Task

- Determine the RE of the language, defined over $\Sigma=\{a, b\}$ of **words beginning with a**.

Solution:

The required RE may be $a(a+b)^*$

- Determine the RE of the language, defined over $\Sigma=\{a, b\}$ of **words beginning with and ending in same letter.**

Solution:

The required RE may be $a(a+b)^*a+b(a+b)^*b$

Task Continued …

- Determine the RE of the language, defined over $\Sigma=\{a, b\}$ of **words ending in b.**

Solution:

The required RE may be

$(a+b)^*b$.

- Determine the RE of the language, defined over $\Sigma=\{a, b\}$ of **words not ending in a.**

Solution: The required RE may be

$(a+b)^*b + \Lambda$ Or $((a+b)^*b)^*$

An important example

The Language EVEN-EVEN :

Language of strings, defined over $\Sigma=\{a, b\}$ having **even number of a's and even number of b's**. *i.e.*

EVEN-EVEN = {Λ, aa, bb, aaaa,aabb,abab, abba, baab, baba, bbaa, bbbb,…} ,

its regular expression can be written as

$(aa+bb+(ab+ba)(aa+bb)^*(ab+ba))^*$

Note

- It is important to be clear about the difference of the following regular expressions

$r_1=a^*+b^*$

$r_2=(a+b)^*$

Here r_1 does not generate any string of concatenation of a and b, while r_2 generates such strings.

Equivalent Regular Expressions

- **Definition:**

Two regular expressions are said to be equivalent if they generate the same language.

Example:

Consider the following regular expressions

$r_1 = (a + b)^* (aa + bb)$

$r_2 = (a + b)^* aa + (a + b)^* bb$ then

both regular expressions define the language of strings **ending in aa or bb**.

Note

- If r_1 $=(aa + bb)$ and $r_2=(a + b)$ then
  1. $r_1+r_2 =(aa + bb) + (a + b)$
  2. $r_1r_2 =(aa + bb)\ (a + b)$
     $=(aaa + aab + bba + bbb)$
  3. $(r_1)^* =(aa + bb)^*$

Regular Languages

- **Definition**:

  The language generated by any regular expression is called a **regular language**.

  It is to be noted that if r_1, r_2 are regular expressions, corresponding to the languages L_1 and L_2 then the languages generated by $r_1 + r_2$, r_1r_2(or r_2r_1) and r_1^*(or r_2^*) are also regular languages.

Note

- It is to be noted that if L_1 and L_2 are expressed by r_1 and r_2, respectively then the language expressed by

  1) $r_1 + r_2$, is the language $L_1 + L_2$ or $L_1 \cup L_2$

  2) $r_1 r_2$, , is the language $L_1 L_2$, of strings obtained by prefixing every string of L_1 with every string of L_2

  3) r_1^*, is the language L_1^*, of strings obtained by concatenating the strings of L, including the null string.

All finite languages are regular.

Example:

Consider the language L, defined over $\Sigma=\{a,b\}$, of strings of length 2, **starting with a**, then

L={aa, ab}, may be expressed by the regular expression aa+ab. Hence L, by definition, is a regular language.

Note

It may be noted that if a language contains even thousand words, its RE may be expressed, placing ‘ + ’ between all the words.

Here the special structure of RE is not important.

Consider the language L={aaa, aab, aba, abb, baa, bab, bba, bbb}, that may be expressed by a RE aaa+aab+aba+abb+baa+bab+bba+bbb, which is equivalent to (a+b)(a+b)(a+b).

Defining Languages (continued)…

- Method (Finite Automaton)

  Definition:

  A Finite automaton (FA), is a collection of the followings

  1) Finite number of states, having one initial and some (maybe none) final states.
  2) Finite set of input letters (Σ) from which input strings are formed.
  3) Finite set of transitions *i.e.* for each state and for each input letter there is a transition showing how to move from one state to another.

Example

- $\Sigma = \{a,b\}$
- **States:** x, y, z where x is an initial state and z is final state.
- **Transitions:**

1. At state **x** reading **a** go to state **z**,
2. At state **x** reading **b** go to state **y**,
3. At state **y** reading **a**, **b** go to state **y**
4. At state **z** reading **a**, **b** go to state **z**

