# 2.2. GRAMÁTICAS REGULARES LINEALES POR LA DERECHA (GRLD)

# **DEFINICIÓN**

$$G = (N, \Sigma, P, S)$$

$$P \subset N \times \Sigma^* (N \cup \varepsilon)$$

$$\begin{array}{ll}
A \to \omega B \\
A \to \omega
\end{array} \right\} \in P \qquad \qquad \begin{array}{ll}
A, B \in N \\
\omega \in \Sigma^*
\end{array}$$

## 2.2.1. UNIÓN

$$\begin{array}{ll} G_{1} = (N_{1}, \Sigma_{1}, P_{1}, S_{1}) \\ G_{2} = (N_{2}, \Sigma_{2}, P_{2}, S_{2}) \end{array} \quad N_{1} \cap N_{2} = \emptyset \Rightarrow G = (N, \Sigma, P, S) \text{ donde } \begin{cases} N = N_{1} \cup N_{2} \cup \{S\} \\ \Sigma = \Sigma_{1} \cup \Sigma_{2} \\ P = P_{1} \cup P_{2} \cup \{S \rightarrow S_{1} \mid S_{2}\} \end{cases}$$

$$L(G) = L(G_1) \cup L(G_2)$$

#### 2.2.2. FORMA NORMAL

$$\begin{split} G &= (N, \Sigma, P, S) \\ P &\subseteq N \times (\Sigma N \cup \epsilon) \\ A &\to \sigma B \\ A &\to \epsilon \end{split} \right\} \in P \qquad \begin{array}{c} A, B \in N \\ \sigma \in \Sigma \end{array}$$

a) Eliminar Producciones Unitarias

Troduce on Cintaria
$$A \to B \qquad A, B \in N$$

$$G = (N, \Sigma, P, S) \Rightarrow G' = (N, \Sigma, P', S)$$

$$A \to B \\ B \to \alpha$$

$$G' = (N, \Sigma, P', S) \Rightarrow A \to \alpha \\ B \to \alpha$$

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$$\begin{split} P' &= \varnothing \\ \forall \ A \in N \\ & \forall \ B \in U(A) \\ & \forall \ B \to \alpha \in P \\ & \text{Si } \alpha \not \in N \text{ entonces} \\ & P' &= P' \cup \{A \to \alpha\} \\ & \text{Fin Si} \end{split}$$

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c) 
$$G = (N, \Sigma, P, S) \\ \Rightarrow G' = (N', \Sigma, P', S) \\ A \rightarrow \sigma_1 D_1 \\ D_1 \rightarrow \sigma_2 D_2 \\ D_2 \rightarrow \sigma_3 D_3 \\ D_3 \rightarrow \sigma_4 D_4 \\ \cdots \\ D_{n-1} \rightarrow \sigma_n D_n \\ D_n \rightarrow \epsilon \\ N' = N \cup \{D_1, D_2, D_3, \cdots, D_n\}$$
 Ejemplo:

$$G = (\{S, A, B\}, \{a, b\}, P, S)$$

$$P = \{$$

$$S \rightarrow aB \mid bA \mid \epsilon$$

$$A \rightarrow abaS$$

$$B \rightarrow babS$$