

Test 2

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1. Write a complete residue system mod 17 only using multiples of 3.

$$\begin{aligned} 0 &\equiv 3 \cdot 0 \pmod{17} \\ 1 &\equiv 3 \cdot 6 \pmod{17} \\ 2 &\equiv 3 \cdot 12 \pmod{17} \\ 3 &\equiv 3 \cdot 18 \pmod{17} \\ 4 &\equiv 3 \cdot 24 \pmod{17} \\ 5 &\equiv 3 \cdot 30 \pmod{17} \\ 6 &\equiv 3 \cdot 36 \pmod{17} \\ 7 &\equiv 3 \cdot 42 \pmod{17} \\ 8 &\equiv 3 \cdot 48 \pmod{17} \\ 9 &\equiv 3 \cdot 54 \pmod{17} \\ 10 &\equiv 3 \cdot 60 \pmod{17} \\ 11 &\equiv 3 \cdot 66 \pmod{17} \\ 12 &\equiv 3 \cdot 72 \pmod{17} \\ 13 &\equiv 3 \cdot 78 \pmod{17} \\ 14 &\equiv 3 \cdot 84 \pmod{17} \\ 15 &\equiv 3 \cdot 90 \pmod{17} \\ 16 &\equiv 3 \cdot 96 \pmod{17} \end{aligned}$$

$\{0, 18, 36, 54, 72, 90, 108, 126, 144, 162, 180, 198, 216, 234, 252, 270, 288\}$ forms a complete residue system mod 17. I generated the table above using the following Python code. For an explanation of Python code in general and the source for `linear_diophantine()`, please read 3.23 in my notebook.

```
1 from tools import linear_diophantine
2 CRS = []
3 for n in range(17):
4     (x_0, y_0), (r_x, r_y) = linear_diophantine(3, -17, n)
5     # now we have 3x_0 - 17y_0 = n
6     # output n ≡ 3 · x_0 (mod 17)
7     print (r'${n} \equiv 3 \cdot {x_0} \pmod{{17}}$ \\' .format(**locals()))
8     CRS.append(3 * x_0)
9
10 # output the whole CRS, seperated by commas
11 print (', '.join(map(str, CRS)))
```

2. Find $2^{100} \pmod{9}$.

All congruence statements are taken mod 9.

$$\begin{aligned}
2^{100} &\equiv ? \\
&\equiv 2^{3 \cdot 33 + 1} \\
&\equiv (2^3)^{33} \cdot 2^1 \\
&\equiv 8^{33} \cdot 2 \\
&\equiv (-1)^{33} \cdot 2 \\
&\equiv -1 \cdot 2 \\
&\equiv 7
\end{aligned}$$

3. **Theorem:** For any CRS mod m , multiplying by a natural coefficient coprime to m produces a new CRS.

Let $\{r_1, r_2, \dots, r_m\}$ form a CRS mod m . Let $a \in \mathbb{N}$ where $\gcd(a, m) = 1$.

Proof: Theorem 3.17 states that m unique (unique mod m) elements describe a CRS. Begin with $\{r_1, r_2, \dots, r_m\}$ which is a CRS, so the elements are unique by definition. $r_i \neq r_j \rightarrow ar_i \neq ar_j$ by the inverse of theorem 1.14 (The inverse is necessarily true since theorem 1.14 is an ‘if-and-only-if’). Thus $\{ar_1, ar_2, \dots, ar_m\}$ are all unique. To reiterate: If they were not ($ar_i \equiv ar_j$) we could use theorem 1.14 to draw a contradiction ($r_i \equiv r_j$) which is not true by the definition of the CRS. Therefore $\{ar_1, ar_2, \dots, ar_m\}$ constitute m unique elements. By theorem 3.17 they must also constitute a CRS mod m .

4. Find the number x where $x \equiv 1 \pmod{4}$ and $x \equiv 2 \pmod{3}$.

$$\begin{aligned}
x &\equiv 1 \pmod{4} && \{\dots, 1, 5, 9, 13, 17, 21, 25, 29, \dots\} \\
x &\equiv 2 \pmod{3} && \{\dots, 2, 5, 8, 11, 14, 17, 20, 23, 26, 29, \dots\} \\
\left. \begin{aligned} x &\equiv 1 \pmod{4} \\ x &\equiv 2 \pmod{3} \end{aligned} \right\} &&& \{\dots, 5, 17, 29, \dots\}
\end{aligned}$$

The first solutions set is $\{1, 5, 9, \dots\}$. By exhaustion, 5 is the first number in the solution set of the first congruence that also satisfies the second congruence. The LCM of the two modulus is 12. Therefore $x \equiv 5 \pmod{12}$.