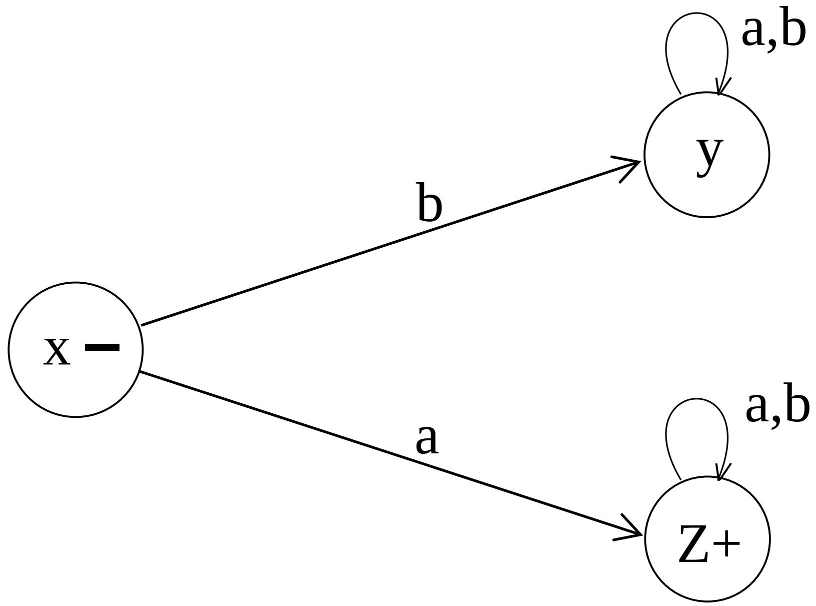
Example Continued ...

- These transitions can be expressed by the following table called transition table

Old States	New States	
	Reading a	Reading b
x -	z	y
y	y	y
z +	z	z

Note

- It may be noted that the information of an FA, given in the previous table, can also be depicted by the following diagram, called the **transition diagram**, of the given FA



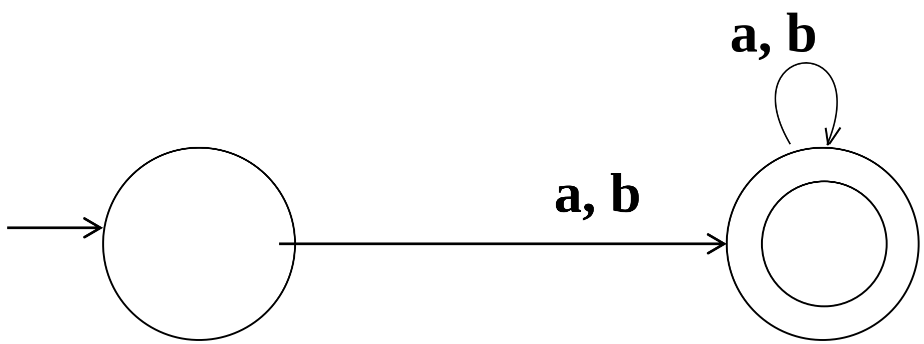
Remark

- The previous transition diagram is an FA accepting the language of strings, defined over $\Sigma=\{a, b\}$, **starting with a**. It may be noted that this language may be expressed by the regular expression

$$a\,(a + b)^*$$

Note

- It may be noted that to indicate the initial state, an arrow head can also be placed before that state and that the final state with double circle, as shown below. It is also to be noted that while expressing an FA by its transition diagram, the labels of states are not necessary.



Example

- $\Sigma = \{a,b\}$

  States: x, y, where x is both initial and final state.

  Transitions:

  1. At state x reading a or b go to state y.
  2. At state y reading a or b go to state x.

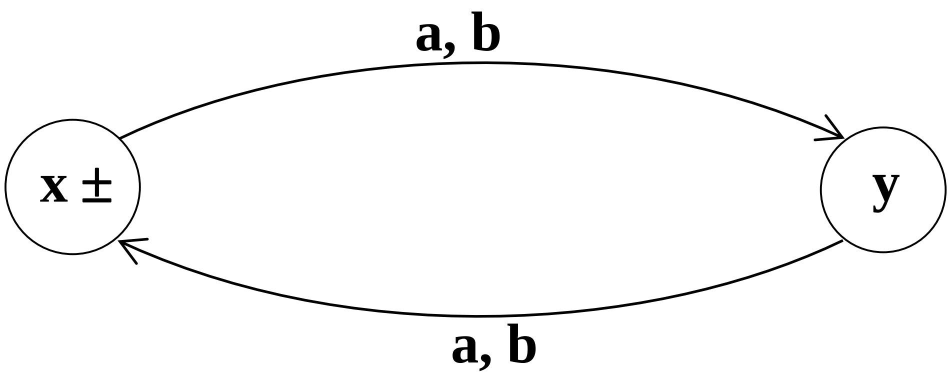
Example Continued …

- These transitions can be expressed by the following transition table

Old States	New States	
	Reading a	Reading b
x ±	y	y
y	x	x

Example Continued …

- It may be noted that the previous transition table may be depicted by the following transition diagram.



Example Continued …

- The previous transition diagram is an FA accepting the language of strings, defined over Σ={a, b} of **even length**. It may be noted that this language may be expressed by the regular expression

$$((a+ b) (a + b))^*$$

TASK

Build an FA for the language L of strings, defined over $\Sigma=\{a, b\}$, **of odd length.**

Solution of Task