Project 3

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Introduction

This project was written in C++ and uses OpenMP

Folders structure

Installation

Requirements

- C++ compiler (g++, clang++, msvc++)
- XMake
- Git
- OpenMP

Installation

```
git clone git@github.com:GlassAlo/CAU_Multicore.git
cd CAU Multicore/proj3
```

Compile

```
xmake f -m release && xmake -y
```

- **-m release** is used to compile the project in release mode, which is faster than debug
- **xmake** -**y** is used to compile the project. The -**y** flag is used to skip the confirmation prompt.
- xmake will create a bin folder with the executables inside.

Usage

```
./bin/prob1 <scheduling type> <number of threads>
./bin/prob2 <number end>
```

Test



- OS: Garuda Linux Broadwing x86_64
- Kernel: 6.13.8-zen1-1-zen
- CPU: AMD Ryzen 9 5900HS with Radeon Graphics (16) @ 4.680GHz
 - Cores 8
 - Uniform core design
 - Threads 16
 - Base clock 3.0GHz
 - Max boost clock up to 4.6GHz
 - L3 cache 16MB
 - Memory PCIe 3.0

- Supports Simultaneous Multithreading (SMT), with each cores supporting two threads
- Integrated GPU: AMD ATI Radeon Vega Series / Radeon Vega Mobile Series
- Discrete GPU: NVIDIA GeForce RTX 3080 Mobile / Max-Q (8GB/16GB)

- RAM: 32GB- Disk: 1TB SSD- Shell: zsh

- Using Arch Linux comes with a cost, the CPU pilots might not be very efficient, stable or up to date.

Problem 1



We measured the execution time and derived performance index using:

- Static load balancing with default chunk size
- Dynamicload balancing with default chunk size
- Static load balancing with a chunk size of 10
- Dynamic load balancing (task size 10)

Each implementation was tested using 1, 2, 4, 6, 8, 10, 12, 14 and 16 threads, with execution times measured in milliseconds (ms).

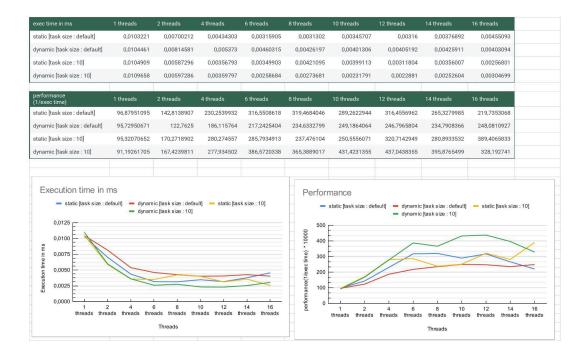
The performance index was defined as:

- Performance Index = (1/Execution Time)

This means that higher values indicate better performance.

- All the tests were done with the same number of threads and the same end number.
- The end number was set to 200000.

Performance Chart



Detailed Observations

- 1. Execution Time (Lower is Better)
 - All four scheduling strategies show significant improvements from 1 to 8 threads.
 - Dynamic (chunk size = 10) shows the fastest execution at most thread counts beyond 4.
 - Static (default chunk size) performs well until 8–10 threads, after which gains taper
 - After 12 threads, all methods plateau or degrade, showing signs of overhead or saturation.
- 2. Performance Index (Higher is Better)
 - Dynamic with chunk size 10 consistently achieves the highest performance from 4 to 12 threads.
 - Static with chunk size 10 also shows strong performance but with more fluctuation.
 - The **default static** and **default dynamic** approaches underperform with high thread counts, likely due to inefficient load balancing or lack of granularity.

- Explanation of Results
- Why Dynamic Scheduling + Chunk Size 10 Performs Best
 - **Better Load Balancing**: Tasks are dynamically pulled by idle threads, reducing idle time.
 - **Fine Granularity**: Smaller chunks allow better utilization of threads, especially in unbalanced workloads like recursive tasks.
- ↑ Why Performance Drops or Plateaus at 14–16 Threads
 - Overhead of too many tasks: With 200,000 elements and chunk size 10, you create 20,000 tasks overhead increases with thread count.
 - **Diminishing returns**: CPU reaches saturation around 8–12 threads, especially with hyperthreading.
 - Contention and synchronization costs (e.g., reduction, task queue contention) rise with more threads.

Conclusion

Dynamic scheduling with small chunk sizes (10) provides the best performance in highly parallel divide-and-conquer tasks on modern CPUs with SMT.

Static scheduling performs reasonably well but lacks adaptability to runtime imbalances. While increasing threads generally improves performance, **beyond 12 threads** there are **diminishing returns** due to overhead, memory bandwidth saturation, and synchronization costs.

The results validate that **task granularity and scheduling policy** are critical to achieving optimal parallel performance in OpenMP.