5. **Theorem:** Let the digits of n be $n = a_k a_{k-1} \dots a_0$. Let $m = a_0 - a_1 + a_2 \dots \pm a_{k-1} \mp a_k$. 11 divides n exactly when 11 divides m .

All congruences are taken mod 11.

Proof: Let $f_n(x) = a_k x^k + a_{k-1} x^{k-1} + \dots + a_0$. We know that $-1 \equiv 10$, which implies $f(-1) \equiv f(10)$ by theorem 3.8. However, $f(-1) = a_k \pm a_{k-1} \mp \dots + a_0 = m$. Similarly $f(10) = a_k 10^k + a_{k-1} 10^{k-1} + \dots + a_0 = n$. Therefore by substitution $m \equiv n$. If $11 \mid n$, then $n \equiv 0 \pmod{11}$, then $m \equiv 0$ by 1.11 (transitivity), which gives $11 \mid m$. And vice versa, If $11 \mid m$, then $m \equiv 0$, then $n \equiv 0$, then $11 \mid n$. Therefore $11 \mid m \leftrightarrow 11 \mid n$. ■

6. Create a test to see if $4 \mid n$. Let the digits of $n = n_m \dots n_1 n_0$.

Algorithm: If n_1 is even, $4 \mid n$ exactly when $4 \mid n_0$. If n_1 is odd, $4 \mid n$ exactly when $4 \mid (n_0 + 2)$.

Lemma: $n \mid (na + b) \leftrightarrow n \mid b$. Justification: $n \mid na$ by definition of divides. $n \mid na$ and $n \mid (na + b)$ imply $n \mid (na + b - na)$ by Theorem 1.2. $n \mid b$ and $n \mid na$ imply $n \mid (na + b)$ by Theorem 1.1.

Lemma: If a_0 is even or odd, then $a_m \dots a_2 a_1 a_0$ is likewise. Justification: By the previous lemma $2 \mid ((a_m \dots a_2 a_1) \cdot 10 + a_0)$ is equivalent to $2 \mid a_0$ since $2 \mid ((a_m \dots a_2 a_1) \cdot 10)$. Therefore if $2 \nmid a_0$, then $2 \nmid (a_m \dots a_2 a_1 a_0)$ and vice versa if a_0 is even.

Proof (odd): If n_1 is even, then $n_m \dots n_2 n_1$. Then $4|(n_m \dots n_2 n_1 n_0)$ is equivalent to $4|(2k \cdot 10 + n_0)$, is equivalent to $4|(4k \cdot 5 + n_0)$, is equivalent to $4|n_0$ (reduced the problem to less than 8).

Proof (even): If n_1 is odd, then $n_m \dots n_2 n_1$ is odd. Then $4|(n_m \dots n_2 n_1 n_0)$ is equivalent to $4|(2k+1) \cdot 10 + n_0$, is equivalent to $4|(20k+10+n_0)$, is equivalent to $4|(10+n_0)$, is equivalent to $4|(8+2+n_0)$, is equivalent to $4|(2+n_0)$ (reduced the problem to less than 11. ■

7. The egg problem.

$$x \equiv 1 \pmod{2}$$

$$x = \{1, 3, 5, 7, 9, 11, 13, 15, 17, 19, 21, 23, \dots\}$$

$$x \equiv 2 \pmod{3} \text{ and all previous equations}$$

$$x = \{5, 11, 17, 23, 29, 35, 41, 47, \dots\}$$

$$x \equiv 3 \pmod{4} \text{ and all previous equations}$$

$$x = \{11, 23, 35, 47, 59, 71, 83, 95, 107, 119, 131, 143, 155, 167, 179, \dots\}$$

$$x \equiv 4 \pmod{5} \text{ and all previous equations}$$

$$x = \{59, 119, 179, \dots\}$$

$$x \equiv 5 \pmod{6} \text{ and all previous equations}$$

$$x = \{59, 119, 179, \dots\}$$

$$x \equiv 0 \pmod{7} \text{ and all previous equations}$$

$$x = \{119, 539, 959, 1379, 1799, 2219, \dots\}$$

Therefore, there are at least 119 eggs in the basket.

