q_1, a, a) = \{(q_3, a)\}$

$$A = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F) \Rightarrow A' = (Q', \Sigma, \Gamma', \delta', q_0', X_0, \emptyset)$$

$$\begin{cases} Q' = Q \cup \{q_0', q_e\} \\ \Gamma' = \Gamma \cup \{X_0\} \\ \delta'(q, a, Z) = \delta(q, a, Z) & \forall q \in (Q - F), a \in (\Sigma \cup \{\epsilon\}), Z \in \Gamma \\ \delta'(q, \sigma, Z) = \delta(q, \sigma, Z) & \forall q \in F, \sigma \in \Sigma, Z \in \Gamma \\ \delta'(q, \epsilon, Z) = \delta(q, \epsilon, Z) \cup \{(q_e, \epsilon)\} & \forall q \in F, Z \in \Gamma \\ \delta'(q, \epsilon, X_0) = \{(q_e, \epsilon)\} & \forall q \in F \\ \delta'(q_e, \epsilon, Z) = \{(q_e, \epsilon)\} & \forall Z \in \Gamma' \end{cases}$$

$3.2.5.2. N(A) \Rightarrow L(A')$

$$\begin{split} A &= (Q, \Sigma, \Gamma, \delta, q_0, Z_0, \emptyset) \Rightarrow A' = (Q', \Sigma, \Gamma', \delta', q_0', X_0, F') \\ & \begin{cases} Q' &= Q \cup \{q_0', q_f\} \\ \Gamma' &= \Gamma \cup \{X_0\} \end{cases} \\ \delta' \begin{cases} \delta'(q_0', \epsilon, X_0) &= \{(q_0, Z_0 X_0)\} \\ \delta'(q, a, Z) &= \delta(q, a, Z) \end{cases} & \forall q \in Q, a \in (\Sigma \cup \{\epsilon\}), Z \in \Gamma \\ \delta'(q, \epsilon, X_0) &= \{(q_f, \epsilon)\} \end{cases} & \forall q \in Q \end{split}$$

Ejemplo:

$$A = (\{q\}, \{i, e\}, \{Z\}, \delta, q, Z, \emptyset)$$

$$\delta(q, i, Z) = \{(q, ZZ)\}$$

$$\delta(q, e, Z) = \{(q, \varepsilon)\}$$