8. **Theorem:** $\{2, 4, 6, \dots, 2m\}$ is a CRS mod m if m is odd.

Proof: Show that any element of the CCRS is congruent to exactly one thing in S where $S = \{2, 4, \dots, 2m\}$. For all x in the CCRS ($0 \leq x < m$),

- if x is even and non-zero, it is in the set S . $x \equiv x \wedge x \in S$. Moreover it is in the set uniquely, the next possible representation is $x + m$, which is not in the set S because an even plus an odd is an odd and the set contains only evens.
- If x is zero, $x \equiv 2m \pmod{m}$ since $m|(2m - 0)$. $0 \equiv 2m \wedge 2m \in S$. (This is why I believe the question should ask about the set $\{0, 2, 4, \dots, 2m - 2\}$, playful nudge). Moreover it is in the set uniquely because no other element is perfectly divisibly by m other than $2m$.
- If x is odd, $x = 2 \cdot \frac{x-1}{2} + 1$ and since m is odd, $m = 2 \cdot \frac{m-1}{2} + 1$. Then $x \equiv x + m$ since $m|(x+m-x)$ and $x+m$ is even, since an odd plus an odd is an even. $x < m - 2$ since x is a smaller odd number than m , therefore $x+m < 2m-2 < 2m$. $x \equiv x+m \wedge x+m \in S$. Moreover it is in the set uniquely since the only other representation of x between 0 and $2m$ is x , and x is odd therefore it is not in the set S .

Since every element in the CCRS is congruent to exactly one number in S , and every integer is congruent to exactly one number in the CCRS, every integer is congruent to exactly one number in S . Therefore S is a CRS. ■

All odd numbers can be written $2k + 1$ for some $k \in \mathbb{Z}$. All even numbers can be written $2k$ for some $k \in \mathbb{Z}$. An even number plus an even number is even since $2k_1 + 2k_2 = 2(k_1 + k_2)$.

An even number plus an odd number is odd since $2k_1 + 2k_2 + 1 = 2(k_1 + k_2) + 1$. An odd number plus an odd number is even since $2k_1 + 1 + 2k_2 + 1 = 2(k_1 + k_2 + 1)$. In general, the sum of two addends is even exactly when both addends are even or both addends are odd. $2|(a + b) \leftrightarrow (2|a \leftrightarrow 2|b)$.

9. **Theorem:** All reduced residue systems modulo m have the same number of elements.

Let $R = \{r_1, r_2, \dots, r_i\}$ be a Reduced Residue System (RRS) modulo m with i elements, and similarly $S = \{s_1, s_2, \dots, s_j\}$ with j elements. All congruence statements are taken mod m .

Proof: Assume for contradiction that $i \neq j$ (without loss of generality $i < j$). By the first tenet of the definition of RRS, $\gcd(s_p, m) = 1$ for all p . By the third tenet of the definition, there exists a q where $s_p \equiv r_q$. Since there are more s s than r s, some of the s s have to be congruent to the same r (in other words the relation can not be injective). $s_p \equiv r_q \equiv s_t$ for some t . By transitivity $s_p \equiv s_t$. This contradicts tenet 2 of the definition.

Therefore $i = j$. ■

10. **Theorem:** $ax \equiv ay \pmod{n} \leftrightarrow a \equiv y \pmod{\frac{n}{\gcd(a,n)}}$

Proof of \rightarrow conditional: Let $ax \equiv ay \pmod{n}$. This is equivalent to $n|(ax - ay)$, which is equivalent to $n|a(x - y)$, which is equivalent to $\frac{n}{\gcd(a,n)} \gcd(a,n) | \frac{a}{\gcd(a,n)}(x - y)$. Since $\gcd(a,n)$ shows up in the divisor and the dividend, they cancel out. $\frac{n}{\gcd(a,n)} | \frac{a}{\gcd(a,n)}(x - y)$. By Test 1 Theorem 4, $\frac{n}{\gcd(a,n)}$ and $\frac{a}{\gcd(a,n)}$ are coprime, which fulfills the condition for Theorem 1.41, which claims that $\frac{n}{\gcd(a,n)} | (x - y)$, which is equivalent to $x \equiv y \pmod{\frac{n}{\gcd(a,n)}}$.

$ab|ac$ exactly when $b|c$. If $b|c$ then $c = ib$, then $ac = iab$, then $ab|ac$. If $ab|ac$ then $iab = ac$, then $ib = c$, then $b|c$.

Proof of \leftarrow conditional: